

Turing machine

$$TM = M = (Q, \Sigma, \Gamma, \delta, q_0, F, B) \quad (\text{7-tuple})$$

Q = set of states

Σ = input alphabet

Γ = tape alphabet (symbols)

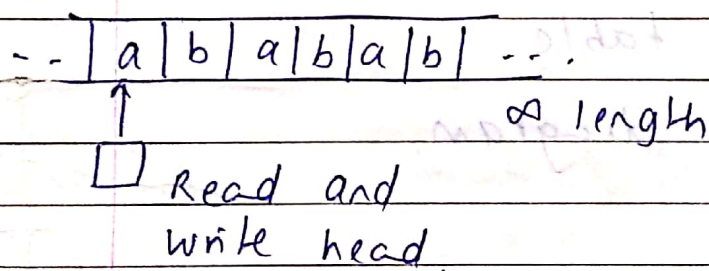
δ : transition function

$$Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R, N\}$$

q_0 = initial state

F = set of final states $\subseteq Q$

B = blank symbol.



e.g.

$$(q_0) \xrightarrow{a/a, R} (q_1)$$

$$\delta(q_0, a) \rightarrow (q_1, a, R)$$

$$(q_0) \xrightarrow{a/x, R} (q_1)$$

$$\delta(q_0, a) \rightarrow (q_1, x, R)$$

(FA/PDA) don't have control over their input.

TM can change its own input.

↳ writing m/c

Tm is most powerful m/c

Any problem that can have
an algorithm and can be solved
can also be solved by Tm.

same power of digital comp.

abstract / mathematical model of
digital comp.

comparision of power

$$\boxed{FA} < \boxed{PDM} < \boxed{Tm}$$

Representation of Tm

(1) Instantaneous Description
(ID)

(2) Transition table

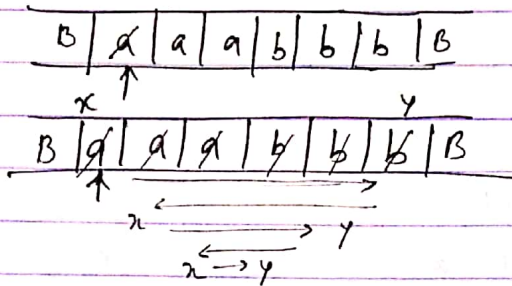
(3) Transition diagram.

After processing a^3b^3 mlc reached to halt state

② T_m accepts strings of $a^n b^n$ $n \geq 1$ rejects other all.

Step I ie $a^3 b^3$

logic



① steps - left most $a \rightarrow x$

② right most $b \rightarrow y$

③ head come back to first a .

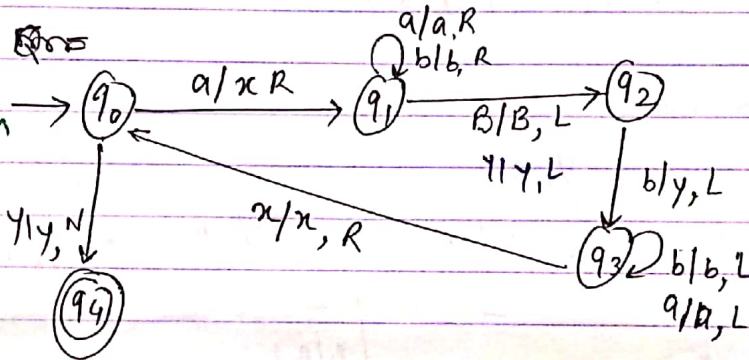
Step II

defn of T_m

$T_m M = (Q, \Sigma, \Gamma, \delta, q_0, B, F)$

Step III

Transition Diagram



$Q = \{q_0, q_1, q_2, q_3, q_4\}$

$q_0 =$ initial state

$\Sigma = \{a, b\}$

$B =$ Blank symbol.

$\Gamma = \{a, b, B\}$
 x, y

$F = \{q_4\}$ Halt state

Step IV

δ is as shown

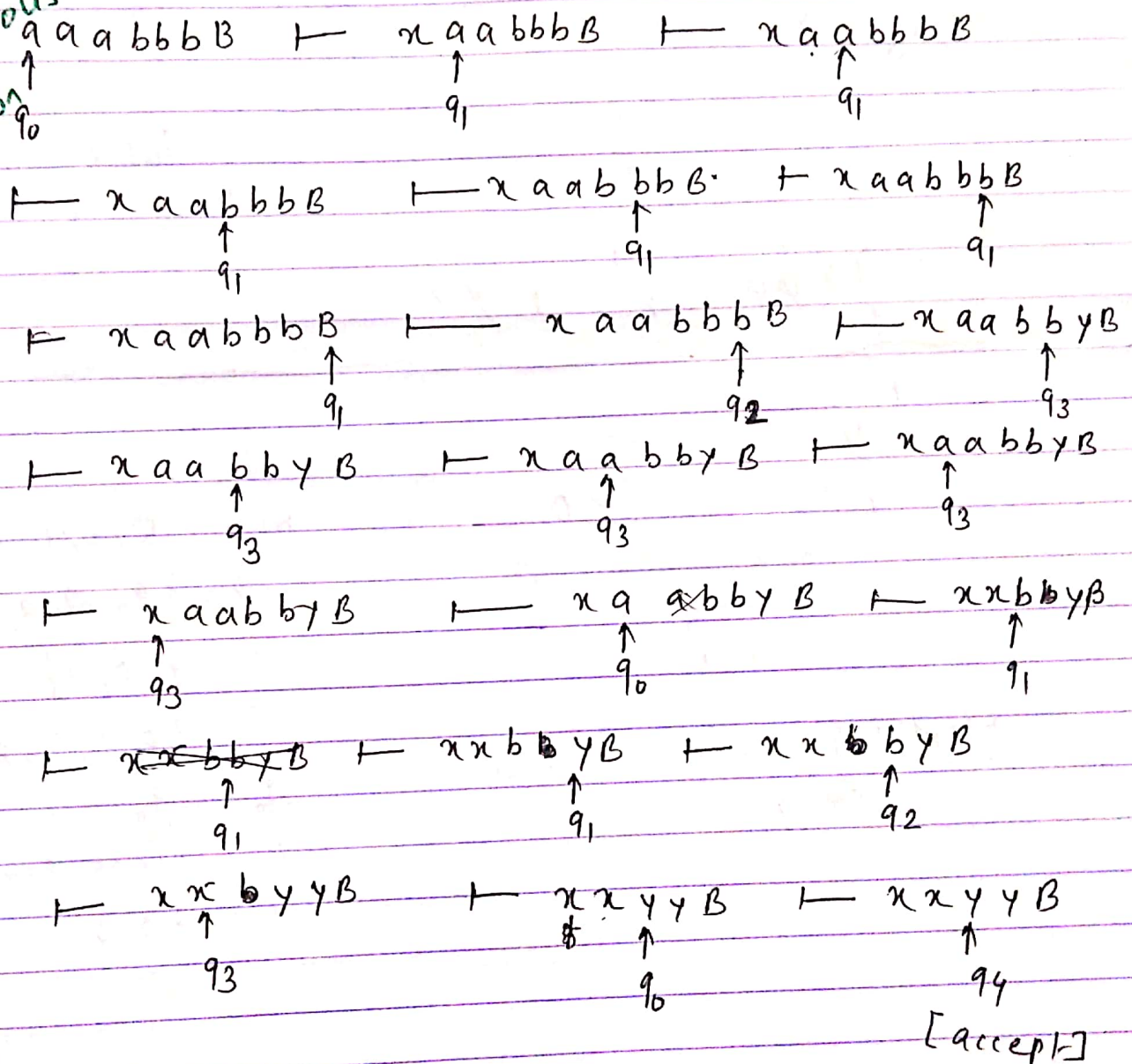
Transition table

	a	b	x	y	B
q_0	$xR q_1$	—	—	$yN q_4$	—
q_1	$aR q_1$	$bR q_1$	—	—	$BL q_2$
q_2	—	$yL q_3$	—	—	—
q_3	$aL q_3$	$bL q_3$	$xR q_0$	—	—
(q_4)	—	—	—	—	—

Halt state

STEP V

Instantaneous
Description



STEP VI

Final
Ans.

$$M = (Q, \Sigma, \Gamma, \delta, q_0, B, F)$$

$$Q = \{q_0, q_1, q_2, q_3, q_4\} \quad \Sigma = \{a, b\} \quad \Gamma = \{a, b, x, y, B\}$$

$$\delta: \text{as per diagram} \quad q_0 = q_0 \quad F = \{q_4\}$$

$B = \text{Blank symbol.}$