Person-In-The-Middle via Arp Spoofing

Execution

(a) What is Kali's main interface's MAC address?

As we can see below (after "ether") it is 26:c3:43:f1:15:92.

```
-(kali⊕kali)-[~]
└─$ ifconfig
eth0: flags=4163<UP,BROADCAST,RUNNING,MULTICAST> mtu 1500
       inet 192.168.64.4 netmask 255.255.255.0 broadcast 192.168.64.255
       inet6 fe80::f254:91db:c5f:f614 prefixlen 64 scopeid 0×20<link>
       inet6 fd66:6c69:1b18:220f:feb7:6634:ec8c:3a1c prefixlen 64 scopeid
0×0<global>
       ether 26:c3:43:f1:15:92 txqueuelen 1000 (Ethernet)
       RX packets 109 bytes 18865 (18.4 KiB)
       RX errors 0 dropped 0 overruns 0 frame 0
       TX packets 51 bytes 8219 (8.0 KiB)
       TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0
lo: flags=73<UP,LOOPBACK,RUNNING> mtu 65536
       inet 127.0.0.1 netmask 255.0.0.0
       inet6 :: 1 prefixlen 128 scopeid 0×10<host>
       loop txqueuelen 1000 (Local Loopback)
       RX packets 4 bytes 240 (240.0 B)
       RX errors 0 dropped 0 overruns 0 frame 0
       TX packets 4 bytes 240 (240.0 B)
       TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0
```

(b) What is Kali's main interface's IP address?

Kali's main interface's IP address is **192.168.64.4**.

(c) What is Metasploitable's main interface's MAC address?

As shown below (after "HWaddr") it is fe:5a:c4:e3:6c:21.

```
msfadmin@metasploitable:~$ ifconfig
          Link encap:Ethernet HWaddr fe:5a:c4:e3:6c:21 inet addr:192.168.64.3 Bcast:192.168.64.255 Mask:255.255.255.0
eth0
          inet6 addr: fd66:6c69:1b18:220f:fc5a:c4ff:fee3:6c21/64 Scope:Global
          inet6 addr: fe80::fc5a:c4ff:fee3:6c21/64 Scope:Link
          UP BROADCAST RUNNING MULTICAST MTU:1500 Metric:1
          RX packets:114 errors:0 dropped:0 overruns:0 frame:0
          TX packets:94 errors:0 dropped:0 overruns:0 carrier:0
          collisions:0 txqueuelen:1000
          RX bytes:18192 (17.7 KB) TX bytes:12291 (12.0 KB)
          Base address:0xc000 Memory:febc0000-febe0000
          Link encap:Local Loopback
lo
          inet addr:127.0.0.1 Mask:255.0.0.0
          inet6 addr: ::1/128 Scope:Host
          UP LOOPBACK RUNNING MTU: 16436 Metric: 1
          RX packets:199 errors:0 dropped:0 overruns:0 frame:0
          TX packets:199 errors:0 dropped:0 overruns:0 carrier:0
          collisions:0 txqueuelen:0
          RX bytes:71137 (69.4 KB) TX bytes:71137 (69.4 KB)
```

(d) What is Metasploitable's main interface's IP address?

Metasploitable's main interface's IP address is 192.168.64.3.

(e) Show Kali's routing table.

```
-(kali⊕kali)-[~]
—$ netstat -r
Kernel IP routing table
                Gateway
Destination
                                                          MSS Window irtt Ifac
                                 Genmask
                                                  Flags
e
default
                192.168.64.1
                                 0.0.0.0
                                                  UG
                                                             0 0
                                                                          0 eth0
192.168.64.0
                0.0.0.0
                                 255.255.255.0
                                                  U
                                                             0 0
                                                                          0 eth0
```

(f) Show Kali's ARP cache.

```
-(kali⊕kali)-[~]
_s arp
Address
                          HWtype
                                  HWaddress
                                                       Flags Mask
                                                                              Iface
192.168.64.3
                          ether
                                  fe:5a:c4:e3:6c:21
                                                       C
                                                                              eth0
192.168.64.1
                          ether
                                  e2:b5:5f:3f:8b:64
                                                                              eth0
```

(g) Show Metasploitable's routing table.

```
msfadmin@metasploitable:~$ netstat -r
Kernel IP routing table
Destination
                 Gateway
                                                    Flags
                                                             MSS Window
                                                                           irtt Iface
                                   Genmask
192.168.64.0
                                   255.255.255.0
                                                    U
                                                               \mathbf{0}
                                                                              0 eth0
default
                 192.168.64.1
                                  0.0.0.0
                                                    UG
                                                               0 0
                                                                              0 eth0
```

(h) Show Metasploitable's ARP cache.

msfadmin@metasploitable:	~Ş arp			
Address	HWtype	HWaddress	Flags Mask	Iface
192.168.64.1	ether	E2:B5:5F:3F:8B:64	С	eth0

(i) Suppose the user of Metasploitable wants to get the CS338 sandbox page via the command "curl http://cs338.jeffondich.com/". To which MAC address should Metasploitable send the TCP SYN packet to get the whole HTTP query started? Explain why.

After testing this on Kali with Wireshark, it looks like the TCP SYN packet is sent to the following MAC address: **e2:b5:5f:3f:8b:64**. It's reasonable to assume that the same would happen for Metasploitable, as it is set up symmetrically to Kali on my system.

That address is the MAC address of my machine, as shown below. This makes sense since the packet must first hop through my physical machine before reaching its destination.

(j) Fire up Wireshark on Kali. Start capturing packets for "tcp port http". On Metasploitable, execute "curl http://cs338.jeffondich.com/". On Kali, stop capturing. Do you see an HTTP response on Metasploitable? Do you see any captured packets in Wireshark on Kali?

After running this command on Metasploitable, I got an HTTP response on Metasploitable:

However, Kali did not capture any packets, as packets did not originate from or arrive at Kali.

(k) Do ARP Poisoning

Got this set up and working!

(l) Show Metasploitable's ARP cache. How has it changed?

Here is Metasploitable's new ARP cache.

```
msfadmin@metasploitable:~$ arp
Address HWtype HWaddress Flags Mask Iface
192.168.64.4 ether 26:C3:43:F1:15:92 C eth0
192.168.64.1 ether 26:C3:43:F1:15:92 C eth0
```

One change is that we now have Kali added to the ARP cache. More importantly, both my MacBook and Kali now have the MAC address of Kali associated with them on Metasploitable (26:c3:43:f1:15:92).

(m) Without actually doing it yet, predict what will happen if you execute "curl http://cs338.jeffondich.com/" on Metasploitable now. Specifically, to what MAC address will Metasploitable send the TCP SYN packet? Explain why.

I guess that Metasploitable will send the TCP SYN packet to Kali's MAC address, as this is now the MAC address associated with my MacBook in Metasploitable's ARP cache. The packet still needs to hop through my machine, so it will need to be addressed to "my machine's" MAC address.

- (n) Start Wireshark capturing "tcp port http" again. Done.
- (o) Execute "curl http://cs338.jeffondich.com/" on Metasploitable. On Kali, stop capturing. Do you see an HTTP response on Metasploitable? Do you see captured packets in Wireshark? Can you tell from Kali what messages went back and forth between Metasploitable and cs338.jeffondich.com?

Again, I do see an HTTP response on Metaspoitable. However, this time, I did capture packets in Wireshark. I can see every message that went back and forth between Metasploitable and cs338.jeffondich.com. Here's a snapshot showing the HTTP response that Metasploitable received (on Kali).

NO.	rime	Jource	Destillation	FIOLOCOL	Length mio
Г	1 0.000000000	192.168.64.3	172.233.221.124	TCP	74 50782 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1
	2 0.005758832	192.168.64.3	172.233.221.124	TCP	74 [TCP Retransmission] 50782 → 80 [SYN] Seq=0
	3 0.024910987	172.233.221.124	192.168.64.3	TCP	66 80 → 50782 [SYN, ACK] Seq=0 Ack=1 Win=64240
	4 0.029757972	172.233.221.124	192.168.64.3	TCP	66 [TCP Retransmission] 80 → 50782 [SYN, ACK]
	5 0.031084956	192.168.64.3	172.233.221.124	TCP	60 50782 → 80 [ACK] Seq=1 Ack=1 Win=5888 Len=0
+	6 0.032078099	192.168.64.3	172.233.221.124	HTTP	212 GET / HTTP/1.1
	7 0.036963077	192.168.64.3	172.233.221.124	TCP	54 50782 → 80 [ACK] Seq=1 Ack=1 Win=5888 Len=0
	8 0.037172042	192.168.64.3	172.233.221.124	TCP	212 [TCP Retransmission] 50782 → 80 [PSH, ACK]
	9 0.055086051	172.233.221.124	192.168.64.3	TCP	60 80 → 50782 [ACK] Seq=1 Ack=159 Win=64128 Le
+	10 0.057355893	172.233.221.124	192.168.64.3	HTTP	789 HTTP/1.1 200 OK (text/html)
	11 0.061179004	172.233.221.124	192.168.64.3	TCP	54 80 → 50782 [ACK] Seq=1 Ack=159 Win=64128 L€
	12 0.061591906	172.233.221.124	192.168.64.3	TCP	789 [TCP Retransmission] 80 → 50782 [PSH, ACK]
	13 0.062474979	192.168.64.3	172.233.221.124	TCP	60 50782 → 80 [ACK] Seq=159 Ack=736 Win=7360 L
	14 0.065227843	192.168.64.3	172.233.221.124	TCP	60 50782 → 80 [FIN, ACK] Seq=159 Ack=736 Win=7
	15 0.069256846	192.168.64.3	172.233.221.124	TCP	54 [TCP Keep-Alive] 50782 → 80 [ACK] Seq=159 A
_		· · · · · · · · · · · · · · · · · · ·			
	<html lang="en">\</html>	.n			0120 65 73 3a 20 62 79 74 65 73 0d 0a 0d 0a <mark>3c</mark>
	<head>\n</head>				0130 4f 43 54 59 50 45 20 68 74 6d 6c 3e 0a 3c
		ırset="utf-8">\n			0140 6d 6c 20 6c 61 6e 67 3d 22 65 6e 22 3e 0a
		338 Sandbox\	.n		0150 20 20 3c 68 65 61 64 3e 0a 20 20 20 20 20
	\n				0160 20 3c 6d 65 74 61 20 63 68 61 72 73 65 74 0170 75 74 66 2d 38 22 3e 0a 20 20 20 20 20 20
	\n				0180 3c 74 69 74 6c 65 3e 43 53 33 33 38 20 53
	<body>\n</body>				0190 64 62 6f 78 3c 2f 74 69 74 6c 65 3e 0a 20
	<h1>CS338 Sandbox</h1> \n				0130 04 02 01 70 30 21 74 03 74 00 03 30 00 20 01a0 20 3c 2f 68 65 61 64 3e 0a 0a 20 20 20 20
	<h2>Fun with security, or maybe insecurity</h2> \n				01b0 6f 64 79 3e 0a 20 20 20 20 20 20 20 20 3c
\n This page should be the page you retrieve for the "Gett			01c0 3e 43 53 33 33 38 20 53 61 6e 64 62 6f 78		
			01d0 68 31 3e 0a 20 20 20 20 20 20 20 20 3c 68		
	assignment. Here's my head, as advertised:\n			01e0 46 75 6e 20 77 69 74 68 20 73 65 63 75 72	
		· <img src="jeff-squar</td><td>e-colorado.jpg" style<="" td=""/> <td>="width:</td> <td>01f0 79 2c 20 6f 72 20 6d 61 79 62 65 20 69 6e</td>	="width:	01f0 79 2c 20 6f 72 20 6d 61 79 62 65 20 69 6e	
	\n				0200 63 75 72 69 74 79 3c 2f 68 32 3e 0a 0a 20
	\n				0210 20 20 20 20 20 3c 70 3e 54 68 69 73 20 70
	\n				0220 65 20 73 68 6f 75 6c 64 20 62 65 20 74 68

(p) Explain in detail what happened. How did Kali change Metasploitable's ARP cache?

We can watch the attack in action by capturing ARP packets:

Source	Destination	Protocol	Length Info
26:c3:43:f1:15:92	fe:5a:c4:e3:6c:21	ARP	42 192.168.64.1 is at 26:c3:43:f1:15:92
26:c3:43:f1:15:92	e2:b5:5f:3f:8b:64	ARP	42 192.168.64.3 is at 26:c3:43:f1:15:92
26:c3:43:f1:15:92	fe:5a:c4:e3:6c:21	ARP	42 192.168.64.1 is at 26:c3:43:f1:15:92
26:c3:43:f1:15:92	e2:b5:5f:3f:8b:64	ARP	42 192.168.64.3 is at 26:c3:43:f1:15:92

As we can see, two ARP announcement packets are sent repeatedly:

- One is sent to my MacBook, saying "The MAC address for Metasploitable is [Kali's MAC address"
- The other is sent to Metasploitable, saying "The MAC address for Varun's MacBook is [Kali's MAC address]"

This way, for any communications passing between Metasploitable and my MacBook, Kali can act as the AITM after these MAC addresses are cached. I believe that these packets are sent repeatedly because ARP caches are periodically cleared.

(q) If you wanted to design an ARP spoofing detector, what would you have your detector do?

Something that I noticed about the ARP poisoning attack was that it requires several ARP packets to be sent since the cache of each machine gets cleared over time. A protective measure would be to flag frequent ARP messages as suspicious. This is susceptible to false positives, however, as even "normal" ARP packets may need to be sent frequently.

A heavy-duty solution could be to hardcode MAC addresses for devices on the same network. Perhaps only a single trusted device could broadcast MAC/IP address updates (like the network's router), using encryption/a digital signature that identifies the source as valid. Then, any ARP packet that conflicts with hardcoded or "trusted" information would be detected as suspicious.

Synthesis

(a) Explain in detail Mal's strategy for intercepting the traffic between Alice and Bob. Use any of your observations from the Execution section to clarify your explanation.

The following picture (also shown above) is fairly illustrative of how Mal acts as an AITM.

Source	Destination	Protocol	Length Info
26:c3:43:f1:15:92	fe:5a:c4:e3:6c:21	ARP	42 192.168.64.1 is at 26:c3:43:f1:15:92
26:c3:43:f1:15:92	e2:b5:5f:3f:8b:64	ARP	42 192.168.64.3 is at 26:c3:43:f1:15:92
26:c3:43:f1:15:92	fe:5a:c4:e3:6c:21	ARP	42 192.168.64.1 is at 26:c3:43:f1:15:92
26:c3:43:f1:15:92	e2:b5:5f:3f:8b:64	ARP	42 192.168.64.3 is at 26:c3:43:f1:15:92

Essentially, Mal can intercept traffic between Alice and Bob by placing herself between Alice and Alice's first hop (on the way to Bob). Mal can accomplish this by sending out two ARP announcement packets (repeatedly):

- One to Alice saying that Alice's first hop's MAC address is [Mal's MAC address]
- One to Alice's first hop saying that Alice's MAC address is [Mal's MAC address]

After Alice and Alice's first hop update their ARP tables to reflect this, Mal can intercept packets sent between the two devices. In other words, communication between Alice and Alice's first hop always passes through Mal, regardless of the direction in which data is traveling.

(b) From Alice's perspective, is this attack detectable? If not, why not? If so, how would Alice's setup need to change to detect the attack?

Probably not: if Alice's ARP table were constantly being cleared and she had no other memory of MAC/IP relationships, then a new suspicious MAC/IP pairing being broadcasted by Mal would be undetectable. To make this attack detectable, Alice would need some form of persistent storage for remembering past MAC/IP pairs to flag newly broadcasted MAC/IP pairs that conflict with older information.

(c) From Bob's perspective, is this attack detectable?

No: Bob is unable to see anything past his first hop (in the direction of Alice). In other words, Mal is protected by all the hops that Bob's packets need to make to get to Alice.

(d) Could Alice or Bob detect and/or prevent this attack if the website in question was using HTTPS instead of HTTP? Explain.

I believe that using HTTPS could provide some protection, as Mal wouldn't be able to decrypt any of the information being shared between Alice and Bob. However, there are still issues that

HTTPS doesn't resolve. For example, Mal could still perform a replay attack by sending Bob copies of Alice's packets. This could potentially cause damage (e.g. imagine Mal adding forum posts in a tight loop to Bob's database and blaming Alice for the damages).