

Greibach Normal Form (GNF)

A CFG $G = (V, T, P, S)$ is said to be in GNF if every production is of the form $A \rightarrow a \alpha$ where $\alpha \in V^*$ [string of zero or more no. of variables]

- ① Lang. should be without ϵ production rules. & simplified.
- ② RHS should start with one terminal symbol followed by zero or more number of variables.

Steps - ① simplify Grammar.

② Bring every production to form $A \rightarrow a \alpha$ or $A \rightarrow \alpha$ where $\alpha \in V^*$

③ Rename variables A_1, A_2, \dots

④ $A_i \rightarrow A_j \alpha$ rule for every $i > j$ convert it to $i \leq j$ by substitution.

⑤ Remove left-recursion.

⑥ get GNF productions by substituting already GNF production.

Q - Find a GNF equivalent of following CFG.

$$\begin{array}{l} S \rightarrow BA \mid ab \\ B \rightarrow AB \mid a \\ A \rightarrow Bb \mid BB \end{array}$$

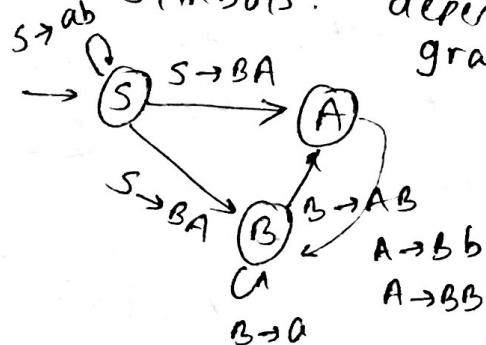
① step 1 - Explain what is GNF.

② step 2 - simplify the grammar.

Grammar does not contain ϵ -production and unit productions.

Generating symbols are (a, b)

To find non reachable symbols. dependency graph.



$B \rightarrow a$
 $S \rightarrow ab$
 (a, b, S, B) are gener.

$A \rightarrow Bb$
 (a, b, S, B, A) all gener.

all the symbols are reachable from S .

No useless symbol is there.

③ step 3 - ~~Rename~~ var bring every production to form $A \rightarrow a\alpha$ or $A \rightarrow \alpha$

where $|\alpha| \geq 2$ then α should contain all variable.

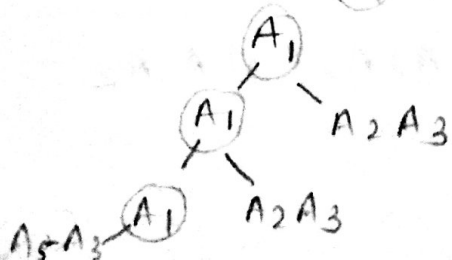
$A \rightarrow Bb \Rightarrow$ convert to $A \rightarrow \alpha$ form by adding $D \rightarrow b$

$$A_4 \rightarrow b$$

② $A_3 \rightarrow A_2$ substitute
Value of A_2

rewrite

$$\frac{A_5 A_3 (A_2 A_3)^n}{B_1}, n \geq 0$$



⑥ step 6 - Remove left recursion

$$A_1 \rightarrow A_2 A_3 \mid a A_4$$

$$A_2 \rightarrow A_3 A_2 \mid a$$

$$A_3 \rightarrow \underline{A_3} A_2 A_4 \mid a A_4 \mid \underline{A_3} A_2 A_2 \mid a A_2$$

$$A_4 \rightarrow b$$

A_3 production has left recursion.

so remove it by adding new variable B_3 .

For $A_3 \rightarrow A_3 A_2 A_4$ →

$A_3 \rightarrow a A_4 \mid a A_2 \mid$

either $a A_4$ or $a A_2$

$\underline{a A_4} (A_2 A_4)^n$
or $\underline{B_3}$
 $\underline{a A_2} (A_2 A_4)^n$
 $\underline{B_3}$

For $\underline{A_3} \rightarrow \underline{A_3} A_2 A_2$

strings of type $\underline{a A_4} (A_2 A_2)^n$
or $\underline{a A_2} (\underline{A_2 A_2})^n$
 $\underline{B_3}$

either $a A_4$ or $a A_2$

$$A_3 \rightarrow a A_4 \mid a A_2 \mid a A_4 B_3 \mid a A_2 B_3$$

$$B_3 \rightarrow A_2 A_4 B_3 \mid A_2 A_4 \mid A_2 A_2 B_3 \mid A_2 A_2$$

Step 7 - convert all productions to CNF.

A_4, A_3 is already in CNF.

A_1, A_2, B_3 can be converted to CNF.

$\checkmark A_4 \rightarrow b$
 $\checkmark A_3 \rightarrow aA_4 \mid aA_2 \mid aA_4B_3 \mid aA_2B_3$ } already in CNF.

$B_3 \rightarrow A_2A_4B_3 \mid A_2A_2B_3 \mid A_2A_4 \mid A_2A_2$.

$\checkmark A_2 \rightarrow aA_4A_2 \mid aA_2A_2 \mid aA_4B_3A_2 \mid aA_2B_3A_2 \mid a$

$A_1 \rightarrow A_2A_3 \mid aA_4$
→ rewrite it as

$\checkmark A_1 \rightarrow aA_4A_2A_3 \mid aA_2A_2A_3 \mid aA_4B_3A_2A_3 \mid$
 $aA_2B_3A_2A_3 \mid aA_3 \mid aA_4$

$B_3 \rightarrow A_2A_4B_3 \mid A_2A_2B_3 \mid A_2A_4 \mid A_2(A_2)$
→ rewrite B_3 by substituting A_2

$\checkmark B_3 \rightarrow aA_4A_2A_4B_3 \mid aA_2A_2A_4B_3 \mid aA_4B_3A_2B_3 \mid$
 $aA_2B_3A_2A_4B_3 \mid aA_4B_3 \mid aA_4A_2A_2B_3 \mid$
 $aA_2A_2A_2B_3 \mid aA_4B_3A_2A_2B_3 \mid aA_2B_3A_2A_2$
 $\mid aA_2B_3 \mid aA_4A_2A_2 \mid aA_2A_2A_2 \mid aA_4B_3A_2A_2$
 $aA_2B_3A_2A_2 \mid aA_2 \mid aA_4A_2A_2 \mid aA_2A_2(A_2) \mid$
 $aA_4B_3A_2(A_2) \mid aA_2B_3A_2(A_2) \mid a(A_2)$

20 terms

Step 8 - write CNF Grammar rules.

$A_4 \rightarrow b$

A_3, A_2, A_1 and B_3 .