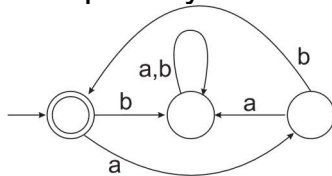


Part III: DFAs and NFAs

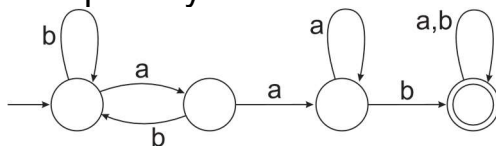
You may find draw.io useful for drawing automata. If you know of a similar and better tool, please let me know. You can insert images from <https://automatonsimulator.com> too as long as the edge annotations are clearly shown (note sometimes they are difficult to read).

1. By writing a regular expression or a grammar, describe the language accepted by the DFA



$(ab)^*$

2. By writing a regular expression or a grammar, describe the language accepted by the DFA



$(b|ab)^*aaa^*b[a,b]^*$

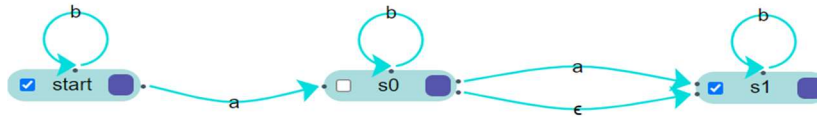
3. Construct NFAs (possibly with ϵ -moves) to recognise the languages on alphabet $\{a,b\}$ such that:

1. $L = \{w \in \{a,b\}^* \mid w \text{ contains at most two } a\text{'s}\}$.

So ϵ and aa and $bbbbabba$ are in L and aaa is not.

```

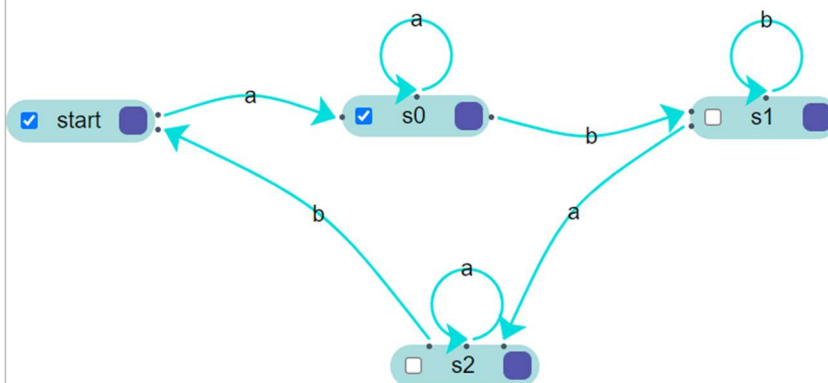
{"type":"NFA","nfa":{"transitions":{"start":{"a":["s0"],"b":["start"]},"s0":{"a":["s1"],"b":["s1"]},"s1":{"b":["s1"]},"startState":"start","acceptStates":["start","s1"],"states":{"start":{"isAccept":true},"s0":{"top":246,"left":249,"displayId":"s0"},"s1":{"isAccept":true,"top":248,"left":542,"displayId":"s1"}},"transitions":[{"stateA":"start","label":"a","stateB":"s0"},{"stateA":"start","label":"b","stateB":"start"},{"stateA":"s0","label":"a","stateB":"s1"},{"stateA":"s0","label":"b","stateB":"s1"},{"stateA":"s1","label":"b","stateB":"s1"}],"bulkTests":{"accept":"\nnaa\nbbbbabba\nbabab","reject":"aaa\nbaaab"}}}
  
```



2. $L = \{w \in \{a,b\}^* \mid w \text{ contains an even number of occurrences of } ab \text{ as a subword}\}$. So ϵ and a and b and $abab$ and $abaaba$ and $baaabab$ are in L , but ab and $ababab$ and $abaabab$ are not.

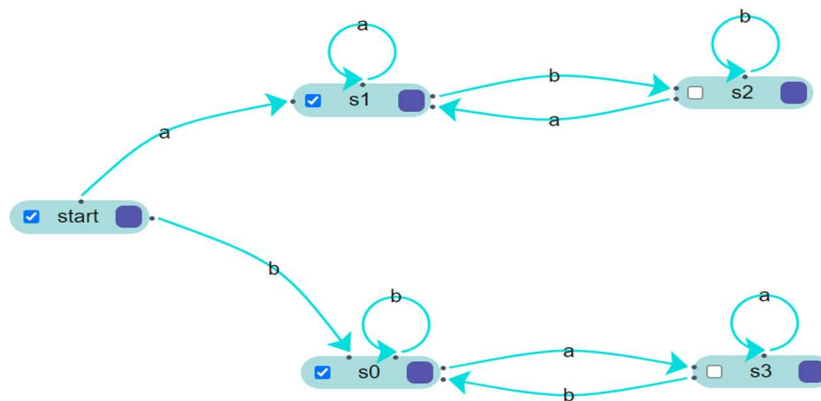
```

{"type":"NFA","nfa":{"transitions":{"start":{"a":["s0"]},"s0":{"b":["s1"],"a":["s0"]},"s1":{"b":["s1"],"a":["s2"]},"s2":{"b":["start"],"a":["s2"]}},"startState":"start","acceptStates":["start","s0"],"states":{"start":{"isAccept":true},"s0":{"isAccept":true,"top":242,"left":237,"displayId":"s0"},"s1":{"top":243,"left":484,"displayId":"s1"},"s2":{"top":419,"left":272,"displayId":"s2"}}, "transitions":[{"stateA":"start","label":"a","stateB":"s0"}, {"stateA":"s0","label":"b","stateB":"s1"}, {"stateA":"s0","label":"a","stateB":"s0"}, {"stateA":"s1","label":"b","stateB":"s1"}, {"stateA":"s1","label":"a","stateB":"s2"}, {"stateA":"s2","label":"b","stateB":"start"}, {"stateA":"s2","label":"a","stateB":"s2"}],"bulkTests":{"accept":"abab\n\nabababab\na","reject":"abaabab\nababab"}}
  
```

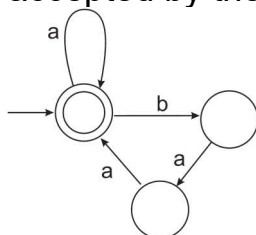


3. $L = \{w \in \{a,b\}^* \mid \text{the first and the last letter of } w \text{ are identical}\}$

```
{
  "type": "NFA",
  "nfa": {
    "transitions": {
      "start": {
        "a": ["s1"],
        "b": ["s0"]
      },
      "s1": {
        "a": ["s1"],
        "b": ["s2"]
      },
      "s2": {
        "a": ["s1"],
        "b": ["s2"]
      },
      "s3": {
        "a": ["s3"],
        "b": ["s0"]
      },
      "s0": {
        "a": ["s3"],
        "b": ["s0"]
      }
    },
    "startState": "start",
    "acceptStates": ["start", "s1", "s0"],
    "states": {
      "start": {
        "isAccept": true,
        "top": 140,
        "left": 213,
        "displayId": "start"
      },
      "s1": {
        "isAccept": true,
        "top": 140,
        "left": 213,
        "displayId": "s1"
      },
      "s0": {
        "isAccept": true,
        "top": 386,
        "left": 220,
        "displayId": "s0"
      },
      "s2": {
        "top": 133,
        "left": 502,
        "displayId": "s2"
      },
      "s3": {
        "top": 385,
        "left": 516,
        "displayId": "s3"
      }
    },
    "transitions": [
      {
        "stateA": "start",
        "label": "a",
        "stateB": "s1"
      },
      {
        "stateA": "start",
        "label": "b",
        "stateB": "s0"
      },
      {
        "stateA": "s1",
        "label": "a",
        "stateB": "s1"
      },
      {
        "stateA": "s1",
        "label": "b",
        "stateB": "s2"
      },
      {
        "stateA": "s2",
        "label": "a",
        "stateB": "s1"
      },
      {
        "stateA": "s2",
        "label": "b",
        "stateB": "s2"
      },
      {
        "stateA": "s3",
        "label": "a",
        "stateB": "s3"
      },
      {
        "stateA": "s3",
        "label": "b",
        "stateB": "s0"
      },
      {
        "stateA": "s0",
        "label": "a",
        "stateB": "s3"
      },
      {
        "stateA": "s0",
        "label": "b",
        "stateB": "s0"
      }
    ],
    "bulkTests": {
      "accept": "",
      "reject": ""
    }
  }
}
```

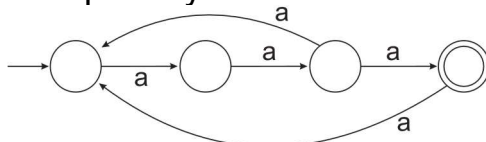


4. By writing a regular expression or a grammar, describe the language accepted by the NFA



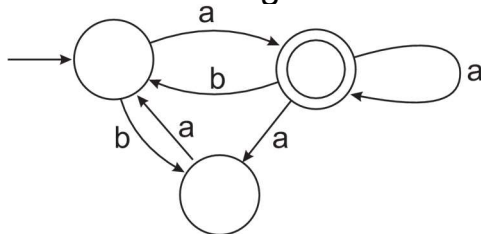
$(a|baa)^*$

5. By writing a regular expression or a grammar, describe the language accepted by the NFA

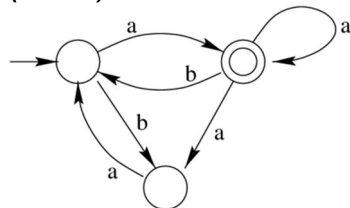


$(aaa|aaaa)^*aaa$

6. (Unmarked) Use the powerset construction (see Lecture 5) to construct a DFA that recognises the same language as the following NFA:



7. (Hard) Let L be the language recognised by the NFA



1. Construct a context-free grammar G that generates language L (a regular grammar is also fine, since every regular grammar is context-free).
 2. State with proof whether G is ambiguous or unambiguous.
8. (Unmarked) Let $G = (\{S, A, B\}, \{a, b\}, P, S)$ be the context-free grammar with productions:

$$S \rightarrow aAbS \mid bBaS \mid \varepsilon$$

$$A \rightarrow aAbA \mid \varepsilon$$

$$B \rightarrow bBaB \mid \varepsilon$$

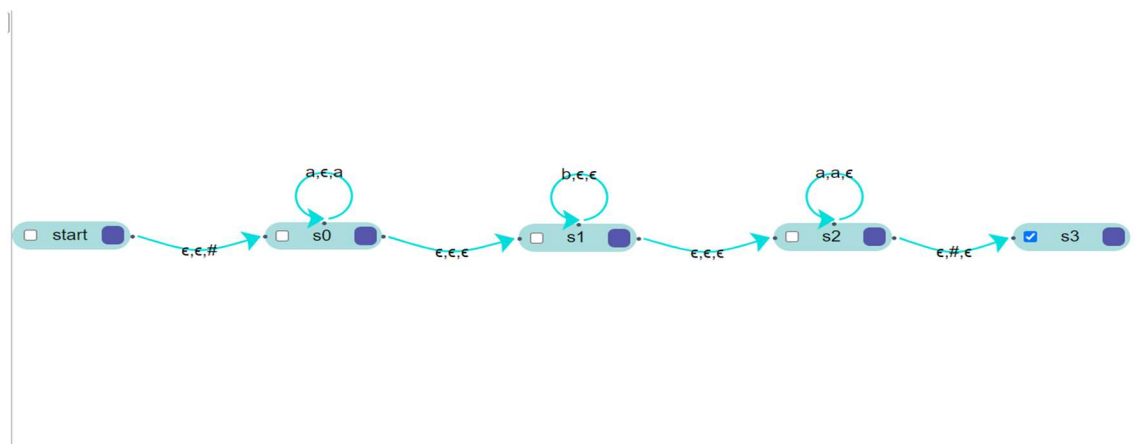
Give a short and intuitive English description of the language determined by G (such a description does exist).

On PDAs

In the questions below, assume the alphabet is $\{a, b\}$. Your answer must state the acceptance mode used. In the questions below, the word *describe* means *draw a precise picture of* in the style I used in lectures or in the style of [this webpage](#) or [this pdf](#), or just search the Internet for [how to draw a PDA](#).

9. Describe a pushdown automaton that recognises $\{a^m b^n a^m \mid m, n \geq 0\}$.

```
{
  "type": "PDA",
  "pda": {
    "transitions": {
      "start": {
        "": {
          "state": "s0",
          "stackPushChar": "#"
        }
      },
      "s0": {
        "a": {
          "state": "s0",
          "stackPushChar": "a"
        },
        "": {
          "state": "s1",
          "stackPushChar": ""
        }
      },
      "s1": {
        "b": {
          "state": "s1",
          "stackPushChar": ""
        },
        "": {
          "state": "s2",
          "stackPushChar": ""
        }
      },
      "s2": {
        "a": {
          "state": "s2",
          "stackPushChar": ""
        },
        "#": {
          "state": "s3",
          "stackPushChar": ""
        }
      },
      "startState": "start",
      "acceptStates": ["s3"],
      "states": {
        "start": {
          "top": 246,
          "left": 225,
          "displayId": "s0"
        },
        "s1": {
          "top": 248,
          "left": 452,
          "displayId": "s1"
        },
        "s2": {
          "top": 247,
          "left": 680,
          "displayId": "s2"
        },
        "s3": {
          "isAccept": true,
          "top": 247,
          "left": 893,
          "displayId": "s3"
        }
      },
      "transitions": [
        {
          "stateA": "start",
          "label": "ε,ε,#",
          "stateB": "s0"
        },
        {
          "stateA": "s0",
          "label": "a,ε,a",
          "stateB": "s0"
        },
        {
          "stateA": "s0",
          "label": "ε,ε,ε",
          "stateB": "s1"
        },
        {
          "stateA": "s1",
          "label": "b,ε,ε",
          "stateB": "s1"
        },
        {
          "stateA": "s1",
          "label": "ε,ε,ε",
          "stateB": "s2"
        },
        {
          "stateA": "s2",
          "label": "a,ε,ε",
          "stateB": "s2"
        },
        {
          "stateA": "s2",
          "label": "ε,ε,ε",
          "stateB": "s3"
        }
      ],
      "bulkTests": {
        "accept": "aabaa\naaaaabbbbbaaaaa\nabab",
        "reject": "aabbbbaaaaa\nabab\nabbbbbaa"
      }
    }
  }
}
```

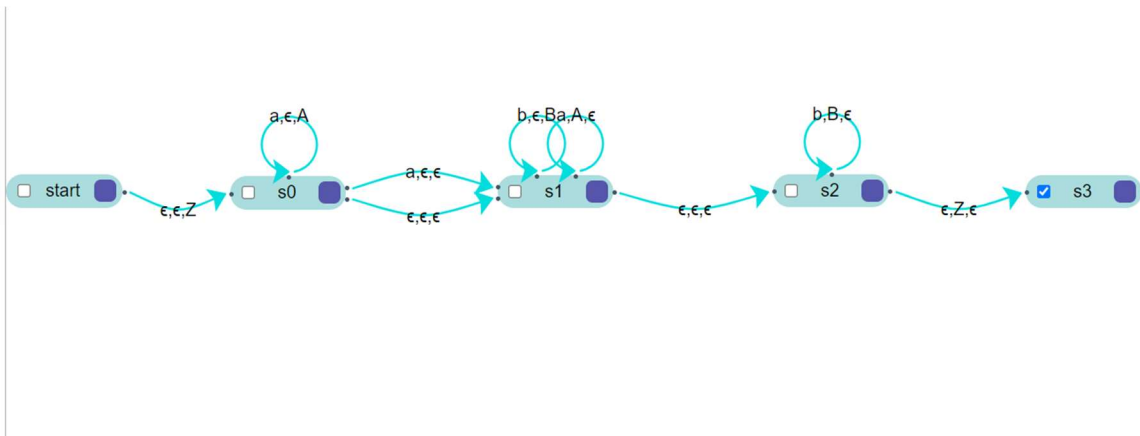


10. Assume $m, n \geq 0$. Describe a DFA that recognises $\{a^m b^n a^m\}$. Explain why this question is different from the previous question.

There are no DFAs that recognise the given language. The reason is that the language given is not regular.

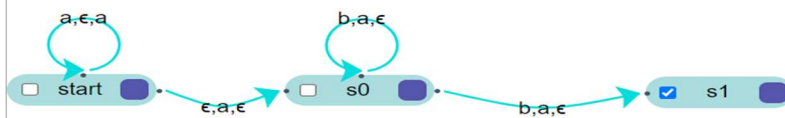
11. Describe a pushdown automaton that recognises $\{a^m b^{2n} \mid m, n \geq 0\}$.

```
{
  "type": "PDA",
  "pda": {
    "transitions": {
      "start": {
        "": {
          [
            {
              "state": "s0",
              "stackPushChar": "Z"
            }
          ]
        },
        "s0": {
          "a": {
            [
              {
                "state": "s0",
                "stackPushChar": "A"
              },
              {
                "state": "s1",
                "stackPushChar": ""
              }
            ]
          },
          "": {
            [
              {
                "state": "s1",
                "stackPushChar": ""
              }
            ]
          },
          "s1": {
            "b": {
              [
                {
                  "state": "s1",
                  "stackPushChar": "B"
                }
              ]
            },
            "a": {
              [
                [
                  [
                    [
                      {
                        "state": "s1",
                        "stackPushChar": ""
                      }
                    ]
                  ],
                  [
                    {
                      "state": "s2",
                      "stackPushChar": ""
                    }
                  ]
                ]
              ],
              "A": [
                {
                  "state": "s1",
                  "stackPushChar": ""
                }
              ]
            },
            "": {
              [
                {
                  "state": "s2",
                  "stackPushChar": ""
                }
              ]
            },
            "s2": {
              "Z": [
                {
                  "state": "s3",
                  "stackPushChar": ""
                }
              ]
            },
            "b": [
              {
                "state": "s2",
                "stackPushChar": ""
              }
            ]
          }
        },
        "s3": {}
      },
      "startState": "start",
      "acceptStates": ["s3"],
      "states": {
        "start": {
          "top": 247,
          "left": 203,
          "displayId": "s0"
        },
        "s0": {
          "top": 246,
          "left": 445,
          "displayId": "s1"
        },
        "s1": {
          "top": 245,
          "left": 695,
          "displayId": "s2"
        },
        "s2": {
          "isAccept": true,
          "top": 246,
          "left": 924,
          "displayId": "s3"
        },
        "s3": {}
      },
      "transitions": [
        {
          "stateA": "start",
          "label": "\u03b5, \u03b5, Z",
          "stateB": "s0"
        },
        {
          "stateA": "s0",
          "label": "a, \u03b5, A",
          "stateB": "s0"
        },
        {
          "stateA": "s0",
          "label": "a, \u03b5, \u03b5",
          "stateB": "s1"
        },
        {
          "stateA": "s0",
          "label": "\u03b5, \u03b5, \u03b5",
          "stateB": "s1"
        },
        {
          "stateA": "s1",
          "label": "b, \u03b5, B",
          "stateB": "s1"
        },
        {
          "stateA": "s1",
          "label": "a, A, \u03b5",
          "stateB": "s1"
        },
        {
          "stateA": "s1",
          "label": "\u03b5, \u03b5, \u03b5",
          "stateB": "s2"
        },
        {
          "stateA": "s2",
          "label": "\u03b5, Z, \u03b5",
          "stateB": "s3"
        },
        {
          "stateA": "s2",
          "label": "b, B, \u03b5",
          "stateB": "s2"
        }
      ],
      "bulkTests": {
        "accept": "abb\naabb\naaaa\naaaaa",
        "reject": "bbbbbb\nbbbbbb\nabbb"
      }
    }
  }
}
```



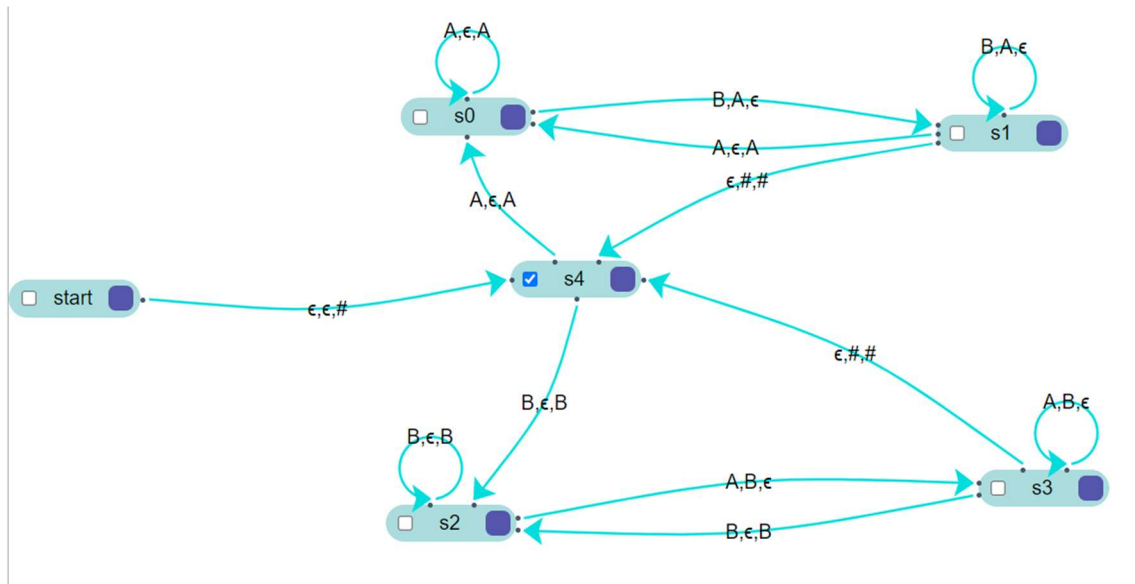
12. Describe a pushdown automaton that recognises $\{a^m b^n \mid m > n > 0\}$.

```
{
  "type": "PDA",
  "pda": {
    "transitions": {
      "start": {
        "a": {
          [
            {
              "state": "start",
              "stackPushChar": "a"
            }
          ]
        },
        "": {
          [
            {
              "state": "s0",
              "stackPushChar": ""
            }
          ]
        },
        "s0": {
          "b": {
            [
              {
                "state": "s1",
                "stackPushChar": ""
              }
            ]
          },
          "a": [
            {
              "state": "s0",
              "stackPushChar": ""
            }
          ]
        },
        "s1": {}
      },
      "startState": "start",
      "acceptStates": ["s1"],
      "states": {
        "start": {
          "top": 246,
          "left": 196,
          "displayId": "s0"
        },
        "s0": {
          "isAccept": true,
          "top": 248,
          "left": 450,
          "displayId": "s1"
        },
        "s1": {}
      },
      "transitions": [
        {
          "stateA": "start",
          "label": "a, \u03b5, a",
          "stateB": "start"
        },
        {
          "stateA": "start",
          "label": "\u03b5, a, \u03b5",
          "stateB": "s0"
        },
        {
          "stateA": "s0",
          "label": "b, a, \u03b5",
          "stateB": "s0"
        },
        {
          "stateA": "s0",
          "label": "b, a, \u03b5",
          "stateB": "s1"
        }
      ],
      "bulkTests": {
        "accept": "aab\naaaaaaaab",
        "reject": "\nabba\naabbbbbb"
      }
    }
  }
}
```



13. Describe a pushdown automaton that recognises $\{w \mid \#_a w = \#_b w\}$, where $\#_a w$ is the number of as appearing in w and $\#_b w$ is the number of bs appearing in w .

```
{
  "type": "PDA",
  "pda": {
    "transitions": {
      "start": {
        "A": {
          "": []
        },
        "B": {
          "": []
        },
        "": {
          "": {
            "state": "s4",
            "stackPushChar": "#"
          }
        }
      },
      "s0": {
        "A": {
          "": {
            "state": "s0",
            "stackPushChar": "A"
          }
        },
        "B": {
          "A": {
            "state": "s1",
            "stackPushChar": ""
          }
        },
        "s1": {
          "B": {
            "A": {
            "state": "s1",
            "stackPushChar": ""
          },
          "#": []
        },
        "A": {
          "": {
            "state": "s0",
            "stackPushChar": "A"
          },
          "": {
            "state": "s4",
            "stackPushChar": "#"
          }
        },
        "s2": {
          "B": {
            "state": "s2",
            "stackPushChar": "B"
          },
          "A": {
            "B": {
            "state": "s3",
            "stackPushChar": ""
          }
        },
        "s3": {
          "B": {
            "state": "s2",
            "stackPushChar": "B"
          },
          "A": {
            "B": {
            "state": "s3",
            "stackPushChar": ""
          },
          "#": []
        },
        "": {
          "state": "s4",
          "stackPushChar": "#"
        }
      },
      "s4": {
        "A": {
          "": {
            "state": "s0",
            "stackPushChar": "A"
          },
          "B": {
            "state": "s2",
            "stackPushChar": "B"
          }
        }
      },
      "startState": "start",
      "acceptStates": ["s4"],
      "states": {
        "start": {
          "isAccept": false,
          "top": 230,
          "left": 402,
          "displayId": "start"
        },
        "s4": {
          "isAccept": true,
          "top": 230,
          "left": 402,
          "displayId": "s4"
        },
        "s0": {
          "top": 100,
          "left": 314,
          "displayId": "s0"
        },
        "s1": {
          "top": 113,
          "left": 744,
          "displayId": "s1"
        },
        "s2": {
          "top": 425,
          "left": 302,
          "displayId": "s2"
        },
        "s3": {
          "top": 397,
          "left": 777,
          "displayId": "s3"
        }
      },
      "transitions": [
        {
          "stateA": "start",
          "label": "ε, ε, #",
          "stateB": "s4"
        },
        {
          "stateA": "s0",
          "label": "A, ε, A",
          "stateB": "s0"
        },
        {
          "stateA": "s0",
          "label": "B, A, ε",
          "stateB": "s1"
        },
        {
          "stateA": "s1",
          "label": "B, A, ε",
          "stateB": "s1"
        },
        {
          "stateA": "s1",
          "label": "A, ε, A",
          "stateB": "s0"
        },
        {
          "stateA": "s1",
          "label": "ε, #, #",
          "stateB": "s4"
        },
        {
          "stateA": "s2",
          "label": "B, ε, B",
          "stateB": "s2"
        },
        {
          "stateA": "s2",
          "label": "A, B, ε",
          "stateB": "s3"
        },
        {
          "stateA": "s3",
          "label": "B, ε, B",
          "stateB": "s2"
        },
        {
          "stateA": "s3",
          "label": "A, B, ε",
          "stateB": "s3"
        },
        {
          "stateA": "s3",
          "label": "ε, #, #",
          "stateB": "s4"
        },
        {
          "stateA": "s4",
          "label": "A, ε, A",
          "stateB": "s0"
        },
        {
          "stateA": "s4",
          "label": "B, ε, B",
          "stateB": "s2"
        }
      ],
      "bulkTests": {
        "accept": "AB\nABAB\nABABAB\nBA\nAAABBB\nAABBBBBBAAA",
        "reject": "A\nB\nABA\nBB\nABABB"
      }
    }
  }
}
```



14. Describe a pushdown automaton that recognises $\{w|_{\#_a}w = 2\#_bw\}$.

15. Describe a pushdown automaton that recognises $\{w|_{\#_a}w \neq \#_bw\}$

```
{
  "type": "PDA",
  "pda": {
    "transitions": {
      "start": {
        "A": {
          "": []
        },
        "B": {
          "": []
        },
        "": {
          [
            {
              "state": "s4",
              "stackPushChar": "#",
              "label": "ε,ε,#"
            }
          ]
        }
      },
      "s0": {
        "A": {
          [
            {
              "state": "s0",
              "stackPushChar": "A",
              "label": "A,ε,A"
            }
          ]
        },
        "B": {
          [
            {
              "state": "s1",
              "stackPushChar": "",
              "label": "B,A,ε"
            }
          ]
        },
        "": {
          [
            {
              "state": "s5",
              "stackPushChar": "A",
              "label": "ε,A,A"
            }
          ]
        }
      },
      "s1": {
        "B": {
          [
            {
              "state": "s1",
              "stackPushChar": "",
              "label": "B,A,ε"
            }
          ]
        },
        "A": {
          [
            {
              "state": "s0",
              "stackPushChar": "A",
              "label": "A,ε,A"
            }
          ]
        },
        "": {
          [
            {
              "state": "s4",
              "stackPushChar": "#",
              "label": "ε,#,#"
            }
          ]
        }
      },
      "s2": {
        "B": {
          [
            {
              "state": "s2",
              "stackPushChar": "B",
              "label": "B,ε,B"
            }
          ]
        },
        "A": {
          [
            {
              "state": "s3",
              "stackPushChar": "",
              "label": "A,B,ε"
            }
          ]
        },
        "": {
          [
            {
              "state": "s5",
              "stackPushChar": "B",
              "label": "ε,B,B"
            }
          ]
        }
      },
      "s3": {
        "B": {
          [
            {
              "state": "s5",
              "stackPushChar": "B",
              "label": "B,ε,B"
            }
          ]
        },
        "A": {
          [
            {
              "state": "s3",
              "stackPushChar": "",
              "label": "A,B,ε"
            }
          ]
        },
        "": {
          [
            {
              "state": "s4",
              "stackPushChar": "#",
              "label": "ε,#,#"
            }
          ]
        }
      },
      "s4": {
        "A": {
          [
            {
              "state": "s0",
              "stackPushChar": "A",
              "label": "A,ε,A"
            }
          ]
        },
        "B": {
          [
            {
              "state": "s2",
              "stackPushChar": "B",
              "label": "B,ε,B"
            }
          ]
        },
        "": {
          [
            {
              "state": "s5",
              "stackPushChar": "A",
              "label": "ε,A,A"
            }
          ]
        }
      }
    },
    "startState": "start",
    "acceptStates": [
      "s5"
    ],
    "states": {
      "start": {
        "top": 219, "left": 282, "displayId": "start"
      },
      "s4": {
        "top": 219, "left": 282, "displayId": "s4"
      },
      "s0": {
        "top": 86, "left": 309, "displayId": "s0"
      },
      "s1": {
        "top": 84, "left": 892, "displayId": "s1"
      },
      "s5": {
        "isAccept": true,
        "top": 234, "left": 908, "displayId": "s5"
      },
      "s2": {
        "top": 435, "left": 216, "displayId": "s2"
      },
      "s3": {
        "top": 434.6000061035156, "left": 900, "displayId": "s3"
      }
    }
  }
}
```