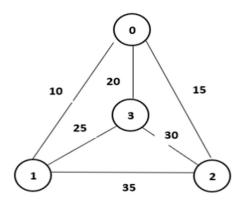
## **EXPERIMENT 9**

# 20CP209P - Design and Analysis of Algorithm Lab

#### Aim:

Given a set of cities and distance between every pair of cities, the problem is to find the shortest possible tour that visits every city exactly once and returns to the starting point. Solve this problem using branch and bound technique.

For example, consider the following graph:



A Travelling Salesman Problem (TSP) tour in the graph is 0 - 1 - 3 - 2 - 0. The cost of the tour is 10 + 25 + 30 + 15 = 80

#### Code:

## **Travelling Salesman Problem (Branch and Bound):**

```
#include <iostream>
#include <vector>
#include <limits> // Use <limits> for numeric_limits
#include <algorithm>
#include <numeric>
#include <vector>
#include <queue> // For priority queue

#define fr(i, a, b) for (int i = a; i < b; i++)
#define pii pair<int,int>

using namespace std;

// Use numeric_limits for infinity for better type safety and clarity const int INF = numeric_limits<int>:::max();

// Structure to represent a node in the state-space tree struct TSPNode {
```

```
// Path taken so far
  vector<int> path;
  vector<bool> visited; // Visited cities marker
                   // Cost of the path so far
  int cost;
  int lower bound;
                        // Estimated lower bound for the total tour cost
                   // Number of cities visited (path.size())
  int level;
  // Constructor
  TSPNode(int n): visited(n, false), cost(0), lower_bound(0), level(0) {}
  // Custom comparator for the priority queue (min-heap based on lower bound)
  bool operator>(const TSPNode& other) const {
    return lower bound > other.lower bound;
  }
};
// Function to calculate the lower bound for a given node
// A simple lower bound: current cost + sum of minimum outgoing edge from each unvisited city
int calculate lower bound(const TSPNode& node, int n, const vector<vector<int>>& graph)
{
  int bound = node.cost;
  int last city = node.path.back();
  for (int i = 0; i < n; ++i)
  {
    // If city 'i' is unvisited OR it's the starting city and we need to return
    if (!node.visited[i])
       // Find the minimum cost edge leaving city 'i' to any *other* city
       int min edge = INF;
       // If it's the last node in the path, find min edge to unvisited or start
       if (i == last city)
         for (int j = 0; j < n; ++j)
         {
            // Must go to an unvisited city OR back to start if it's the last hop
            if (j == node.path[0] \&\& node.level == n)
            { // Check if it's the last step
               if (graph[i][j] != INF)
               {
                 min_edge = min(min_edge, graph[i][j]);
              }
            else if (!node.visited[j] && i != j && graph[i][j] != INF)
              min_edge = min(min_edge, graph[i][j]);
            }
         }
       }
       else
       { // For other unvisited cities, find the absolute minimum outgoing edge
```

```
for(int j=0; j<n; ++j)
        if (i != j && graph[i][j] != INF)
           min_edge = min(min_edge, graph[i][j]);
        }
      }
    }
    // If an unvisited city has no way out (shouldn't happen in complete graph)
    // or if it's the last node with no path back to start
    if (min edge == INF)
      return INF; // This path is infeasible or bound calculation failed
    bound += min edge;
 }
// A slightly simpler version (often used, might be less tight but still valid):
// Sum minimum edge cost for all unvisited nodes, regardless of where they go
int simpler_bound = node.cost;
for(int i=0; i<n; ++i) {
  if(!node.visited[i]) {
    int min edge = INF;
    for(int j=0; j<n; ++j) {
       if(i != j && graph[i][j] != INF) {
         min_edge = min(min_edge, graph[i][j]);
    if(min_edge == INF) return INF; // Infeasible
    simpler_bound += min_edge;
  }
// Add cost from last visited node back to start if it's the last leg needed
if (node.level == n - 1 \&\& !node.visited[0]) { // Check if only start is left}
  int last = node.path.back();
  if(graph[last][0] != INF) {
     simpler_bound += graph[last][0]; // Already partially counted, tricky
  } else {
     return INF; // No path back to start
  }
}
return simpler bound; // Let's use the simpler bound for clarity
*/
// Let's stick to the slightly more accurate one calculated first.
// If the path is nearly complete, the last leg must go back to start
if (node.level == n)
{
```

```
int return cost = graph[last city][node.path[0]];
     if (return_cost == INF) return INF; // No path back
     // The bound IS the final cost here
     return node.cost + return_cost;
   }
   return bound;
}
// --- Simplified Lower Bound Implementation ---
// Calculate the lower bound: cost so far + sum of min edges from unvisited nodes
int calculate simplified lower bound(int current cost, int current city, int n, const vector<bool>&
visited, const vector<vector<int>>& graph)
{
  int bound = current_cost;
  for (int i = 0; i < n; ++i)
     // Consider unvisited cities
     if (!visited[i])
       int min_val = INF;
       for (int j = 0; j < n; ++j)
       {
         // Find the cheapest edge leaving this unvisited city 'i'
         // The destination 'j' can be any other city (visited or unvisited)
         // or the starting city if 'i' is the last city visited in a partial tour
         if (i != j && graph[i][j] != INF)
         {
           // Simple version: min edge to *any* other node
           min_val = min(min_val, graph[i][j]);
         }
       if(min val == INF) return INF; // This city is stranded
       bound += min val;
    }
  }
   // Special handling for the edge FROM the current_city
   // Find the minimum edge from current city to an UNVISITED city
   int min from current = INF;
   bool can_reach_unvisited = false;
   for(int j = 0; j < n; ++j) {
    if (!visited[j] && graph[current_city][j] != INF)
       min_from_current = min(min_from_current, graph[current_city][j]);
       can reach unvisited = true;
   }
```

```
// If we can't reach any unvisited node from current city (and tour not complete)
  // it might mean we only need to go back to start
  if (!can_reach_unvisited && count(visited.begin(), visited.end(), false) == 1 && !visited[0])
    if (graph[current_city][0] != INF)
      // The only unvisited node is the start node
      // The lower bound should include the cost back to start
      // Note: The loop above already added min cost *from* start (if unvisited)
      // Let's refine the loop logic slightly:
      // bound = current cost
      // sum min outgoing edge for all nodes k (!visited)
      // sum min incoming edge for all nodes k (!visited)
      // return bound = current_cost + (sum_min_out + sum_min_in) / 2 (Held-Karp idea - complex)
      // Stick to simpler: Sum of min outgoing from unvisited
      // Let's recalculate simpler bound:
      bound = current cost;
      for(int i=0; i<n; ++i)
        // For the current node, find min edge to UNVISITED node
        if(i == current_city)
           int min edge curr = INF;
           for(int j=0; j<n; ++j)
             if(!visited[j] && graph[i][j] != INF)
             {
                min_edge_curr = min(min_edge_curr, graph[i][j]);
             }
           }
           // If only start node is left unvisited
           if (min_edge_curr == INF && count(visited.begin(), visited.end(), false) == 1 &&
!visited[0])
              if(graph[i][0] != INF)
                min_edge_curr = graph[i][0];
              else
              {
                return INF; // Cannot return to start
           else if (min edge curr == INF && count(visited.begin(), visited.end(), false) > 0)
              return INF; // Cannot reach any other unvisited node
           }
```

```
if(min edge curr != INF) bound += min edge curr; // Add cost from current node
onwards
         }
         // For OTHER unvisited nodes, find the absolute minimum outgoing edge
         else if (!visited[i])
            int min_edge_other = INF;
            for(int j=0; j<n; ++j)
               if(i != j && graph[i][j] != INF)
                 min_edge_other = min(min_edge_other, graph[i][j]);
            if(min_edge_other == INF) return INF; // Stranded node
            bound += min edge other;
      return bound;
     }
     else
       return INF; // Can't get back to start
     }
  }
  return bound; // Return the calculated simple bound
}
int main(void)
  // Example Graph (Cost Matrix)
  // graph[i][j] = cost from city i to city j
  // Use INF if no direct path (or for diagonal i == j)
   vector<vector<int>> graph = {
    {INF, 10, 8, 9, 7},
    {10, INF, 10, 5, 6},
    {8, 10, INF, 8, 9},
    {9, 5, 8, INF, 6},
    {7, 6, 9, 6, INF}
  int n = graph.size(); // Number of cities
  // Use graph2 for testing
  // vector<vector<int>> graph2 = {
  // {INF, 10, 15, 20},
  // {5, INF, 9, 10},
      {6, 13, INF, 12}, // Changed 9 to 13 based on common examples for cost 35
```

```
// {8, 8, 9, INF}
// };
// vector<vector<int>> graph = graph2; // Uncomment to use graph2
// n = graph.size();
// Priority Queue for B&B nodes (min-heap based on lower bound)
priority queue<TSPNode, vector<TSPNode>, greater<TSPNode>> pq;
// Global minimum cost found so far
int min cost = INF;
vector<int> final path;
// Create the root node (starting at city 0)
TSPNode root(n);
root.level = 1;
root.path.push_back(0); // Start at city 0
root.visited[0] = true;
root.cost = 0;
// Calculate initial lower bound (sum of min edges from all nodes) - Adjusted calculation
root.lower_bound = calculate_simplified_lower_bound(root.cost, 0, n, root.visited, graph);
// root.lower bound = 0; // Initial cost is 0
// for(int i=0; i<n; ++i) {
// int min_row = INF;
    for (int j=0; j<n; ++j) {
//
//
       if (i != j && graph[i][j] != INF) {
         min_row = min(min_row, graph[i][j]);
//
      }
// }
     if(min_row == INF && n > 1) { // Check for infeasible graphs early
       cout << "Graph is not strongly connected or has isolated nodes." << endl;</pre>
//
//
       return 1;
//
     if (min_row != INF) root.lower_bound += min_row;
//
//}
if (root.lower_bound != INF) {
  pq.push(root);
} else if (n > 0) {
   cout << "Initial lower bound is INF. Problem might be infeasible." << endl;
}
cout << "Starting Branch and Bound TSP..." << endl;
if (n <= 1) {
  cout << "Minimum Cost: 0" << endl;
   cout << "Path: 0" << endl;
  return 0;
}
// Branch and Bound main loop
while (!pq.empty()) {
```

```
// Get the node with the lowest lower bound
TSPNode current node = pq.top();
pq.pop();
// --- Pruning Step ---
// If the current node's lower bound is already worse than the best solution found, prune it.
if (current node.lower bound >= min cost) {
  // cout << "Pruning node. Lower bound (" << current node.lower bound
  // << ") >= min_cost (" << min_cost << ")" << endl;
  continue;
// --- Goal Check ---
// If all cities have been visited (path length = n)
if (current node.level == n) {
  // We need to add the cost of returning to the starting city (city 0)
  int last city = current node.path.back();
  int return cost = graph[last city][0]; // Cost from last city back to start
  if (return cost != INF) {
    int total cost = current node.cost + return cost;
    // Found a new best solution
    if (total cost < min cost) {
       min cost = total cost;
      final_path = current_node.path;
      // Add the starting city to the end to show the cycle
      final path.push back(0);
       cout << "Found new best solution: Cost = " << min cost << ", Path = ";
       for(int city: final path) cout << city << " "; cout << endl;
    }
  // No further branching needed from a complete tour node
  continue;
}
// --- Branching Step ---
// Explore neighbors (unvisited cities)
int current city = current node.path.back();
for (int next_city = 0; next_city < n; ++next_city) {
  // If the next city is not visited and there's an edge
  if (!current_node.visited[next_city] && graph[current_city][next_city] != INF)
    // Create a child node representing the extended path
    TSPNode child node = current node; // Copy parent state
    child node.level++;
    child node.path.push back(next city);
    child node.visited[next city] = true;
    child_node.cost = current_node.cost + graph[current_city][next_city];
```

```
// Calculate the lower bound for the child node
        child_node.lower_bound = calculate_simplified_lower_bound(child_node.cost, next_city, n,
child_node.visited, graph);
         // cout << " Exploring edge " << current_city << "->" << next_city
              << ", Cost: " << child node.cost
         // << ", Lower Bound: " << child node.lower bound << endl;
        // --- Pruning before adding to queue ---
        // Only add the child to the queue if its bound is promising
        if (child node.lower bound < min cost) {
           pq.push(child_node);
        } else {
          // cout << " Pruning child path (bound " << child_node.lower_bound << " >= min_cost "
<< min_cost << ")" << endl;
        }
      }
    }
  // Output the result
  if (min_cost == INF) {
    cout << "\nNo feasible solution found." << endl;</pre>
    cout << "\n-----" << endl;
    cout << "Optimal Minimum Cost: " << min_cost << endl;</pre>
    cout << "Optimal Path: ";
    for (int i = 0; i < final path.size(); ++i) {
      cout << final path[i] << (i == final path.size() - 1 ? "" : " -> ");
    cout << endl;
    cout << "-----" << endl;
  }
  return 0;
}
```

# **Output:**

```
PS 8:\sem4\23bcp153_daa\lab9> g++ tspbnbgem.cpp -o tspbnbgem
PS 8:\sem4\23bcp153_daa\lab9> ./tspbnbgem
Starting Branch and Bound TSP...
Found new best solution: Cost = 34, Path = 0 4 1 3 2 0

Optimal Minimum Cost: 34
Optimal Path: 0 -> 4 -> 1 -> 3 -> 2 -> 0

PS 8:\sem4\23bcp153_daa\lab9> [
```