

L.J INSTITUTE OF ENGINEERING AND TECHNOLOGY, Ahmedabad
THEORY OF COMPUTATION(TOC)
SEM-IV CE/IT DEPARTMENT

| subj ect_ code | unit _nu mbe r | chapt er_n umbe r | gro up_ id | Keywo rds | question_text | Type | ans wer_ tex t | m ar ks | difficult y_level | pr ev io us_ y ea r | opti on1 | opti on2 | optio n3 | opti on4 |
|----------------------|-------------------------|----------------------------|------------------|------------------------------|---|------------|---|---------------|----------------------|---------------------------------------|-------------------|---------------------------------|---------------------|---|
| 0170 1349 2 | 1 | 1 | 1 | Basics of Automa ta | Give True or False: If $A \subseteq B$ and $B \subseteq C$ then $A \subseteq C$ | True/False | True | 1 | Easy | | | | | |
| 0170 1349 2 | 1 | 1 | 2 | Basics of Automa ta | Give True or False: If $x \in A$ and $A \subseteq B$ then $x \in B$ | True/False | True | 1 | Easy | | | | | |
| 0170 1349 2 | 1 | 1 | 3 | Basics of Automa ta | If $n(A) = 110$, $n(B) = 300$, $n(A - B) = 50$ then find $n(A \cup B)$ | MCQ | 350 | 1 | Easy | | 300 | 360 | 110 | 350 |
| 0170 1349 2 | 1 | 1 | 4 | Basics of Automa ta | RR* can be expressed in which of the forms: | MCQ | a) R+ | 1 | Easy | | a) R+ | b) R- | c) R+ U R- | d) R |
| 0170 1349 2 | 1 | 1 | 5 | Basics of Automa ta | If $\Sigma = \{0,1\}$, then Φ^* will result to: | MCQ | a) ϵ | 1 | Easy | | a) ϵ | b) Φ | c) Σ | d) Non e of the ment ion |
| 0170 1349 2 | 1 | 1 | 6 | Basics of Automa ta | Which among the following is not an associative operation? | MCQ | d) Non e of the ment ione d | 1 | Easy | | a) Uni on | b) Con cate nati on | c) Dot | d) Non e of the ment ione d |
| 0170 1349 2 | 1 | 1 | 7 | Regular Express ion | Which of the following statements about Regular Expression Is/are | MCQ | c) The itera | 1 | Medium | | a) The unio | b) The conc | c) The iterat | d) All |

| | | | | | | | | | | | | | | |
|-------------------|---|---|----|--------------------|---|-----|---|---|--------|--|---|---|--|--------------------------|
| | | | | | incorrect? | | tion of a regular expression is also a regular expression | | | | n of two regular expressions is also a regular expression | aten atio n of two regular expressions is also a regular expression | ion of a regular expression is also a regular expression | |
| 0170 1349 2 | 1 | 1 | 8 | Basics of Automata | Dot operator in regular expression resembles which of the following? | MCQ | a) Expressions are juxtaposed | 1 | Medium | | a) Expressions are juxtaposed | b) Expressions are multiplied | c) Cross operation | d) None of the mentioned |
| 0170 1349 2 | 1 | 1 | 9 | Basics of Automata | Which of the following regular expressions represents the set of strings which do not contain a substring 'rt' if $\Sigma = \{r, t\}$ | MCQ | d) (t^*r^*) | 1 | Hard | | a) $(rt)^*$ | b) $(tr)^*$ | c) (r^*t^*) | d) (t^*r^*) |
| 0170 1349 2 | 1 | 1 | 10 | Basics of Automata | If $\Sigma = \{0,1\}$, then Φ^* will result to: | MCQ | a) ϵ | 1 | Medium | | a) ϵ | b) Φ | c) Σ | d) None of the mentioned |

| | | | | | | | | | | | | | | |
|-------------------|---|---|----|------------------------------|---|-------------|----------------------------------|---|--------|--|--|-----------------------------------|----------------------------------|----------------------------|
| 0170 1349 2 | 1 | 1 | 11 | Basics of Automa ta | Finite automata requires minimum _____ number of stacks. | Short Que. | 0 | 1 | Easy | | | | | |
| 0170 1349 2 | 1 | 1 | 12 | Basics of Automa ta | FSM with output capability can be used to add two given integer in binary representation. | True /False | True | 1 | Easy | | | | | |
| 0170 1349 2 | 1 | 1 | 13 | Basics of Automa ta | Differentiate between (a.b) and (a+b)? | Short Que. | | 1 | Easy | | | | | |
| 0170 1349 2 | 1 | 1 | 14 | Basics of Automa ta | Define: Transition Diagram | Short Que. | | 1 | Easy | | | | | |
| 0170 1349 2 | 1 | 1 | 15 | Basics of Automa ta | What are the applications of automata theory? | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 1 | 1 | 16 | Basics of Automa ta | What is a Regular language ? | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 1 | 1 | 17 | Basics of Automa ta | Differentiate L* and L+ | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 1 | 1 | 18 | Basics of Automa ta | What is the closure property of regular sets? | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 1 | 1 | 19 | Basics of Automa ta | The appropriate precedence order of operations over a Regular Language is | MCQ | Klee ne, Dot, Unio n | 1 | Easy | | Klee ne, Uni on, Con cate nate | Klee ne, Star, Uni on | Klee ne, Dot, Unio n | star, Unio n, Dot |

| | | | | | | | | | | | | | | |
|-------------------|---|---|----|------------------------------|---|------------|--|---|--------|--|--|---|---|----------------------|
| 0170 1349 2 | 2 | 2 | 20 | Regular Express ion | RR* can be expressed in which of the forms: | MCQ | a) R+ | 1 | Medium | | a) R+ | b) R- | c) R+ U R- | d) R |
| 0170 1349 2 | 2 | 2 | 21 | Basics of Automa ta | What Automata Theory ? | Short Ques | | 1 | Easy | | | | | |
| 0170 1349 2 | 2 | 2 | 22 | Basics of Automa ta | What is the difference between the strings and the words of a language? | Short Que. | | 1 | Easy | | | | | |
| 0170 1349 2 | 2 | 2 | 23 | Regular Express ion | Define Regular Expression? | Short Que. | | 1 | Easy | | | | | |
| 0170 1349 2 | 2 | 2 | 24 | Basics of Automa ta | Define FA | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 25 | Regular Express ion | Regular expression for all strings starts with ab and ends with bba is? | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 26 | Basics of Automa ta | The basic limitation of finite automata is that. | MCQ | It can't remember arbitrary large amount of information. | 1 | Medium | | It can't remember arbitrary large amount of information. | It sometimes recognize grammar that are not regular |) It sometimes fails to recognize regular grammar | All of the mentioned |
| 0170 1349 2 | 2 | 2 | 27 | Regular Express ion | Which of the following is not a regular expression? | MCQ | $[(0+1)-(0b+al)^*]$ | 1 | Medium | | $[(a+b)^*-(aa+bb)]$ | $[(0+1)-(0b+al)^*]$ | $(01+11+10)^*$ | $(1+2+0)^* (1+2)^*$ |

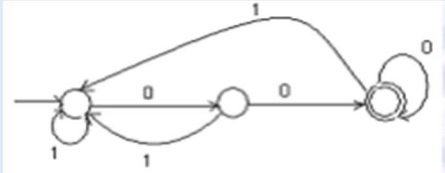
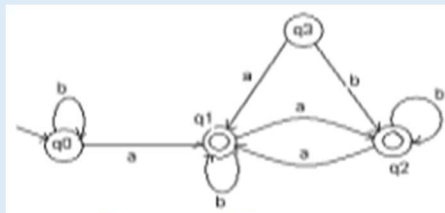
| | | | | | | | | | | | | | | |
|-------------------|---|---|----|---------------------------|---|------------|---|---|--------|--|---|---|---|---|
| | | | | | | | (a+b)]* | | | | | (a+b)]* | | |
| 0170 1349 2 | 2 | 2 | 28 | Regular Express ion | The language described by the regular expression $(0+1)^*0(0+1)^*0(0+1)^*$ over the alphabet $\{0, 1\}$ is the set of | MCQ | All strings containing at least two 0's | 1 | Hard | | All strings containing at least two 1's | All strings that begin and end with either 0's or 1's | All strings containing at least two 0's | All strings containing the substring 00 |
| 0170 1349 2 | 2 | 2 | 29 | Regular Express ion | What is a Regular language? | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 30 | Regular Express ion | Construct a r.e for the language which accepts all strings with atleast two c's over the set $\Sigma=\{c,b\}$ | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 31 | Regular Express ion | Give a regular expression for the following regular language: The set of all strings containing 00. | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 32 | Regular Express ion | for which of the following applications regular expressions can be used ? | MCQ | All of these | 1 | Easy | | Designing compilers | Developing text editors | Simulating sequential circuits | All of these |
| 0170 1349 2 | 2 | 2 | 33 | Regular Express ion | Regular Expression R and the language it describes can be represented as: | MCQ | c) R, L(R) | 1 | Easy | | a) R, R(L) | b) L(R), R(L) | c) R, L(R) | d) All of the mentioned |

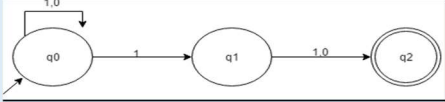
| | | | | | | | | | | | | | | |
|-------------------|---|---|----|---------------------------|--|------------|-------------------------|---|--------|--|--------------------------------|-----------------------|--------------------------|------------------------|
| 0170 1349 2 | 2 | 2 | 34 | Regular Express ion | If R is a regular expression, then which of the followings is FALSE? | MCQ | d) (R*) * = R+ | 1 | Easy | | a) R*= R*+ ϵ | b) R*R *= R* | c) ϵ R= R | d) (R*) *= R+ |
| 0170 1349 2 | 2 | 2 | 35 | Regular Express ion | Define Dead-End State with Example | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 36 | Regular Express ion | Write Regular Expressions for the following languages of all strings in $\{0,1\}^*$ (i) Strings that do not end with 01. (ii) Strings with odd numbers of 1's (Ones) | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 37 | Regular Express ion | Define regular language and regular expressions | Short Que. | | 1 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 38 | Regular Express ion | Find regular expression for the following: • Describe the language corresponding to following: $(1+01)^*(0+01)^*$ | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 39 | Regular Express ion | Write Regular Expressions for the following languages of all strings in $\{0,1\}^*$ (i) Strings that start with 1 and do not end with 10. (ii) Strings with length 6 or less | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 40 | Regular Express ion | Find a regular expression corresponding to each of the following subsets of $\{0, 1\}^*$. i. The language of all strings that do not contain the substring 110. ii. The language of all strings containing both 101 and 010 as substrings. iii. The language of all strings in which both the number of 0's and the number of 1's are odd | Short Que. | | 2 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 41 | Regular Express ion | Write a regular expression for language L over $\{0,1\}$ such that every string in L i) Begins with 00 and ends with 11. ii) Contains alternate 0 and 1 | Short Que. | | 2 | Hard | | | | | |

| | | | | | | | | | | | | | | |
|-------------------|---|---|----|---------------------------|---|------------|--|---|------|--|--|--|--|--|
| 0170 1349 2 | 2 | 2 | 42 | Regular Express ion | Prove the formula $(00^*1)^*1 = 1+0(0+10)^*11$ | Short Que. | | 2 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 43 | Regular Express ion | Write regular expression for the following languages i) $L1 = \{x \in (0,1)^* \mid x \text{ do not ends with } 11\}$ ii) $L2 = \{x \in (0,1)^* \mid x \text{ contains both } 101 \text{ and } 110\}$ | Short Que. | | 2 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 44 | Regular Express ion | Find regular expression for following i. Language of all strings containing exactly two 0's. ii. Language of all strings that begins or ends with 00 or 11. iii. Language of all strings in which every 0 is followed immediately by 11 | Short Que. | | 2 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 45 | Regular Express ion | Write Regular Expressions for the following languages of all strings in $\{0,1\}^*$ (i) Strings that do not end with 01. (ii) The language of all strings containing both 101 and 010 as substrings | Short Que. | | 2 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 46 | Regular Express ion | Write RE for the languages of all Strings that do not end with 01. | Short Que. | | 1 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 47 | Regular Express ion | Show that for any language L, $L^* = (L^*)^* = (L^*)^+ = (L^+)^*$ | Short Que. | | 1 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 48 | Regular Express ion | Write Regular Expression corresponding to each of the following subsets of $\{0, 1\}^*$ a. The language of all strings containing both 101 and 010 as substrings. b. The language of all strings in which both the number of 0's and the number of 1's are even | Short Que. | | 2 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 49 | Regular Express ion | L1 and L2 are two languages: $L1 = \{x \mid 11 \text{ is not a substring of } x\}$ $L2 = \{x \mid x \text{ starts with } 0 \text{ and ends with } 0\}$ Draw FA for both L1 and L2 and construct FA for $L3 = L2 - L1$ | Short Que. | | 2 | Hard | | | | | |

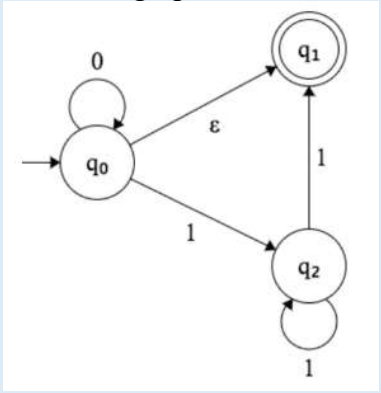
| | | | | | | | | | | | | | | |
|-------------------|---|---|----|---------------------------|---|------------|--|---|--------|--|--|--|--|--|
| 0170 1349 2 | 2 | 2 | 50 | Regular Express ion | Find a regular expression corresponding to each of the following subsets of $\{0, 1\}^*$ 1. The language of all strings that begin or end with 00 or 11. 2. The language of all strings containing both 11 and 010 as substrings | Short Que. | | 2 | Hard | | | | | |
| 0170 1349 2 | 2 | 2 | 51 | Closure properti es | What are the closure properties of regular languages? | Short Que. | | 1 | Easy | | | | | |
| 0170 1349 2 | 2 | 2 | 52 | Regular Express ion | Find a regular expression corresponding to each of the following subsets of $\{0, 1\}^*$ (i) The language of all strings that begin or end with 00 or 11. (ii) The language of all strings beginning with 1 and ending with 0 | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 53 | Regular Express ion | What are the applications of regular expressions and finite automata? | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 54 | Regular Express ion | Write Regular Expression over the alphabets $\{a, b\}$ consisting strings: Second last character as 'a' Starting with 'a' and ending with 'b' | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 55 | Regular Express ion | Define Regular Expression. Find Regular Expression corresponding to each of the following subsets of $\{0,1\}^*$ 1) The Language of all strings containing exactly two 0's 2) The Language of all strings that end with 01 3) The Language of all strings that begin or end with 00 or 11 | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 56 | Regular Express ion | Write regular expressions for the following languages defined over $\Sigma = \{0, 1\}$: (i) The language of all the strings that do not end with 01. (ii) The language of all the strings containing even number of 0's and even number of 1's | Short Que. | | 2 | Medium | | | | | |

| | | | | | | | | | | | | | | |
|-------------------|---|---|----|---------------------------|---|---|--|---|--------|--|--|--|--|---------------------------------------|
| 0170 1349 2 | 2 | 2 | 57 | Regular Express ion | Find a regular expression of following subsets of $\{0, 1\}^*$ 1. The language of all strings that begin or end with 00 or 11. 2. The language of all strings ending with 1 and not containing 00. | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 58 | Regular Express ion | Find a regular expression of following subsets of $\{0, 1\}^*$ 1. The language of all strings that contain odd number of 1's 2. The language of all strings with next to last symbol 0 | Short Que. | | 2 | Medium | | | | | |
| 0170 1349 2 | 2 | 2 | 59 | Regular Express ion | Find a regular expression corresponding to each of the following subsets of $\{0,1\}^*$ (i). the language of all strings that do not end with 01 (ii). the language of all strings that begin with or end with 00 or 11 | Short Que. | | 2 | Medium | | | | | |
| 1701 3492 | 3 | 3 | 60 | FA-1 | Finite automata requires minimum _____ number of stacks. | Rememberi ng- Multiple_Ch oice | 0 | 1 | Easy | | 0 | 1 | 2 | 3 |
| 1701 3492 | 3 | 3 | 61 | FA-1 | FSM with output capability can be used to add _____ given integer in binary representation. | Understandi ng- Multiple_Ch oice | 2 | 1 | Easy | | 3 | 1 | 0 | 5 |
| 1701 3492 | 3 | 3 | 62 | FA-1 | We can represent one language in more one FSMs, true or false? | Understandi ng- Multiple_Ch oice | TRUE | 1 | Easy | | May be true | TR UE | Cann ot be said | FAL SE |
| 1701 3492 | 3 | 3 | 63 | FA-1 | "The basic limitations of Finite Automata is | Understandi ng- Multiple_Ch oice | It can't reme mber arbit rary | 1 | Easy | | It som etim e reco gniz ed | It can' t rem emb er arbit | It some time fails to recog nize | All of the ment ione d |

| | | | | | | | | | | | | | | |
|----------|---|---|----|------|---|-----------------------------|------------------------------|---|------|--|------------------------------|---|----------------------------|--------------------------|
| | | | | | | | large amount of information. | | | | grammar that are not regular | regular language amount of information. | the Regular grammar. | |
| 17013492 | 3 | 3 | 64 | FA-1 | A language is regular if and only if | Remembering-Multiple_Choice | Accepted by DFA | 1 | Easy | | Accepted by PDA | Accepted by LBA | Accepted by Turing Machine | Accepted by DFA |
| 17013492 | 4 | 4 | 65 | FA-2 | <p>"The DFA shown below accepts the set of all strings over {0,1} that</p>  | Evaluating-Multiple_Choice | End with 00 | 1 | Hard | | End with 0 | End with 00 | Begin either with 0 or 1 | Contain the substring 00 |
| 17013492 | 4 | 4 | 66 | FA-2 | <p>Which one of the following is true for the automata</p>  | Evaluating-Multiple_Choice | $b^*a(a+b)^*$ | 1 | Hard | | $b^*ab^*ab^*$ | $b^*ab^*ab^*$ | $b^*a(a+b)^*$ | $(a+b)^*$ |
| 17013492 | 3 | 3 | 67 | FA-1 | There are _____ tuples in finite state machine. | Remembering-Multiple_Choice | 5 | 1 | Easy | | 4 | 3 | 5 | 0 |
| 17013492 | 3 | 3 | 68 | FA-1 | Number of states require to accept string ends with 10. | Evaluating-Multiple_Choice | 3 | 1 | Easy | | 4 | 3 | 5 | 0 |

| | | | | | | | | | | | | | | |
|----------|---|---|----|------|--|---|---|---|------|--|--|---|---|---|
| 17013492 | 3 | 3 | 69 | FA-1 | <p>Statement:</p>  <p>The DFA shown represents all strings which has 1 at second last position.</p> | Evaluating-Multiple_Ch oice | Wro ng prop osi ti on | 1 | Hard | | May be corr ect | Corr ect | Wro ng prop osi ti on | Inco rrect, Inco mple te DFA |
| 17013492 | 3 | 3 | 70 | FA-1 | Language of finite automata is. | Rememberi ng- Multiple_Ch oice | Typ e 3 | 1 | Easy | | Typ e 0 | Typ e 1 | Type 3 | Type 4 |
| 17013492 | 3 | 3 | 71 | FA-1 | Number of final state require to accept Φ in minimal finite automata. | Understandi ng- Multiple_Ch oice | No Fina l State Req uire d | 1 | Easy | | 1 | 2 | No Final State Req uired | 4 |
| 17013492 | 4 | 4 | 72 | FA-2 | Regular expression for all strings starts with ab and ends with bba is. | Applying- Multiple_Ch oice | ab(a +b)* bba | 1 | Easy | | aba* b*b ba | ab(a b)*b ba | ab(a +b)* bba | All of the ment ione d |
| 17013492 | 4 | 4 | 73 | FA-2 | NFA, in its name has 'non-deterministic' because of : | Rememberi ng- Multiple_Ch oice | The choi ce of path is non- deter mini stic | 1 | Easy | | The resul t is unde term ined | The state to be trans ited next is non- deter min istic | None of the ment ione d | The choi ce of path is non- deter mini stic |
| 17013492 | 3 | 3 | 74 | FA-1 | Which of the following option is correct? | Understandi ng- Multiple_Ch oice | NFA is slow er to proc ess | 1 | Easy | | NF A is slow er to proc ess | DF A is faste r to proc ess | NFA is slow er to proc ess | DFA is slow er to proc ess |

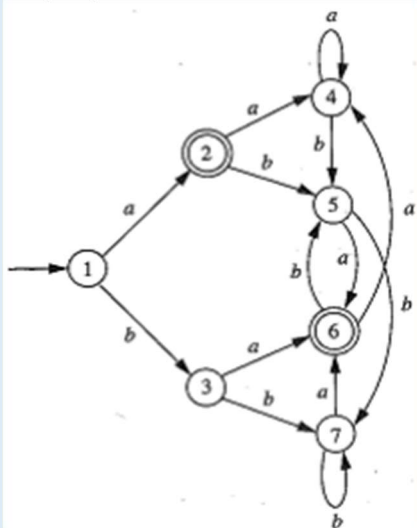
| | | | | | | | | | | | | | | |
|----------|---|---|----|------|---|-----------------------------|---|---|------|--|--|---|---|--|
| | | | | | | | and its representation uses less memory than DFA | | | | and its representation uses more memory than DFA | and its representation uses less memory than NFA | and its representation uses less memory than DFA | and its representation uses less memory than NFA |
| 17013492 | 3 | 3 | 75 | FA-1 | What is wrong in the given definition? Def: ($\{q_0, q_1, q_2\}$, $\{0,1\}$, δ , q_3 , $\{q_3\}$) | Evaluating-Multiple_Choice | Initial and Final states do not belong to the Graph | 1 | Hard | | Initial and final states can't be same | The definition does not satisfy 5 Tuple definition of NFA | Initial and Final states do not belong to the Graph | There are no transition definition |
| 17013492 | 3 | 3 | 76 | FA-1 | What is the relation between DFA and NFA on the basis of computational power? | Remembering-Multiple_Choice | Equal | 1 | Easy | | DFA > NFA | NFA > DFA | Equal | Can't be said |

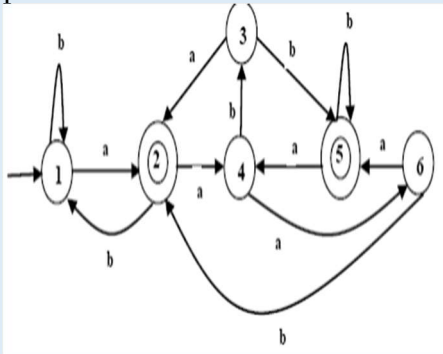
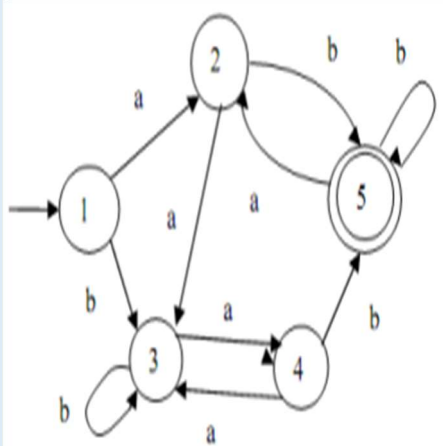
| | | | | | | | | | | | | | | |
|--------------|---|---|----|------|---|--------------------------------|--|---|--------|--|---|--|---|------------------------------------|
| 1701 3492 | 3 | 3 | 77 | FA-1 | <p>ϵ- closure of q_1 in the given transition graph:</p>  | Evaluating-Multiple_Ch oice | {q0, q1} | 1 | Easy | | {q1} | {q1, q2} | {q0, q1, q2} | {q0, q1} |
| 1701 3492 | 3 | 3 | 78 | FA-1 | <p>According to the given transitions, which among the following are the epsilon closures of q_1 for the given NFA? $\Delta(q_1, \epsilon) = \{q_2, q_3, q_4\}$ $\Delta(q_4, 1) = q_1$ $\Delta(q_1, \epsilon) = q_1$</p> | Evaluating-Multiple_Ch oice | {q1, q2, q3, q4} | 1 | Easy | | {q4} | {q1, q2, q3, q4} | {q1, q3} | {q1, q3,q 4} |
| 1701 3492 | 3 | 3 | 79 | FA-1 | <p>Given: $\Sigma = \{a, b\}$ $L = \{x \in \Sigma^* x \text{ is a string combination}\}$ Σ^4 represents which among the following?</p> | Evaluating-Multiple_Ch oice | {aaa a, abab , ϵ , abaa , aabb } | 1 | Easy | | {aa, ab, ba, bb} | {aaa a, abab , ϵ , abaa , aabb } | {aaa , aab, aba, bbb} | All of the ment ione d |
| 1701 3492 | 3 | 3 | 80 | FA-1 | <p>Which of the following pairs of regular expression are not equivalent?</p> | Evaluating-Multiple_Ch oice | none of the abov e | 1 | Medium | | $(a+b)^*$ and $(a^*+b)^*$ | $(a^*+b)^*$ and $(a+b)^*$ | $(ab)^*$ a and $a(ba)^*$ | none of the abov e |
| 1701 3492 | 3 | 3 | 81 | FA-1 | <p>We have two statements S1 and S2 whose definition are as follows: S1 – $\{0^2n \mid n \geq 1\}$ is a regular language. S2 – $\{0^m 1^n 0^l \mid m+n+l=1 \text{ and } n \geq 1\}$ is a regular language.</p> | Evaluating-Multiple_Ch oice | Only S1 is corre ct | 1 | Medium | | Both S1 and S2 are corr ect | Onl y S1 is corr ect | Neit her S1 nor S2 is corre ct | Only S2 is corre ct |

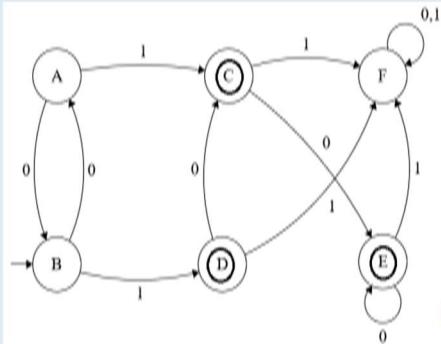
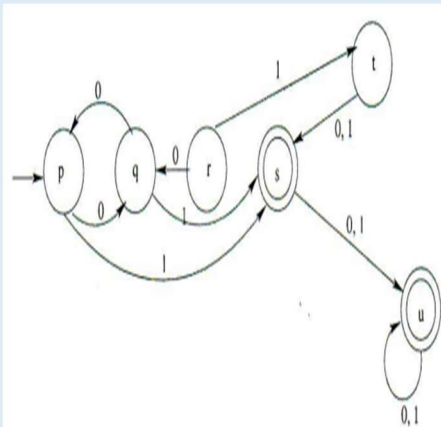
| | | | | | | | | | | | | | | |
|--------------|---|---|----|------|---|---|--------------------------|---|--------|--|----------------------|--------------------------|-------------------------|----------------------|
| | | | | | Which one of the following statements is correct? | | | | | | | | | |
| 1701 3492 | 3 | 3 | 82 | FA-1 | The minimum number of states in any DFA accepting the regular language $L = (111+11111)^*$ is | Evaluating- Multiple_Ch oice | 9 | 1 | Hard | | 7 | 5 | 9 | 11 |
| 1701 3492 | 3 | 3 | 83 | FA-1 | How many states are present in the smallest finite automaton which accepts the language $\{x \mid \text{length of } x \text{ is divisible by } 3\}$? | Evaluating- Multiple_Ch oice | 4 | 1 | Medium | | 4 | 5 | 3 | 2 |
| 1701 3492 | 3 | 3 | 84 | FA-1 | Which of the following regular expression identities are true ? | Rememberi ng- Multiple_Ch oice | $(r + s)^* = (r^*s^*)^*$ | 1 | Medium | | $(r + s)^* = r^*s^*$ | $(r + s)^* = (r^*s^*)^*$ | $(r + s)^* = r^* + s^*$ | $r^*s^* = r^* + s^*$ |
| 1701 3492 | 4 | 4 | 85 | FA-2 | Draw FA for Regular Expression: $(111 + 100)^*0$ | Applying- Short_Answ er_2_5_Min | | 2 | Hard | | | | | |
| 1701 3492 | 4 | 4 | 86 | FA-2 | Draw FA for Regular Expression: $(11 + 100)^*1$ | Applying- Short_Answ er_2_5_Min | | 2 | Hard | | | | | |
| 1701 3492 | 4 | 4 | 87 | FA-2 | Draw FA for Regular Expression: $0(10 + 01)^* + 1(00 + 01)^*$ | Applying- Short_Answ er_2_5_Min | | 2 | Medium | | | | | |
| 1701 3492 | 4 | 4 | 88 | FA-2 | Draw FA for the string 1) The string with next to last symbol as 0. 2) The string with number of 0's odd and numbers of 1's odd. | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Medium | | | | | |
| 1701 3492 | 4 | 4 | 89 | FA-2 | Draw FA for 1) $(11 + 110)^*0$ 2) $\{11\}^*\{00\}^*$ | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Hard | | | | | |
| 1701 3492 | 4 | 4 | 90 | FA-2 | Draw FA for the string 1) The string in $\{0,1\}^*$ ending in 10 or 11. | Evaluating- Long_Answ | | 4 | Hard | | | | | |

| | | | | | | | | | | | | | | |
|--------------|---|---|----|------|--|---|--|---|--------|--|--|--|--|--|
| | | | | | 2) The string end with 1 and does not contain substring 00 | er_5_10_Mi n | | | | | | | | |
| 1701 3492 | 4 | 4 | 91 | FA-2 | Draw FA for the strings: 1) The string in $\{a,b\}^*$ ending in aba. 2) The string corresponding to regular expression $(111 + 100)^*0$. | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Hard | | | | | |
| 1701 3492 | 4 | 4 | 92 | FA-2 | Draw FA for the strings: 1) The string in $\{0,1\}^*$ ending in 00 or 01. 2) The string corresponding to regular expression $(10 + 110)^*1$. | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Hard | | | | | |
| 1701 3492 | 4 | 4 | 93 | FA-2 | Draw FA for the corresponding language 1) $1(01 + 10)^* + 0(11 + 10)^*$ 2) $(010 + 00)^* (10)^*$ | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Hard | | | | | |
| 1701 3492 | 4 | 4 | 94 | FA-2 | Draw FA for the corresponding language 1) $(1+110)^*0$ 2) $(1+10+110)^*0$ | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Hard | | | | | |
| 1701 3492 | 4 | 4 | 95 | FA-2 | Draw FA for the corresponding language 1) $1(1+10)^*+10(0+01)^*$ 2) $0+(10)^*+01^*0$ | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Hard | | | | | |
| 1701 3492 | 4 | 4 | 96 | FA-2 | Draw FA for the string 1) The string with number of 0's odd and numbers of 1's even. 2) $(0+1)(01)^*(011)^*$ | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Medium | | | | | |
| 1701 3492 | 4 | 4 | 97 | FA-2 | Draw FA for accepting: 1) The string in $\{0,1\}^*$ ending in 1 and not containing substring 00. 2) The string with even numbers of 0's and even numbers of 1's. | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Medium | | | | | |
| 1701 3492 | 4 | 4 | 99 | FA-2 | Draw DFA for the following languages | Evaluating- Long_Answ | | 4 | Medium | | | | | |

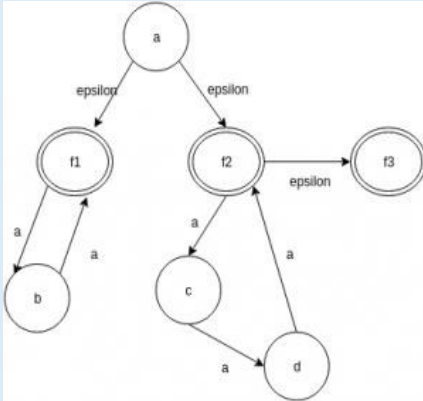
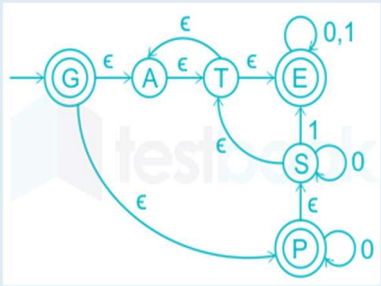
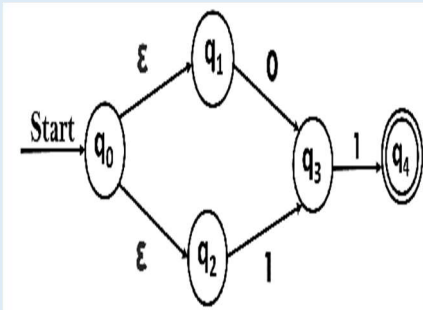
| | | | | | 1) $L1 = \{x \in (0,1)^* x \text{ contains } 110111\}$ 2) $L2 = \{x \in (0,1)^* x \text{ contains odd numbers of 1's and even numbers of 0's}\}$ | er_5_10_Min | | | | | | | | | | | | | | | | | | | | | | | | | | |
|--------------|---------------|---------------|-----|------|--|---------------------------------|---------------|---------------|--------|-------|-----|---|-----|-----|---|-----|-----|---|-----|-------------|---|-------------|-----|---------------------------------|--|---|--------|--|--|--|--|--|
| 1701 3492 | 4 | 4 | 100 | FA-2 | Draw DFA for the following languages 1) $L1 = \{x \in (0,1)^* x \text{ do not contains } 110\}$ 2) $L2 = \{x \in (0,1)^* x \text{ do not contain } 00 \text{ as a substring}\}$ | Evaluating-Long_Answer_5_10_Min | | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 101 | FA-2 | Draw FA for each of the following RE: 1) $(0+1)^*(1+00)(0+1)^*$ 2) $(0+1)^*(01+110)$ | Evaluating-Long_Answer_5_10_Min | | 4 | Hard | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 102 | FA-2 | Draw DFA for the following languages 1) $L1 = \{x \in (0,1)^* x \text{ end with } 01\}$ 2) $(0+1)^*(10+11)$ | Evaluating-Long_Answer_5_10_Min | | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 103 | FA-2 | Draw FA for the following languages 1) $L1 = \{x \in (0,1)^* \text{ends with } 11\}$ 2) $L2 = \{x \in (0,1)^* x \text{ contains both } 101 \text{ and } 110\}$ | Evaluating-Long_Answer_5_10_Min | | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 104 | FA-2 | An NFA with states 1-5 and input alphabets {a,b} has following transition table Q-1 Draw its transition diagram Q-2 Calculate $\delta^*(1,a)$ Q-3 Calculate $\delta^*(1,aaabaab)$ <table border="1"><thead><tr><th>q</th><th>$\delta(q,a)$</th><th>$\delta(q,b)$</th></tr></thead><tbody><tr><td>1</td><td>{1,2}</td><td>{1}</td></tr><tr><td>2</td><td>{3}</td><td>{3}</td></tr><tr><td>3</td><td>{4}</td><td>{4}</td></tr><tr><td>4</td><td>{5}</td><td>\emptyset</td></tr><tr><td>5</td><td>\emptyset</td><td>{5}</td></tr></tbody></table> | q | $\delta(q,a)$ | $\delta(q,b)$ | 1 | {1,2} | {1} | 2 | {3} | {3} | 3 | {4} | {4} | 4 | {5} | \emptyset | 5 | \emptyset | {5} | Evaluating-Long_Answer_5_10_Min | | 4 | Medium | | | | | |
| q | $\delta(q,a)$ | $\delta(q,b)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 1 | {1,2} | {1} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 2 | {3} | {3} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 3 | {4} | {4} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 4 | {5} | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 5 | \emptyset | {5} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

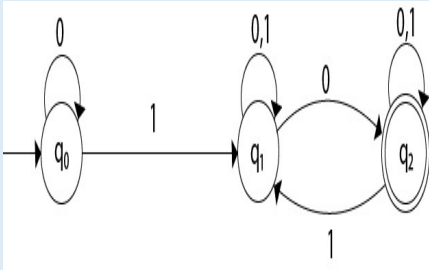
| | | | | | | | | | | | | | | |
|--------------|---|---|-----|------|--|---|---|--------|--|--|--|--|--|--|
| 1701 3492 | 4 | 4 | 105 | FA-2 | Draw FA for each of the following RE 1) $(a+b)^*baaa$ 2) $(bbb + baa)^*a$ | Evaluating- Long_Answ er_5_10_Mi n | 4 | Medium | | | | | | |
| 1701 3492 | 4 | 4 | 106 | FA-2 | For following NFA, find minimum FA accepting same language  | Evaluating- Long_Answ er_5_10_Mi n | 4 | Medium | | | | | | |
| 1701 3492 | 4 | 4 | 107 | FA-2 | For the following RE, draw an NFA 1) $(a+b)^*(abba^* + (ab)^*ba)$ 2) $(aa+aab)^*b$ | Evaluating- Long_Answ er_5_10_Mi n | 4 | Medium | | | | | | |
| 1701 3492 | 4 | 4 | 108 | FA-2 | For the following RE, draw an NFA 1) $((0+1)^*10 + (00)^*(11)^*)^*$ 2) $(0+1)^*1(0+1)$ | Evaluating- Long_Answ er_5_10_Mi n | 4 | Medium | | | | | | |
| 1701 3492 | 4 | 4 | 109 | FA-2 | For the following RE, draw an NFA 1) $(0+1)^*(011+01010)(0+1)^*$ 2) $(0+1)(01)^*(011)^*$ | Evaluating- Long_Answ er_5_10_Mi n | 4 | Medium | | | | | | |
| 1701 3492 | 4 | 4 | 110 | FA-2 | For the following RE, draw an NFA 1) $(0+1)^*(10+110)^*1$ 2) $0^*(01)^*1+1^*0$ | Evaluating- Long_Answ | 4 | Medium | | | | | | |

| | | | | | | er_5_10_Min | | | | | | | | | | | | | | | | | | | | | | | | | | |
|--------------|----------------|----------------|-----|------|--|---------------------------------|----------------|----------------|---|----------|-------|---|-------|-------|---|-------|-------|---|-------|-------------|---|-------------|-------|---------------------------------|---|--------|--|--|--|--|--|--|
| 1701 3492 | 4 | 4 | 111 | FA-2 | <p>Minimize the following DFA if possible:</p>  | Evaluating-Long_Answer_5_10_Min | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 112 | FA-2 | <p>Minimize the following DFA if possible:</p>  | Evaluating-Long_Answer_5_10_Min | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 113 | FA-2 | <p>Calculate $\delta^*(1,abbaa)$ & $\delta^*(1,abaabba)$ form following transition table.</p> <table border="1" data-bbox="619 1224 961 1485"><thead><tr><th>q</th><th>$\delta(q, a)$</th><th>$\delta(q, b)$</th></tr></thead><tbody><tr><td>1</td><td>{ 1, 2 }</td><td>{ 1 }</td></tr><tr><td>2</td><td>{ 3 }</td><td>{ 3 }</td></tr><tr><td>3</td><td>{ 4 }</td><td>{ 4 }</td></tr><tr><td>4</td><td>{ 5 }</td><td>\emptyset</td></tr><tr><td>5</td><td>\emptyset</td><td>{ 5 }</td></tr></tbody></table> | q | $\delta(q, a)$ | $\delta(q, b)$ | 1 | { 1, 2 } | { 1 } | 2 | { 3 } | { 3 } | 3 | { 4 } | { 4 } | 4 | { 5 } | \emptyset | 5 | \emptyset | { 5 } | Evaluating-Long_Answer_5_10_Min | 4 | Medium | | | | | | |
| q | $\delta(q, a)$ | $\delta(q, b)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 1 | { 1, 2 } | { 1 } | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 2 | { 3 } | { 3 } | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 3 | { 4 } | { 4 } | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 4 | { 5 } | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 5 | \emptyset | { 5 } | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

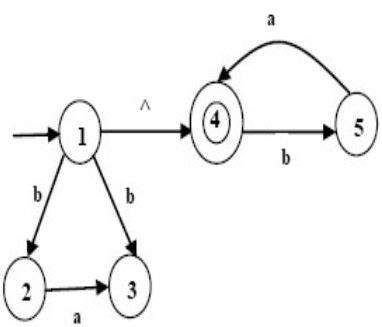
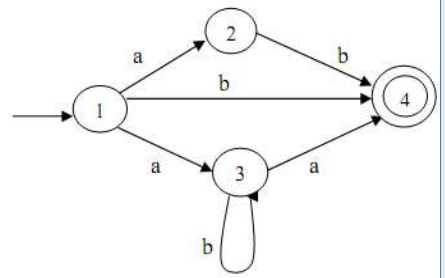
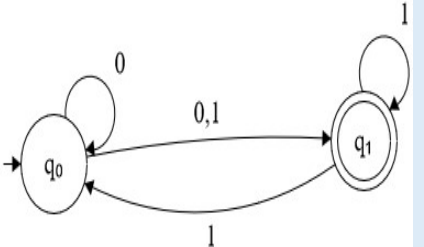
| 1701 3492 | 4 | 4 | 114 | FA-2 | <p>Minimize the following DFA if possible:</p> <table> <tr> <th>Q</th> <th>$\delta(q,a)$</th> <th>$\delta(q,b)$</th> </tr> <tr> <td>- +1</td> <td>{3}</td> <td>{2}</td> </tr> <tr> <td>2</td> <td>{4}</td> <td>{1}</td> </tr> <tr> <td>3</td> <td>{5}</td> <td>{4}</td> </tr> <tr> <td>4</td> <td>{4}</td> <td>{4}</td> </tr> <tr> <td>5</td> <td>{3}</td> <td>{2}</td> </tr> </table> | Q | $\delta(q,a)$ | $\delta(q,b)$ | - +1 | {3} | {2} | 2 | {4} | {1} | 3 | {5} | {4} | 4 | {4} | {4} | 5 | {3} | {2} | Evaluating-Long_Answer_5_10_Min | 4 | Easy | | | | | | |
|--------------|---------------|---------------|-----|------|--|---------------------------------|---------------|---------------|------|-----|-----|---|-----|-----|---|-----|-----|---|-----|-----|---|-----|-----|---------------------------------|---|------|--|--|--|--|--|--|
| Q | $\delta(q,a)$ | $\delta(q,b)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| - +1 | {3} | {2} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 2 | {4} | {1} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 3 | {5} | {4} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 4 | {4} | {4} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 5 | {3} | {2} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 115 | FA-2 | <p>Minimize the following DFA if possible:</p>  | Evaluating-Long_Answer_5_10_Min | 4 | Easy | | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 116 | FA-2 | <p>Minimize the following DFA if possible:</p>  | Evaluating-Long_Answer_5_10_Min | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | | |

| 1701 3492 | 4 | 4 | 117 | FA-2 | <p>Find \wedge-closure of the each of the states in following NFA-\wedge</p> <table><tr><th>q</th><th>$\delta(q,a)$</th><th>$\delta(q,b)$</th><th>$\delta(q,\Lambda)$</th></tr><tr><td>1</td><td>\emptyset</td><td>\emptyset</td><td>$\{2\}$</td></tr><tr><td>2</td><td>$\{3\}$</td><td>\emptyset</td><td>$\{5\}$</td></tr><tr><td>3</td><td>\emptyset</td><td>$\{4\}$</td><td>\emptyset</td></tr><tr><td>4</td><td>$\{4\}$</td><td>\emptyset</td><td>$\{1\}$</td></tr><tr><td>5</td><td>\emptyset</td><td>$\{6,7\}$</td><td>\emptyset</td></tr><tr><td>6</td><td>$\{5\}$</td><td>\emptyset</td><td>\emptyset</td></tr><tr><td>7</td><td>\emptyset</td><td>\emptyset</td><td>$\{1\}$</td></tr></table> | q | $\delta(q,a)$ | $\delta(q,b)$ | $\delta(q,\Lambda)$ | 1 | \emptyset | \emptyset | $\{2\}$ | 2 | $\{3\}$ | \emptyset | $\{5\}$ | 3 | \emptyset | $\{4\}$ | \emptyset | 4 | $\{4\}$ | \emptyset | $\{1\}$ | 5 | \emptyset | $\{6,7\}$ | \emptyset | 6 | $\{5\}$ | \emptyset | \emptyset | 7 | \emptyset | \emptyset | $\{1\}$ | Evaluating- Short_Answ er_2_5_Min | | 2 | Medium | | | | | | |
|-----------------|---------------|---------------|---------------------|-------------|--|---|---------------|---------------|---------------------|-------|-----------------|-------------|---------|---------|---------|-------------|---------|---------|-------------|-------------|---|---------|---------|-------------|---------|---|-------------|-----------|-------------|---|---------|-------------|-------------|---|-------------|-------------|---------|---|--|---|--------|--|--|--|--|--|--|
| q | $\delta(q,a)$ | $\delta(q,b)$ | $\delta(q,\Lambda)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 1 | \emptyset | \emptyset | $\{2\}$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 2 | $\{3\}$ | \emptyset | $\{5\}$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 3 | \emptyset | $\{4\}$ | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 4 | $\{4\}$ | \emptyset | $\{1\}$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 5 | \emptyset | $\{6,7\}$ | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 6 | $\{5\}$ | \emptyset | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 7 | \emptyset | \emptyset | $\{1\}$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 118 | FA-2 | <p>Consider the NFA-\wedge depicted in following table</p> <p>1) Compute the \wedge-closure of each states.</p> <p>2) Find $\delta^*(q_0,1111)$</p> <p>3) Find $\delta^*(q_0,0011)$</p> <p>4) Find $\delta^*(q_0,1001)$</p> <p>5) Find $\delta^*(q_0,0111)$</p> <table><tr><th>q</th><th>$\delta(q,0)$</th><th>$\delta(q,1)$</th></tr><tr><td>q_0</td><td>q_0</td><td>q_0, q_1</td></tr><tr><td>q_1</td><td>q_2</td><td>q_2</td></tr><tr><td>q_2</td><td>q_3</td><td>q_3</td></tr><tr><td>q_3</td><td>\emptyset</td><td>\emptyset</td></tr></table> | q | $\delta(q,0)$ | $\delta(q,1)$ | q_0 | q_0 | q_0, q_1 | q_1 | q_2 | q_2 | q_2 | q_3 | q_3 | q_3 | \emptyset | \emptyset | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | |
| q | $\delta(q,0)$ | $\delta(q,1)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| q_0 | q_0 | q_0, q_1 | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| q_1 | q_2 | q_2 | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| q_2 | q_3 | q_3 | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| q_3 | \emptyset | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 1701 3492 | 4 | 4 | 119 | FA-2 | <p>Consider the NFA-\wedge depicted in following table</p> <p>1) Compute the \wedge-closure of each states.</p> <p>2) Find $\delta^*(q_0,abab)$</p> <p>3) Find $\delta^*(q_0,aaabbb)$</p> <table><tr><th></th><th>Λ</th><th>a</th><th>b</th><th>c</th></tr><tr><td>$\rightarrow p$</td><td>\emptyset</td><td>$\{p\}$</td><td>$\{q\}$</td><td>$\{r\}$</td></tr><tr><td>q</td><td>$\{p\}$</td><td>$\{q\}$</td><td>$\{r\}$</td><td>\emptyset</td></tr><tr><td>$*r$</td><td>$\{q\}$</td><td>$\{r\}$</td><td>\emptyset</td><td>$\{p\}$</td></tr></table> | | Λ | a | b | c | $\rightarrow p$ | \emptyset | $\{p\}$ | $\{q\}$ | $\{r\}$ | q | $\{p\}$ | $\{q\}$ | $\{r\}$ | \emptyset | $*r$ | $\{q\}$ | $\{r\}$ | \emptyset | $\{p\}$ | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Medium | | | | | | | | | | | | | | | | | | |
| | Λ | a | b | c | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $\rightarrow p$ | \emptyset | $\{p\}$ | $\{q\}$ | $\{r\}$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| q | $\{p\}$ | $\{q\}$ | $\{r\}$ | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $*r$ | $\{q\}$ | $\{r\}$ | \emptyset | $\{p\}$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

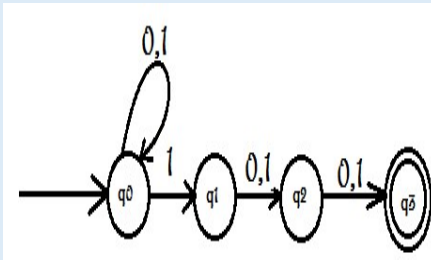
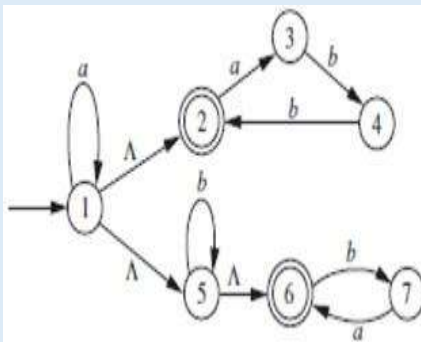
| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|--|---|----------------------------|-----------------|---|--------|--|-----------------|--------------|--------------|-------------|
| 0170 1349 4 | 05 | 05 | 120 | Finite Automata with Epsilon Transition | <p>Which of the following belongs to the epsilon closure set of a?</p>  | Evaluating Multiple-Choice | {a, f1, f2, f3} | 1 | Medium | | {a, f1, f2, f3} | {a, f1, f2} | {f1, f2, f3} | {a, f2, f3} |
| 0170 1349 4 | 05 | 05 | 121 | Finite Automata with Epsilon Transition | <p>Which of the following belongs to the epsilon closure set of S?</p>  | Evaluating Multiple-Choice | {S, T, A, E} | 1 | Medium | | {S, T} | {S, T, A, E} | {S} | {S, P, T} |
| 0170 1349 4 | 05 | 05 | 122 | Conversion from finite automata with Epsilon Transition to Deterministic finite automata | <p>While converting NFA with null to DFA, what will be $\delta'(A, 0)$ for the following NFA?</p>  | Evaluating Multiple-Choice | {q3} | 1 | Hard | | {q0, q1, q2} | {q1, q2} | {q3} | {q0, q1} |

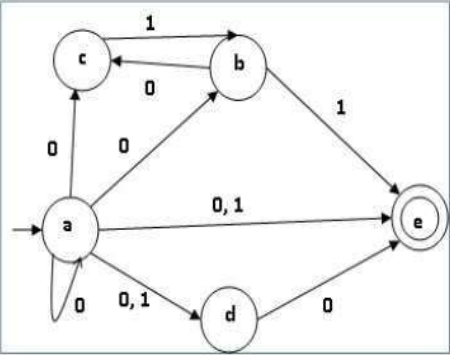
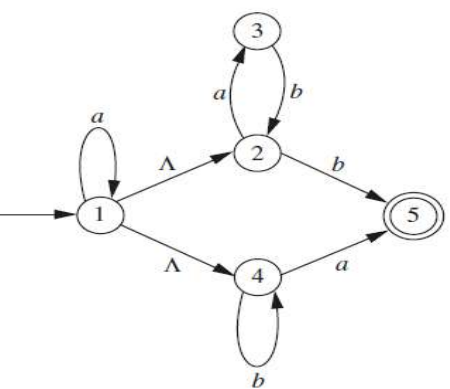
| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|--|---|-------------------------------|----------------------------------|---|------|--|---------------|----------------------------------|-----------|----------------------|
| 0170 1349 4 | 05 | 05 | 123 | Conversion from Non-Deterministic finite automata to Deterministic finite automata | <p>Which new state is generated while converting NFA to DFA and finding $\delta'([q1], 0)$?</p>  | Evaluating-Multiple-Choice | [q1, q2] | 1 | Hard | | [q0, q1, q2] | [q1, q2] | [q2] | [q1] |
| 0170 1349 4 | 05 | 05 | 124 | Conversion from Deterministic finite automata to Non-Deterministic finite automata | Conversion of a DFA to an NFA? | Understanding-Multiple-Choice | Requires the subset construction | 1 | Easy | | Is impossible | Requires the subset construction | Is chancy | Is non-deterministic |
| 0170 1349 4 | 05 | 05 | 125 | Non-Deterministic finite automata | If we consider an arbitrary NFA (non-deterministic finite automaton) with N states in total, the maximum number of states that are there in an equivalent DFA (minimized) is at least: | Understanding-Multiple-Choice | 2^N | 1 | Easy | | N! | $2N$ | 2^N | N^2 |
| 0170 1349 4 | 05 | 05 | 126 | Deterministic finite automata, Non-Deterministic finite | The total time needed to run any input string in DFA is than time required in NFA. | Understanding-Multiple-Choice | less | 1 | Easy | | more | less | equal | None of these |

| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
|-------------------|---------------------|----------------|----------------|---|--|---------------------------------------|---------------------|----------------|----------------|---|-----|---------|-------------|--------------------|-----|-----|-------------|---|-------------|-------------|-----|---|-------------|-----|-------------|---|--|---|------|--|--|--|--|--|
| | | | | automat a | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 127 | Determi nstic finite automat a ,Non- Determi nstic finite automat a | Which of the following cannot use Empty String transition? | Understandi ng-Multiple- Choice | DFA | 1 | Easy | | FA | NF A | DFA | All of these | | | | | | | | | | | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 128 | Non- Determi nstic finite automat a | Which of the following can use Empty String transition? | Understandi ng-Multiple- Choice | NFA | 1 | Easy | | FA | NF A | DFA | All of these | | | | | | | | | | | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 129 | Convers ion from Non - Determi nstic finite automat a with epsilon to Determi nstic finite automat a and Non- Determi nstic finite automat a | Convert NFA-^ to NFA and DFA. Initial State: A, Final State: D <table><tr><td>Q</td><td>$\delta(q, \wedge)$</td><td>$\delta(q, 0)$</td><td>$\delta(q, 1)$</td></tr><tr><td>A</td><td>{B}</td><td>{A}</td><td>\emptyset</td></tr><tr><td>B</td><td>{D}</td><td>{C}</td><td>\emptyset</td></tr><tr><td>C</td><td>\emptyset</td><td>\emptyset</td><td>{B}</td></tr><tr><td>D</td><td>\emptyset</td><td>{D}</td><td>\emptyset</td></tr></table> | Q | $\delta(q, \wedge)$ | $\delta(q, 0)$ | $\delta(q, 1)$ | A | {B} | {A} | \emptyset | B | {D} | {C} | \emptyset | C | \emptyset | \emptyset | {B} | D | \emptyset | {D} | \emptyset | Evaluating- Long_Answ er_5_10_Mi n | | 5 | Hard | | | | | |
| Q | $\delta(q, \wedge)$ | $\delta(q, 0)$ | $\delta(q, 1)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| A | {B} | {A} | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| B | {D} | {C} | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| C | \emptyset | \emptyset | {B} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| D | \emptyset | {D} | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

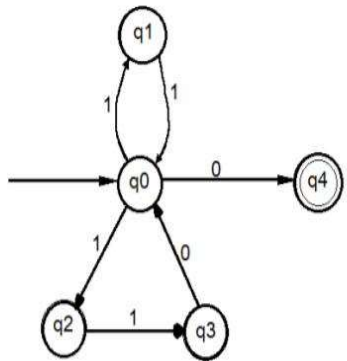
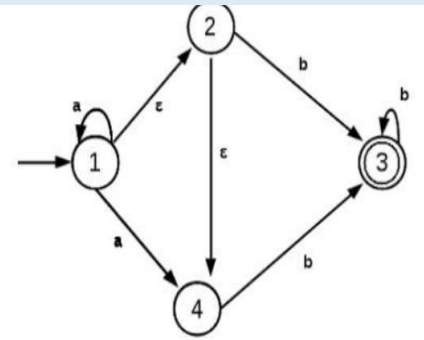
| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|---|--|---------------------------------|---|--------|--|--|--|--|--|--|
| 0170 1349 4 | 05 | 05 | 130 | Conversion from Non-Deterministic finite automata with epsilon to Deterministic finite automata | Convert the following NFA- Λ into FA.  | Evaluating-Long_Answer_5_10_Min | 5 | Hard | | | | | | |
| 0170 1349 4 | 05 | 05 | 131 | Conversion from Non-Deterministic finite automata to Deterministic finite automata | Convert the following NFA into FA.  | Evaluating-Long_Answer_5_10_Min | 4 | Medium | | | | | | |
| 0170 1349 4 | 05 | 05 | 132 | Conversion from Non-Deterministic finite automata to Deterministic finite | Convert the following NFA to DFA  | Evaluating-Short_Answer_2_5_Min | 4 | Medium | | | | | | |

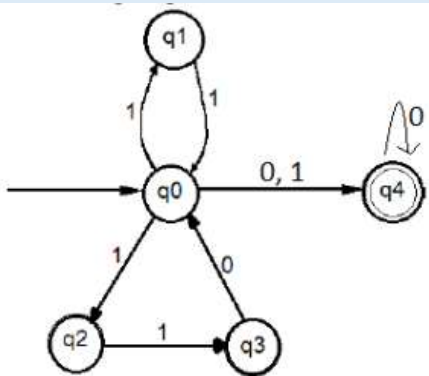
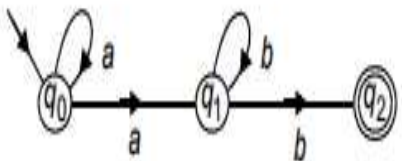
| | | | | automat a | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
|-----------------------------|----------------|---------------------------------|---------|---|--|----------------------|--------------|---------|---------|---|-----------------|----------------|---------------------------------|----------------|----------------|----------------|-----------------------------|----------------|-----|---------------------------------|-----|---|------|-----|-----|---|---|---|---|---------------------------------|--|---|------|--|--|--|--|--|
| 0170 1349 4 | 05 | 05 | 133 | Conversion from Non - Deterministic finite automata to Deterministic finite automata | Convert the NFA given in Table below to its corresponding DFA and draw the DFA. ----- <table><tr><th rowspan="2">Current State</th><th colspan="2">Input symbol</th></tr><tr><th>0</th><th>1</th></tr><tr><td>→Q₀</td><td>Q₁</td><td>Q₀, Q₂</td></tr><tr><td>Q₁</td><td>Q₂</td><td>Q₀</td></tr><tr><td>Q₂[*]</td><td>Q₀</td><td>---</td></tr></table> | Current State | Input symbol | | 0 | 1 | →Q ₀ | Q ₁ | Q ₀ , Q ₂ | Q ₁ | Q ₂ | Q ₀ | Q ₂ [*] | Q ₀ | --- | Evaluating-Short_Answer_2_5_Min | | 4 | Hard | | | | | | | | | | | | | | | |
| Current State | Input symbol | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | 0 | 1 | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| →Q ₀ | Q ₁ | Q ₀ , Q ₂ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| Q ₁ | Q ₂ | Q ₀ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| Q ₂ [*] | Q ₀ | --- | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 134 | Conversion from Non - Deterministic finite automata with epsilon to Deterministic finite automata and Non-Deterministic finite automata | Convert NFA-^ to NFA and DFA. Initial State: A , Final State: E <table><tr><th>Q</th><th>δ(q, ^)</th><th>δ(q, 0)</th><th>δ(q, 1)</th></tr><tr><td>A</td><td>{B,D}</td><td>{A}</td><td>∅</td></tr><tr><td>B</td><td>∅</td><td>{C}</td><td>{E}</td></tr><tr><td>C</td><td>∅</td><td>∅</td><td>{B}</td></tr><tr><td>D</td><td>∅</td><td>{E}</td><td>{D}</td></tr><tr><td>E</td><td>∅</td><td>∅</td><td>∅</td></tr></table> | Q | δ(q, ^) | δ(q, 0) | δ(q, 1) | A | {B,D} | {A} | ∅ | B | ∅ | {C} | {E} | C | ∅ | ∅ | {B} | D | ∅ | {E} | {D} | E | ∅ | ∅ | ∅ | Evaluating-Long_Answer_5_10_Min | | 5 | Hard | | | | | |
| Q | δ(q, ^) | δ(q, 0) | δ(q, 1) | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| A | {B,D} | {A} | ∅ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| B | ∅ | {C} | {E} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| C | ∅ | ∅ | {B} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| D | ∅ | {E} | {D} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| E | ∅ | ∅ | ∅ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 135 | Conversion from Non - | Consider the NFA-Λ depicted in following table: | Evaluating-Long_Answ | | 4 | Hard | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

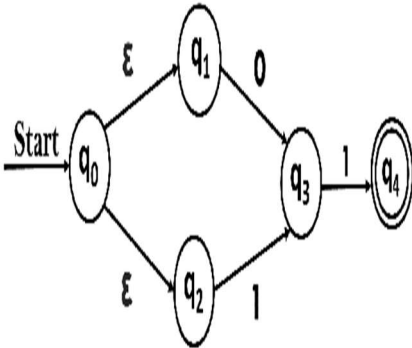
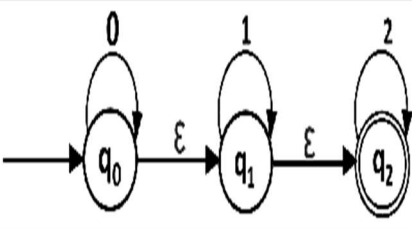
| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
|-----------------|-----------|---------|---------|---|---|---------------------------------|-----------|--------|---|---|-----------------|--------|---------|---------|---------|---|---------|---------|---------|--------|----|---------|---------|--------|---------|-------------|--|--|--|--|--|--|--|--|
| | | | | Deterministic finite automata with epsilon to Deterministic finite automata | <table><tr><td></td><td>Λ</td><td>a</td><td>b</td><td>c</td></tr><tr><td>$\rightarrow p$</td><td>Φ</td><td>$\{p\}$</td><td>$\{q\}$</td><td>$\{r\}$</td></tr><tr><td>q</td><td>$\{p\}$</td><td>$\{q\}$</td><td>$\{r\}$</td><td>Φ</td></tr><tr><td>*r</td><td>$\{q\}$</td><td>$\{r\}$</td><td>Φ</td><td>$\{p\}$</td></tr></table> <div><div>i)</div><div>Compute the Λ-closure of each state.</div></div> <div><div>ii)</div><div>Convert the NFA-Λ to a DFA</div></div> | | Λ | a | b | c | $\rightarrow p$ | Φ | $\{p\}$ | $\{q\}$ | $\{r\}$ | q | $\{p\}$ | $\{q\}$ | $\{r\}$ | Φ | *r | $\{q\}$ | $\{r\}$ | Φ | $\{p\}$ | er_5_10_Min | | | | | | | | |
| | Λ | a | b | c | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $\rightarrow p$ | Φ | $\{p\}$ | $\{q\}$ | $\{r\}$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| q | $\{p\}$ | $\{q\}$ | $\{r\}$ | Φ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| *r | $\{q\}$ | $\{r\}$ | Φ | $\{p\}$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 017013494 | 05 | 05 | 136 | Conversion from Non - Deterministic finite automata to Deterministic finite automata | Convert this NFA to FA <div></div> | Evaluating-Long_Answer_5_10_Min | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 017013494 | 05 | 05 | 137 | Conversion from Non - Deterministic finite automata with epsilon to Deterministic finite automata | Figure shows NFA- Λ . Draw an FA accepting the same language. <div></div> | Evaluating-Long_Answer_5_10_Min | 5 | Medium | | | | | | | | | | | | | | | | | | | | | | | | | | |

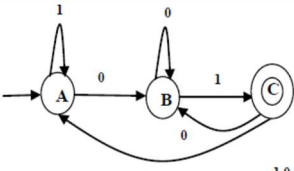
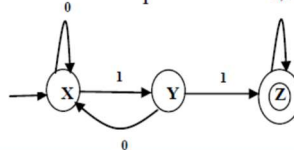
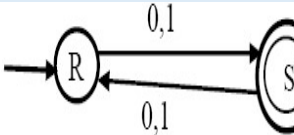
| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|---|---|---------------------------------|---|--------|--|--|--|--|--|--|
| 0170 1349 4 | 05 | 05 | 138 | Conversion from Non-Deterministic finite automata to Deterministic finite automata | <p>Convert the following NDFA to DFA.</p>  | Evaluating-Long_Answer_5_10_Min | 4 | Hard | | | | | | |
| 0170 1349 4 | 05 | 05 | 139 | Conversion from Non-Deterministic finite automata with epsilon to Deterministic finite automata | <p>Convert the following NFA - Λ into its equivalent DFA that accepts the same language:</p>  | Evaluating-Long_Answer_5_10_Min | 5 | Hard | | | | | | |
| 0170 1349 4 | 05 | 05 | 140 | Conversion from Non-Deterministic finite automata with epsilon to Deterministic | <p>Convert NFA- Λ to FA for following figure.</p> | Evaluating-Long_Answer_5_10_Min | 4 | Medium | | | | | | |

| | | | | | | | | | | | | | | | |
|-------------------|----|----|-----|--|---|---|--|---|------|--|--|--|--|--|--|
| | | | | nstic finite automat a | | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 141 | Convers ion from Non - Determi nistic finite automat a to Determi nistic finite automat a | Convert the following NFA into its equivalent DFA | Evaluating- Long_Answ er_5_10_Mi n | | 5 | Hard | | | | | | |
| 0170 1349 4 | 05 | 05 | 142 | Convers ion from Non - Determi nistic finite automat a to Determi nistic finite automat a | Convert the Given NFAwith null into its equivalent NFA | Evaluating- Short_Answ er_2_5_Min | | 4 | Easy | | | | | | |
| 0170 1349 4 | 05 | 05 | 143 | Convers ion from Non - Determi | Convert the given NFA to FA. | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Easy | | | | | | |

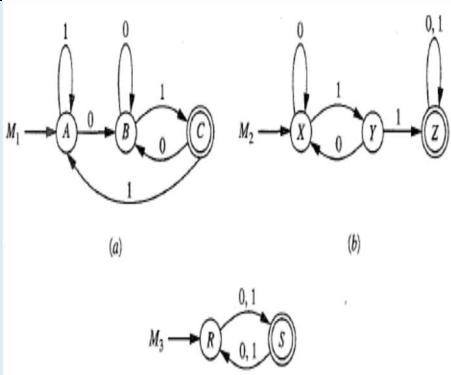
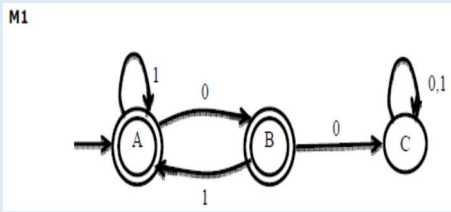
| | | | | nistic finite automat a to Determi nistic finite automat a |  | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
|-------------------|---------------|---------------|---------------------|---|--|---|---------------|---------------|---------------------|---|-------------|-------------|-----|---|-----|-------------|-----|---|-------------|-----|-------------|---|-----|-------------|-----|---|-------------|-------|-------------|---|-----|-------------|-------------|---|-------------|-------------|-----|---|--|---|------|--|--|--|--|--|--|
| 0170 1349 4 | 05 | 05 | 144 | Convers ion from Non - Determi nistic finite automat a with epsilon to Determi nistic finite automat a | Find Λ -Closure for each of the states in following NFA- Λ . And convert it into NFA and FA. <table><tr><th>q</th><th>$\delta(q,a)$</th><th>$\delta(q,b)$</th><th>$\delta(q,\Lambda)$</th></tr><tr><td>1</td><td>\emptyset</td><td>\emptyset</td><td>{2}</td></tr><tr><td>2</td><td>{3}</td><td>\emptyset</td><td>{5}</td></tr><tr><td>3</td><td>\emptyset</td><td>{4}</td><td>\emptyset</td></tr><tr><td>4</td><td>{4}</td><td>\emptyset</td><td>{1}</td></tr><tr><td>5</td><td>\emptyset</td><td>{6,7}</td><td>\emptyset</td></tr><tr><td>6</td><td>{5}</td><td>\emptyset</td><td>\emptyset</td></tr><tr><td>7</td><td>\emptyset</td><td>\emptyset</td><td>{1}</td></tr></table> | q | $\delta(q,a)$ | $\delta(q,b)$ | $\delta(q,\Lambda)$ | 1 | \emptyset | \emptyset | {2} | 2 | {3} | \emptyset | {5} | 3 | \emptyset | {4} | \emptyset | 4 | {4} | \emptyset | {1} | 5 | \emptyset | {6,7} | \emptyset | 6 | {5} | \emptyset | \emptyset | 7 | \emptyset | \emptyset | {1} | Evaluating- Long_Answ er_5_10_Mi n | | 5 | Hard | | | | | | |
| q | $\delta(q,a)$ | $\delta(q,b)$ | $\delta(q,\Lambda)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 1 | \emptyset | \emptyset | {2} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 2 | {3} | \emptyset | {5} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 3 | \emptyset | {4} | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 4 | {4} | \emptyset | {1} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 5 | \emptyset | {6,7} | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 6 | {5} | \emptyset | \emptyset | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 7 | \emptyset | \emptyset | {1} | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 145 | Convers ion from Non - Determi nistic finite automat a with epsilon to Non- Determi nistic finite | Convert the following ϵ -NFA into NFA.  | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Medium | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

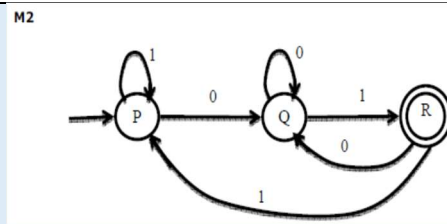
| | | | | | | | | | | | | | | | | |
|-------------------|----|----|-----|--|--|------------------------------------|--|---|--------|--|--|--|--|--|--|--|
| | | | | automat a | | | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 146 | Conversion from Non - Deterministic finite automata with null to Non - Deterministic finite automata | Convert NFA- Λ to NFA for following figure  | Evaluating-Long_Answer_5_10_Min | | 4 | Hard | | | | | | | |
| 0170 1349 4 | 05 | 05 | 147 | Conversion from Non - Deterministic finite automata to Deterministic finite automata | Determine the equivalent DFA for the above given NFA.  | Understanding-Short_Answer_2_5_Min | | 3 | Easy | | | | | | | |
| 0170 1349 4 | 05 | 05 | 148 | Conversion from Non - Deterministic finite automata with epsilon | Convert the NFA with ϵ into its equivalent DFA. | Evaluating-Long_Answer_5_10_Min | | 4 | Medium | | | | | | | |

| | | | | | | | | | | | | | | | |
|-------------------|----|----|-----|--|---|--|--|----|--------|--|--|--|--|--|--|
| | | | | to Determi nistic finite automat a |  | | | | | | | | | | |
| 0170 1349 4 | 05 | 05 | 149 | Convers ion from Non - Determi nistic finite automat a to Determi nistic finite automat a | Convert the given NFA into its equivalent DFA.  | Evaluating- Long_Answ er_5_10_Mi n | | 4 | Medium | | | | | | |
| 0170 1349 2 | 06 | 06 | 150 | kleene's Theore m | Using kleene's Theorem Draw NFA- Λ for $((01)^*10 + (00)^*)^*$ | Understandi ng- Short_Answ er_2_5_Min | | 03 | Medium | | | | | | |
| 0170 1349 2 | 06 | 06 | 151 | kleene's Theore m | Using kleene's Theorem Draw NFA- Λ for $((0+1)^*10 + (00)^*)^*$ | Understandi ng- Short_Answ er_2_5_Min | | 03 | Medium | | | | | | |
| 0170 1349 2 | 06 | 06 | 152 | kleene's Theore m | Draw NFA recognizing the language $(\{0,1\}^*10 \cup \{00\}^*\{11\}^*)^*$ using kleene's theorem part 1, where $\Sigma = \{0,1\}$ | Understandi ng- Short_Answ er_2_5_Min | | 03 | Medium | | | | | | |
| 0170 1349 2 | 06 | 06 | 153 | kleene's Theore m | Using kleene's Theorem Draw NFA- Λ for $((0+1)(01)^*)$ | Understandi ng- | | 03 | Medium | | | | | | |

| | | | | | | | | | | | | | | | |
|-----------|----|----|-----|---------------------------------------|--|------------------------------------|--|----|--------|--|--|--|--|--|--|
| | | | | | | Short_Answer_2_5_Min | | | | | | | | | |
| 017013492 | 06 | 06 | 154 | Union, Intersection, Difference of FA | <p>Let M1, M2 and M3 be the FAs pictured in Figure, recognizing languages L1, L2 and L3, respectively.</p> <div><p>M₁ = </p><p>M₂ = </p><p>M₃ = </p><p>Draw FAs recognizing the following languages.</p><p>a. $L_1 \cup L_2$ b. $L_1 \cap L_3$.</p></div> | Evaluating-Long_Answer_5_10_Min | | 05 | Hard | | | | | | |
| 017013492 | 06 | 06 | 155 | Union, Intersection, Difference of FA | <p>L1 is a language over $\{0, 1\}^*$ that accepts strings ending in 11. L2 is a language over $\{0, 1\}^*$ that accepts strings containing 101 as sub-string. Write the regular expressions, draw FA for L1 and L2 and derive FA for $L_1 \cup L_2$</p> | Understanding-Short_Answer_2_5_Min | | 05 | Hard | | | | | | |
| 017013492 | 06 | 06 | 156 | Union, Intersection, Difference of FA | <p>Draw FA for follow.</p> <p>□ $L_1 = \{w \mid 00 \text{ is not substring of } w\}$</p> <p>□ $L_2 = \{w \mid w \text{ ends in } 01\}$</p> <p>languages:</p> <p>Find FA accepting languages (i)$L_1 \cup L_2$ and (ii)$L_1 \cap L_2$</p> | Understanding-Short_Answer_2_5_Min | | 05 | Medium | | | | | | |

| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|--|---|--|----|--------|--|--|--|--|--|--|
| 0170 1349 2 | 06 | 06 | 157 | Union, Intersec tion, Differ ence of FA | <p>Let M1 and M2 be the FAs pictured in Figure, recognizing languages L1 and L2 respectively.</p> <p>Draw FAs recognizing the following languages.</p> <p>a. $L1 \cup L2'$</p> <p>b. $L2 - L1$</p> | Evaluating- Long_Answ er_5_10_Mi n | 04 | Medium | | | | | | |
| 0170 1349 2 | 06 | 06 | 158 | Union, Intersec tion, Differ ence of FA | <p>Suppose that Languages L1 and L2 are the subsets given below. Where $\Sigma = \{0, 1\}$ $L1 = \{x \mid 00 \text{ is not a substring of } x\}$ $L2 = \{x \mid x \text{ ends with } 01\}$ Draw FAs recognizing the following languages (1) $L1 - L2$ (2) $L1 \cap L2$</p> | Understandi ng- Short_Answ er_2_5_Min | 05 | Medium | | | | | | |
| 0170 1349 2 | 06 | 06 | 159 | Union, Intersec tion, Differ ence of FA | <p>Suppose that L1 and L2 are the subsets:</p> | Understandi ng- Short_Answ er_2_5_Min | 05 | Medium | | | | | | |

| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|--|---|--|----|--------|--|--|--|--|--|--|
| | | | | |  <p>Draw FAs recognizing the following languages:</p> <ul style="list-style-type: none"> i. $L_1 \cup L_2$ ii. $L_1 \cap L_2$ iii. $L_1 - L_2$ iv. $L_1 \cap L_3$ v. $L_3 - L_2$ | | | | | | | | | |
| 0170 1349 2 | 06 | 06 | 162 | Union, Intersec tion, Differ ence of FA | <p>There are 2 languages over $\Sigma = \{a, b\}$</p> <p>L_1 = all strings with a double “a”</p> <p>L_2 = all strings with an even number of “a”</p> <p>Find a regular expression and an FA that define $L_1 \cap L_2$.</p> | Understandi ng- Short_Answ er_2_5_Min | 05 | Medium | | | | | | |
| 0170 1349 2 | 06 | 06 | 163 | Union, Intersec tion, Differ ence of FA | <p>Let M_1 and M_2 be the two FAs as given below.</p>  | Evaluating- Long_Answ er_5_10_Mi n | 05 | Medium | | | | | | |



Draw FA recognizing $(L_1 \cup L_2)$ and $(L_1 - L_2)$ where L_1 and L_2 correspond to M_1 and M_2 respectively.

0170
1349
2

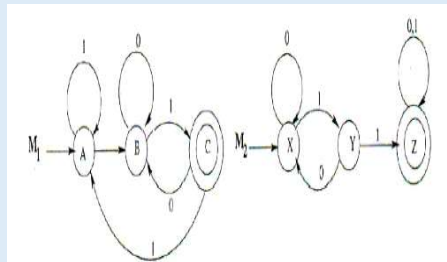
06

06

164

Union, Intersection, Difference of FA

Fig. shows two DFAs M1 and M2, to accept languages L1 and L2, respectively. Determine DFAs to recognize $L1 \cup L2$.



Understanding-
Short_Answer_2_5_Min

04

Medium

$$\begin{array}{r} 0170 \\ 1349 \\ \hline 2 \end{array}$$

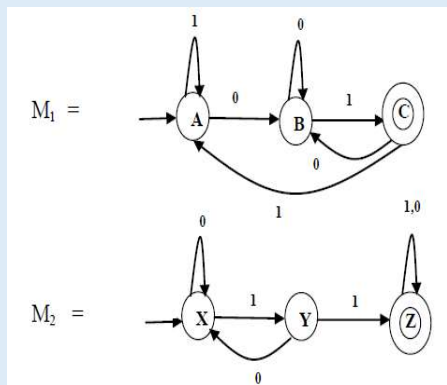
06

06

165

Union,
Intersec
tion,
Differe
nce of
FA

Let M_1 and M_2 be the
FAs pictured below,
recognizing languages
 L_1 and L_2 respectively



| |
|---|
| Evaluating- Long_Answ er_5_10_Mi n |
|---|

| |
|----|
| 04 |
|----|

Medium

| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|---------------------------------------|---|------------------------------------|--|----|--------|--|--|--|--|--|
| | | | | | Draw the FAs recognizing the following languages: $L_1 \cap L_2, L_2 - L_1$ | | | | | | | | | |
| 0170 1349 2 | 06 | 06 | 166 | Union, Intersection, Difference of FA | <p>L_1 and L_2 are two languages:</p> <p>$L_1 = \{x \mid 11 \text{ is not a substring of } x\}$</p> <p>$L_2 = \{x \mid x \text{ starts with 0 and ends with 0}\}$</p> <p>Draw FA for both L_1 and L_2 and construct FA for $L_3 = L_2 - L_1$</p> | Evaluating-Long_Answer_5_10_Min | | 04 | Medium | | | | | |
| 0170 1349 2 | 06 | 06 | 167 | Union, Intersection, Difference of FA | <p>Let M_1 and M_2 be the FAs pictured in Fig. (i) and Fig. (ii) accept the languages L_1 and L_2, respectively</p> <div style="text-align: center;"> </div> <p>Draw FAs accepting the following languages: (i) $L_1 \cup L_2$ (ii) L_2</p> | Understanding-Short_Answer_2_5_Min | | 04 | Medium | | | | | |
| 0170 1349 2 | 06 | 06 | 168 | Union, Intersection, Difference of FA | <p>Let FA_1 and FA_2 be the FAs as shown in the figure recognizing the languages L_1 and L_2 respectively. Draw an FA recognizing the language, $L_1 \cup L_2$.</p> <p>FA₁:</p> <div style="text-align: center;"> </div> | Understanding-Short_Answer_2_5_Min | | 04 | Medium | | | | | |

| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|----------------------|--|--|-------------------|----|--------|--|--|--|--|--|
| 0170 1349 2 | 06 | 06 | 175 | Pumpin g Lemma | <p>Is the following set regular</p> $\{0^{2n} \mid n \geq 1\}$ <p>If yes, write down the corresponding regular expression. Else, prove that the language is not regular.</p> <p>Answer: Yes. $(00)^+$</p> | Understandi ng- Short_Answ er_2_5_Min | Yes. (00) + | 01 | Medium | | | | | |
| 0170 1349 2 | 06 | 06 | 176 | Pumpin g Lemma | <p>Use the pumping lemma to show that following language is not regular.</p> <p>$L = \{xy \mid x, y \in \{0, 1\}^* \text{ and } y \text{ is either } x \text{ or } x^r\}.$</p> | Understandi ng- Short_Answ er_2_5_Min | | 03 | Medium | | | | | |
| 0170 1349 2 | 06 | 06 | 177 | Pumpin g Lemma | <p>Use the pumping lemma to show that following language is not regular: $L = \{ww \mid w \in \{0, 1\}^*\}.$</p> | Understandi ng- Short_Answ er_2_5_Min | | 04 | Medium | | | | | |
| 0170 1349 2 | 06 | 06 | 178 | Pumpin g Lemma | <p>What do you mean by Regular Language? Explain the application of the Pumping Lemma to show a Language is Regular or Not.</p> | Understandi ng- Short_Answ er_2_5_Min | | 03 | Medium | | | | | |
| 0170 1349 2 | 06 | 06 | 179 | Pumpin g Lemma | <p>Define Pumping Lemma for Regular Languages. Prove that the language $L = \{a^n \mid n \text{ is a prime number}\}$ is not regular.</p> | Understandi ng- Short_Answ er_2_5_Min | | 04 | Medium | | | | | |
| 0170 1349 2 | 06 | 06 | 180 | Pumpin g Lemma | <p>Define Pumping Lemma. Use the Pumping Lemma to show that the following languages are not regular:</p> <ul style="list-style-type: none"> $L = \{0^n 1 0^{2n} \mid n \geq 0\}$ $L = \{0^i 1^j 0^k \mid k > i+j\}$ | Understandi ng- Short_Answ er_2_5_Min | | 04 | Hard | | | | | |

| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|----------------------|--|--|--------------------|----|--------|--|--------------------|--------------------|-------------------|----------------------|
| 0170 1349 2 | 06 | 06 | 181 | Pumpin g Lemma | Use Pumping Lemma to show that $L = \{x \in \{0,1\}^* \mid x \text{ is a palindrome}\}$ is not a regular language. | Understandi ng- Short_Answ er_2_5_Min | | 04 | Medium | | | | | |
| 0170 1349 2 | 06 | 06 | 182 | Pumpin g Lemma | While applying Pumping lemma over a language, we consider a string w that belong to L and fragment it into _____ parts. | Understandi ng- Multiple_Ch oice | 3 | 01 | Easy | | 3 | 5 | 2 | Non e of these |
| 0170 1349 2 | 06 | 06 | 183 | Pumpin g Lemma | If we select a string w such that $w \in L$, and $w=xyz$. Which of the following portions cannot be an empty string? | Understandi ng- Multiple_Ch oice | y | 01 | Easy | | x | y | z | Non e of these |
| 0170 1349 2 | 06 | 06 | 184 | Pumpin g Lemma | Let $w=xyz$ and y refers to the middle portion and $ y >0$. What do we call the process of repeating y 0 or more times before checking that they still belong to the language L or not? | Understandi ng- Multiple_Ch oice | Pum ping | 01 | Easy | | Gen erati ng | Pum ping | Prod ucin g | Non e of these |
| 0170 1349 2 | 06 | 06 | 185 | Pumpin g Lemma | here exists a language L. We define a string w such that $w \in L$ and $w=xyz$ and $ w \geq n$ for some constant integer n. What can be the maximum length of the substring xy i.e. $ xy \leq ?$ | Understandi ng- Multiple_Ch oice | n | 01 | Easy | | n | $ y $ | $ x $ | Non e of these |
| 0170 1349 2 | 06 | 06 | 186 | Pumpin g Lemma | Fill in the blank in terms of p, where p is the maximum string length in L. Statement: Finite languages trivially satisfy the pumping lemma by having $n =$ _____ | Understandi ng- Multiple_Ch oice | $p+1$ | 01 | Easy | | p^*1 | $p+1$ | $p-1$ | Non e of these |
| 0170 1349 2 | 06 | 06 | 187 | Pumpin g Lemma | Answer in accordance to the third and last statement in pumping lemma: For all _____ $xy^iz \in L$ | Understandi ng- Multiple_Ch oice | $i \geq 0$ | 01 | Easy | | $i > 0$ | $i < 0$ | $i \leq 0$ | Non e of these |
| 0170 1349 2 | 06 | 06 | 188 | Pumpin g Lemma | Which of the following one can relate to the given statement: Statement: If n items are put into | Understandi ng- | Pige on Hole | 01 | Easy | | Pum ping | Pige on Hole | Coun t | Non e of these |

| | | | | | | | | | | | | | | |
|-------------------|----|----|-----|--|--|---|---|----|------|--|---|--------------------------------|---|----------------------|
| | | | | | m containers, with $n > m$, then atleast one container must contain more than one item. | Multiple_Ch oice | prin ciple | | | | lem ma | prin ciple | princ iple | |
| 0170 1349 2 | 06 | 06 | 189 | Pumpin g Lemma | <p>f d is a final state, which of the following is correct according to the given diagram?</p> <pre> graph LR start(()) --> a((a)) a -- p --> b((b)) b -- q --> c((c)) c -- r --> b b -- s --> d(((d))) </pre> | Understandi ng- Multiple_Ch oice | x=p, y=qr , z=s | 01 | Easy | | x=p, y=qr , z=s | x=p, z=qr s | x=pr , y=r, z=s | Non e of these |
| 0170 1349 2 | 06 | 06 | 190 | Pumpin g Lemma | Relate the following statement: Statement: All sufficiently long words in a regular language can have a middle section of words repeated a number of times to produce a new word which also lies within the same language. | Understandi ng- Multiple_Ch oice | Pum ping Lem ma | 01 | Easy | | Turi ng Mac hine | Pum ping Lem ma | Arde n's theor em | Non e of these |
| 0170 1349 2 | 06 | 06 | 191 | Union, Intersec tion, Differe nce of FA | Regular sets are closed under union,concatenation and kleene closure. | Understandi ng- Multiple_Ch oice | True | 01 | Easy | | True | Fals e | Depe nds on regul ar set | Non e of these |
| 0170 1349 2 | 06 | 06 | 192 | Union, Intersec tion, Differe nce of FA | Explanation: Regular sets are closed under these three operation. | Understandi ng- Multiple_Ch oice | maki ng final state s non- final and non- final to final | 01 | Easy | | maki ng start ing state as final state | no triva l met hod | maki ng final state s non- final and non- final to final | Non e of these |

