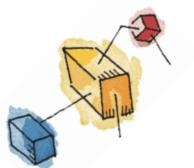
Readers/Writers Problem

- A data area is shared among many processes
 - Some processes only read the data area, some only write to the area
- Conditions to satisfy:
 - Multiple readers may read the file at once.
 - 2. Only one writer at a time may write
 - 3. If a writer is writing to the file, no reader may read it.



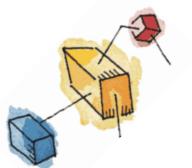




Readers have Priority

- The writer process is simple. The semaphore wsem is used to enforce mutual exclusion.
- As long as one writer is accessing the shared data area, no other writers and no readers may access it.
- The reader process also makes use of wsem to enforce mutual exclusion.
- However, to allow multiple readers, we require that, when there are no readers reading, the first reader that attempts to read should wait on wsem.



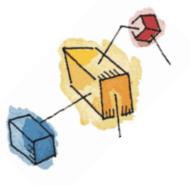


Readers have Priority

- When there is already at least one reader reading, subsequent readers need not wait before entering.
- The global variable readcount is used to keep track of the number of readers.
- The semaphore x is used to assure that readcount is updated properly.





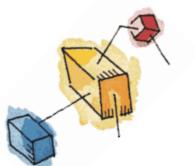


Readers have Priority

```
/* program readersandwriters */
int readcount;
semaphore x = 1, wsem = 1;
void reader()
   while (true) {
     semWait (x);
     readcount++;
     if (readcount == 1) semWait (wsem);
     semSignal (x);
     READUNIT();
     semWait (x);
     readcount --;
     if (readcount == 0) semSignal (wsem);
     semSignal (x);
void writer()
   while (true) {
     semWait (wsem);
     WRITEUNIT();
     semSignal (wsem);
void main()
   readcount = 0;
   parbegin (reader, writer);
```

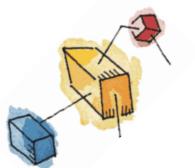






- This solution guarantees that no new readers are allowed access to the data area once at least one writer has declared a desire to write.
- For writers, the following semaphores and variables are added to the ones already defined:
 - A semaphore rsem that inhibits all readers while there is at least one writer desiring access to the data area
 - A variable writecount that controls the setting of rsem
 - A semaphore y that controls the updating of writecount





- For readers, one additional semaphore z is needed.
- A long queue must not be allowed to build up on rsem;
 otherwise writers will not be able to jump the queue.
- Therefore, only one reader is allowed to queue on rsem, with any additional readers queuing on semaphore z



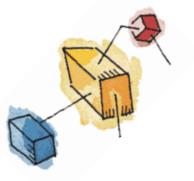


```
/* program readersandwriters */
int readcount, writecount;
semaphore x = 1, y = 1, z = 1, wsem = 1, rsem = 1;
void reader()
   while (true) {
     semWait (z);
          semWait (rsem);
               semWait (x);
                    readcount++;
                    if (readcount == 1) semWait (wsem);
               semSignal (x);
          semSignal (rsem);
     semSignal (z);
     READUNIT();
     semWait (x);
          readcount--:
          if (readcount == 0) semSignal (wsem);
     semSignal (x);
```

```
void writer ()
   while (true) {
     semWait (y);
          writecount++;
          if (writecount == 1) semWait (rsem);
     semSignal (y);
     semWait (wsem);
     WRITEUNIT();
     semSignal (wsem);
     semWait (y);
          writecount--;
          if (writecount == 0) semSignal (rsem);
     semSignal (y);
void main()
   readcount = writecount = 0;
   parbegin (reader, writer);
```







Readers only in the system	 wsem set no queues
Writers only in the system	wsem and rsem setwriters queue on wsem
Both readers and writers with read first	 wsem set by reader rsem set by writer all writers queue on wsem one reader queues on rsem other readers queue on z
Both readers and writers with write first	 wsem set by writer rsem set by writer writers queue on wsem one reader queues on rsem other readers queue on z

