

# Chapter 6: Database Design Using the E-R Model

**Database System Concepts, 7th Ed.** 

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#### **Design Phases**

- Initial phase -- characterize fully the data needs of the prospective database users.
- Second phase -- choosing a data model
  - Applying the concepts of the chosen data model
  - Translating these requirements into a conceptual schema of the database.
  - A fully developed conceptual schema indicates the functional requirements of the enterprise.
    - Describe the kinds of operations (or transactions) that will be performed on the data.



# **Design Phases (Cont.)**

- Final Phase -- Moving from an abstract data model to the implementation of the database
  - Logical Design Deciding on the database schema.
    - Database design requires that we find a "good" collection of relation schemas.
    - Business decision What attributes should we record in the database?
    - Computer Science decision What relation schemas should we have and how should the attributes be distributed among the various relation schemas?
  - Physical Design Deciding on the physical layout of the database



#### **Design Alternatives**

- In designing a database schema, we must ensure that we avoid two major pitfalls:
  - Redundancy: a bad design may result in repeat information.
    - Redundant representation of information may lead to data inconsistency among the various copies of information
  - Incompleteness: a bad design may make certain aspects of the enterprise difficult or impossible to model.
- Avoiding bad designs is not enough. There may be a large number of good designs from which we must choose.



#### **Design Approaches**

- Entity Relationship Model (covered in this chapter)
  - Models an enterprise as a collection of entities and relationships
    - Entity: a "thing" or "object" in the enterprise that is distinguishable from other objects
      - Described by a set of attributes
    - Relationship: an association among several entities
  - Represented diagrammatically by an entity-relationship diagram:
- Normalization Theory (Chapter 7)
  - Formalize what designs are bad, and test for them



#### **Outline of the ER Model**



#### **Entity Sets**

- An entity is an object that exists and is distinguishable from other objects.
  - Example: specific person, company, event, plant
- An entity set is a set of entities of the same type that share the same properties.
  - Example: set of all persons, companies, trees, holidays
- An entity is represented by a set of attributes; i.e., descriptive properties possessed by all members of an entity set.
  - Example:

```
instructor = (ID, name, salary )
course= (course_id, title, credits)
```

 A subset of the attributes form a primary key of the entity set; i.e., uniquely identifying each member of the set.



# Representing Entity sets in ER Diagram

- Entity sets can be represented graphically as follows:
  - Rectangles represent entity sets.
  - Attributes listed inside entity rectangle
  - Underline indicates primary key attributes

instructor

<u>ID</u>

name

salary

ID
name
tot cred



#### **Relationship Sets**

A relationship is an association among several entities

Example:

44553 (Peltier) <u>advisor</u> 22222 (<u>Einstein</u>) student entity relationship set instructor entity

A relationship set is a mathematical relation among n ≥ 2 entities, each taken from entity sets

$$\{(e_1, e_2, \dots e_n) \mid e_1 \in E_1, e_2 \in E_2, \dots, e_n \in E_n\}$$

where  $(e_1, e_2, ..., e_n)$  is a relationship

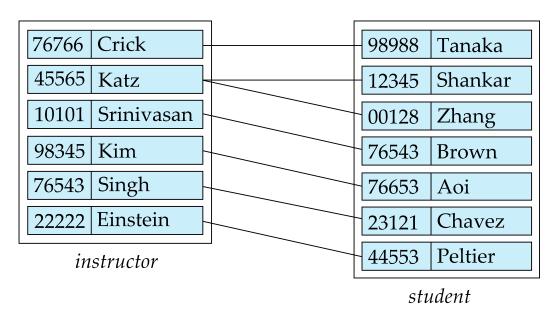
Example:

 $(44553,22222) \in advisor$ 



# Relationship Sets (Cont.)

- Example: we define the relationship set advisor to denote the associations between students and the instructors who act as their advisors.
- Pictorially, we draw a line between related entities.





#### Representing Relationship Sets via ER Diagrams

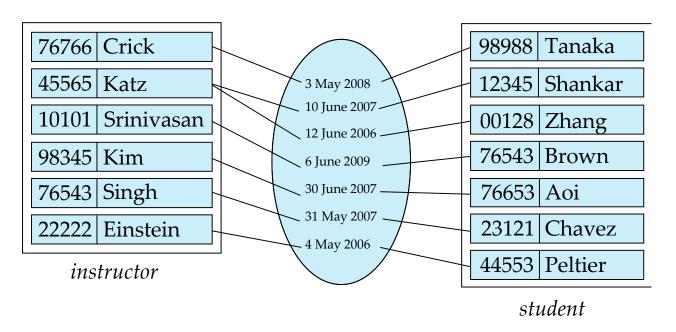
Diamonds represent relationship sets.





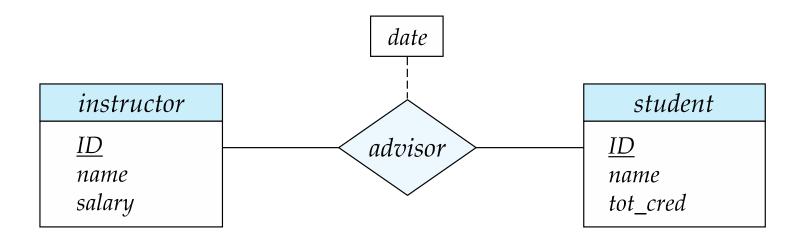
### **Relationship Sets (Cont.)**

- An attribute can also be associated with a relationship set.
- For instance, the advisor relationship set between entity sets instructor and student may have the attribute date which tracks when the student started being associated with the advisor





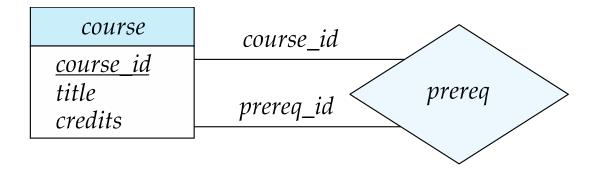
# **Relationship Sets with Attributes**





#### Roles

- Entity sets of a relationship need not be distinct
  - Each occurrence of an entity set plays a "role" in the relationship
- The labels "course\_id" and "prereq\_id" are called roles.





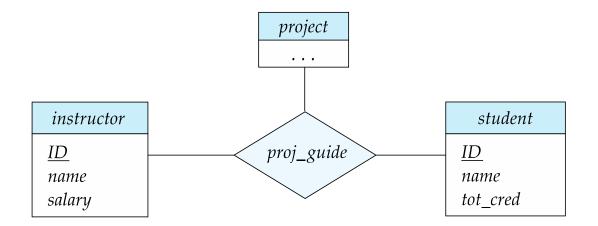
#### Degree of a Relationship Set

- Binary relationship
  - involve two entity sets (or degree two).
  - most relationship sets in a database system are binary.
- Relationships between more than two entity sets are rare. Most relationships are binary. (More on this later.)
  - Example: students work on research projects under the guidance of an instructor.
  - relationship proj\_guide is a ternary relationship between instructor, student, and project



### **Non-binary Relationship Sets**

- Most relationship sets are binary
- There are occasions when it is more convenient to represent relationships as non-binary.
- E-R Diagram with a Ternary Relationship





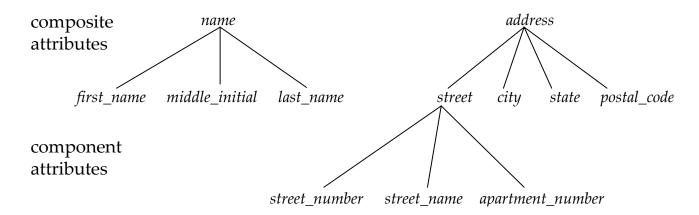
#### **Complex Attributes**

- Attribute types:
  - Simple and composite attributes.
  - Single-valued and multivalued attributes
    - Example: multivalued attribute: phone\_numbers
  - Derived attributes
    - Can be computed from other attributes
    - Example: age, given date\_of\_birth
- Domain the set of permitted values for each attribute



#### **Composite Attributes**

Composite attributes allow us to divided attributes into subparts (other attributes).





#### Representing Complex Attributes in ER Diagram

#### instructor

```
ID
name
  first_name
   middle_initial
   last_name
address
   street
      street_number
      street_name
     apt_number
   city
   state
   zip
{ phone_number }
date_of_birth
age()
```

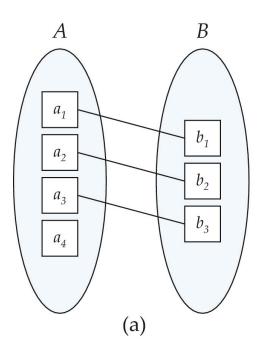


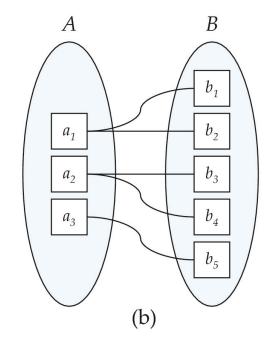
#### **Mapping Cardinality Constraints**

- Express the number of entities to which another entity can be associated via a relationship set.
- Most useful in describing binary relationship sets.
- For a binary relationship set the mapping cardinality must be one of the following types:
  - One to one
  - One to many
  - Many to one
  - Many to many



# **Mapping Cardinalities**





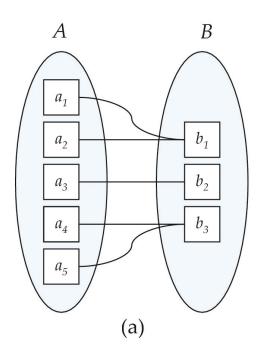
One to one

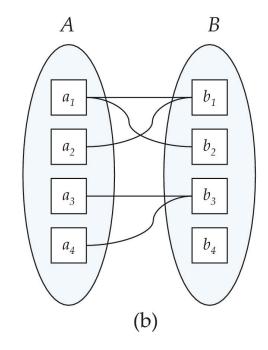
One to many

Note: Some elements in A and B may not be mapped to any elements in the other set



### **Mapping Cardinalities**





Many to one

Many to many

Note: Some elements in A and B may not be mapped to any elements in the other set



#### Representing Cardinality Constraints in ER Diagram

- We express cardinality constraints by drawing either a directed line (→), signifying "one," or an undirected line (—), signifying "many," between the relationship set and the entity set.
- One-to-one relationship between an instructor and a student :
  - A student is associated with at most one instructor via the relationship advisor
  - A student is associated with at most one department via stud dept





#### **One-to-Many Relationship**

- one-to-many relationship between an *instructor* and a *student* 
  - an instructor is associated with several (including 0) students via advisor
  - a student is associated with at most one instructor via advisor,





### Many-to-One Relationships

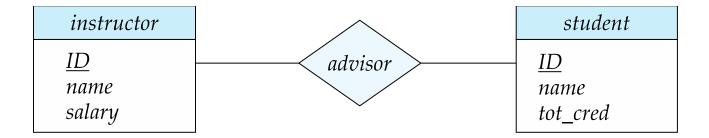
- In a many-to-one relationship between an instructor and a student,
  - an instructor is associated with at most one student via advisor,
  - and a student is associated with several (including 0) instructors via advisor





#### Many-to-Many Relationship

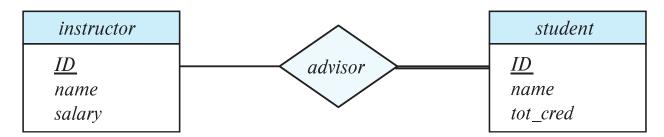
- An instructor is associated with several (possibly 0) students via advisor
- A student is associated with several (possibly 0) instructors via advisor





### **Total and Partial Participation**

• **Total participation** (indicated by double line): every entity in the entity set participates in at least one relationship in the relationship set



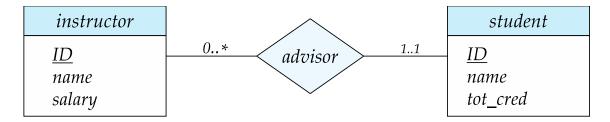
participation of *student* in *advisor r*elation is total

- every student must have an associated instructor
- Partial participation: some entities may not participate in any relationship in the relationship set
  - Example: participation of instructor in advisor is partial



#### **Notation for Expressing More Complex Constraints**

- A line may have an associated minimum and maximum cardinality, shown in the form *l..h*, where *l* is the minimum and *h* the maximum cardinality
  - A minimum value of 1 indicates total participation.
  - A maximum value of 1 indicates that the entity participates in at most one relationship
  - A maximum value of \* indicates no limit.
- Example



 Instructor can advise 0 or more students. A student must have 1 advisor; cannot have multiple advisors



# **Cardinality Constraints on Ternary Relationship**

- We allow at most one arrow out of a ternary (or greater degree)
   relationship to indicate a cardinality constraint
- For example, an arrow from proj\_guide to instructor indicates each student has at most one guide for a project
- If there is more than one arrow, there are two ways of defining the meaning.
  - For example, a ternary relationship R between A, B and C with arrows to B and C could mean
    - 1. Each A entity is associated with a unique entity from B and C or
    - 2. Each pair of entities from (A, B) is associated with a unique C entity, and each pair (A, C) is associated with a unique B
  - Each alternative has been used in different formalisms
  - To avoid confusion we outlaw more than one arrow



# **Primary Key**

- Primary keys provide a way to specify how entities and relations are distinguished. We will consider:
  - Entity sets
  - Relationship sets.
  - Weak entity sets



# **Primary key for Entity Sets**

- By definition, individual entities are distinct.
- From database perspective, the differences among them must be expressed in terms of their attributes.
- The values of the attribute values of an entity must be such that they can uniquely identify the entity.
  - No two entities in an entity set are allowed to have exactly the same value for all attributes.
- A key for an entity is a set of attributes that suffice to distinguish entities from each other



# **Primary Key for Relationship Sets**

- To distinguish among the various relationships of a relationship set we use the individual primary keys of the entities in the relationship set.
  - Let R be a relationship set involving entity sets E1, E2, .. En
  - The primary key for R is consists of the union of the primary keys of entity sets E1, E2, ..En
  - If the relationship set *R* has attributes a1, a2, .., am associated with it, then the primary key of *R* also includes the attributes a1, a2, .., am
- Example: relationship set "advisor".
  - The primary key consists of instructor.ID and student.ID
- The choice of the primary key for a relationship set depends on the mapping cardinality of the relationship set.



# **Choice of Primary key for Binary Relationship**

- Many-to-Many relationships. The preceding union of the primary keys is a minimal superkey and is chosen as the primary key.
- One-to-Many relationships. The primary key of the "Many" side is a minimal superkey and is used as the primary key.
- Many-to-one relationships. The primary key of the "Many" side is a minimal superkey and is used as the primary key.
- One-to-one relationships. The primary key of either one of the participating entity sets forms a minimal superkey, and either one can be chosen as the primary key.



#### **Weak Entity Sets**

- Consider a section entity, which is uniquely identified by a course\_id, semester, year, and sec\_id.
- Clearly, section entities are related to course entities. Suppose we create
  a relationship set sec\_course between entity sets section and course.
- Note that the information in sec\_course is redundant, since section
  already has an attribute course\_id, which identifies the course with which
  the section is related.
- One option to deal with this redundancy is to get rid of the relationship sec\_course; however, by doing so the relationship between section and course becomes implicit in an attribute, which is not desirable.



# Weak Entity Sets (Cont.)

- An alternative way to deal with this redundancy is to not store the attribute course\_id in the section entity and to only store the remaining attributes section\_id, year, and semester.
  - However, the entity set section then does not have enough attributes to identify a particular section entity uniquely
- To deal with this problem, we treat the relationship sec\_course as a special relationship that provides extra information, in this case, the course\_id, required to identify section entities uniquely.
- A weak entity set is one whose existence is dependent on another entity, called its identifying entity
- Instead of associating a primary key with a weak entity, we use the identifying entity, along with extra attributes called **discriminator** to uniquely identify a weak entity.



# Weak Entity Sets (Cont.)

- An entity set that is not a weak entity set is termed a strong entity set.
- Every weak entity must be associated with an identifying entity; that is, the weak entity set is said to be existence dependent on the identifying entity set.
- The identifying entity set is said to own the weak entity set that it identifies.
- The relationship associating the weak entity set with the identifying entity set is called the identifying relationship.
- Note that the relational schema we eventually create from the entity set section does have the attribute course\_id, for reasons that will become clear later, even though we have dropped the attribute course\_id from the entity set section.



# **Expressing Weak Entity Sets**

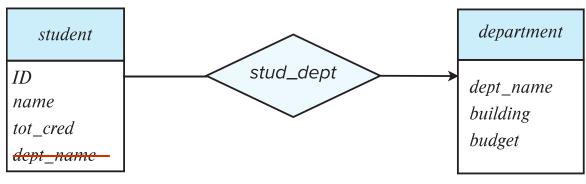
- In E-R diagrams, a weak entity set is depicted via a double rectangle.
- We underline the discriminator of a weak entity set with a dashed line.
- The relationship set connecting the weak entity set to the identifying strong entity set is depicted by a double diamond.
- Primary key for section (course\_id, sec\_id, semester, year)





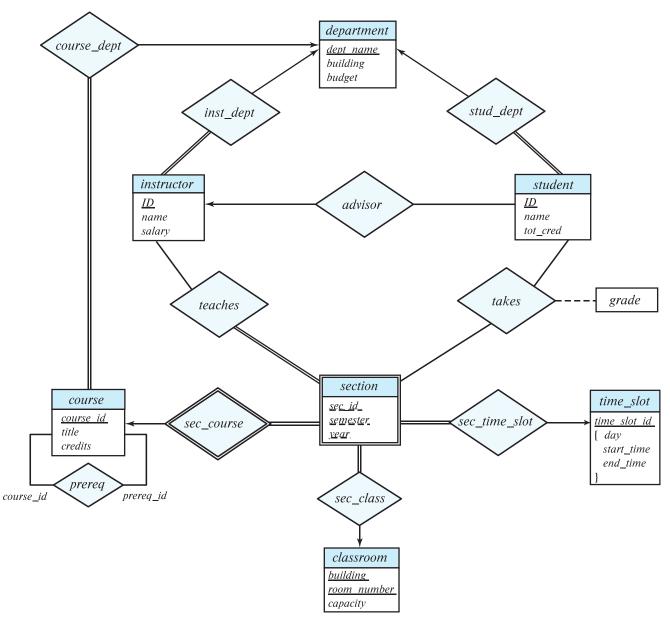
#### **Redundant Attributes**

- Suppose we have entity sets:
  - student, with attributes: ID, name, tot\_cred, dept\_name
  - department, with attributes: dept\_name, building, budget
- We model the fact that each student has an associated department using a relationship set stud\_dept
- The attribute dept\_name in student below replicates information present in the relationship and is therefore redundant
  - and needs to be removed.
- BUT: when converting back to tables, in some cases the attribute gets reintroduced, as we will see later.





# E-R Diagram for a University Enterprise





#### **Reduction to Relation Schemas**



#### **Reduction to Relation Schemas**

- Entity sets and relationship sets can be expressed uniformly as relation schemas that represent the contents of the database.
- A database which conforms to an E-R diagram can be represented by a collection of schemas.
- For each entity set and relationship set there is a unique schema that is assigned the name of the corresponding entity set or relationship set.
- Each schema has a number of columns (generally corresponding to attributes), which have unique names.



### **Representing Entity Sets**

A strong entity set reduces to a schema with the same attributes

student(<u>ID</u>, name, tot\_cred)

 A weak entity set becomes a table that includes a column for the primary key of the identifying strong entity set

section ( course\_id, sec\_id, sem, year )

Example





#### Representation of Entity Sets with Composite Attributes

#### instructor

```
ID
name
  first name
  middle initial
  last name
address
  street
     street number
     street name
     apt number
  city
  state
  zip
{ phone_number }
date of birth
age()
```

- Composite attributes are flattened out by creating a separate attribute for each component attribute
  - Example: given entity set instructor with composite attribute name with component attributes first\_name and last\_name the schema corresponding to the entity set has two attributes name\_first\_name and name\_last\_name
    - Prefix omitted if there is no ambiguity (name\_first\_name could be first\_name)
- Ignoring multivalued attributes, extended instructor schema is
  - instructor(ID, first\_name, middle\_initial, last\_name, street\_number, street\_name, apt\_number, city, state, zip\_code, date of birth)



#### Representation of Entity Sets with Multivalued Attributes

- A multivalued attribute M of an entity E is represented by a separate schema EM
- Schema EM has attributes corresponding to the primary key of E and an attribute corresponding to multivalued attribute M
- Example: Multivalued attribute phone\_number of instructor is represented by a schema: inst\_phone= (<u>ID</u>, phone\_number)
- Each value of the multivalued attribute maps to a separate tuple of the relation on schema EM
  - For example, an *instructor* entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples:

     (22222, 456-7890) and (22222, 123-4567)



# Representing Relationship Sets

- A many-to-many relationship set is represented as a schema with attributes for the primary keys of the two participating entity sets, and any descriptive attributes of the relationship set.
- Example: schema for relationship set advisor

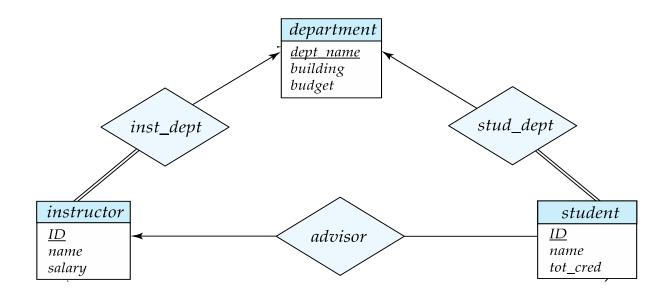
$$advisor = (\underline{s}\underline{id}, \underline{i}\underline{id})$$





### **Redundancy of Schemas**

- Many-to-one and one-to-many relationship sets that are total on the manyside can be represented by adding an extra attribute to the "many" side, containing the primary key of the "one" side
- Example: Instead of creating a schema for relationship set inst\_dept, add an attribute dept\_name to the schema arising from entity set instructor
- Example





## Redundancy of Schemas (Cont.)

- For one-to-one relationship sets, either side can be chosen to act as the "many" side
  - That is, an extra attribute can be added to either of the tables corresponding to the two entity sets
- If participation is partial on the "many" side, replacing a schema by an extra attribute in the schema corresponding to the "many" side could result in null values



## Redundancy of Schemas (Cont.)

- The schema corresponding to a relationship set linking a weak entity set to its identifying strong entity set is redundant.
- Example: The section schema already contains the attributes that would appear in the sec\_course schema





### **Extended E-R Features**



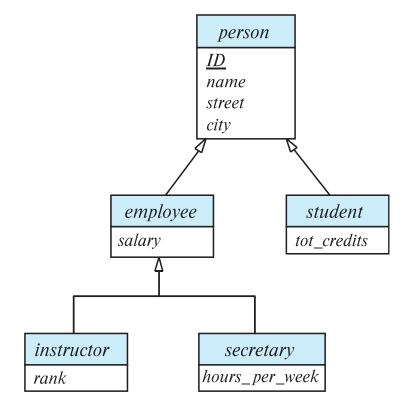
## **Specialization**

- Top-down design process; we designate sub-groupings within an entity set that are distinctive from other entities in the set.
- These sub-groupings become lower-level entity sets that have attributes or participate in relationships that do not apply to the higher-level entity set.
- Depicted by a triangle component labeled ISA (e.g., instructor "is a" person).
- Attribute inheritance a lower-level entity set inherits all the attributes and relationship participation of the higher-level entity set to which it is linked.



## **Specialization Example**

- Overlapping employee and student
- Disjoint instructor and secretary
- Total and partial





## Representing Specialization via Schemas

- Method 1:
  - Form a schema for the higher-level entity
  - Form a schema for each lower-level entity set, include primary key
    of higher-level entity set and local attributes

schema	attributes
person	ID, name, street, city
student	ID, tot_cred
employee	ID, salary

 Drawback: getting information about, an employee requires accessing two relations, the one corresponding to the low-level schema and the one corresponding to the high-level schema



# Representing Specialization as Schemas (Cont.)

#### Method 2:

Form a schema for each entity set with all local and inherited attributes

ty
ty, tot_cred
ty, salary

 Drawback: name, street and city may be stored redundantly for people who are both students and employees



#### **Generalization**

- A bottom-up design process combine a number of entity sets that share the same features into a higher-level entity set.
- Specialization and generalization are simple inversions of each other;
   they are represented in an E-R diagram in the same way.
- The terms specialization and generalization are used interchangeably.



## **Completeness constraint**

- Completeness constraint -- specifies whether or not an entity in the higher-level entity set must belong to at least one of the lower-level entity sets within a generalization.
  - total: an entity must belong to one of the lower-level entity sets
  - partial: an entity need not belong to one of the lower-level entity sets



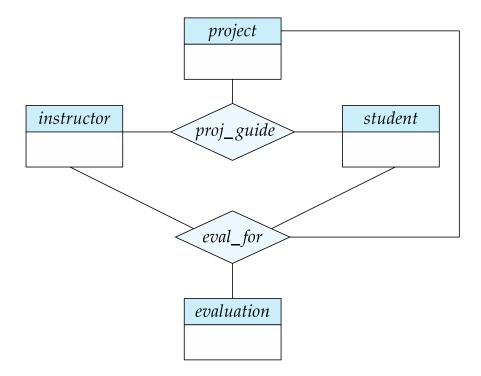
# **Completeness constraint (Cont.)**

- Partial generalization is the default.
- We can specify total generalization in an ER diagram by adding the keyword total in the diagram and drawing a dashed line from the keyword to the corresponding hollow arrow-head to which it applies (for a total generalization), or to the set of hollow arrow-heads to which it applies (for an overlapping generalization).
- The student generalization is total: All student entities must be either graduate or undergraduate. Because the higher-level entity set arrived at through generalization is generally composed of only those entities in the lower-level entity sets, the completeness constraint for a generalized higher-level entity set is usually total



## **Aggregation**

- Consider the ternary relationship proj\_guide, which we saw earlier
- Suppose we want to record evaluations of a student by a guide on a project





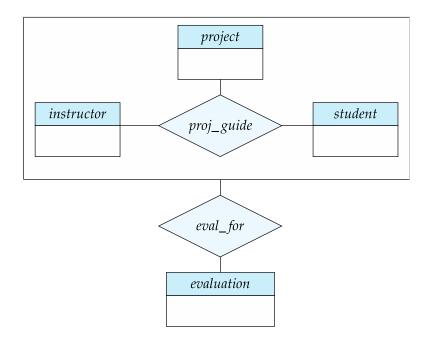
### **Aggregation (Cont.)**

- Relationship sets eval\_for and proj\_guide represent overlapping information
  - Every eval\_for relationship corresponds to a proj\_guide relationship.
  - However, some proj\_guide relationships may not correspond to any eval\_for relationships
    - So we can't discard the proj\_guide relationship
- Eliminate this redundancy via aggregation
  - Treat relationship as an abstract entity
  - Allows relationships between relationships
  - Abstraction of relationship into new entity



# **Aggregation (Cont.)**

- Eliminate this redundancy via aggregation without introducing redundancy, the following diagram represents:
  - A student is guided by a particular instructor on a particular project
  - A student, instructor, project combination may have an associated evaluation





#### **Reduction to Relational Schemas**

- To represent aggregation, create a schema containing
  - Primary key of the aggregated relationship,
  - The primary key of the associated entity set
  - Any descriptive attributes
- In our example:
  - The schema eval\_for is:
     eval\_for (s\_ID, project\_id, i\_ID, evaluation\_id)
  - The schema proj\_guide is redundant.

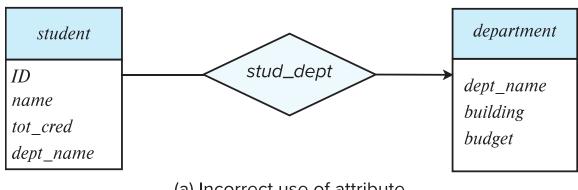


# **Design Issues**

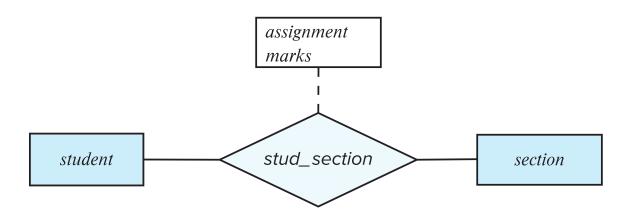


## **Common Mistakes in E-R Diagrams**

Example of erroneous E-R diagrams



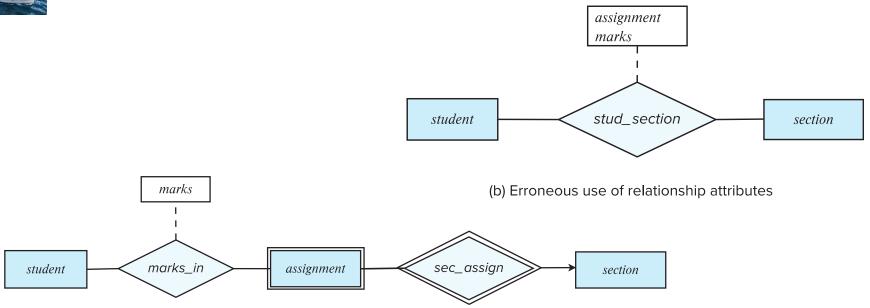
(a) Incorrect use of attribute



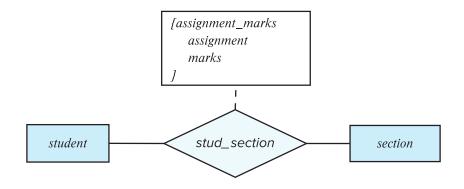
(b) Erroneous use of relationship attributes



# **Common Mistakes in E-R Diagrams (Cont.)**



(c) Correct alternative to erroneous E-R diagram (b)

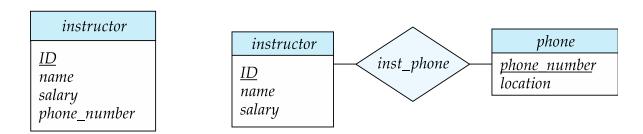


(d) Correct alternative to erroneous E-R diagram (b)



#### **Entities vs. Attributes**

Use of entity sets vs. attributes



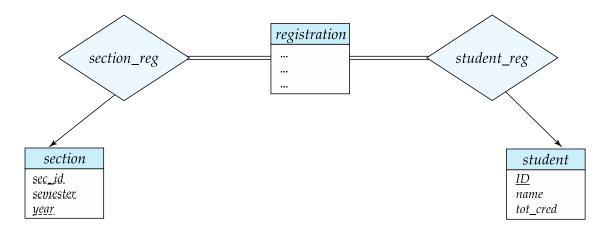
 Use of phone as an entity allows extra information about phone numbers (plus multiple phone numbers)



#### **Entities vs. Relationship sets**

#### Use of entity sets vs. relationship sets

Possible guideline is to designate a relationship set to describe an action that occurs between entities



#### Placement of relationship attributes

For example, attribute date as attribute of advisor or as attribute of student



# **Binary Vs. Non-Binary Relationships**

- Although it is possible to replace any non-binary (n-ary, for n > 2)
  relationship set by a number of distinct binary relationship sets, a n-ary
  relationship set shows more clearly that several entities participate in a
  single relationship.
- Some relationships that appear to be non-binary may be better represented using binary relationships
  - For example, a ternary relationship parents, relating a child to his/her father and mother, is best replaced by two binary relationships, father and mother
    - Using two binary relationships allows partial information (e.g., only mother being known)
  - But there are some relationships that are naturally non-binary
    - Example: proj\_guide

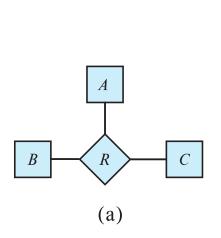


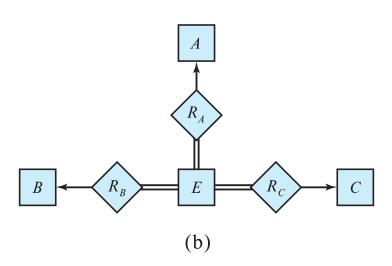
#### Converting Non-Binary Relationships to Binary Form

- In general, any non-binary relationship can be represented using binary relationships by creating an artificial entity set.
  - Replace R between entity sets A, B and C by an entity set E, and three relationship sets:
    - 1.  $R_A$ , relating E and A 2.  $R_B$ , relating E and B
    - 3.  $R_{\rm C}$ , relating E and C
  - Create an identifying attribute for *E* and add any attributes of *R* to *E*
  - For each relationship  $(a_i, b_i, c_i)$  in R, create
    - 1. a new entity  $e_i$  in the entity set E 2. add  $(e_i, a_i)$  to  $R_{\Delta}$

3. add  $(e_i, b_i)$  to  $R_B$ 

4. add  $(e_i, c_i)$  to  $R_C$ 







# **Converting Non-Binary Relationships (Cont.)**

- Also need to translate constraints
  - Translating all constraints may not be possible
  - There may be instances in the translated schema that cannot correspond to any instance of R
    - Exercise: add constraints to the relationships R<sub>A</sub>, R<sub>B</sub> and R<sub>C</sub> to ensure that a newly created entity corresponds to exactly one entity in each of entity sets A, B and C
  - We can avoid creating an identifying attribute by making E a weak entity set (described shortly) identified by the three relationship sets

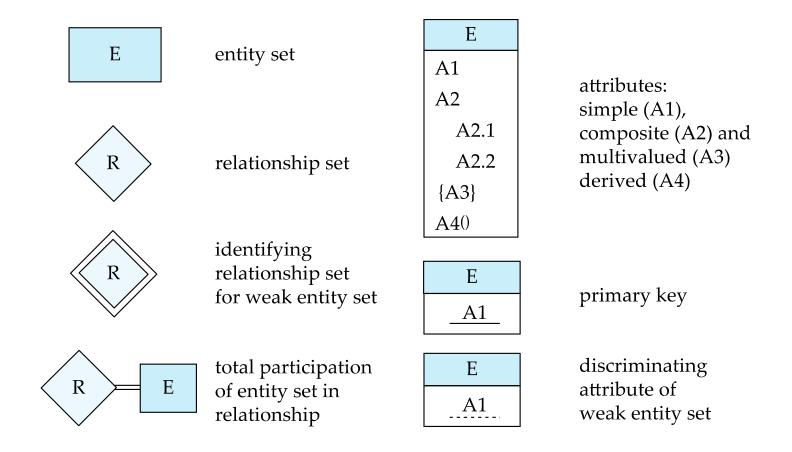


## **E-R Design Decisions**

- The use of an attribute or entity set to represent an object.
- Whether a real-world concept is best expressed by an entity set or a relationship set.
- The use of a ternary relationship versus a pair of binary relationships.
- The use of a strong or weak entity set.
- The use of specialization/generalization contributes to modularity in the design.
- The use of aggregation can treat the aggregate entity set as a single unit without concern for the details of its internal structure.

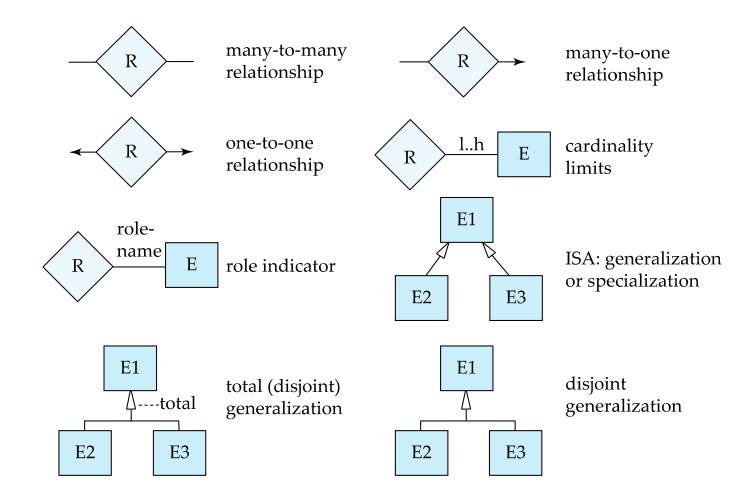


## **Summary of Symbols Used in E-R Notation**





# Symbols Used in E-R Notation (Cont.)

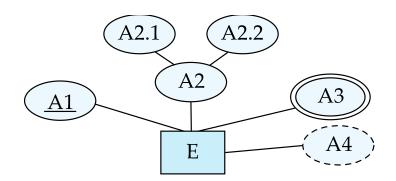




#### **Alternative ER Notations**

• Chen, IDE1FX, ...

entity set E with simple attribute A1, composite attribute A2, multivalued attribute A3, derived attribute A4, and primary key A1



weak entity set



generalization



total generalization

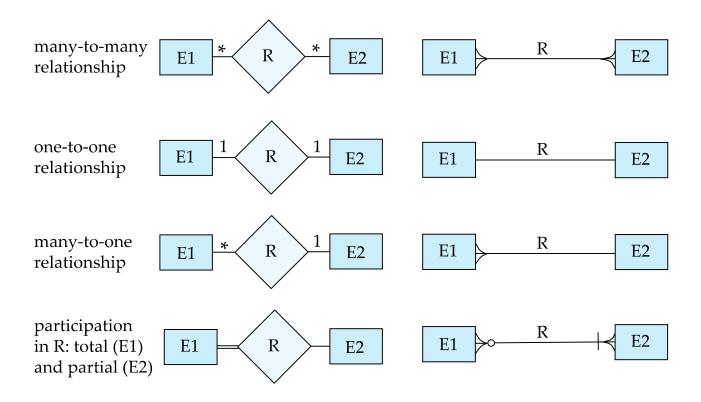




#### **Alternative ER Notations**

#### Chen

#### **IDE1FX (Crows feet notation)**





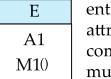
#### **UML**

- UML: Unified Modeling Language
- UML has many components to graphically model different aspects of an entire software system
- UML Class Diagrams correspond to E-R Diagram, but several differences.

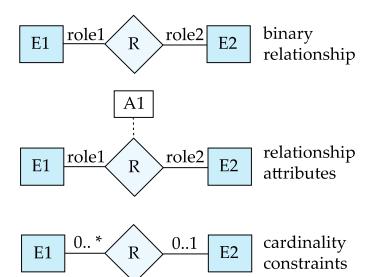


### **ER vs. UML Class Diagrams**

#### **ER Diagram Notation**

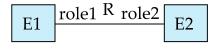


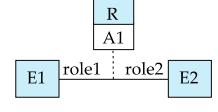
entity with attributes (simple, composite, multivalued, derived)

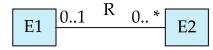


#### **Equivalent in UML**









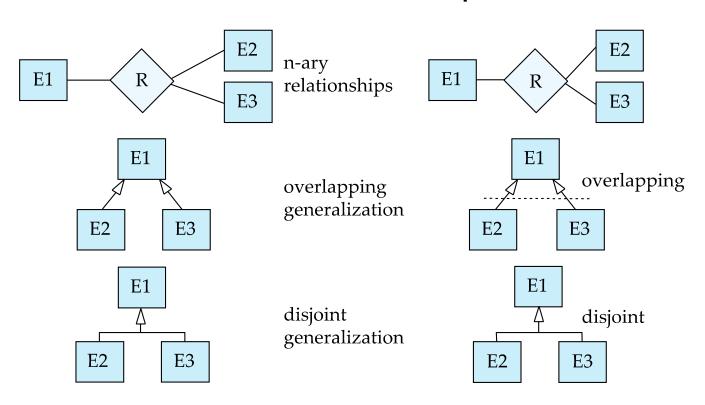
\* Note reversal of position in cardinality constraint depiction



### **ER vs. UML Class Diagrams**

#### **ER Diagram Notation**

#### **Equivalent in UML**



\* Generalization can use merged or separate arrows independent of disjoint/overlapping



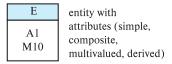
## **UML Class Diagrams (Cont.)**

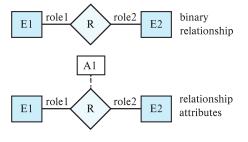
- Binary relationship sets are represented in UML by just drawing a line connecting the entity sets. The relationship set name is written adjacent to the line.
- The role played by an entity set in a relationship set may also be specified by writing the role name on the line, adjacent to the entity set.
- The relationship set name may alternatively be written in a box, along with attributes of the relationship set, and the box is connected, using a dotted line, to the line depicting the relationship set.

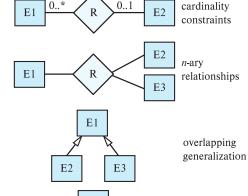


# **ER vs. UML Class Diagrams**

#### **ER Diagram Notation**









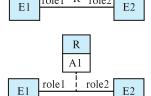
disjoint



#### **Equivalent in UML**

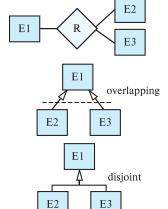


class with simple attributes and methods (attribute prefixes: + = public, - = private, # = protected)



role1 R role2









# Other Aspects of Database Design

- Functional Requirements
- Data Flow, Workflow
- Schema Evolution



# **End of Chapter 6**