

# Uzhhorod National University

# UzhNU Machata

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Geometry	
10 Misc. algorithms	19
$\underline{\text{Contest}}$ (1)	
template.cpp	101 lines
// #pragma GCC optimize("Ofast") // #pragma GCC optimize("O3, unroll—loops") // #pragma GCC target("sse, sse2, sse3, ssse3, sse4") // #pragma GCC target("avx2, bmi, bmi2, popcnt, lzcnt") // #define _GLIBCXX_DEBUG // #define _GLIBCXX_DEBUG_PEDANTIC	
<pre>#include <cassert> #include <iomanip> #include <iostream> #include <algorithm> #include <map> #include <set> #include <functional> #include <array> #include <numeric> #include <queue> #include <deque> #include &lt;<cathodox< pre=""></cathodox<></deque></queue></numeric></array></functional></set></map></algorithm></iostream></iomanip></cassert></pre>	
using namespace std;	
<pre>const int MOD = 998244353; const long double PI = 3.141592653589793; using ll = long long; const ll INF = le18;</pre>	
// #define int ll	
//> sashko123's defines:	
<pre>#define itn int //Vasya sorry :( #define p_b push_back #define fi first</pre>	

1 Contest

```
1 | #define se second
   #define pii std::pair<int, int>
   #define oo LLONG MAX
   #define big INT_MAX
   #define elif else if
   int input()
   {
       int x:
       cin>>x:
       return x;
   // -----> end of sashko123's defines (thank you Vasya <3)
   template<typename T>
   using graph = vector<vector<T>>;
   template<typename T>
   istream& operator>>(istream& in, vector<T>& a) {
       for (auto& i: a) {
          in >> i;
       return in;
   11 fast_pow(ll a, ll b, ll mod) {
       if (b == 0)
          return 1;
       if (b % 2) {
           return (111 * a * fast_pow(a, b - 1, mod)) % mod;
       11 k = fast_pow(a, b / 2, mod);
       return (111 * k * k) % mod;
   11 fast_pow(ll a, ll b) {
       if (b == 0)
          return 1:
       if (b % 2) {
           return (111 * a * fast pow(a, b - 1));
       11 k = fast pow(a, b / 2);
       return (111 * k * k);
   void solve() {
   int32_t main(int32_t argc, const char * argv[]) {
       cin.tie(0);
       cout.tie(0);
       ios_base::sync_with_stdio(0);
       // insert code here...
       int tt= 1;
       // std::cin >> tt;
       while (tt--) {
           solve();
       return 0;
   fast-input.h
   double readNumber() {
       const int BSIZE = 4096;
```

static char buffer[BSIZE];

```
static char* bptr = buffer + BSIZE;
    auto getChar = []() {
       if (bptr == buffer + BSIZE) {
           memset (buffer, 0, BSIZE);
            cin.read(buffer, BSIZE);
           bptr = buffer;
       return *bptr++;
   char c = getChar();
   while (c && (c < '0' || c > '9') && c != '-')
       c = getChar();
    bool minus = false;
   if (c == '-') minus = true, c = getChar();
    double res = 0;
   while (c >= '0' \&\& c <= '9') {
       res = res * 10 + c - '0';
       c = getChar();
     if (c == '.') {
       c = getChar();
       double cur = 0.1;
       while (c >= '0' \&\& c <= '9') {
           res = res + (c - '0') * cur;
           c = getChar();
           cur /= 10.0;
    return minus ? -res : res;
.bashrc
alias c='q++ -Wall -Wconversion -Wfatal-errors -q -std=c++20 \
 -fsanitize=undefined, address'
xmodmap -e 'clear lock' -e 'keycode 66=less greater' #caps = <>
.vimrc
                                                          6 lines
set cin aw ai is ts=4 sw=4 tm=50 nu noeb bg=dark ru cul
sy on | im jk <esc> | im kj <esc> | no;:
" Select region and then type : Hash to hash your selection.
" Useful for verifying that there aren't mistypes.
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space:]' \
\| md5sum \| cut -c-6
hash.sh
# Hashes a file, ignoring all whitespace and comments. Use for
# verifying that code was correctly typed.
cpp -dD -P -fpreprocessed | tr -d '[:space:]' | md5sum |cut -c-6
```

### Combinatorial (2)

### 2.1 Formulas

### 2.1.1 Hockey-stick identity

$$\sum_{i=k}^{n} C(i,k) = C(n+1,k+1)$$

### template fast-input .bashrc .vimrc hash

### 2.2 Permutations

### **2.2.1** Cycles

Let  $g_S(n)$  be the number of *n*-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

### 2.2.2 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

### 2.2.3 Burnside's lemma

Given a group G of symmetries and a set X, the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where  $X^g$  are the elements fixed by g(g.x = x).

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using  $G = \mathbb{Z}_n$  to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k).$$

### 2.3 Partitions and subsets

### 2.3.1 Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

### 2.3.2 Lucas' Theorem

Let n, m be non-negative integers and p a prime. Write  $n=n_kp^k+\ldots+n_1p+n_0$  and  $m=m_kp^k+\ldots+m_1p+m_0$ . Then  $\binom{n}{m}\equiv\prod_{i=0}^k\binom{n_i}{m_i}\pmod{p}$ .

### 2.4 General purpose numbers

### 2.1 General purpose i

2.4.1 Bernoulli numbers

EGF of Bernoulli numbers is  $B(t) = \frac{t}{e^t - 1}$  (FFT-able).  $B[0, \ldots] = [1, -\frac{1}{2}, \frac{1}{6}, 0, -\frac{1}{30}, 0, \frac{1}{42}, \ldots]$ 

Sums of powers:

$$\sum_{i=1}^{n} n^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} \cdot (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\sum_{i=m}^{\infty} f(i) = \int_{m}^{\infty} f(x)dx - \sum_{k=1}^{\infty} \frac{B_k}{k!} f^{(k-1)}(m)$$

$$\approx \int_{m}^{\infty} f(x)dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m))$$

### 2.4.2 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$
$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)\dots(x+n-1)$$

c(8, k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1 $c(n, 2) = 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots$ 

### 2.4.3 Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \geq j$ , k j:s s.t.  $\pi(j) > j$ .

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n, n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} \binom{n+1}{j} (k+1-j)^{n}$$

### 2.4.4 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^n$$

### 2.4.5 Bell numbers

Total number of partitions of n distinct elements. B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, . . . . For <math>p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

### 2.4.6 Labeled unrooted trees

# on n vertices:  $n^{n-2}$ 

# on k existing trees of size  $n_i$ :  $n_1 n_2 \cdots n_k n^{k-2}$ 

# with degrees  $d_i$ :  $(n-2)!/((d_1-1)!\cdots(d_n-1)!)$ 

### 2.4.7 Catalan numbers

$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2}C_n, \ C_{n+1} = \sum C_i C_{n-i}$$

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$ 

- sub-diagonal monotone paths in an  $n \times n$  grid.
- $\bullet$  strings with n pairs of parenthesis, correctly nested.
- binary trees with with n+1 leaves (0 or 2 children).
- ordered trees with n+1 vertices.
- ways a convex polygon with n + 2 sides can be cut into triangles by connecting vertices with straight lines.
- permutations of [n] with no 3-term increasing subseq.

# Mathematics (3)

### 3.1 Equations

$$ax^2 + bx + c = 0 \Rightarrow x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

The extremum is given by x = -b/2a.

$$ax + by = e$$

$$cx + dy = f \Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

In general, given an equation Ax = b, the solution to a variable  $x_i$  is given by

$$x_i = \frac{\det A_i'}{\det A}$$

where  $A'_i$  is A with the *i*'th column replaced by b.

### 3.2 Recurrences

If  $a_n = c_1 a_{n-1} + \cdots + c_k a_{n-k}$ , and  $r_1, \dots, r_k$  are distinct roots of  $x^k - c_1 x^{k-1} - \cdots - c_k$ , there are  $d_1, \dots, d_k$  s.t.

$$a_n = d_1 r_1^n + \dots + d_k r_k^n.$$

Non-distinct roots r become polynomial factors, e.g.  $a_n = (d_1n + d_2)r^n$ .

### Trigonometry

$$\sin(v + w) = \sin v \cos w + \cos v \sin w$$
$$\cos(v + w) = \cos v \cos w - \sin v \sin w$$

$$\tan(v+w) = \frac{\tan v + \tan w}{1 - \tan v \tan w}$$
$$\sin v + \sin w = 2\sin\frac{v+w}{2}\cos\frac{v-w}{2}$$
$$\cos v + \cos w = 2\cos\frac{v+w}{2}\cos\frac{v-w}{2}$$

$$(V+W)\tan(v-w)/2 = (V-W)\tan(v+w)/2$$

where V, W are lengths of sides opposite angles v, w.

$$a\cos x + b\sin x = r\cos(x - \phi)$$
$$a\sin x + b\cos x = r\sin(x + \phi)$$

where  $r = \sqrt{a^2 + b^2}$ ,  $\phi = \operatorname{atan2}(b, a)$ .

### 3.4 Geometry

### 3.4.1 Triangles

Side lengths: a, b, c

Semiperimeter: 
$$p = \frac{a+b+c}{2}$$

Area: 
$$A = \sqrt{p(p-a)(p-b)(p-c)}$$

Circumradius: 
$$R = \frac{abc}{4A}$$

Inradius: 
$$r = \frac{A}{p}$$

Length of median (divides triangle into two equal-area triangles):  $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$ 

Length of bisector (divides angles in two):

$$s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c}\right)^2\right]}$$

Law of sines: 
$$\frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$$
  
Law of cosines:  $a^2 = b^2 + c^2 - 2bc \cos \alpha$ 

Law of tangents: 
$$\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$$

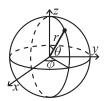
### 3.4.2 Quadrilaterals

With side lengths a, b, c, d, diagonals e, f, diagonals angle  $\theta$ , area A and magic flux  $F = b^2 + d^2 - a^2 - c^2$ :

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180°, ef = ac + bd, and  $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$ .

### 3.4.3 Spherical coordinates



$$\begin{array}{ll} x = r \sin \theta \cos \phi & r = \sqrt{x^2 + y^2 + z^2} \\ y = r \sin \theta \sin \phi & \theta = \arccos(z/\sqrt{x^2 + y^2 + z^2}) \\ z = r \cos \theta & \phi = \operatorname{atan2}(y, x) \end{array}$$

### Derivatives/Integrals

$$\frac{d}{dx}\arcsin x = \frac{1}{\sqrt{1-x^2}} \qquad \frac{d}{dx}\arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}\tan x = 1 + \tan^2 x \qquad \frac{d}{dx}\arctan x = \frac{1}{1+x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x\sin ax = \frac{\sin ax - ax\cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \qquad \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1)$$

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

#### Sums 3.6

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

### 3.7

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^{2}}{2} + \frac{x^{3}}{3} - \frac{x^{4}}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^{2}}{8} + \frac{2x^{3}}{32} - \frac{5x^{4}}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$$
$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

### 3.8 Probability theory

assuming the value x. It will then have an expected value (mean)  $\mu = \mathbb{E}(X) = \sum_{x} x p_X(x)$  and variance  $\sigma^2 = V(X) = \mathbb{E}(X^2) - (\mathbb{E}(X))^2 = \sum_x (x - \mathbb{E}(X))^2 p_X(x)$  where  $\sigma$ is the standard deviation. If X is instead continuous it will have a probability density function  $f_X(x)$  and the sums above will

instead be integrals with  $p_X(x)$  replaced by  $f_X(x)$ .

Let X be a discrete random variable with probability  $p_X(x)$  of

Expectation is linear:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

For independent X and Y,

$$V(aX + bY) = a^2V(X) + b^2V(Y).$$

### 3.8.1 Discrete distributions Binomial distribution

The number of successes in n independent yes/no experiments, each which yields success with probability p is  $Bin(n, p), n = 1, 2, ..., 0 \le p \le 1.$ 

$$p(k) = \binom{n}{k} p^k (1-p)^{n-k}$$

$$\mu = np, \, \sigma^2 = np(1-p)$$

Bin(n, p) is approximately Po(np) for small p.

### First success distribution

The number of trials needed to get the first success in independent yes/no experiments, each which yields success with probability p is Fs(p),  $0 \le p \le 1$ .

$$p(k) = p(1-p)^{k-1}, k = 1, 2, ...$$
  
 $\mu = \frac{1}{p}, \sigma^2 = \frac{1-p}{p^2}$ 

### Poisson distribution

The number of events occurring in a fixed period of time t if these events occur with a known average rate  $\kappa$  and independently of the time since the last event is  $Po(\lambda)$ ,  $\lambda = t\kappa$ .

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, k = 0, 1, 2, \dots$$
$$\mu = \lambda, \sigma^2 = \lambda$$

fft.h

### 3.8.2 Continuous distributions Uniform distribution

If the probability density function is constant between a and band 0 elsewhere it is U(a, b), a < b.

$$f(x) = \begin{cases} \frac{1}{b-a} & a < x < b \\ 0 & \text{otherwise} \end{cases}$$

$$\mu = \frac{a+b}{2}, \, \sigma^2 = \frac{(b-a)^2}{12}$$

### Exponential distribution

The time between events in a Poisson process is  $\text{Exp}(\lambda), \lambda > 0.$ 

$$f(x) = \begin{cases} \lambda e^{-\lambda x} & x \ge 0\\ 0 & x < 0 \end{cases}$$
$$\mu = \frac{1}{\lambda}, \, \sigma^2 = \frac{1}{\lambda^2}$$

### Normal distribution

Most real random values with mean  $\mu$  and variance  $\sigma^2$  are well described by  $\mathcal{N}(\mu, \sigma^2)$ ,  $\sigma > 0$ .

$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

If  $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$  and  $X_2 \sim \mathcal{N}(\mu_2, \sigma_2^2)$  then

$$aX_1 + bX_2 + c \sim \mathcal{N}(\mu_1 + \mu_2 + c, a^2\sigma_1^2 + b^2\sigma_2^2)$$

### 3.9 Markov chains

A Markov chain is a discrete random process with the property that the next state depends only on the current state. Let  $X_1, X_2, \ldots$  be a sequence of random variables generated by the Markov process. Then there is a transition matrix  $\mathbf{P} = (p_{ij})$ , with  $p_{ij} = \Pr(X_n = i | X_{n-1} = j)$ , and  $\mathbf{p}^{(n)} = \mathbf{P}^n \mathbf{p}^{(0)}$  is the probability distribution for  $X_n$  (i.e.,  $p_i^{(n)} = \Pr(X_n = i)$ ), where  $\mathbf{p}^{(0)}$  is the initial distribution.

 $\pi$  is a stationary distribution if  $\pi = \pi \mathbf{P}$ . If the Markov chain is irreducible (it is possible to get to any state from any state), then  $\pi_i = \frac{1}{\mathbb{E}(T_i)}$  where  $\mathbb{E}(T_i)$  is the expected time between two visits in state i.  $\pi_i/\pi_i$  is the expected number of visits in state j between two visits in state i.

For a connected, undirected and non-bipartite graph, where the transition probability is uniform among all neighbors,  $\pi_i$  is proportional to node i's degree.

A Markov chain is *ergodic* if the asymptotic distribution is independent of the initial distribution. A finite Markov chain is ergodic iff it is irreducible and aperiodic (i.e., the gcd of cycle lengths is 1).  $\lim_{k\to\infty} \mathbf{P}^k = \mathbf{1}\pi$ .

A Markov chain is an A-chain if the states can be partitioned into two sets A and G, such that all states in A are absorbing  $(p_{ii}=1)$ , and all states in **G** leads to an absorbing state in **A**. The probability for absorption in state  $i \in \mathbf{A}$ , when the initial state is j, is  $a_{ij} = p_{ij} + \sum_{k \in \mathbf{G}} a_{ik} p_{kj}$ . The expected time until absorption, when the initial state is i, is  $t_i = 1 + \sum_{k \in C} p_{ki} t_k$ .

### Algebra (4)

xor-basis.h

Description: Xor basis, all elements in the main set can be constructed using xor operation and elements in the basis

**Time:** insert per element -  $\mathcal{O}(log(A_{max}))$ 

```
hash46 lines
template<typename T = int, int max_bit = 31>
struct xor basis
 std::vector<T> basis; // basis[i] \rightarrow element with smallest
       set bit equal to i
 int sz; // Current size of the basis
 xor_basis() {
   basis.assign(max_bit);
 bool insert(T val) {
    for (int i = 0; i < max_bit; i++) {
      if (((val >> i)&1) == 0)
        continue;
      if (!basis[i]) {
        basis[i] = val;
        return true;
      val ^= basis[i];
    return false;
 bool contains(T val) {
    for (int i = 0; i < max_bit; i++) {</pre>
      if (((val >> i) &1) == 0)
       continue;
      if (!basis[i])
       return false;
      val ^= basis[i];
   return true;
 T max_element() { // not-sure
   T \text{ val} = 0;
    for (int i = max_bit - 1; i >= 0; i--) {
      if (basis[i] && !((val>>i)&1)) {
        val ^= basis[i];
    return val;
};
```

```
Description: FFT implementation
Time: \mathcal{O}((n+m)*\log(n+m))
<vector>, <complex>
                                                         hash56 lines
using cd = std::complex<double>;
void fft(std::vector<cd> & a, bool invert) {
    int n = a.size();
    for (int i = 1, j = 0; i < n; i++) {
        int bit = n >> 1;
        for (; j & bit; bit >>= 1)
            j ^= bit;
        j ^= bit;
        if (i < j)
            swap(a[i], a[j]);
    for (int len = 2; len <= n; len <<= 1) {
        double ang = 2 * PI / len * (invert ? -1 : 1);
        cd wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i += len) {
            cd w(1);
            for (int j = 0; j < len / 2; j++) {
                cd u = a[i+j], v = a[i+j+len/2] * w;
                a[i+j] = u + v;
                a[i+j+len/2] = u - v;
                 w \star= wlen;
    if (invert) {
        for (cd & x : a)
            x /= n;
template<typename T>
std::vector<T> multiply(const std::vector<T>& a, const std::
     vector<T>& b) {
    std::vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end
         ());
    int n = 1;
    while (n < a.size() + b.size())</pre>
        n <<= 1;
    fa.resize(n);
    fb.resize(n);
    fft(fa, false);
    fft(fb, false);
    for (int i = 0; i < n; i++)
        fa[i] \star = fb[i];
    fft(fa, true);
    std::vector<T> result(n);
    for (int i = 0; i < n; i++)
        result[i] = round(fa[i].real());
    result.resize(a.size() + b.size() - 1);
    return result:
nnt.h
Description: NNT implementation by modulo 998244353
Time: \mathcal{O}((n+m)*\log(n+m))
                                                         hash68 lines
```

const int root = 31; // primitive root of module

### floor-sum berlekamp-massey primitive-root

return A;

```
const int root_1 = fast_pow(root, MOD - 2, MOD); // (primitive
const int root pw = 1 << 23; // max power of 2 in MOD - 1
inline int inverse(int n, int mod) {
 return fast pow(n, mod - 2, mod);
void nnt(std::vector<int> &a, bool invert, int mod) {
    int n = a.size();
    for (int i = 1, j = 0; i < n; i++) {
       int bit = n >> 1;
        for (; j & bit; bit >>= 1)
          j ^= bit;
       j ^= bit;
       if (i < j)
            std::swap(a[i], a[j]);
   for (int len = 2; len <= n; len <<= 1) {
       int wlen = invert ? root 1 : root;
        for (int i = len; i < root_pw; i <<= 1)</pre>
           wlen = (int)(1LL * wlen * wlen % mod);
        for (int i = 0; i < n; i += len) {
           int w = 1;
            for (int j = 0; j < len / 2; j++) {
                int u = a[i+j], v = (int)(1LL * a[i+j+len/2] *
                a[i+j] = u + v < mod ? u + v : u + v - mod;
                a[i+j+len/2] = u - v >= 0 ? u - v : u - v + mod
                w = (111 * w * wlen % mod);
       }
   if (invert) {
       int n_1 = inverse(n, mod);
        for (int& x : a)
           x = (111 * x * n_1 % mod);
std::vector<int> multiply(const std::vector<int>& a, const std
    ::vector<int>& b, int mod) {
  int n = a.size(), m = b.size();
  int k = n + m - 1;
  int 1 = 1;
  while (1 < k)
   1 <<= 1;
  std::vector<int> A(1), B(1);
  for (int i = 0; i < n; i++) {
   A[i] = a[i];
  for (int i = 0; i < m; i++) {
   B[i] = b[i];
  nnt(A, false, mod);
  nnt(B, false, mod);
  for (int i = 0; i < B.size(); i++) {
   A[i] = (111 * A[i] * B[i]) % mod;
 nnt(A, true, mod);
  A.resize(k);
```

```
floor-sum.h
Description: sum_{i=0}^{n-1} floor((a * i + b) / m)
Time: log(n)
<utility>
                                                           hash24 lines
using ull = unsigned long long;
ull floor_sum_unsigned(ull n, ull m, ull a, ull b) {
    ull ans = 0;
    while (true) {
        if (a >= m) {
             ans += n * (n - 1) / 2 * (a / m);
        if (b >= m) {
             ans += n * (b / m);
            b %= m;
        ull y_max = a * n + b;
        if (v max < m) break;
        // y_{-}max < m * (n + 1)
        // floor(y_max / m) \le n
        n = (ull) (y_max / m);
        b = (ull) (y max % m);
        std::swap(m, a);
    return ans:
berlekamp-massev.h
Description: For given n first elements of sequence a, return array c, a[i] =
sum(j = 1 \dots -c-) a[i-j] * c[j])
Time: \mathcal{O}\left(n^2\right)
                                                           hash45 lines
template<typename T>
vector<T> berlekamp_massey(const vector<T> &s) {
    vector<T> c:
    vector<T> oldC:
    int f = -1;
    for (int i=0; i<(int)s.size(); i++) {</pre>
        T delta = s[i];
        for (int j=1; j<=(int)c.size(); j++)</pre>
            delta -= c[j-1] * s[i-j];
        if (delta == 0)
             continue;
        if (f == -1) {
            c.resize(i + 1);
            mt19937 rng(chrono::steady_clock::now().
                 time_since_epoch().count());
             for (T &x : c)
                x = rnq();
            f = i;
        } else {
            vector<T> d = oldC;
            for (T &x : d)
                x = -x;
            d.insert(d.begin(), 1);
```

T df1 = 0;

assert (df1 != 0);

 $x \star = coef;$ 

for (T &x : d)

T coef = delta / df1;

vector<T> zeros(i - f - 1);

for (int j=1; j<=(int)d.size(); j++)</pre>

df1 += d[j-1] \* s[f+1-j];

### $\underline{\text{Numeric}}$ (5)

```
primitive-root.h
```

```
Description: Primitive root of n
```

```
hash52 lines
```

```
constexpr long long safe mod(long long x, long long m) {
    if (x < 0) x += m;
    return x;
long long pow_mod(long long x, long long n, int m) {
    if (m == 1) return 0;
    unsigned int m = (unsigned int)(m);
    unsigned long long r = 1;
    unsigned long long y = safe_mod(x, m);
    while (n) {
        if (n \& 1) r = (r * y) % _m;
        y = (y * y) % _m;
        n >>= 1;
    return r;
int primitive root(int m) {
    if (m == 2) return 1;
    if (m == 167772161) return 3;
    if (m == 469762049) return 3;
    if (m == 754974721) return 11;
    if (m == 998244353) return 3;
    int divs[20] = \{\};
    divs[0] = 2;
    int cnt = 1;
    int x = (m - 1) / 2;
    while (x % 2 == 0) x /= 2;
    for (int i = 3; (long long) (i) *i <= x; i += 2) {
        if (x % i == 0) {
            divs[cnt++] = i;
            while (x \% i == 0) {
                x /= i;
    if (x > 1) {
        divs[cnt++] = x;
    for (int g = 2;; g++) {
        bool ok = true;
        for (int i = 0; i < cnt; i++) {
            if (pow_mod(g, (m - 1) / divs[i], m) == 1) {
                ok = false;
                break;
```

```
if (ok) return q;
pollard-rho.h
Description: Finds divider of n.
Time: \mathcal{O}\left(n^{1/4} * \log(n)\right)
long long mult(long long a, long long b, long long mod) {
    return ( int128)a * b % mod;
long long f(long long x, long long c, long long mod) {
    return (mult(x, x, mod) + c) % mod;
long long rho(long long n, long long x0=2, long long c=1) {
    long long x = x0;
    long long v = x0:
    long long g = 1;
    while (q == 1) {
        x = f(x, c, n);
        y = f(y, c, n);
        y = f(y, c, n);
        q = std::gcd(abs(x - y), n);
    return q;
miller-rabin.h
Description: checks whether given number (up to 1e18) is prime
```

Time:  $+-\mathcal{O}\left(\log^3(n)\right)$ hash43 lines using u128 = \_\_uint128\_t; 11 fast\_pow(ll a, ll b, ll mod) { if (b == 0)return 1; if (b % 2) { return ((u128) a \* fast\_pow(a, b - 1, mod)) % mod;  $11 k = fast_pow(a, b / 2, mod);$ return ((u128) k \* k) % mod; bool check\_composite(ll n, ll a, ll d, int s) {  $11 x = fast_pow(a, d, n);$ if (x == 1 | | x == n - 1)return false; for (int r = 1; r < s; r++) { x = (u128) x \* x % n;if (x == n - 1)return false; return true; bool miller\_rabin(ll n) { if (n < 2) return false; int r = 0: 11 d = n - 1;while ((d & 1) == 0) { d >>= 1;

r++;

```
for (int a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37})
   if (n == a)
       return true;
   if (check_composite(n, a, d, r))
       return false:
return true;
```

euclidean-algorithm.h

**Description:** Find x and y, s.t.  $x^*a + y^*b = gcd(a, b)$  // extended Euclidean algorithm

Time:  $\mathcal{O}(\log(\min(a,b)))$ 

```
<tuple>
                                                        hash12 lines
template<typename T>
std::pair<T, T> euclidean_algorithm(T a, T b) {
   T x = 1, y = 0;
   T \times 1 = 0, y1 = 1, a1 = a, b1 = b;
   while (b1) {
       T q = a1 / b1;
       std::tie(x, x1) = std::make\_tuple(x1, x - q * x1);
       std::tie(y, y1) = std::make\_tuple(y1, y - q * y1);
        std::tie(a1, b1) = std::make\_tuple(b1, a1 - q * b1);
   return {x, y};
```

### chinese-remainder-theorem.h

**Description:** Solves chinese remainder theorem.

**Time:**  $\mathcal{O}(n)$ , n - number of equations.

hash19 lines

```
struct Congruence {
    long long a, m;
long long chinese_remainder_theorem(vector<Congruence> const&
    congruences) {
    long long M = 1;
    for (auto const& congruence : congruences) {
        M *= congruence.m;
    long long solution = 0;
    for (auto const& congruence : congruences) {
        long long a_i = congruence.a;
        long long M_i = M / congruence.m;
       long long N_i = mod_inv(M_i, congruence.m);
        solution = (solution + a_i * M_i % M * N_i) % M;
    return solution;
mod-sart.h
Description: for given n finds all a, s.t. a*a = n \pmod{p}
```

Time:  $\mathcal{O}\left(log^2(p)\right)$ 

```
hash45 lines
unsigned xrand() {
 static unsigned x = 314159265, y = 358979323, z = 846264338,
      w = 327950288;
 unsigned t = x ^ x << 11; x = y; y = z; z = w; return w = w ^
       w >> 19 ^ t ^ t >> 8;
int jacobi(ll a, ll m) {
 int s = 1;
 if (a < 0) a = a % m + m;
```

```
for (; m > 1; ) {
   a %= m;
   if (a == 0) return 0;
   const int r = __builtin_ctz(a);
   if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
   a >>= r:
   if (a \& m \& 2) s = -s;
   std::swap(a, m);
 return s:
vector<ll> mod_sqrt(ll a, ll p) {
 if (p == 2) return {a & 1};
 const int j = jacobi(a, p);
 if (j == 0) return {0};
 if (j == -1) return {};
 11 b, d;
 for (; ; ) {
   b = xrand() % p;
   d = (b * b - a) % p;
   if (d < 0) d += p;
   if (jacobi(d, p) == -1) break;
 11 f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
 for (11 e = (p + 1) >> 1; e; e >>= 1) {
   if (e & 1) {
     tmp = (g0 * f0 + d * ((g1 * f1) % p)) % p;
     q1 = (q0 * f1 + q1 * f0) % p;
   tmp = (f0 * f0 + d * ((f1 * f1) % p)) % p;
   f1 = (2 * f0 * f1) % p;
   f0 = tmp;
 return (g0  ? vector<11>{g0, p - g0} : vector<11>{p
      - q0, q0};
```

### Data structures (6)

fenwick-tree.h

**Description:** Fenwick Tree, update(+=) at element, sum at segment. R is excluded.

```
Time: update - \mathcal{O}(\log N), get - \mathcal{O}(\log N)
```

```
hash32 lines
<cassert>, <vector>
template <class T> struct fenwick_tree {
 public:
    fenwick\_tree() : \_n(0) {}
    fenwick_tree(int n) : _n(n), data(n) {}
    void add(int p, T x) {
        assert(0 <= p && p < _n);
        p++;
        while (p \le \underline{n}) \{
            data[p - 1] += T(x);
            p += p \& -p;
    T sum(int 1, int r) {
        assert(0 <= 1 && 1 <= r && r <= _n);
        return sum(r) - sum(l);
 private:
    std::vector<T> data;
```

```
T sum(int r) {
       T s = 0;
       while (r > 0) {
           s += data[r - 1];
           r -= r & -r;
       }
        return s;
};
```

### fenwick-tree-2d.h

**Description:** 2d fenwick tree, update(+=) at element, sum at 2d segment. R is excluded.

**Time:** update -  $\mathcal{O}(\log N * \log M)$ , get -  $\mathcal{O}(\log N * \log M)$ 

```
hash53 lines
template <class T> struct fenwick_tree_2d{
    struct fenwick tree {
       int n:
       unordered_map<int, T> data;
        fenwick tree(): n(0) {};
        fenwick_tree(int n): n(n) {};
       void add(int p, T x) {
           assert (0 <= p && p < n);
           p++;
            while (p \le n) {
                data[p - 1] += T(x);
                p += p \& -p;
       T pref sum(int r){
           T s = 0;
            while (r > 0) {
               s += data[r - 1];
                r -= r & -r;
            return s;
    };
  int n, m;
  std::vector<fenwick tree> data;
  fenwick_tree_2d(int n, int m): n(n), m(m), data(n,
      fenwick tree(m)) {};
  void add(int x, int y, T val) {
       assert (0 <= x \&\& x < n);
   x++;
       while (x \le n) {
           data[x - 1].add(y, val);
           x += x & -x;
 T pref_sum(int xr, int yr){
       T s = 0;
       while (xr > 0) {
           s += data[xr - 1].pref_sum(yr);
           xr -= xr \& -xr;
        return s;
   T sum(int xl, int yl, int xr, int yr) {
        return pref_sum(xr, yr) - pref_sum(xr, yl) - pref_sum(
            xl, yr) + pref_sum(xl, yl);
```

```
};
segment-tree.h
Description: Segment tree, update(+=) at element, sum at segment. R is
Time: update - \mathcal{O}(\log N), get - \mathcal{O}(\log N)
                                                          hash38 lines
<array>
template<typename T, int N = (1 << 18)> struct segment_tree {
 std::array<T, 2 * N> tree;
 segment_tree() {
   tree.fill(T());
 void update(int pos, T val) {
    pos += N;
    tree[pos] += val;
    pos >>= 1;
    while (pos > 0) {
     tree[pos] = tree[pos << 1] + tree[(pos << 1) | 1];
     pos >>= 1;
 T get_sum(int 1, int r) {
   1 += N;
   r += N;
   T ans = T();
    while (1 < r) {
      if (1 & 1) {
        ans += tree[1++];
      if (r & 1) {
       ans += tree[--r];
     1 >>= 1;
     r >>= 1;
   return ans;
 T get(int pos) {
    return tree[N + pos];
};
lazy-segment-tree.h
Description: Segment tree, update(+=) at segment, sum at segment. R is
excluded.
Time: update - \mathcal{O}(\log N), get - \mathcal{O}(\log N)
                                                          hash82 lines
<array>, <vector>
template<typename T> struct lazy_segment_tree {
 struct node{
   T sum = 0;
   T promise = 0;
 };
 int n:
 std::vector<node> tree;
 lazy_segment_tree(int n_, const vector<T>& init): n(n_) {
   tree.assign(4 * n, node{});
   build(1, 0, n, init);
 void build(int v, int 1, int r, const vector<T>& init) {
   if (1 + 1 == r) {
```

```
tree[v].sum = init[l];
     return:
   if (1 >= r)
     return:
    int mid = (1 + r) / 2;
   build(2 * v, 1, mid, init);
   build(2 * v + 1, mid, r, init);
   tree[v].sum = (tree[2 * v].sum + tree[2 * v + 1].sum);
 void update(int 1, int r, T value) {
   update(1, 0, n, 1, r, value);
 T get(int l, int r) {
   return get (1, 0, n, 1, r);
 T get(int pos) {
   return get (1, 0, n, pos, pos + 1);
 void push(int v, int l, int r) {
   if (tree[v].promise == 0)
     return;
   tree[v].sum += tree[v].promise * (r - 1);
   if (1 + 1 < r) {
     tree[2 * v].promise += tree[v].promise;
     tree[2 * v + 1].promise += tree[v].promise;
   tree[v].promise = 0;
 void update(int v, int tl, int tr, int l, int r, T value) {
   push(v, tl, tr);
   if (1 >= tr || t1 >= r)
     return:
    if (1 <= t1 && tr <= r) {
     tree[v].promise += value;
     push(v, tl, tr);
     return;
   int mid = (tl + tr) / 2;
   update(2 * v, tl, mid, l, r, value);
   update (2 * v + 1, mid, tr, 1, r, value);
   tree[v].sum = tree[2 * v].sum + tree[2 * v + 1].sum;
 T get(int v, int tl, int tr, int l, int r) {
   push(v, tl, tr);
   if (1 >= tr || t1 >= r)
     return 0:
   if (1 <= t1 && tr <= r) {
     return tree[v].sum;
   int mid = (tl + tr) / 2;
   auto left = get(2 * v, tl, mid, l, r);
   auto right = get(2 * v + 1, mid, tr, 1, r);
    return left + right;
};
```

```
persistent-segment-tree.h
```

**Description:** Persistent segment tree, update(+=) at segment, sum at segment. R is excluded.

Time: update -  $\mathcal{O}(\log N)$ , get -  $\mathcal{O}(\log N)$ 

```
hash83 lines
template<typename T> struct persistent_segment_tree {
  struct node {
   T sum;
    int left, right;
   node(T val = T()): sum(val), left(-1), right(-1) {}
   node(int left_, int right_): sum(T()), left(left_), right(
        right_) {}
  std::vector<node> tree;
  node create_node(T val = 0) {
   return node(val);
  node create_node(int left, int right) {
   auto v = node(left, right);
   v.sum = (left != -1 ? tree[left].sum : 0) + (right != -1 ?
        tree[right].sum : 0);
    return v;
  int LAST = 0;
  template<typename... params>
  int new_node(params... args) {
   tree[++LAST] = create_node(args...);
    return LAST;
  persistent_segment_tree(int n_, int sz_, const vector<T>& a,
      int& first_root): n(n_) {
    tree.resize(sz_);
    first_root = build(0, n, a);
  int build(int 1, int r, const vector<T>& a) {
   if (1 + 1 == r) {
     return new_node(a[1]);
    int mid = (r + 1) / 2;
    int left = build(1, mid, a);
    int right = build(mid, r, a);
    return new_node(left, right);
  int update(int root, int pos, T val) {
   return update(root, 0, n, pos, val);
  T get(int root, int 1, int r) {
    return get (root, 0, n, 1, r);
 T get(int root, int pos) {
   return get(root, 0, n, pos, pos + 1);
  int update(int v, int l, int r, int pos, T val) {
   if (1 + 1 == r) {
     return new_node(val + tree[v].sum);
   int mid = (1 + r) / 2;
   if (pos < mid) {
     return new_node(
```

```
update(tree[v].left, 1, mid, pos, val),
       tree[v].right
     );
   } else {
     return new_node(
       tree[v].left,
       update(tree[v].right, mid, r, pos, val)
     );
 T get(int v, int tl, int tr, int l, int r) {
    if (tr <= 1 || r <= t1)
     return 0;
   if (1 <= t1 && tr <= r) {
     return tree[v].sum;
   int mid = (tl + tr) / 2;
   return get(tree[v].left, tl, mid, l, r) + get(tree[v].right
        , mid, tr, l, r);
};
```

### segment-tree-2d.h

**Description:** 2D segment tree, update(+=) at element, sum at segment. R is excluded.

```
Time: update - \mathcal{O}(\log N * \log M), get - \mathcal{O}(\log N * \log M)
                                                          hash76 lines
template<tvpename T>
struct segment_tree_2d {
    int n, m;
    vector<vector<T>> tree;
    segment tree 2d(const vector<vector<T>>& arr) {
        n = arr.size():
        m = arr[0].size();
        tree.assign(2 * n, vector<T > (2 * m));
        for (int i = 2 * n - 1; i > 0; i--) {
            for (int j = 2 * m - 1; j > 0; j--) {
                if (i < n) {
                     tree[i][j] = tree[2 * i][j] + tree[2 * i +
                         1][j];
                } else if (j < m) {</pre>
                     tree[i][j] = tree[i][2 * j] + tree[i][2 * j]
                           + 1];
                } else {
                     tree[i][j] = arr[i - n][j - m];
    void update_point(int x, int y, T newval) { // arr[x][y] :=
          nennal
        x += n;
        y += m;
        int curx = x;
        while (curx > 0) {
            int cury = y;
            while (cury > 0) {
                if (curx < n) {
                     tree[curx][cury] = tree[2 * curx][cury] +
                          tree[2 * curx + 1][cury];
                } else if (cury < m) {
                     tree[curx][cury] = tree[curx][2 * cury] +
                          tree[curx][2 * cury + 1];
                     tree[curx][cury] = newval;
```

```
cury >>= 1;
            curx >>= 1;
    T find_sum(int lx, int rx, int ly, int ry) { // [lx, rx] *
         [ly, ry)
        1x += n:
        rx += n;
       T ans = 0;
        while (lx < rx) {</pre>
            int curly = ly + m;
            int curry = ry + m;
            while (curly < curry) {
                if (curly & 1) {
                    if (lx & 1) {
                        ans += tree[lx][curly];
                    if (rx & 1) {
                        ans += tree[rx - 1][curly];
                if (curry & 1) {
                    if (lx & 1) {
                        ans += tree[lx][curry - 1];
                    if (rx & 1) {
                        ans += tree[rx - 1][curry - 1];
                curly = (curly + 1) >> 1;
                curry >>= 1;
            1x = (1x + 1) >> 1;
            rx >>= 1:
        return ans;
};
li-chao-tree.h
```

**Description:** Li-Chao tree, online convex hull for maximizing f(x) = k \* x+ b, for minimization use (-k) \* x + (-b)**Time:** add -  $\mathcal{O}(\log N)$ , get -  $\mathcal{O}(\log N)$ hash74 lines

```
template<typename T> struct li_chao_tree {
 const T MX = 1e9 + 1;
 struct line {
   T k = 0;
    T b = -INF;
    T f(T x) const {
      return k * x + b;
 };
 struct node {
   line ln:
    node* left = nullptr;
    node* right = nullptr;
  node* new_node() {
    const int N = 100000;
    static node* block;
    static int count = N;
```

```
if (count == N) {
      block= new node[N];
      count = 0;
    return (block + count++);
  node* root = new_node();
  T get(T x) {
    return get (root, -MX, MX, x);
    void add(line ln) {
        add(root, -MX, MX, ln);
  T get(node*& v, T 1, T r, T x) {
    if (!v || 1 > r) {
      return -INF:
    T ans = v->ln.f(x);
    if (r == 1) {
     return ans;
    T \text{ mid} = (r + 1) / 2;
    if (x <= mid) {</pre>
      return max(ans, get(v->left, 1, mid, x));
             return max(ans, get(v->right, mid + 1, r, x));
    void add(node*& v, T l, T r, line ln) {
    if (1 > r)
      return:
        if (!v) {
             v = new_node();
    T m = (r + 1) / 2;
        bool left = v \rightarrow ln.f(1) < ln.f(1);
    bool md = v \rightarrow ln.f(m) < ln.f(m);
        if (md)
             swap (v->ln, ln);
        if (1 == r) {
             return;
        if (left != md) {
            add(v->left, 1, m, ln);
             add(v->right, m + 1, r, ln);
};
```

**Description:** Rope data structure **Time:** all get queries  $\mathcal{O}(\log N)$ 

using namespace \_\_gnu\_cxx; rope<int> v; for (int i = 1;  $i \le n$ ; ++i) v.push\_back(i); rope<int> cur = v.substr(1, r - 1 + 1); // start, lengthv.erase(1, r - 1 + 1); // start, length v.insert(v.mutable\_begin(), cur); for(rope <int>::iterator it = v.mutable\_begin(); it != v. mutable end(); ++it)

```
cout << *it << " ";
ordered-set.h
Description: A red-black tree with the ability to get an element by index
(find_by_order) and index of a specific element (order_of_key)
Time: get - \mathcal{O}(\log N), segment tree is 2 times faster
<ext/pb_ds/assoc_container.hpp>, <ext/pb_ds/tree_policy.hpp>
                                                             hash4 lines
using namespace __gnu_pbds;
template<typename T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
sparse-table.h
Description: Min sparse table.
Time: build - \mathcal{O}(N * \log N), get - \mathcal{O}(1)
                                                            hash29 lines
<cassert>, <array>, <vector>
template<typename T>
struct sparse_table {
    static const int K = 20;
    std::array<std::vector<T>, K> ar;
    std::vector<int> lg;
    int n;
    sparse_table(int n, const vector<T>& a): n(n) {
        lg.resize(n + 1);
        lq[1] = 0;
        for (int i = 2; i <= n; i++)
             lq[i] = lq[i >> 1] + 1;
        ar[0] = a;
        for (int k = 0; k + 1 < K; k++) {
             ar[k + 1].resize(n);
             for (int i = 0; i + (1 << k) < n; i++) {
                 ar[k + 1][i] = min(ar[k][i], ar[k][i + (1 << k)
    T get(int 1, int r) {
        assert(0 <= 1 < n && 1 < r && r <= n);
        int power = lg[r - 1];
        return min(ar[power][1], ar[power][r - (1 << power)]);</pre>
};
```

hash11 lines

Description: Binary t for integer numbers. Get finds maximum xor of two numbers.

Time: add -  $\mathcal{O}(\log A)$ , get -  $\mathcal{O}(\log A)$ 

```
hash35 lines
<cassert>, <array>
struct xor_trie_node {
 int cnt = 0;
 std::array<xor_trie_node*, 2> mp = {nullptr, nullptr};
 void add(int mask, int k = 30) {
   cnt++;
   if (k == -1)
      return;
    int bit = (mask>>k) &1;
   if (!mp[bit])
      mp[bit] = new xor_trie_node();
    mp[bit] \rightarrow add(mask, k - 1);
```

```
void remove(int mask, int k = 30) {
    cnt--:
    if (k == -1)
      return;
    int bit = (mask>>k)&1;
    assert(mp[bit] && mp[bit]->cnt > 0);
    mp[bit]->remove(mask, k - 1);
  int get(int mask, int k = 30) {
    if(k == -1)
      return 0;
    int bit = (mask >> k) &1;
    int cur= bit;
    if (mp[!bit] && mp[!bit]->cnt)
      cur = !bit;
    return ((cur^bit) << k) | mp[cur]->get(mask, k - 1);
};
cartesian-tree.h
Description: Cartesian tree
Time: all get queries \mathcal{O}(\log N)
                                                        hash63 lines
struct Node {
    int val;
    int v;
    int 1. r:
    int cnt;
    Node(int val = -1) {
        this->val = val;
        v = rand();
        1 = -1;
        r = -1;
        cnt = 1;
array<Node, 2 * N> tree;
int getCnt(int v) {
    if (v == -1)
        return 0;
    return tree[v].cnt;
void upd(int v) {
    tree[v].cnt = getCnt(tree[v].1) + getCnt(tree[v].r) + 1;
pair<int, int> split(int v, int k) {
    if (v == -1)
        return {-1, -1};
    if (k == 0)
        return {-1, v};
    pair<int, int> res;
    if (getCnt(tree[v].1) >= k) {
        res = split(tree[v].1, k);
        tree[v].1 = res.second;
        res.second = v;
        res = split(tree[v].r, k - getCnt(tree[v].l) - 1);
        tree[v].r = res.first;
        res.first = v;
```

upd(v);

return res;

```
int merge(int u, int v) {
    if (u == -1)
        return v;
    if (v == -1)
       return u;
    int res = -1;
   if (tree[u].y > tree[v].y) {
        res = merge(tree[u].r, v);
        tree[u].r = res;
        res = u;
    } else {
        res = merge(u, tree[v].1);
       tree[v].l = res;
        res = v;
    upd(res);
    return res;
```

### link-cut-tree.h

Description: Represents a forest of unrooted trees. You can add and remove edges (as long as the result is still a forest), and check whether two nodes are in the same tree.

```
Time: All operations take amortized O (log N ).
                                                        hash90 lines
struct Node { // Splay tree. Root's pp contains tree's parent.
 Node *p = 0, *pp = 0, *c[2];
  bool flip = 0;
  Node() { c[0] = c[1] = 0; fix(); }
  void fix() {
    if (c[0]) c[0] -> p = this;
    if (c[1]) c[1] -> p = this;
    // (+ update sum of subtree elements etc. if wanted)
  void pushFlip() {
    if (!flip) return;
    flip = 0; swap(c[0], c[1]);
   if (c[0]) c[0]->flip ^= 1;
   if (c[1]) c[1]->flip ^= 1;
  int up() { return p ? p \rightarrow c[1] == this : -1; }
  void rot(int i, int b) {
   int h = i ^ b;
   Node *x = c[i], *y = b == 2 ? x : x -> c[h], *z = b ? y : x;
   if ((y->p = p)) p->c[up()] = y;
    c[i] = z->c[i ^ 1];
    if (b < 2) {
     x->c[h] = y->c[h ^ 1];
     y -> c[h ^1] = x;
    z->c[i ^1] = this;
    fix(); x->fix(); y->fix();
    if (p) p->fix();
    swap(pp, y->pp);
  void splay() {
    for (pushFlip(); p; ) {
     if (p->p) p->p->pushFlip();
     p->pushFlip(); pushFlip();
     int c1 = up(), c2 = p->up();
      if (c2 == -1) p -> rot(c1, 2);
      else p->p->rot(c2, c1 != c2);
  Node* first() {
   pushFlip();
    return c[0] ? c[0]->first() : (splay(), this);
```

```
};
struct LinkCut {
 vector<Node> node;
 LinkCut(int N) : node(N) {}
 void link(int u, int v) { // add an edge (u, v)
   assert(!connected(u, v));
   makeRoot (&node[u]);
   node[u].pp = &node[v];
 void cut (int u, int v) { // remove an edge (u, v)
   Node *x = &node[u], *top = &node[v];
   makeRoot(top); x->splay();
   assert(top == (x-pp ?: x-c[0]));
   if (x->pp) x->pp = 0;
   else {
     x->c[0] = top->p = 0;
     x->fix();
 bool connected(int u, int v) { // are u, v in the same tree?
   Node* nu = access(&node[u])->first();
   return nu == access(&node[v])->first();
 void makeRoot(Node* u) {
   access(u);
   u->splay();
   if(u->c[0]) {
     u - c[0] - p = 0;
     u - c[0] - flip ^= 1;
     u - c[0] - pp = u;
     u - > c[0] = 0;
      u->fix();
 Node* access(Node* u) {
   u->splay();
    while (Node* pp = u->pp) {
     pp \rightarrow splay(); u \rightarrow pp = 0;
     if (pp->c[1]) {
       pp->c[1]->p = 0; pp->c[1]->pp = pp; }
     pp -> c[1] = u; pp -> fix(); u = pp;
   return u:
```

## Graphs (7)

```
articulation-points.h
Description: Finds all articulation points.
Time: \mathcal{O}(N+M)
<set>, <vector>, <functional>
                                                         hash35 lines
template<typename T>
using graph = std::vector<std::vector<T>>;
std::set<int> find_articulation_points(int n, graph<int> g) {
    std::vector<int> used(n), tin(n), fup(n);
    int T = 0;
    std::set<int> nodes;
    std::function < void(int, int) > dfs = [&](int v, int p = -1)
        used[v] = true;
        tin[v] = fup[v] = T++;
        int cnt = 0;
        for (auto to : g[v]) {
            if (to == p)
                continue;
            if (used[to]) {
                fup[v] = std::min(fup[v], tin[to]);
            } else {
                dfs(to, v);
                 fup[v] = std::min(fup[v], fup[to]);
                if (fup[to] >= tin[v] \&\& p != -1) {
                    nodes.insert(v);
                cnt++;
        if (cnt > 1 \&\& p == -1)
            nodes.insert(v);
    for (int i = 0; i < n; i++) {
        if (!used[i])
            dfs(i, -1);
    return nodes;
Description: Finds all bridges in the undirected graph.
Time: \mathcal{O}(N+M)
<vector>, <functional>
                                                         hash30 lines
template<typename T>
using graph = std::vector<std::vector<T>>;
std::vector<std::pair<int, int>> find_bridges(int n, graph<int>
    std::vector<int> used(n), tin(n), fup(n);
    int T = 0;
    std::vector<std::pair<int, int>> edges;
    std::function<void(int, int)> dfs = [\&] (int v, int p = -1)
        used[v] = true;
        tin[v] = fup[v] = T++;
        for (auto to : g[v]) {
            if (to == p)
                continue;
```

if (used[to]) {

dfs(to, v);

} else {

fup[v] = std::min(fup[v], tin[to]);

fup[v] = std::min(fup[v], fup[to]);

edges.push\_back({v, to});

if (fup[to] == tin[to]) {

```
};
    for (int i = 0; i < n; i++) {
        if (!used[i])
            dfs(i, -1);
    return edges;
Description: Disjoint Set Union
Time: \mathcal{O}(n) time (amortized) by sizes
<algorithm>, <cassert>, <vector>
                                                        hash59 lines
struct dsu {
 public:
    dsu() : _n(0) {}
    dsu(int n) : _n(n), parent_or_size(n, -1) {}
    int merge(int a, int b) {
        assert(0 <= a && a < n);
        assert(0 <= b && b < _n);
        int x = leader(a), y = leader(b);
        if (x == y) return x;
        if (-parent_or_size[x] < -parent_or_size[y]) std::swap(</pre>
        parent_or_size[x] += parent_or_size[y];
        parent_or_size[y] = x;
        return x;
   bool same(int a, int b) {
       assert(0 <= a && a < n);
        assert (0 \le b \& b \le n);
        return leader(a) == leader(b);
   int leader(int a) {
        assert(0 <= a && a < _n);
        if (parent_or_size[a] < 0) return a;</pre>
        return parent_or_size[a] = leader(parent_or_size[a]);
    int size(int a) {
        assert(0 <= a && a < _n);
        return -parent_or_size[leader(a)];
    std::vector<std::vector<int>> groups() {
        std::vector<int> leader_buf(_n), group_size(_n);
        for (int i = 0; i < _n; i++) {
            leader_buf[i] = leader(i);
            group_size[leader_buf[i]]++;
        std::vector<std::vector<int>> result(_n);
        for (int i = 0; i < _n; i++) {
            result[i].reserve(group_size[i]);
        for (int i = 0; i < _n; i++) {
            result[leader_buf[i]].push_back(i);
        result.erase(
            std::remove_if(result.begin(), result.end(),
                            [&] (const std::vector<int>& v) {
                                 return v.empty(); }),
            result.end());
        return result;
```

```
private:
   int _n;
    // root node: -1 * component size
    // otherwise: parent
    std::vector<int> parent_or_size;
dynamic-connectivity-problem.h
Description: Disjoint Set Union with roolbacks
Time: \mathcal{O}(q * log(q) * log(n)), q - number of queries, n - number of vertices
<array>, <vector>
struct Query {
    int v, u;
    bool united:
    Query(int _v, int _u) : v(_v), u(_u) {}
};
struct DSU
 struct DSU_save
    int v, rang_v;
    int u, rang u;
  std::vector<int> pred, rang;
  std::vector<DSU_save> saves;
  DSU(int n_): n(n_) {
    pred.resize(n);
    rang.resize(n);
    for (int i = 0; i < n; i++) {
      pred[i] = i;
      rang[i] = 0;
 int get(int v) {
    if (pred[v] == v) {
      return v;
    return get(pred[v]);
  bool merge(int u, int v) {
    u = get(u);
    v = get(v);
    if (u == v)
      return false;
    if (rang[u] < rang[v])</pre>
      std::swap(u, v);
    saves.push_back({v, rang[v], u, rang[u]});
    pred[v] = u;
    if (rang[u] == rang[v]) {
      rang[u]++;
    return true;
  void rollback() {
    if (saves.empty())
      return:
    auto [v, rang_v, u, rang_u] = saves.back();
    rang[v] = rang_v;
    rang[u] = rang_u;
    pred[v] = v;
    pred[u] = u;
```

```
saves.pop_back();
};
struct dynamic_connectivity_problem {
  std::vector<std::vector<Query>> tree;
 DSU dsu:
  dynamic_connectivity_problem(int q_, int n_): q(q_), n(n_),
      dsu(DSU(n_)) {
   tree.resize(4 * q);
  void add(Query a, int 1, int r) {
   add(1, 0, q, 1, r, a);
  void add(int v, int tl, int tr, int l, int r, const Query& a)
    if (1 <= t1 && tr <= r) {
     tree[v].push_back(a);
     return;
    if (1 >= tr || t1 >= r)
     return;
    int mid = (tr + t1) / 2;
   add(2 * v, t1, mid, 1, r, a);
   add(2 * v + 1, mid, tr, 1, r, a);
  void dfs(int v, int l, int r) {
   if (1 >= r)
     return;
    for (auto& q: tree[v]) {
     q.united = dsu.merge(q.u, q.v);
    if (1 + 1 == r) {
     int x = dsu.qet(0);
     // do something
    } else {
     int mid = (r + 1) / 2;
     dfs(2 * v, 1, mid);
     dfs(2 * v + 1, mid, r);
    for (auto& q: tree[v]) {
     if (q.united)
       dsu.rollback();
     q.united = false;
};
```

### kuhn-matching.h

**Description:** Fast pair matching algorithm. To find the minimum vertex cover start dfs from each vertice in the left that is not in maximum matching, from the left side choice unvisited vertices and from right chose visited.

#### Time: $\mathcal{O}(n*(n+m))$

```
used[v] = 1;
    for (auto to : q[v]) {
       if (mt[to] == -1) {
           rev_mt[v] = to;
           mt[to] = v;
           return true;
    for (auto to : g[v]) {
       if (dfs(mt[to], g, mt, rev_mt, used)) {
           rev_mt[v] = to;
           mt[to] = v;
           return true;
   return false;
pair<int, vector<int>> pair_matching(int n, int m, graph<int> q
    vector<int> mt(m, -1), used, rev_mt(n, -1);
   int cnt = 0;
    for (int it = 0; ; it++) {
       bool found = false;
       used.assign(n, 0);
       for (int i = 0; i < n; i++) {
            if (rev_mt[i] == -1 && dfs(i, g, mt, rev_mt, used))
                cnt++;
                found = true;
        if (!found) {
           break;
    return {cnt, mt};
```

#### lca h

 $\bf Description:$  Finds lowest common ancestor of two vertices using sparce table

```
Time: \mathcal{O}\left(n * log(n)\right) - build, \mathcal{O}\left(1\right) - get
```

hash30 lir

```
struct LCA {
 vector<pair<int, int>> traversal;
 vector<int> pos;
 graph<int> g;
 int n:
 sparse_table<pair<int, int>> st;
 LCA(int n_, qraph<int> q_, int root = 0): n(n_), q(q_) {
   pos.resize(n);
   dfs(root, -1, 0);
   st = sparse_table<pair<int, int>>(traversal.size(),
        traversal);
 void dfs(int v, int pred, int depth) {
   pos[v] = traversal.size();
   traversal.push_back({depth, v});
   for (auto to : q[v]) {
     if (to != pred) {
       dfs(to, v, depth + 1);
       traversal.push_back({depth, v});
```

```
int get(int a, int b) {
    if (a == b)
      return a;
    return st.get(min(pos[a], pos[b]), max(pos[a], pos[b]) + 1)
};
two-sat.h
Description: 2-sat implementation
Time: \mathcal{O}(n)
<vector>
                                                        hash74 lines
template<typename T>
using graph = std::vector<std::vector<T>>;
struct two_sat {
    graph<int> g, rev;
    std::vector<int> used, order, comp, ans;
    two sat(int n): n(n) {
        g.assign(2 * n, {});
        rev.assign(2 * n, \{\});
    void add edge(int u, int v) {
        q[u].push_back(v);
        rev[v].push_back(u);
    void add clause or (int a, bool val a, int b, bool val b) {
        add_edge(a + val_a * n, b + !val_b * n);
        add_edge(b + val_b * n, a + !val_a * n);
    void add_clause_xor(int a, bool val_a, int b, bool val_b) {
        add clause or (a, val a, b, val b);
        add_clause_or(a, !val_a, b, !val_b);
    void add_clause_and(int a, bool val_a, int b, bool val_b) {
        add_clause_xor(a, !val_a, b, val_b);
    void top_sort(int v) {
        used[v] = 1;
        for (auto to : g[v]) {
            if (!used[to])
                top_sort(to);
        order.push_back(v);
    void compress(int v, int id) {
        comp[v] = id;
        for (auto to : rev[v]) {
            if (comp[to] == -1)
                compress(to, id);
    bool satisfiable() {
        order.clear();
        used.assign(2 * n, 0);
        comp.assign(2 * n, -1);
        ans.assign(n, 0);
        for (int i = 0; i < 2 * n; i++) {
            if (!used[i])
```

```
top sort(i);
        reverse(order.begin(), order.end());
        int id = 0:
        for (auto v : order) {
            if (comp[v] == -1)
                 compress(v, id++);
        for (int i = 0; i < n; i++) {
             if (comp[i] == comp[i + n])
                 return false;
            ans[i] = (comp[i + n] < comp[i]);
        }
        return true;
};
eulerian-path.h
Description: Finds eulerian path and cycle for undirected graph
Time: \mathcal{O}\left(m * \log(m)\right)
<vector>, <set>
                                                         hash101 lines
using namespace std;
struct eulerian_path
    vector<multiset<int>>graph;
    eulerian_path(vector<vector<int>>g)
        graph.resize(g.size());
        for(int i=0;i<q.size();i++)</pre>
             for(int j:q[i])
                 graph[i].insert(j);
    void dfs(int u, vector<int>&cvcle)
        auto p = graph[u];
        for(int j:p)
             if(graph[u].find(j) != graph[u].end())
                 graph[u].erase(graph[u].find(j));
                graph[j].erase(graph[j].find(u));
                 dfs(j, cycle);
        cycle.push_back(u);
    vector<int> find_cycle(int v = 0)
        for(int i = 0;i < graph.size();i++)</pre>
            if(graph[i].size() % 2)
                return {};
        vector<int>cycle;
        dfs(v, cycle);
        for(auto x:graph)
             if(x.size())
            return {};
        return cycle;
```

```
vector<int> find path()
        int st = -1, fi = -1;
        int mx = 0:
        for(int i = 0;i < graph.size();i++)</pre>
             if(graph[i].size() % 2)
                 if(st == -1)
                     st = i;
                 else
                 if(fi == -1)
                      fi = i;
                 else
                     return {};
             if (graph[mx].size() < graph[i].size()) mx = i;</pre>
        if(fi == -1)
        {
             auto cycle = find_cycle(mx);
             return cycle;
        graph[st].insert(fi);
        graph[fi].insert(st);
        auto cycle = find_cycle(st);
        if(!cycle.size())
        return {};
        cycle.pop_back();
        if(cycle[0] == st and cycle.back() == fi or cycle[0] == fi
              and cycle.back() == st)
             return cycle;
        vector<int>path;
        for(int i=0;;i++)
             if(cvcle[i] == st and cvcle[i+1] == fi or cvcle[i]
                  == fi and cycle[i+1] == st)
                 for(int j = i + 1; j < cycle.size(); j++)
                      path.push_back(cycle[j]);
                 for (int j = 0; j \le i; j++)
                      path.push_back(cycle[j]);
                 break;
        return path;
};
max-flow.h
Description: Maximum flow problem
Time: \mathcal{O}\left(\min(n^{2/3}*m,m^{3/2})\right) - all capacities are 1, \mathcal{O}\left(\min(n^2,m)\right) -
template<typename Cap>
struct mf graph {
    struct mf_edge {
        int from, to;
        Cap cap, flow;
        int back_id;
        int _id;
```

```
};
  int n:
  vector<vector<mf edge>> q;
bool need_clear = false;
  mf_graph(int n): n(n) {
      g.assign(n, {});
  void add_edge(int from, int to, Cap cap, int id = -1) {
      int id1 = g[from].size();
      int id2 = g[to].size();
      g[from].push_back(mf_edge {
          .from = from,
          .to = to,
          .cap = cap,
          .flow = Cap(),
          .back_id = id2,
          .\_id = id
     });
      g[to].push_back(mf_edge {
          .from = to,
          .to = from,
          .cap = Cap(),
          .flow = Cap(),
          .back_id = id1,
          ._{id} = -1
      });
  }
void clear() {
  for (int i = 0; i < n; i++) {
    for (mf_edge& e: g[i])
      e.flow = Cap();
  Cap flow(int s, int t, Cap limit) {
  if (need_clear)
    clear();
      vector<int> dist(n, n + 1);
      auto bfs = [\&] (int s, int t) {
          dist.assign(n, n + 1);
          dist[s] = 0;
          deque<int> q;
          q.push_back(s);
          while (!q.empty())
              int v = q.front();
              q.pop_front();
              for (mf_edge& e : g[v]) {
                  if (e.flow < e.cap && dist[e.to] == n + 1)
                      dist[e.to] = dist[v] + 1;
                      q.push_back(e.to);
          return dist[t] != n + 1;
      };
      vector<int> lst(n);
      function<Cap(int, int, Cap)> dfs = [&](int v, int
          target, Cap F) {
          if (v == target)
```

### min-cut min-cost-flow ford-bellman spfa

```
return F;
            Cap pushed = Cap();
            for (; lst[v] < g[v].size(); lst[v]++) {</pre>
                mf_edge& e = g[v][lst[v]];
                if (dist[e.to] == dist[v] + 1 && e.flow < e.cap</pre>
                    Cap x = dfs(e.to, target, min(F, e.cap - e.
                         flow));
                    if (x) {
                        e.flow += x;
                        g[e.to][e.back_id].flow -= x;
                        pushed += x;
                        F -= x;
                }
            return pushed;
        };
        Cap flow = 0;
        while (bfs(s, t)) {
            lst.assign(n, 0);
            while (Cap f = dfs(s, t, limit)) {
                flow += f;
    need_clear = true;
        return flow;
};
```

### min-cut.h

**Description:** Start dfs from source and using edges with capacity - flow > 0, mark visited vertices. Min cut is edges between marked and unmarked vertices.

Time:  $\mathcal{O}(n+m)$ 

#### min-cost-flow.h

 $\bf Description:$  Solves minimum-cost flow problem for specific flow value (or infinity).

**Time:**  $\mathcal{O}(F * (n+m) * log(n+m))$ , where F is the amount of the flow and m is the number of added edges.

hash106 lines template<typename Cap, typename Cost> struct mcf\_graph struct mcf\_edge { int from, to: Cap cap, flow; Cost cost; int back\_id; int \_id; int n; vector<vector<mcf\_edge>> g; mcf graph(int n): n(n) { q.assiqn(n, {}); void add\_edge(int from, int to, Cap cap, Cost cost, int \_id = -1) { int id1 = q[from].size(); int id2 = g[to].size(); g[from].push\_back(mcf\_edge{ .from = from,.to = to,

```
.cap = cap,
        .flow = Cap(),
        .cost = cost,
        .back_id = id2,
        ._id = _id
    });
    g[to].push_back(mcf_edge{
        .from = to,
        .to = from,
        .cap = Cap(),
        .flow = Cap(),
        .cost = -cost,
        .back_id = id1,
        ._{id} = -1
    });
}
pair<Cap, Cost> flow(int s, int t, Cap target_flow) {
    Cap flow = Cap();
    vector<pair<int, int>> pred(n);
    vector<Cost> dist(n), dual(n, 0);
    auto shortests_path = [&](int s, int t) {
        pred.assign(n, \{-1, -1\});
        dist.assign(n, numeric_limits<Cost>::max());
        dist[s] = Cost();
        set<pair<int, int>> q;
        q.insert({dist[s], s});
        while (!q.empty()) {
            int v = q.begin()->second;
            q.erase(q.begin());
            for (int i = 0; i < g[v].size(); i++) {
                mcf_edge& e = g[v][i];
                if (e.flow >= e.cap)
                    continue;
                Cost cost = e.cost - dual[e.to] + dual[v];
                if (dist[e.to] > dist[v] + cost) {
                    q.erase({dist[e.to], e.to});
                    dist[e.to] = dist[v] + cost;
                    pred[e.to] = \{v, i\};
                    g.insert({dist[e.to], e.to});
        if (dist[t] == numeric_limits<Cost>::max())
            return false;
        for (int v = 0; v < n; v++) {
            if (dist[v] == numeric_limits<Cost>::max())
                continue;
            dual[v] -= dist[t] - dist[v];
        return true;
   Cost total cost = {};
    while (flow < target flow) {
        if (!shortests_path(s, t))
            break;
        Cap f = target_flow - flow;
        int cur = t;
        while (cur != s) {
            auto [p, id] = pred[cur];
            mcf_edge& e = g[p][id];
            f = min(f, e.cap - e.flow);
            cur = p;
```

```
cur = t;
while (cur != s) {
    auto [p, id] = pred[cur];
    mcf_edge& e = g[p][id];
    e.flow += f;
    g[e.to][e.back_id].flow -= f;
    cur = p;
}
Cost d = -dual[s];
flow += f;
total_cost += f * d;
}
return {flow, total_cost};
}
```

### ford-bellman.h

 $\bf Description:$  Single source shortest path with negative weight edges in directed graph

Time:  $\mathcal{O}(n*m)$ 

hash16 line

```
vector<int> d(n, INF);
d[v] = 0;
vector<int> p(n, -1);

for (;;) {
   bool any = false;
   for (Edge e : edges)
    if (d[e.a] < INF)
        if (d[e.b] > d[e.a] + e.cost) {
        d[e.b] = d[e.a] + e.cost;
        p[e.b] = e.a;
        any = true;
    }
   if (!any)
        break;
}
```

#### spfa.h

**Description:** Single source shortest path with negative weight edges in directed graph

Time:  $\mathcal{O}\left(n*m\right)$ 

hash36 lines

```
const int INF = 10000000000;
vector<vector<pair<int, int>>> adj;
bool spfa(int s, vector<int>& d) {
    int n = adj.size();
    d.assign(n, INF);
    vector<int> cnt(n, 0);
    vector<bool> inqueue(n, false);
    queue<int> q;
    d[s] = 0;
    q.push(s);
    inqueue[s] = true;
    while (!q.empty()) {
        int v = q.front();
        q.pop();
        inqueue[v] = false;
        for (auto edge : adj[v]) {
            int to = edge.first;
            int len = edge.second;
            if (d[v] + len < d[to]) {</pre>
                d[to] = d[v] + len;
                if (!inqueue[to]) {
                    q.push(to);
```

### directed-mst z-function prefix-function suffix-array

```
inqueue[to] = true;
               cnt[to]++;
               if (cnt[to] > n)
                   return false; // negative cycle
   }
return true;
```

### directed-mst.h

Description: finds directed mst from given root

```
Time: \mathcal{O}(n * \log(m))
                                                        hash82 lines
struct RollbackUF {
  vector<int> e; vector<pair<int, int>> st;
  RollbackUF(int n) : e(n, -1) {}
  int size(int x) { return -e[find(x)]; }
  int find(int x) { return e[x] < 0 ? x : find(e[x]); }
  int time() { return (int)st.size(); }
  void rollback(int t) {
    for (int i = time(); i --> t;)
     e[st[i].first] = st[i].second;
    st.resize(t);
  bool join(int a, int b) {
   a = find(a), b = find(b);
   if (a == b) return false;
   if (e[a] > e[b]) swap(a, b);
   st.push_back({a, e[a]});
    st.push_back({b, e[b]});
   e[a] += e[b]; e[b] = a;
    return true;
};
struct Edge { int a, b; ll w; };
struct Node {
  Edge key;
 Node *1, *r;
  ll delta;
  void prop() {
   key.w += delta;
   if (1) 1->delta += delta;
   if (r) r->delta += delta;
    delta = 0;
  Edge top() { prop(); return key; }
Node *merge(Node *a, Node *b) {
  if (!a || !b) return a ?: b;
  a->prop(), b->prop();
  if (a->key.w > b->key.w) swap(a, b);
  swap(a->1, (a->r = merge(b, a->r)));
  return a;
void pop(Node*\& a) { a->prop(); a = merge(a->1, a->r); }
pair<11, vector<int>> dmst(int n, int r, vector<Edge>& g) {
  RollbackUF uf(n);
  vector<Node*> heap(n);
  for (Edge e : q) heap[e.b] = merge(heap[e.b], new Node{e});
  11 \text{ res} = 0;
  vector<int> seen(n, -1), path(n), par(n);
  seen[r] = r;
  vector<Edge> Q(n), in(n, \{-1,-1\}), comp;
  deque<tuple<int, int, vector<Edge>>> cycs;
  for (int s = 0; s < n; s + +) {
```

```
int u = s, qi = 0, w;
  while (seen[u] < 0) {
    if (!heap[u]) return {-1,{}};
    Edge e = heap[u]->top();
    heap[u]->delta -= e.w, pop(heap[u]);
    Q[qi] = e, path[qi++] = u, seen[u] = s;
    res += e.w, u = uf.find(e.a);
    if (seen[u] == s) {
     Node \star cyc = 0;
      int end = qi, time = uf.time();
      do cyc = merge(cyc, heap[w = path[--qi]]);
      while (uf.join(u, w));
     u = uf.find(u), heap[u] = cyc, seen[u] = -1;
      cycs.push_front({u, time, {&Q[qi], &Q[end]}});
  for (int i = 0; i < qi; i++) in [uf.find(Q[i].b)] = Q[i];
for (auto& [u,t,comp] : cycs) { // restore sol (optional)
  uf.rollback(t);
  Edge inEdge = in[u];
  for (auto& e : comp) in[uf.find(e.b)] = e;
  in[uf.find(inEdge.b)] = inEdge;
for(int i=0;i<n;i++) par[i] = in[i].a;</pre>
return {res, par};
```

### Strings (8)

### z-function.h

**Description:** Z-functions, z[i] equal to the length of largest common prefix of string s and suffix of s starting at i.

**Time:**  $\mathcal{O}(N)$ , N - size of string s

hash18 lines

```
vector<int> z_function(const string& s) {
 int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
 for (int i = 1; i < n; i++) {
   if(i < r) {
     z[i] = min(r - i, z[i - 1]);
   while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]]) {
     z[i]++;
   if(i + z[i] > r) {
     1 = i;
     r = i + z[i];
   }
 return z;
```

### prefix-function.h

**Description:** The prefix function for this string is defined as an array p of length n, where p[i] is the length of the longest valid prefix of the substring s[0...i], which is also a suffix of this substring.

**Time:**  $\mathcal{O}(N)$ , N - size of string s

```
vector<int> prefix_function(string s) {
   int n = (int)s.length();
   vector<int> pi(n);
   for (int i = 1; i < n; i++) {
       int j = pi[i-1];
       while (j > 0 \&\& s[i] != s[j])
            j = pi[j-1];
```

```
if (s[i] == s[j])
        j++;
    pi[i] = j;
return pi;
```

#### suffix-array.h

**Description:** Suffix array will contain integers that represent the starting indexes of the all the suffixes of a given string, after the aforementioned suffixes are sorted.

```
Time: \mathcal{O}(N * log(N)), N - size of string s
 vector<int> suffix_arrays(string s) {
     s = s + "$";
     int n = s.size();
     std::vector<int> p(n);
     vector<vector<int>> c(20, vector<int>(n));
     int alphabet = 256;
     auto set_classes = [&](int k) {
          int classes = 0;
           c[k][p[0]] = classes++;
           for (int i = 1; i < n; i++) {
                 auto cur = pair\{c[k - 1][p[i]], c[k - 1][(p[i] + (1 << (k - 1))][(p[i] + (k << (k - 1))][(p[i] + (k << (k - 1))][(p[i] + (k << (k 
                              -1))) % n]};
                 auto prev = pair\{c[k-1][p[i-1]], c[k-1][(p[i-1]+
                                  (1 << (k-1)) % n]};
                 if (cur == prev) {
                      c[k][p[i]] = c[k][p[i-1]];
                      c[k][p[i]] = classes++;
     };
     vector<int> cnt(alphabet);
     for (int i = 0; i < n; i++) {
           cnt[s[i]]++;
     for (int i = 1; i < alphabet; i++) {</pre>
           cnt[i] += cnt[i - 1];
     for (int i = n - 1; i >= 0; i--) {
           p[cnt[s[i]] - 1] = i;
           cnt[s[i]]--;
     int classes = 0;
     c[0][p[0]] = classes++;
     for (int i = 1; i < n; i++) {
          if (s[p[i]] == s[p[i-1]]) {
                 c[0][p[i]] = c[0][p[i - 1]];
          } else {
                 c[0][p[i]] = classes++;
     for (int k = 0; (1<<k) < n; k++) {
           vector<int> pn(n), cnt(n);
           for (int i = 0; i < n; i++) {
                 pn[i] = (p[i] - (1 << k) + n) % n;
                 cnt[c[k][pn[i]]]++;
           for (int i = 1; i < n; i++)
                 cnt[i] += cnt[i - 1];
```

### lcp-array hash trie aho-corasick

```
for (int i = n - 1; i >= 0; i--) {
   p[cnt[c[k][pn[i]]] - 1] = pn[i];
   cnt[c[k][pn[i]]]--;
  set classes(k + 1);
p.erase(p.begin());
return p;
```

### lcp-array.h

Description: Largest common prefix of substrings. Given a string s of length n, it returns the LCP array of s. Here, the LCP array of s is the array of length n-1, such that the i-th element is the length of the LCP (Longest Common Prefix) of s[sa[i]..n) and s[sa[i+1]..n).

Time:  $\mathcal{O}(N)$ , N - size of string s

```
template <class T>
std::vector<int> lcp_array(const std::vector<T>& s,
                          const std::vector<int>& sa) {
   int n = int(s.size());
   assert (n >= 1);
   std::vector<int> rnk(n);
   for (int i = 0; i < n; i++) {
       rnk[sa[i]] = i;
   std::vector<int> lcp(n - 1);
   int h = 0;
   for (int i = 0; i < n; i++) {
       if (h > 0) h--;
       if (rnk[i] == 0) continue;
       int j = sa[rnk[i] - 1];
       for (; j + h < n && i + h < n; h++) {
           if (s[j+h] != s[i+h]) break;
       lcp[rnk[i] - 1] = h;
   return lcp;
std::vector<int> lcp array(const std::string& s, const std::
    vector<int>& sa) {
   int n = int(s.size());
   std::vector<int> s2(n);
   for (int i = 0; i < n; i++) {
       s2[i] = s[i];
   return lcp_array(s2, sa);
```

### hash.h

Description: Polynomial hashes for strings

**Time:**  $\mathcal{O}(n * log(m))$ , n - size of string s, m - module

hash47 lines

```
template<int P, int MOD>
struct hash_st {
    int n;
    vector<int> hash_, rev_, p, rev_p;
   hash_st(string s) {
       n = s.size();
       hash_.resize(n);
        rev .resize(n);
        p.resize(n);
        rev_p.resize(n);
        p[0] = 1;
        rev_p[0] = 1;
```

```
for (int i = 1; i < n; i++) {
           p[i] = (p[i - 1] * P) % MOD;
           rev_p[i] = fast_pow(p[i], MOD - 2, MOD);
       int last = 0;
       for (int i = 0; i < n; i++) {
           hash_[i] = (last + p[i] * (s[i] - 'a' + 1)) % MOD;
           last = hash_[i];
       last = 0:
       for (int i = n - 1; i >= 0; i--) {
           rev_[i] = (last + p[n - 1 - i] * (s[i] - 'a' + 1))
                % MOD;
           last = rev_[i];
   int get(int 1, int r) {
       r--;
       if (1 == 0)
           return hash_[r];
       int x = (MOD + hash_[r] - hash_[1 - 1]) % MOD;
       return (x * rev_p[1]) % MOD;
   int get_rev(int 1, int r) {
       r--;
       if (r == n - 1)
           return rev_[1];
       int x = (MOD + rev_[1] - rev_[r + 1]) % MOD;
       int st = n - 1 - r;
       return (x * rev_p[st]) % MOD;
};
```

### trie.h

Description: Trie implementation

Time:  $\mathcal{O}(n)$ 

```
hash63 lines
template<int N = (int) 1e6 + 1>
struct trie {
 struct node {
   int cnt = 0;
   array<int, 27> links;
   node() {
     links.fill(-1);
     cnt = 0;
 };
 array<node, N> tree;
 int sz = 0;
 int root = 0;
 trie() {
   root = sz++;
 void add(string s) {
   int cur = root;
   tree[cur].cnt++;
   for (int i = 0; i < s.size(); i++) {</pre>
     int nxt = tree[cur].links[s[i] - 'a'];
     if (nxt == -1) {
       nxt = sz++;
       tree[cur].links[s[i] - 'a'] = nxt;
```

```
tree[nxt].cnt++;
     cur = nxt;
 void remove(string s) {
   int cur = root;
    tree[cur].cnt--;
    for (int i = 0; i < s.size(); i++) {
     int nxt = tree[cur].links[s[i] - 'a'];
     assert (nxt !=-1);
     tree[nxt].cnt--;
      cur = nxt;
 int get_cnt_of_str(const string& s) {
   int cur = root;
    for (int i = 0; i < s.size(); i++) {
     int nxt = tree[cur].links[s[i] - 'a'];
     if (nxt == -1) {
       return 0;
     cur = nxt;
   return tree[cur].cnt;
 void clear() {
   for (int i = 0; i < sz; i++)
     tree[i] = node();
};
```

### aho-corasick.h

Description: Creates suffix automaton

```
Time: \mathcal{O}(|alphabet| * \sum |s_i|)
                                                         hash103 lines
const int K = 26;
struct Vertex {
    int next[K];
    bool output = false;
    int p = -1;
    char pch;
    int link = -1;
    int ans_link = -1;
    int go[K];
    Vertex(int p=-1, char ch='\$') : p(p), pch(ch) {
        fill(begin(next), end(next), -1);
        fill (begin (go), end (go), -1);
};
vector<Vertex> t(1);
void add_string(string const& s) {
    int v = 0;
    for (char ch : s) {
        int c = ch - 'a';
        if (t[v].next[c] == -1) {
             t[v].next[c] = t.size();
            t.emplace_back(v, ch);
        v = t[v].next[c];
    t[v].output = true;
```

```
int go(int v, char ch);
int get_link(int v) {
    if (t[v].link == -1) {
       if (v == 0 || t[v].p == 0)
           t[v].link = 0;
        else
            t[v].link = go(get_link(t[v].p), t[v].pch);
    return t[v].link;
int go(int v, char ch) {
    int c = ch - 'a';
    if (t[v].go[c] == -1) {
        if (t[v].next[c] != -1)
            t[v].qo[c] = t[v].next[c];
            t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
    return t[v].go[c];
int get_ans_link(int v)
    if(v == 0)
    return 0;
    if(t[v].ans_link == -1)
        if(t[get_link(v)].output)
            t[v].ans_link = get_link(v);
        t[v].ans_link = get_ans_link(get_link(v));
    return t[v].ans_link;
// returns all strings from dictionary that is suffix of given
vector<string> get_ans(int v)
    vector<string>ans;
    if(t[v].output)
        string cur = "";
       int v1 = v;
        while(v1)
            cur += t[v1].pch;
            v1 = t[v1].p;
        reverse(cur.begin(), cur.end());
        ans.push_back(cur);
    v = get_ans_link(v);
    while(v)
        string cur = "";
        int v1 = v;
        while(v1)
            cur += t[v1].pch;
            v1 = t[v1].p;
        reverse(cur.begin(), cur.end());
        ans.push_back(cur);
```

```
v = get ans link(v);
return ans;
```

```
Geometry (9)
point.h
Description: Geometry formuls
                                                       hash124 lines
<complex>, <iostream>
class point : public std::complex<long double> {
 using ld = long double;
 static constexpr long double PI = 3.141592653589793;
public:
 point() : std::complex<long double>() {}
 point(ld x, ld y) : std::complex<long double>(x, y) {}
 point(std::complex<long double> obj) : std::complex<long</pre>
      double > (obi) {}
 ld x() {
   return this->real();
 }
 ld v() {
   return this->imag();
 }
 ld x() const {
   return this->real();
 ld v() const {
   return this->imag();
  // a_{-}x * b_{-}x + a_{-}y * b_{-}y
 static ld dot_product(const point& a, const point& b) {
    return (conj(a) * b).real();
  // a_{-}x * b_{-}y - a_{-}y * b_{-}x
 static ld cross_product(const point& a, const point& b) {
   return (conj(a) * b).imag();
 static ld squared_distance(const point& a, const point& b) {
   return norm(a - b);
 static ld distance (const point & a, const point & b) {
   return abs(a - b);
 // angle\_of\_elevation of line (a, b) to oX
 static ld angle_of_elevation(const point& a, const point& b)
    return arg(b - a);
 // k \text{ from } y = k * x + b
 static ld slope_of_line(const point& a, const point& b) {
   return tan(arg(b - a));
```

```
static point from_polar(ld r, ld theta) {
  return std::polar(r, theta);
static point rotate_above_pivot(const point& a, const ld
     theta, const point& pivot = point(0, 0)){
  return (a - pivot) * std::polar<ld>(1.0, theta) + pivot;
point& rotate(const ld theta, const point& pivot = point(0,
  *this = point::rotate_above_pivot(*this, theta, pivot);
  return *this;
// angle of ABC
static ld angle (const point& a, const point& b, const point&
  return abs(remainder(arg(a-b) - arg(c-b), 2.0 * PI));
static point project_on_vector(const point& p, const point& v
  return v * dot_product(p, v) / norm(v);
static point project_on_line(const point& p, const point& a,
     const point & b) {
  return a + (b - a) * dot_product(p - a, b - a) / norm(b - a
static point reflect_accros(const point& p, const point& a,
     const point& b) {
  return a + conj((p - a) / (b - a)) * (b - a);
// intersection of lines (a, b) and (p, q). if parallel
     returns {false, ...} else {true, intersection}.
friend std::pair<bool, point> intersection_of_lines(const
    point& a, const point& b, const point& p, const point& q
  1d c1 = cross\_product(p - a, b - a), c2 = cross\_product(q -
       a, b - a);
  if (c1 == c2) {
    return {false, {}};
  return {true, (c1 * q - c2 * p) / (c1 - c2)}; // undefined
       if parallel
// returns a, b, c from a * x + b * y + c = 0 by two points
friend std::tuple<ld, ld, ld> get_line(const point& p, const
    point& q) {
  1d a = (p.y() - q.y());
  1d b = -(p.x() - q.x());
  ld c = p.y() * (p.x() - q.x()) - p.x() * (p.y() - q.y());
  return {a, b, c};
friend ld distance_from_point_to_line(const point& p, const
     point& a, const point& b) {
  point q = project_on_line(p, a, b);
  return point::distance(p, q);
friend ld distance_from_point_to_segment(const point& p,
     const point& a, const point& b) {
```

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```
point q = project_on_line(p, a, b);
    if (std::min(a.x(), b.x()) \le q.x() \&\& q.x() \le std::max(a.
        x(), b.x())
      return point::distance(p, q);
      return std::min(distance(p, a), distance(p, b));
  friend std::istream& operator>> (std::istream& in, point& p)
    ld x, y;
   in >> x >> y;
    p = point(x, y);
    return in;
};
faces-of-planar-graph.h
```

**Description:** Finds faces of planar graph. Rreturns a vector of vertices for each face, outer face goes first. Inner faces are returned in counter-clockwise orders and the outer face is returned in clockwise order.

```
Time: \mathcal{O}(n * \log(n))
<vector>
template<typename T>
using graph = vector<vector<T>>;
struct Point {
    int64 t x, v;
    Point(int64_t x_, int64_t y_): x(x_), y(y_) {}
    Point operator - (const Point & p) const {
        return Point(x - p.x, y - p.y);
    int64_t cross (const Point & p) const {
        return x * p.y - y * p.x;
    int64_t cross (const Point & p, const Point & q) const {
        return (p - *this).cross(q - *this);
    int half () const {
        return int (y < 0 \mid | (y == 0 \&\& x < 0));
};
std::vector<std::vector<size_t>> find_faces(std::vector<Point>
     vertices, graph<int> adj) {
    size_t n = vertices.size();
    std::vector<std::vector<char>> used(n);
    for (size_t i = 0; i < n; i++) {</pre>
        used[i].resize(adj[i].size());
        used[i].assign(adj[i].size(), 0);
        auto compare = [&](size_t l, size_t r) {
            Point pl = vertices[1] - vertices[i];
            Point pr = vertices[r] - vertices[i];
            if (pl.half() != pr.half())
                 return pl.half() < pr.half();</pre>
            return pl.cross(pr) > 0;
        std::sort(adj[i].begin(), adj[i].end(), compare);
    std::vector<std::vector<size_t>> faces;
    for (size_t i = 0; i < n; i++) {
        for (size_t edge_id = 0; edge_id < adj[i].size();</pre>
             edge_id++) {
            if (used[i][edge_id]) {
```

```
continue;
            std::vector<size t> face;
            size t v = i;
            size_t e = edge_id;
            while (!used[v][e]) {
                used[v][e] = true;
                face.push_back(v);
                size_t u = adj[v][e];
                size_t e1 = std::lower_bound(adj[u].begin(),
                     adj[u].end(), v, [&](size_t l, size_t r) {
                     Point pl = vertices[1] - vertices[u];
                     Point pr = vertices[r] - vertices[u];
                    if (pl.half() != pr.half())
                         return pl.half() < pr.half();</pre>
                     return pl.cross(pr) > 0;
                }) - adj[u].begin() + 1;
                if (e1 == adj[u].size()) {
                     e1 = 0;
                v = u;
                e = e1;
             std::reverse(face.begin(), face.end());
            int sign = 0;
            for (size_t j = 0; j < face.size(); j++) {</pre>
                size_t j1 = (j + 1) % face.size();
                size_t j2 = (j + 2) % face.size();
                int64_t val = vertices[face[j]].cross(vertices[
                      face[j1]], vertices[face[j2]]);
                if (val > 0) {
                     sign = 1;
                     break;
                } else if (val < 0) {</pre>
                     sign = -1;
                     break;
             if (sign <= 0) {
                 faces.insert(faces.begin(), face);
                 faces.emplace back(face);
    return faces;
minimum-enclosing-circle.h
Description: Minimum enclosing circle
Time: \mathcal{O}(n)
<algorithm>, <assert.h>, <iostream>, <math.h>, <vector>, <iomanip> hash89 lines
using namespace std;
const double INF = 1e18;
struct Point {
 double X, Y;
struct Circle {
 Point C;
 double R;
double dist(const Point& a, const Point& b)
  return sqrt (pow(a.X - b.X, 2)
        + pow(a.Y - b.Y, 2));
bool is_inside(const Circle& c, const Point& p)
```

return dist(c.C, p) <= c.R;

```
Point get_circle_center(double bx, double by, double cx, double
 double B = bx * bx + by * by;
 double C = cx * cx + cy * cy;
 double D = bx * cy - by * cx;
 return { (cy * B - by * C) / (2 * D),
      (bx * C - cx * B) / (2 * D) };
Circle circle_from(const Point& A, const Point& B, const Point&
 Point I = get_circle_center(B.X - A.X, B.Y - A.Y, C.X - A.X,
      C.Y - A.Y);
 I.X += A.X;
 I.Y += A.Y;
 return { I, dist(I, A) };
Circle circle_from(const Point& A, const Point& B)
 Point C = \{ (A.X + B.X) / 2.0, (A.Y + B.Y) / 2.0 \};
 return { C, dist(A, B) / 2.0 };
bool is_valid_circle(const Circle& c, const vector<Point>& P)
 for (const Point& p : P)
   if (!is_inside(c, p))
     return false;
 return true;
Circle min circle trivial (vector < Point > & P)
  assert (P.size() <= 3);
 if (P.empty()) {
    return { { 0, 0 }, 0 };
 else if (P.size() == 1) {
    return { P[0], 0 };
 else if (P.size() == 2) {
    return circle_from(P[0], P[1]);
 for (int i = 0; i < 3; i++) {
    for (int j = i + 1; j < 3; j++) {
      Circle c = circle_from(P[i], P[j]);
      if (is valid circle(c, P))
        return c;
 return circle_from(P[0], P[1], P[2]);
Circle welzl_helper(vector<Point>& P, vector<Point> R, int n)
 if (n == 0 || R.size() == 3) {
   return min_circle_trivial(R);
 int idx = rand() % n;
 Point p = P[idx];
  swap(P[idx], P[n - 1]);
  Circle d = welzl_helper(P, R, n - 1);
  if (is_inside(d, p)) {
    return d;
 R.push_back(p);
 return welzl_helper(P, R, n - 1);
```

UzhNU Machata primes sos-dp

```
Circle welzl(const vector<Point>& P)
{
  vector<Point> P_copy = P;
  random_shuffle(P_copy.begin(), P_copy.end());
  return welzl_helper(P_copy, {}, P_copy.size());
}
```

# Misc. algorithms (10)

```
primes.txt
```

3 lines

### sos-dp.h

**Description:** Sum over submasks

Time:  $\mathcal{O}\left(n*2^n\right)$ 

hash7 lines

```
for (int i = 0; i < n; i++) {
  for (int mask = 0; mask < (1<<n); mask++) {
    if (mask&(1<<i)) {
        dp[mask] += dp[mask^(1<<i)];
    }
  }
}</pre>
```

primes sos-dp

# Techniques (A)

### techniques.txt

Bitonic cycle

161 lines

Recursion Divide and conquer Finding interesting points in N log N Algorithm analysis Master theorem Amortized time complexity Greedy algorithm Scheduling Max contiquous subvector sum Invariants Huffman encoding Graph theory Dynamic graphs (extra book-keeping) Breadth first search Depth first search \* Normal trees / DFS trees Dijkstra's algorithm MST: Prim's algorithm Bellman-Ford Konig's theorem and vertex cover Min-cost max flow Lovasz toggle Matrix tree theorem Maximal matching, general graphs Hopcroft-Karp Hall's marriage theorem Graphical sequences Floyd-Warshall Euler cycles Flow networks \* Augmenting paths \* Edmonds-Karp Bipartite matching Min. path cover Topological sorting Strongly connected components Cut vertices, cut-edges and biconnected components Edge coloring \* Trees Vertex coloring \* Bipartite graphs (=> trees) \* 3^n (special case of set cover) Diameter and centroid K'th shortest path Shortest cycle Dynamic programming Knapsack Coin change Longest common subsequence Longest increasing subsequence Number of paths in a dag Shortest path in a dag Dynprog over intervals Dynprog over subsets Dynprog over probabilities Dynprog over trees 3^n set cover Divide and conquer Knuth optimization Convex hull optimizations Slope trick Aliens trick RMQ (sparse table a.k.a 2^k-jumps)

Log partitioning (loop over most restricted) Combinatorics Computation of binomial coefficients Pigeon-hole principle Inclusion/exclusion Catalan number Pick's theorem Number theory Integer parts Divisibility Euclidean algorithm Modular arithmetic \* Modular multiplication \* Modular inverses \* Modular exponentiation by squaring Chinese remainder theorem Fermat's little theorem Euler's theorem Phi function Frobenius number Quadratic reciprocity Pollard-Rho Miller-Rabin Hensel lifting Vieta root jumping Game theory Combinatorial games Game trees Mini-max Nim Games on graphs Games on graphs with loops Grundy numbers Bipartite games without repetition General games without repetition Alpha-beta pruning Probability theory Optimization Binary search Ternary search Unimodality and convex functions Binary search on derivative Numerical methods Numeric integration Newton's method Root-finding with binary/ternary search Golden section search Matrices Gaussian elimination Exponentiation by squaring Sorting Radix sort Geometry Coordinates and vectors \* Cross product \* Scalar product Convex hull Polygon cut Closest pair Coordinate-compression Quadtrees KD-trees All segment-segment intersection Discretization (convert to events and sweep) Angle sweeping Line sweeping Discrete second derivatives Strings

Longest common substring Palindrome subsequences Knuth-Morris-Pratt Tries Rolling polynomial hashes Suffix array Suffix tree Aho-Corasick Manacher's algorithm Letter position lists Combinatorial search Meet in the middle Brute-force with pruning Best-first (A\*) Bidirectional search Iterative deepening DFS / A\* Data structures LCA (2^k-jumps in trees in general) Pull/push-technique on trees Heavy-light decomposition Centroid decomposition Lazy propagation Self-balancing trees Convex hull trick (wcipeg.com/wiki/Convex\_hull\_trick) Monotone queues / monotone stacks / sliding queues Sliding queue using 2 stacks Persistent segment tree

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