



Modélisation & Résolution

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Plan

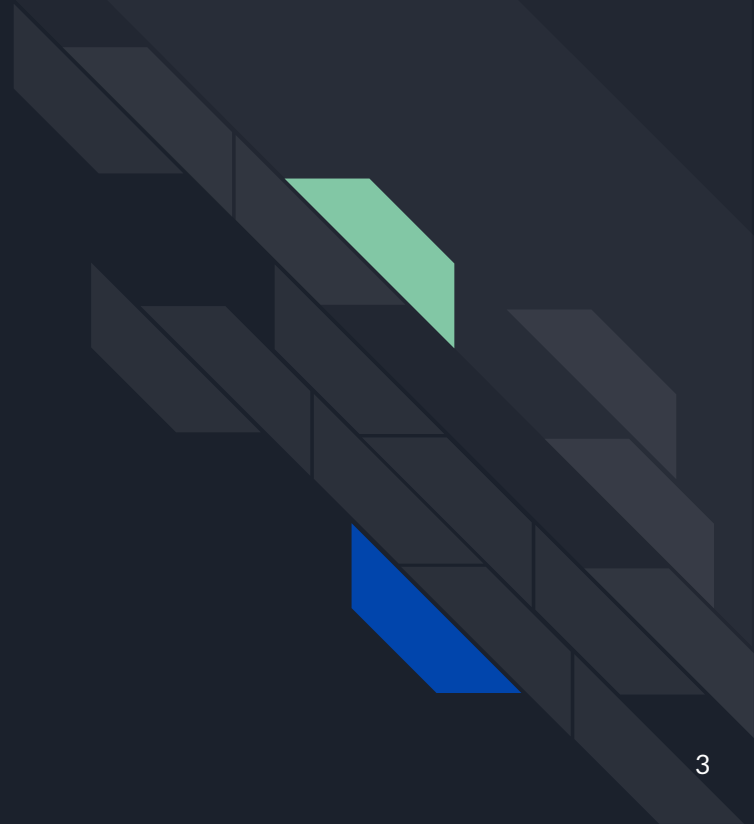
I/ Les n-reines

- 1) M1 & R1
- 2) M2 & R2
- 3) M3 pseudo booléen
- 4) M3 booléen

II/ Garam

III/ Rikudo

I/ Les n-reines





1) M1

- Variables : $X_i, i \in [1..4]$
 - Reines sur l'échiquier
- Domaines (valeurs possible) : $\{1, \dots, n\}$
 - Index des colonnes de chaque reines placées
- Contraintes (requirements):
 - Une et une seule reine par colonne
 - Déjà dans la formulation des variables et du domaine
 - Une et une seule reine par ligne
 - $\forall i, j \in [1..4], i \neq j$
 - Pas deux reines sur une diagonale
 - Voir R1



R1

Chaque reine à sa colonne.

Nous voulons savoir quelle ligne elle occupe.

Nous avons appliqué des contraintes finies pour 4 reines :

```
% ligne
constraint R1 != R2;
constraint R1 != R3;
constraint R1 != R4;
constraint R2 != R3;
constraint R2 != R4;
constraint R3 != R4;
```

```
% Diagonales montantes
constraint R1 != R2 + 1;
constraint R1 != R3 + 2;
constraint R1 != R4 + 3;
constraint R2 != R3 + 1;
constraint R2 != R4 + 2;
constraint R3 != R4 + 1;
```

```
% Diagonales descendantes
constraint R1 != R2 - 1;
constraint R1 != R3 - 2;
constraint R1 != R4 - 3;
constraint R2 != R3 - 1;
constraint R2 != R4 - 2;
constraint R3 != R4 - 1;
```

R1	R2	R3	R4
		Q	
Q			
			Q
	Q		



2) M2

- Variables : $q = \text{array}[1..n]$ of var $1..n$
 - Reines sur l'échiquier
- Domaines (valeurs possible) : $\{1, \dots, n\}$
 - Index des colonnes de chaque reines placées
- Contraintes (requirements):
 - Une et une seule reine par colonne
 - Déjà dans la formulation des variables et du domaine
 - Une et une seule reine par ligne
 - $\text{alldifferent}(q)$
 - Pas deux reines sur une diagonale
 - Voir R2



R2

```
include "alldifferent.mzn";

int: n;
array[1..n] of var 1..n:q;

% Une et une seule reine par colonne
constraint alldifferent(q);
% Diagonale
constraint alldifferent([q[i] + i | i in 1..n]);
constraint alldifferent([q[i] - i | i in 1..n]);

constraint q[1] <= q[n];

solve satisfy;

output [ if fix(q[j]) == i then "|Q" else "|" endif ++
        | if j == n then "|\\n" else "" endif | i, j in 1..n];
```



3) M3 pseudo booléen

Première modélisation

```
% lines
constraint forall(i in 1..n)(
|   sum(k in 1..n)(bool2int(q[k,i])) = 1
);
```

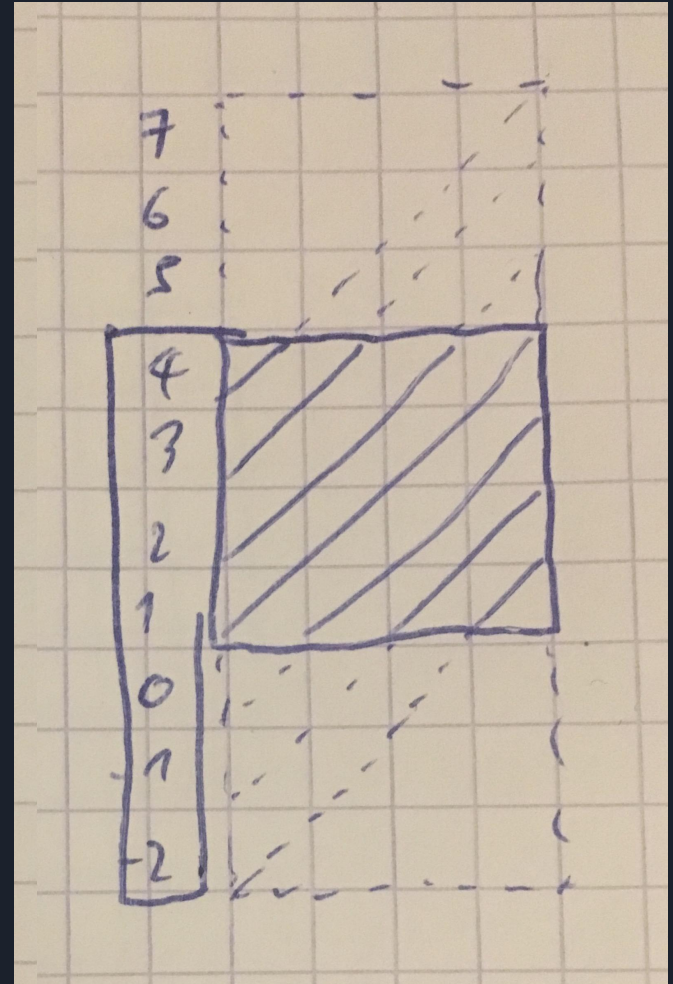
Variables booléennes

Contrainte sur les colonnes

```
% columns
constraint forall(i in 1..n)(
|   sum(k in 1..n)(bool2int(q[i,k])) = 1
);
```


3) M3 pseudo booléen

```
% increasing diagonales
% decreasing diagonales
constraint forall(i in 1..(2 * n - 1))(
  sum(k in 1..n)(
    if((i - k + 1 > n) \vee (i - k + 1 < 1)) then 0
    else bool2int(q[i - k + 1, k])
    endif
  ) < 2
);
```



4) M3

P

```
% lines
constraint forall(i in 1..n) (
  forall(j in 1..n) (
    orall([q[i, k] | k in 1..n where k != j]) != q[i, j]
  )
);
```

else"

=

```
% columns
constraint forall(i in 1..n) (
  forall(j in 1..n) (
    orall([q[k, j] | k in 1..n where k != i]) != q[i, j]
  )
);
```

```
function var bool: orall(array[int] of var bool: tab) =
  if length(tab) = 1
  then tab[1]
  else tab[1] /\ orall([tab[i] | i in 1..length(tab) where i > 1])
endif;
```

4) M3

Néanmoins, pour ce qui est des diagonales, le travail est plus fastidieux.

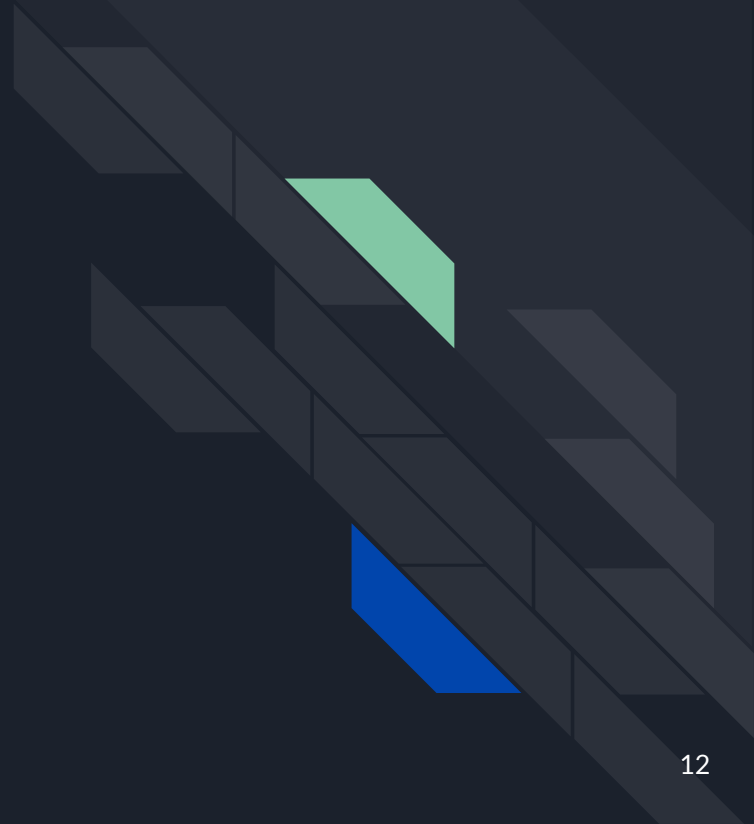
Pour chaque case :

```
constraint
for i in lines do
  for j in columns do
    if not out of bounds && queen[i,j] == true
    then orall( AllOtherCases ) == false
    else true
    endif
  endfor
endfor
```

```
% decreasing diagonales
constraint forall(i in 1..n) (
  forall(j in 1..n) (
    if q[i, j] = true /\ length([q[i + k, j + k] | k in (-n)..n where i + k > 0 /\ i + k < n /\ j + k > 0 /\ j + k < n /\ k != 0]) != 0
    then orall([q[i + k, j + k] | k in (-n)..n where i + k > 0 /\ i + k < n /\ j + k > 0 /\ j + k < n /\ k != 0]) != true
    else true
    endif
  )
);

% increasing diagonales
constraint forall(i in 1..n) (
  forall(j in 1..n) (
    if q[i, j] = true /\ length([q[i - k, j + k] | k in (-n)..n where i - k > 0 /\ i - k < n /\ j + k > 0 /\ j + k < n /\ k != 0]) != 0
    then orall([q[i - k, j + k] | k in (-n)..n where i - k > 0 /\ i - k < n /\ j + k > 0 /\ j + k < n /\ k != 0]) != true
    else true
    endif
  )
);
```

II/ Garam



II/ Garam

On travaille sur des indices au début de bas gauche

% 1st line

o10 = "-";

o90 = "+";

x100 = 1;

% 2nd line

x01 = 2;

x81 = 2;

x121 = 4;



```
numbers = [  
    _,_,_,1,_,_  
    2,_,2,4,_,_  
    _,_,2,_,9,_,_  
    _,_,9,4,_,_,_  
    2,2,_,_,_,_  
    _,_,2,_,_,_,_  
    1,1,1,1,_,_  
    7,_,1,_,9,_,_  
    _,3,_,_,_,_,_  
];  
  
operators = [  
    "-", "+",  
    "+",  
    "*", "+", "*", *,  
    "+", "+",  
    "+", "-",  
    "-", "-",  
    "+",  
    "*", "+", "*", +,  
    "+", +  
];
```

```

predicate operation(string: operator, var 0..9: x1, var 0..9: x2, var 0..90: x3) =
    (operator = "+" /\ x1 + x2 = x3) \/
    (operator = "-" /\ x1 - x2 = x3) \/
    (operator = "*" /\ x1 * x2 = x3) \/
    (operator = "/" /\ x1 / x2 = x3);

```

```

constraint operation(operators[1], numbers[1], numbers[2], numbers[3]);

```

```

constraint numbers[7] > 0;

```

```

constraint operation(operators[15], numbers[39], numbers[34], numbers[30] * 10 + numbers[24]);

```

Garam : niveau Facile (01/11/2019)

$\begin{array}{r} +3= \\ \times \\ 7 \\ \hline 1 \end{array}$	$\begin{array}{r} + \\ + \\ 1= \\ \hline 1 \end{array}$	$\begin{array}{r} + \\ \times \\ 1 \\ \hline 1 \end{array}$	$\begin{array}{r} + \\ \times \\ 9 \\ \hline 1 \end{array}$
$\begin{array}{r} - \\ =2 \end{array}$	$\begin{array}{r} + \\ 2= \\ \hline \end{array}$	$\begin{array}{r} - \\ = \end{array}$	$\begin{array}{r} 2 \\ \hline \end{array}$
$\begin{array}{r} + \\ =9 \end{array}$	$\begin{array}{r} + \\ +2= \\ \hline \end{array}$	$\begin{array}{r} 4+ \\ \times \\ 9 \\ \hline 4 \end{array}$	$\begin{array}{r} 2 \\ \hline \end{array}$
$\begin{array}{r} \times \\ 2 \\ \hline \end{array}$	$\begin{array}{r} - \\ = \end{array}$	$\begin{array}{r} +1= \\ \hline \end{array}$	$\begin{array}{r} 2 \\ \hline \end{array}$

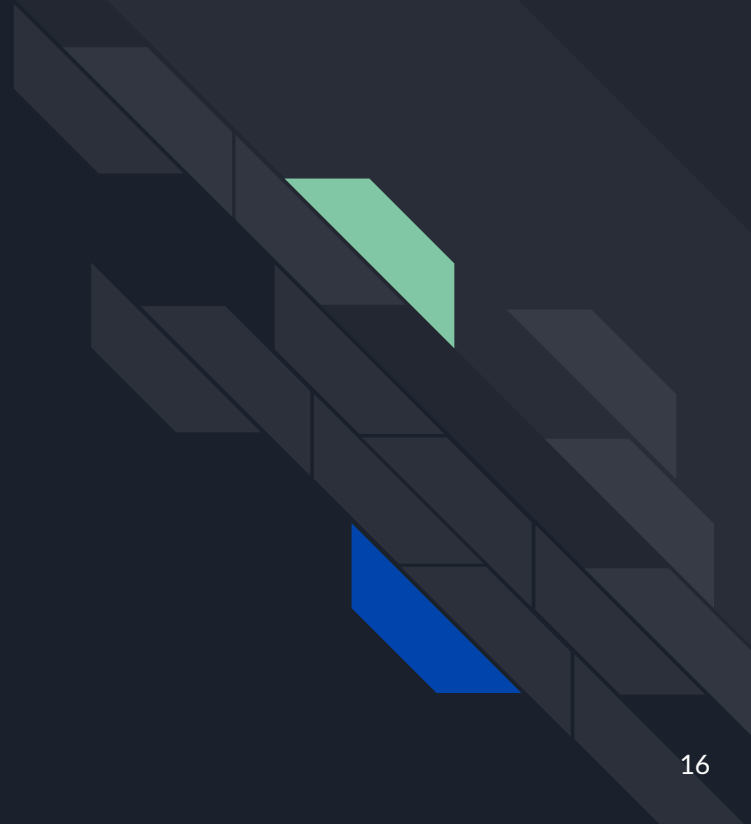
1 2 3 4 5 6
7 8 9 0



$2+3=5$	$2+2=4$		
$*$	$+$	$*$	$+$
7	$7+1=8$	9	
$=$	$=$	$=$	$=$
1	1	1	1
$4-2=2$	$6-3=3$		
$+$	$-$		
2	2		
$=$	$=$		
$5+4=9$	$4+1=5$		
$*$	$+$	$*$	$*$
5	$4+2=6$	9	
$=$	$=$	$=$	$=$
2	1	2	4
$5-2=3$	$4+1=5$		

\times	$+3=$		$+$	$=$	\times	$+$	$=$
7		$+$	$+1=$		9		
$=$	$=$	$=$	$=$	$=$	$=$	$=$	$=$
1	1				1	1	
	$-$	$=2$			$-$	$=$	
	$+$	$=$			$-$	$=$	
	2				2		
	$=$				$=$		
\times	$+$	$=9$		$+$	$=$		
		$+$	$+2=$			\times	9
$=$	$=$	$=$	$=$	$=$	$=$	$=$	$=$
2					2		4
	$-$	$=$			$+$	$+1=$	

III/ Rikudo

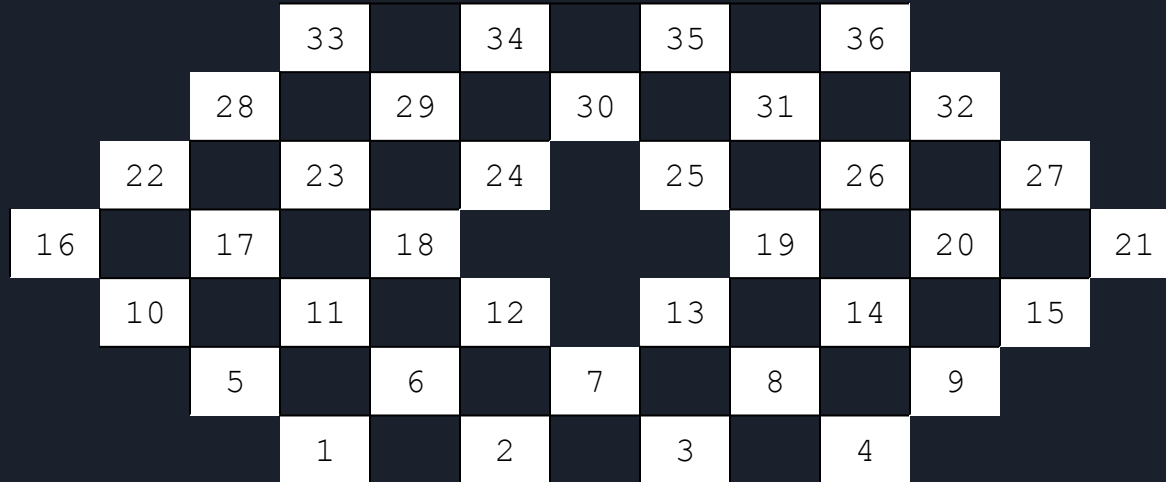


III/ Rikudo



-1	-1	-1	0	-1	0	-1	13	-1	0	-1	-1	-1
-1	-1	7	-1	8	-1	0	-1	0	-1	0	-1	-1
-1	5	-1	0	-1	0	-1	0	-1	0	-1	0	-1
3	-1	0	-1	30	-1	-1	-1	34	-1	0	-1	18
-1	0	-1	0	-1	0	-1	0	-1	0	-1	0	-1
-1	-1	1	-1	26	-1	0	-1	36	-1	0	-1	-1
-1	-1	-1	0	-1	0	-1	0	-1	22	-1	-1	-1

III/ Rikudo

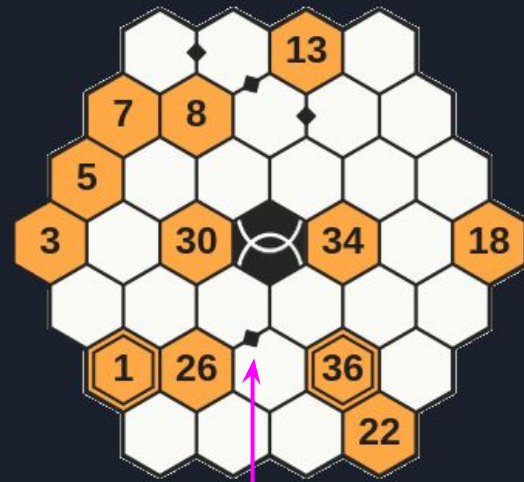


III/ Rikudo

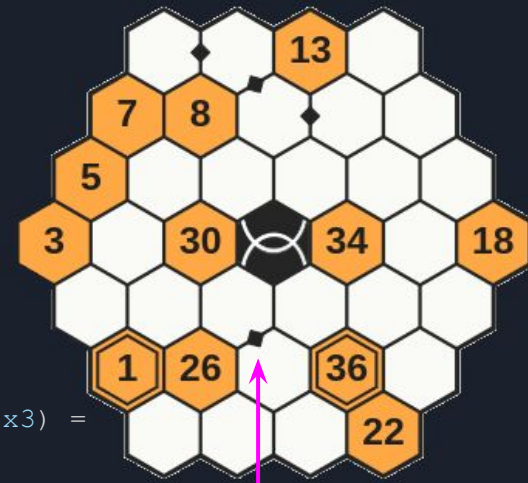
Sens de
lecture

```
rikudo = [  
  _,_,_,22,  
  1,26,_,36,_,  
  _,'_','_','_','_','_'  
  3,_,30,34,_,18,  
  5,_,_,_,_,_,  
  7,8,_,_,_,  
  _,_,13,_  
];
```

```
links_size = 4;  
links = array2d(1..links_size,1..2,  
  [ 7,12,  
    30,31,  
    33,34,  
    34,30  
  ],  
  );
```



III/ Rikudo



```
% for boxes with 3 links
predicate links(var 1..36 : x, var 1..36 : x1, var 1..36 : x2, var 1..36 : x3) =
    x != 36 -> x1 = x + 1 \/ x2 = x + 1 \/ x3 = x + 1;
    3-6 links ...
```

```
% first line
constraint links(rikudo[1], rikudo[2], rikudo[5], rikudo[6]);
constraint links(rikudo[2], rikudo[1], rikudo[3], rikudo[6], rikudo[7]);
constraint links(rikudo[3], rikudo[2], rikudo[4], rikudo[7], rikudo[8]);
constraint links(rikudo[4], rikudo[3], rikudo[8], rikudo[9]);
    1-7 line ...
```

```
% for links
    rikudo[links[1]] = rikudo[links[2]] + 1
    OU
    rikudo[links[2]] = rikudo[links[1]] + 1
```



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Merci
Avez-vous des questions ?