

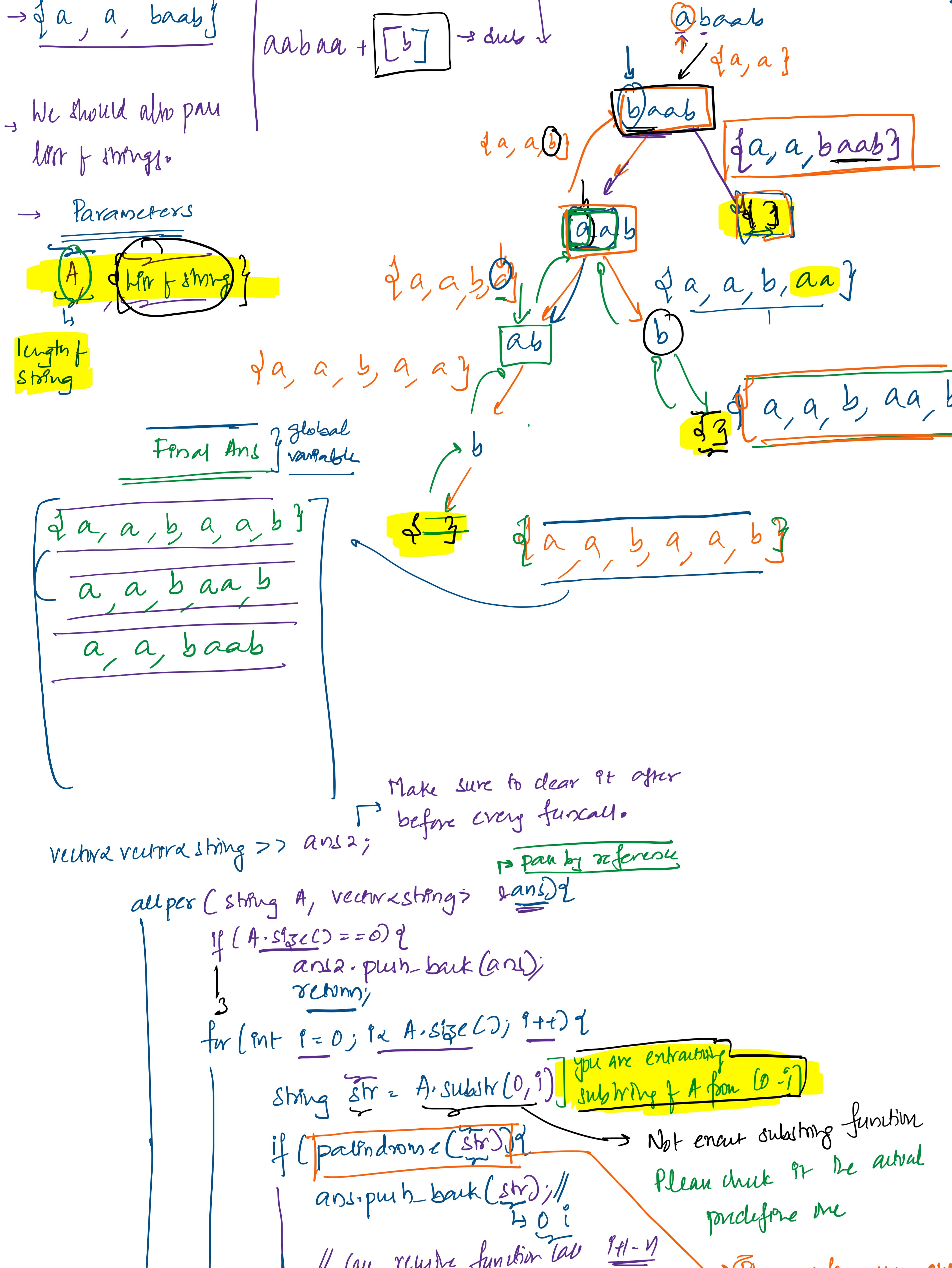
- 9) SISTER ✓
- 10) Palindrome Partitioning
- 11) Combinations Sum
- 12) Remove Invalid Parenthesis → YES
- 13) Max Consecutive ones
- 14) Max Rectangular area
- 15) Quick Select
- 16) Count Divers/SquareSum

Ques: Number of Subsets of size B having ≤ 1000
 Given N array elements find no. of subsets having ≤ 1000

Time: Generate all subset sums & check ≤ 1000 → If ($\text{sum} > 1000$) return 0;
 if ($\text{cnt} == B$) { return sum ≤ 1000 ; }

Backtracking: How many functions? // parameters → 2
 arr[], i, N, sum, cnt, B
 array p[i] index N number of subsets requested subset B
 return

l
 1001 - 2
 1001
 q-2
 q 1001-2y
 }
 If i == N {
 If sum ≤ 1000 & cnt == B {
 return 1;
 return 0;
 }
 func(arr, i+1, N, sum + arr[i], cnt+1, B) +
 func(arr, i+1, N, sum, cnt+1, B)



$$T(n) = T(n-1) + T(n-2) \dots T(1)$$

```

allper(string A, vector<string> &ans) {
    if (A.size() == 0) {
        ans.push_back("");
        return;
    }
    for (int i = 0; i < A.size(); i++) {
        string str = A.substr(0, i);
        if (palindrome(str)) {
            ans.push_back(str);
            if (i < A.size() - 1) {
                string str2 = A.substr(i + 1, A.size() - i - 1);
                allper(str2, ans);
            }
            ans.pop_back();
        }
    }
}
    
```

- Ques: Given a parenthesis at max we can remove B, $\underline{\underline{C}}$
 get all balanced parentheses of max length you can get
- Constraints
 len = 20 → Total Subsets $\approx 10^6$
- Important Constraints
 → Max Iterations = $\frac{10^6}{10^7}$
- len = 2
 → () a ()
 → () a ()
 → Note: If we are at a i^{th} character which is not '(' or ')' in that case you cannot remove that character.
- Global Variables:
 len of string
 vector<string> ans;
 Set<string> ans2;
 No1:
- BruteForce:
 Simply generate all subsets & check if balanced or not
- Parameters:
 A, i, cdif, temp str, ans
 current string
 current index position
 Diff of count
 + nos of open -
 nos of closed
 Ans string where we are constructing our balanced parentheses
- 4 steps:
- 1) Use backtracking we are generating all balanced parentheses of all lengths & storing in a set<string> to avoid duplicates, & also getting len of max balanced parentheses
 - 2) Iterate on set of all strings having len = (max length) in that can add it in the final ans
 - 3) Return final ans

All solutions :: solve(string A) → len of string
 ans2.clear(); Initializing Global Variables
 len = 0;
 string a = " ";
 bal(A, 0, 0, a); Before function call
 // we have data in ans2 & also len
 vector<string> ans;
 // iterate on set string by string, if a string length == len, we
 can simply add this to our final answer ans
 for (auto it = ans2.begin(); it != ans2.end(); it++) {
 if ((*it).size() == len) {
 ans.push_back(*it);
 }
 }
 return ans;