



SCHOOL OF  
COMPUTING

M M VASANTH

CH.SC.U4CSE24227

Week - 5

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Design and Analysis of Algorithm(23CSE211)

## 1. QUICK SORT

**Code:**

```
#include <stdio.h>

void swap(int* a, int* b) {
    int t = *a;
    *a = *b;
    *b = t;
}

int partition(int arr[], int low, int high) {
    int pivot = arr[high];
    int i = (low - 1);

    for (int j = low; j <= high - 1; j++) {
        if (arr[j] < pivot) {
            i++;
            swap(&arr[i], &arr[j]);
        }
    }
    swap(&arr[i + 1], &arr[high]);
    return (i + 1);
}

void quickSort(int arr[], int low, int high) {
    if (low < high) {
        int pi = partition(arr, low, high);

        quickSort(arr, low, pi - 1);
        quickSort(arr, pi + 1, high);
    }
}
```

```
int main() {
    int arr[100], n, i;
    printf("Name: M M VASANTH\n");
    printf("Roll No: CH.SC.U4CSE24227\n\n");
    printf("Enter number of elements: ");
    scanf("%d", &n);
    printf("Enter %d elements:\n", n);
    for(i = 0; i < n; i++) {
        scanf("%d", &arr[i]);
    }

    quickSort(arr, 0, n - 1);

    printf("\nSorted array:\n");
    for(i = 0; i < n; i++) {
        printf("%d ", arr[i]);
    }
    return 0;
}
```

## Output:

```
vasanth@VasanthMM:/mnt/c/Users/vasan$ nano quick.c
vasanth@VasanthMM:/mnt/c/Users/vasan$ gcc quick.c -o quick
vasanth@VasanthMM:/mnt/c/Users/vasan$ ./quick
Name: M M VASANTH
Roll No: CH.SC.U4CSE24227

Enter number of elements: 5
Enter 5 elements:
34
56
63
92
35

Sorted array:
34 35 56 63 92 vasanth@VasanthMM:/mnt/c/Users/vasan$ |
```

### ▣ Time Complexity: $O(N^2)$

The array is partitioned  $\log N$  times (divide step). In each level of recursion, we iterate through the array segments to compare with the pivot ( $O(N)$ ). Thus, the total time is  $N * \log N$ . However, in the worst case (already sorted array), it degrades to  $O(N^2)$ .

### ▣ Space Complexity: $O(N)$

Quick Sort is an in-place sorting algorithm (it doesn't create a new array for sorting). However, it uses stack memory for recursion. In the worst case, the height of the recursion tree is  $N$ .

## 2. MERGE SORT

### Code:

```
#include <stdio.h>

void merge(int arr[], int l, int m, int r) {

    int i, j, k;
    int n1 = m - l + 1;
    int n2 = r - m;

    int L[50], R[50];

    for (i = 0; i < n1; i++)
        L[i] = arr[l + i];
    for (j = 0; j < n2; j++)
        R[j] = arr[m + 1 + j];

    i = 0; j = 0; k = l;
    while (i < n1 && j < n2) {
        if (L[i] <= R[j]) {
            arr[k] = L[i];
            i++;
        } else {
            arr[k] = R[j];
            j++;
        }
    }
}
```

```
    k++;

}

while (i < n1) {

    arr[k] = L[i];

    i++; k++;

}

while (j < n2) {

    arr[k] = R[j];

    j++; k++;

}

void mergeSort(int arr[], int l, int r) {

    if (l < r) {

        int m = l + (r - 1) / 2;

        mergeSort(arr, l, m);

        mergeSort(arr, m + 1, r);

        merge(arr, l, m, r);

    }

}
```

```
int main() {  
    int arr[100], n, i;  
  
    printf("Name: M M VASANTH\n");  
    printf("Roll No: CH.SC.U4CSE24227\n\n");  
    printf("Enter number of elements: ");  
    scanf("%d", &n);  
    printf("Enter %d elements:\n", n);  
    for(i = 0; i < n; i++) {  
        scanf("%d", &arr[i]);  
    }  
  
    mergeSort(arr, 0, n - 1);  
  
    printf("\nSorted array:\n");  
    for(i = 0; i < n; i++) {  
        printf("%d ", arr[i]);  
    }  
    return 0;  
}
```

## **Output:**

```
Name: M M VASANTH
Roll No: CH.SC.U4CSE24227
```

```
Enter number of elements: 7
```

```
Enter 7 elements:
```

```
23
34
64
54
79
92
95
```

```
Sorted array:
```

```
23 34 54 64 79 92 95 vasanth@VasanthMM:/mnt/c/Users/vasan$ |
```

### **② Time Complexity: $O(N \log N)$**

Similar to Quick Sort, the array is divided  $\log N$  times. The merge operation takes linear time  $O(N)$  to combine the subarrays. Since the split is always perfectly balanced (halved), Merge Sort guarantees  $O(N \log N)$  time complexity even in the worst case.

### **③ Space Complexity: $O(N)$**

Unlike Quick Sort, Merge Sort is not in-place. The merge function requires temporary arrays (L and R) to store the split elements before combining them. This consumes auxiliary space proportional to the input size N.

### 3. BINARY SEARCH TREE (BST)

**Code:**

```
#include <stdio.h>
#include <stdlib.h>

struct Node {
    int key;
    struct Node *left, *right;
};

struct Node* newNode(int item) {
    struct Node* temp = (struct Node*)malloc(sizeof(struct Node));
    temp->key = item;
    temp->left = temp->right = NULL;
    return temp;
}

struct Node* insert(struct Node* node, int key) {
    if (node == NULL) return newNode(key);

    if (key < node->key)
        node->left = insert(node->left, key);
    else if (key > node->key)
        node->right = insert(node->right, key);
}
```

```
    return node;

}

void inorder(struct Node* root) {

    if (root != NULL) {

        inorder(root->left);

        printf("%d ", root->key);

        inorder(root->right);

    }

}

struct Node* search(struct Node* root, int key) {

    if (root == NULL || root->key == key)

        return root;

    if (root->key < key)

        return search(root->right, key);

    return search(root->left, key);

}

int main() {

    struct Node* root = NULL;

    int n, i, val, key;

    printf("Name: M M VASANTH\n");
```

```
printf("Roll No: CH.SC.U4CSE24227\n\n");

printf("Enter number of elements to insert: ");
scanf("%d", &n);

printf("Enter elements:\n");
for(i=0; i<n; i++) {
    scanf("%d", &val);
    root = insert(root, val);
}

printf("Inorder Traversal (Sorted): ");
inorder(root);
printf("\n");

printf("Enter element to search: ");
scanf("%d", &key);
if(search(root, key) != NULL)
    printf("Element Found\n");
else
    printf("Element Not Found\n");

return 0;
}
```

## **Output:**

```
Name: M M VASANTH
Roll No: CH.SC.U4CSE24227

Enter number of elements to insert: 5
Enter elements:
24
43
56
95
94
Inorder Traversal (Sorted): 24 43 56 94 95
Enter element to search: 94
Element Found
vasanth@VasanthMM:/mnt/c/Users/vasan$ |
```

### **② Time Complexity: O(N)**

- Insertion/Search: In a balanced tree, we discard half the nodes at each step, making the height  $\log N$ .
- Worst Case: If the tree becomes "skewed" (like a linked list, e.g., inserting 1, 2, 3, 4, 5 in order), the height becomes  $N$ , degrading operations to  $O(N)$ .

### **③ Space Complexity: O(N)**

The space is required to store the  $N$  nodes of the tree. Additionally, the recursion stack for insertion or traversal can grow up to  $O(N)$  in the worst-case (skewed tree).