

Minimization of DFA

One important result on finite automata, both theoretically and practically, is that for any regular language there is a unique DFA having the smallest number of states that accepts it. Let $M = \langle Q, \Sigma, q_0, \delta, A \rangle$ be a DFA that accepts a language L . Then the following algorithm produces the DFA, denote it by M_1 , that has the smallest number of states among the DFAs that accept L .

Minimization Algorithm for DFA

Construct a partition $\Pi = \{ A, Q - A \}$ of the set of states Q ;

$\Pi_{\text{new}} := \text{new_partition}(\Pi)$;

while ($\Pi_{\text{new}} \neq \Pi$)

$\Pi := \Pi_{\text{new}}$;

$\Pi_{\text{new}} := \text{new_partition}(\Pi)$

$\Pi_{\text{final}} := \Pi$;

function **new_partition**(Π)

for each set S of Π **do**

partition S into subsets such that two states p and q of S are in the same subset of S

if and only if for each input symbol, p and q make a transition to (states of) the same set of Π .

The subsets thus formed are sets of the output partition in place of S .

If S is not partitioned in this process, S remains in the output partition.

end

Minimum DFA M_1 is constructed from Π_{final} as follows:

- Select one state in each set of the partition Π_{final} as the representative for the set. These representatives are states of minimum DFA M_1 .

- Let p and q be representatives i.e. states of minimum DFA M_1 . Let us also denote by p and q the sets of states of the original DFA M represented by p and q , respectively. Let s be a state in p and t a state in q . If a transition from s to t on symbol a exists in M , then the minimum DFA M_1 has a transition from p to q on symbol a .
- The start state of M_1 is the representative which contains the start state of M .
- The accepting states of M_1 are representatives that are in A .

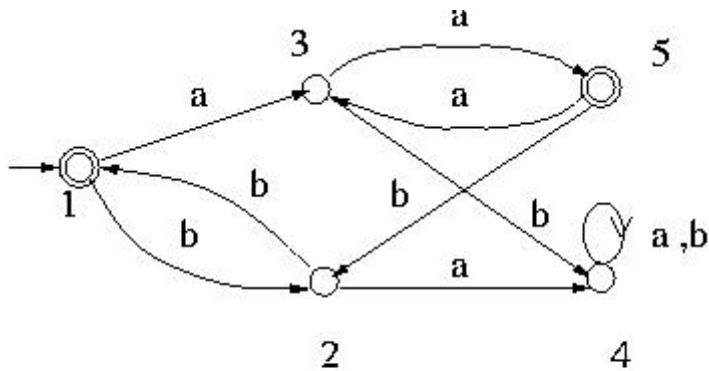
Note that the sets of Π_{final} are either a subset of A or disjoint from A .

Remove from M_1 the dead states and the states not reachable from the start state, if there are any. Any transitions to a dead state become undefined.

A state is a **dead state** if it is not an accepting state and has no out-going transitions except to itself.

Example 1 :

Let us try to minimize the number of states of the following DFA.



DFA Example 1

Initially $\Pi = \{ \{ 1, 5 \}, \{ 2, 3, 4 \} \}$.

new_partition is applied to Π .

Since on b state 2 goes to state 1, state 3 goes to state 4 and 1 and 4 are in different sets in Π , states 2 and 3 are going to be separated from each other in Π_{new} .

Also since on a state 4 goes to state 4, state 3 goes to state 5 and 4 and 5 are in different sets in Π , states 3 and 4 are going to be separated from each other in Π_{new} .

Further, since on b 2 goes to 1, 4 goes to 4 and 1 and 4 are in different sets in Π , 2 and 4 are separated from each other in Π_{new} .

On the other hand 1 and 5 make the same transitions. So they are not going to be split.

Thus the new partition is $\{ \{ 1, 5 \}, \{ 2 \}, \{ 3 \}, \{ 4 \} \}$. This becomes the Π in the second iteration.

When new_partition is applied to this new Π , since 1 and 5 do the same transitions, Π remains unchanged.

Thus $\Pi_{\text{final}} = \{ \{ 1, 5 \}, \{ 2 \}, \{ 3 \}, \{ 4 \} \}$.

Select 1 as the representative for $\{ 1, 5 \}$. Since the rest are singletons, they have the obvious representatives.

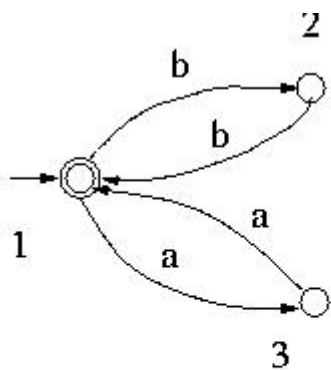
Note here that state 4 is a dead state because the only transition out of it is to itself.

Thus the set of **states** for the minimized DFA is $\{ 1, 2, 3 \}$.

For the **transitions**, since 1 goes to 3 on a, and to 2 on b in the original DFA, in the minimized DFA transitions are added from 1 to 3 on a, and 1 to 2 on b. Also since 2 goes to 1 on b, and 3 goes to 1 on a in the original DFA, in the minimized DFA transitions are added from 2 to 1 on b, and from 3 to 1 on a.

Since the rest of the states are singletons, all transitions between them are inherited for the minimized DFA.

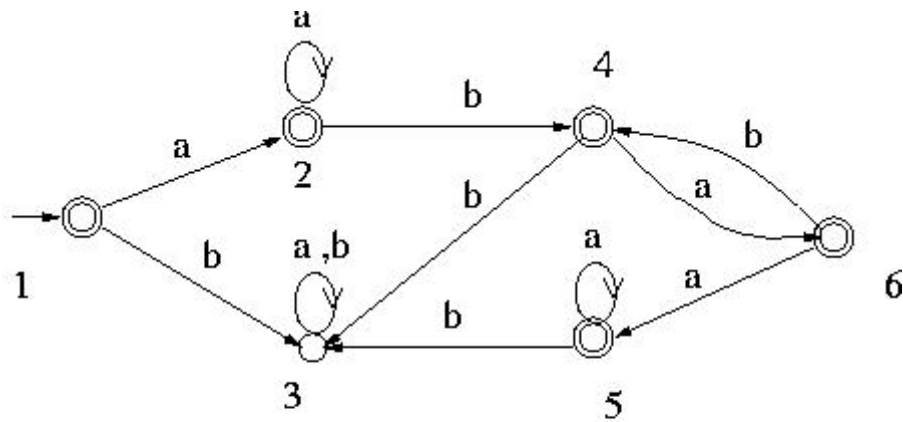
Thus the minimized DFA is as given in the following figure:



Minimized DFA

Example 2 :

Let us try to minimize the number of states of the following DFA.



DFA Example 2

Initially $\Pi = \{ \{ 3 \}, \{ 1, 2, 4, 5, 6 \} \}$.

By applying new_partition to this Π , $\Pi_{\text{new}} = \{ \{ 3 \}, \{ 1, 4, 5 \}, \{ 2, 6 \} \}$ is obtained.

Applying new_partition to this Π , $\Pi_{\text{new}} = \{ \{ 3 \}, \{ 1, 4 \}, \{ 5 \}, \{ 2 \}, \{ 6 \} \}$ is obtained.

Applying new_partition again, $\Pi_{\text{new}} = \{ \{ 1 \}, \{ 2 \}, \{ 3 \}, \{ 4 \}, \{ 5 \}, \{ 6 \} \}$ is obtained.

Thus the number of states of the given DFA is already minimum and it can not be reduced any further.

Test Your Understanding of Minimization of DFA

Indicate which of the following statements are correct and which are not.
Click True or Fals , then Submit.