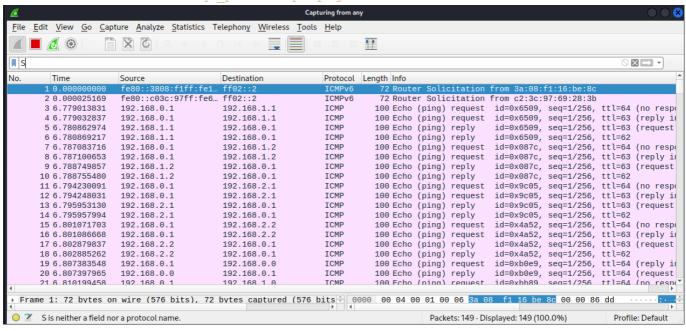
## Question 1

(a) The network topology called NetworkTopo() has been defined inside the q1.py file using the python api of mininet. pingall command inside mininet give the following output.

```
mininet> pingall
   Ping: testing ping reachability
         h3 h4 h5 h6 ra rb rc
          h3 h4
                h5
                    h6
                       ra rb
             h4
          h2
                 h5
                    h6
          h2
             h3
                h5
                    h6
                       ra
          h2
             h3
                 h4
                    h6
                       ra
          h2
            h3
                h4
                    h5
                       ra
             h3
          h2
                 h4
                    h5
                       h6
          h2
             h3
                 h4
          h2
             h3
                h4
                    h5
                       h6
                           ra
    Results: 0% dropped (72/72 received)
```

**(b)** Running the wireshark on router ra using the command ra wireshark &, and then running the pingall command inside mininet captures the following packets. The entire TCP Dump has been captured and saved inside the file called gl packets.pcapng.



**(c)** Updated the routing table in both ra and rc so that the packets are only sent to the rb, where they are again forwarded.

Comparing latency using the ping command and iperf tool gave the following

```
1. For path h1 -> ra -> rc -> h6
      mininet> h1 ping -c 5 h6
      PING 192.168.2.2 (192.168.2.2) 56(84) bytes of data.
      64 bytes from 192.168.2.2: icmp_seq=1 ttl=62 time=6.56 ms
      64 bytes from 192.168.2.2: icmp seq=2 ttl=62 time=0.493 ms
      64 bytes from 192.168.2.2: icmp_seq=3 ttl=62 time=0.143 ms
      64 bytes from 192.168.2.2: icmp_seq=4 ttl=62 time=0.123 ms
      64 bytes from 192.168.2.2: icmp_seq=5 ttl=62 time=0.112 ms
       — 192.168.2.2 ping statistics -
      5 packets transmitted, 5 received, 0% packet loss, time 4067ms
      rtt min/avg/max/mdev = 0.112/1.486/6.563/2.542 ms
     mininet> h6 iperf -s -i 1 -p 80 &
mininet> h1 iperf -c h6 -t 10 -i 1 -p 80
      Client connecting to 192.168.2.2, TCP port 80
      TCP window size: 85.3 KByte (default)
        1] local 192.168.0.1 port 33176 connected with 192.168.2.2 port 80 (icwnd/mss/irtt=14/1448/8773)
      [ ID] Interval
                       Transfer
                                  Bandwidth
        1 0.0000-1.0000 sec 5.14 GBytes 44.1 Gbits/sec 1 1.0000-2.0000 sec 5.10 GBytes 43.8 Gbits/sec
   2. For the path h1 -> ra -> rb -> rc -> h6
      mininet> h1 ping -c 5 h6
      PING 192.168.2.2 (192.168.2.2) 56(84) bytes of data.
      64 bytes from 192.168.2.2: icmp_seq=1 ttl=61 time=6.04 ms
      64 bytes from 192.168.2.2: icmp_seq=2 ttl=61 time=1.19 ms
      64 bytes from 192.168.2.2: icmp seq=3 ttl=61 time=0.140 ms
      64 bytes from 192.168.2.2: icmp_seq=4 ttl=61 time=0.224 ms
      64 bytes from 192.168.2.2: icmp_seq=5 ttl=61 time=0.116 ms

    192.168.2.2 ping statistics

      5 packets transmitted, 5 received, 0% packet loss, time 4189ms
      rtt min/avg/max/mdev = 0.116/1.541/6.037/2.283 ms
     mininet> h6 iperf -s -i 1 -p 80 &
      mininet> h1 iperf -c h6 -t 10 -i 1 -p 80
      Client connecting to 192.168.2.2, TCP port 80
      TCP window size: 85.3 KByte (default)
        1] local 192.168.0.1 port 43302 connected with 192.168.2.2 port 80 (icwnd/mss/irtt=14/1448/10838)
       IDl Interval
                       Transfer
                                  Bandwidth
        1] 0.0000-1.0000 sec 4.55 GBytes 39.1 Gbits/sec
1] 1.0000-2.0000 sec 4.58 GBytes 39.3 Gbits/sec
1] 2.0000-3.0000 sec 4.69 GBytes 40.3 Gbits/sec
        1] 3.0000-4.0000 sec 4.56 GBytes 39.1 Gbits/sec
1] 4.0000-5.0000 sec 4.66 GBytes 40.1 Gbits/sec
        1] 5.0000-6.0000 sec 4.43 GBytes 38.1 Gbits/sec
        1] 6.0000-7.0000 sec 4.64 GBytes 39.8 Gbits/sec
It can be observed that the latency is less for the path h1 -> ra -> rb -> rc -> h6 from both ping's avg
time and irtt in iperf tool.
```

(d) The routing tables are printed using the code. The routing tables are in order ra, rb, rc.

For (a),							
***Routing ta	bles***				1 - 1		
Kernel IP rou							
Destination	Gateway	Genmask	Flags	Metri	c Ref	Use	Iface
192.168.0.0	0.0.0.0	255.255.255.0	U	0	0	0	ra-eth1
192.168.1.0	192.168.3.1	255.255.255.0	UG	0	0	0	ra-eth2
192.168.2.0	192.168.5.0	255.255.255.0	UG	0	0	0	ra-eth3
192.168.3.0	0.0.0.0	255.255.255.0	U	0	0		ra-eth2
192.168.5.0	0.0.0.0	255.255.255.0	U	0	0	0	ra-eth3
Kernel IP rou	ting table						
Destination	Gateway	Genmask	Flags	Metri	c Ref	Use	Iface
192.168.0.0	192.168.3.0	255.255.255.0	UG	0	0	0	rb-eth2
192.168.1.0	0.0.0.0	255.255.255.0	U	0	0	0	rb-eth1
192.168.2.0	192.168.4.1	255.255.255.0	UG	0	0		rb-eth3
192.168.3.0	0.0.0.0	255.255.255.0	U	0	0	0	rb-eth2
192.168.4.0	0.0.0.0	255.255.255.0	U	0	0	0	rb-eth3
Kernel IP rou							
Destination	Gateway	Genmask	Flags	Metri	c Ref	Use	Iface
192.168.0.0	192.168.5.1	255.255.255.0	UG	0	0	0	rc-eth3
192.168.1.0	192.168.4.0	255.255.255.0	UG	0	0	0	rc-eth2
192.168.2.0	0.0.0.0	255.255.255.0	U	0	0	0	rc-eth1
192.168.4.0	0.0.0.0	255.255.255.0	U	0	0		rc-eth2
192.168.5.0	0.0.0.0	255.255.255.0	U	0	0		rc-eth3
For (c),							
***Routing tab							
Kernel IP rout	_				5 6	A.,	- 6
Destination	Gateway	Genmask	_	Metric			Iface
192.168.0.0	0.0.0.0	255.255.255.0		0	0		ra-eth1
192.168.1.0	192.168.3.1	255.255.255.0		0	0		ra-eth2
192.168.2.0	192.168.3.1	255.255.255.0		0	0		ra-eth2
192.168.3.0	0.0.0.0	255.255.255.0		0	0		ra-eth2
192.168.5.0	0.0.0.0	255.255.255.0	U	0	0	Ø	ra-eth3
Kernel IP rout		C  -	F1		D-6		T.C
Destination	Gateway	Genmask		Metric			Iface
192.168.0.0	192.168.3.0	255.255.255.0		0	0		rb-eth2
192.168.1.0	0.0.0.0	255.255.255.0		0	0		rb-eth1
192.168.2.0	192.168.4.1	255.255.255.0		0	0		rb-eth3
192.168.3.0	0.0.0.0	255.255.255.0		0	0		rb-eth2
192.168.4.0	0.0.0.0	255.255.255.0	U	0	0	V	rb-eth3
Kernel IP rout	_	Conmode	F1 2 4 2	Motric	Dof	llaa	T. C. C. C.
Destination	Gateway	Genmask		Metric			Iface
192.168.0.0	192.168.4.0	255.255.255.0		0	0		rc-eth2
192.168.1.0	192.168.4.0	255.255.255.0		0	0		rc-eth2
192.168.2.0	0.0.0.0	255.255.255.0		0	0		rc-eth1
192.168.4.0 192.168.5.0	0.0.0.0	255.255.255.0 255.255.255.0		0 0	0		rc-eth2 rc-eth3
192.108.3.0	0.0.0.0	233.233.233.0	0	V	U	- U	rc-etiis
The routing table	s bays also book say	ad to the files of		+**+ 200	01 ro	11±00 (	+ +++ For
The routing tuble	is nave also been sav	ed to the files q1 ro	utes a.	LXLand	QI IC	utes t	. LXLIUI

the respective questions (a) and (c)

## Question 2

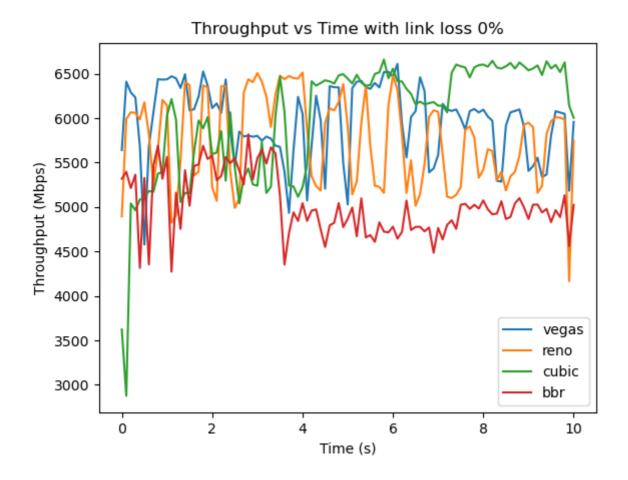
(a) The network topology given in the image is defined in NetworkTopo () class, and the class takes the linkloss parameter, which sets the value of the \_s1-s2 link packet loss percentage. The file q2.py can be run using python with sudo permissions. The file can be run using the below command. Where the values of config(required) can be b (or) c, values of congestion control(not required) can be any one of {Vegas, Cubic, Reno, BBR}, and values of link loss(not required) can be any integer between 0 and 100

python q2.py --config[-c] <config> --link-loss[-ll] <link loss> -congestion-control[-cc] <congestion control>

## Note:

- 1. If no congestion-control option is given for config b then the program plots Throughput for all values of all the above given congestion control mechanism, for the client h1.
- 2. If a congestion-control option is given for config c then the program plots Throughput for all hosts h1, h2, and h3, for the given value of congestion control in the same graph. If it is not given then it plots Throughput for all values of the congestion control mechanism for each given host {h1, h2, h3} in a different graph.
- 3. If no link-loss option is mentioned then the throughput analysis is done using 0 link loss for s1-s2 link.
- 4. The program does not enter into mininet CLI.
- **(b)** The client is run for 10sec and the plots are obtained by running the file using the below command.

sudo python q2.py --config b



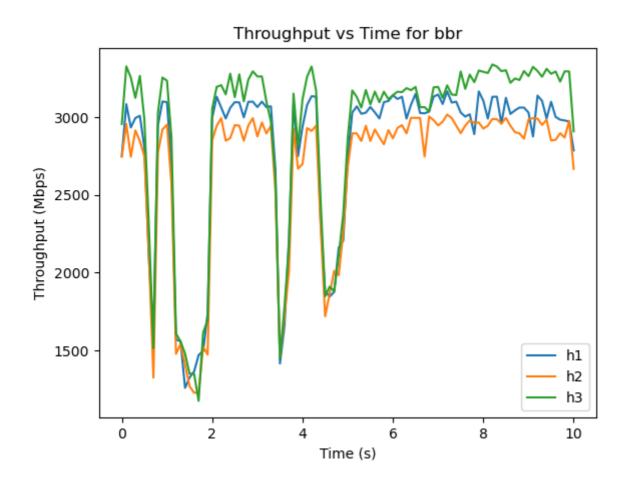
It can be observed that BBR congestion control gives less throughput, vegas shows a steady throughput after some time. Cubic, and Reno congestion control shows a decreasing throughput after some time, and

cubic also increases after the decrease. Reno and BBR have slow start, whereas Vegas and Cubic have fast start.

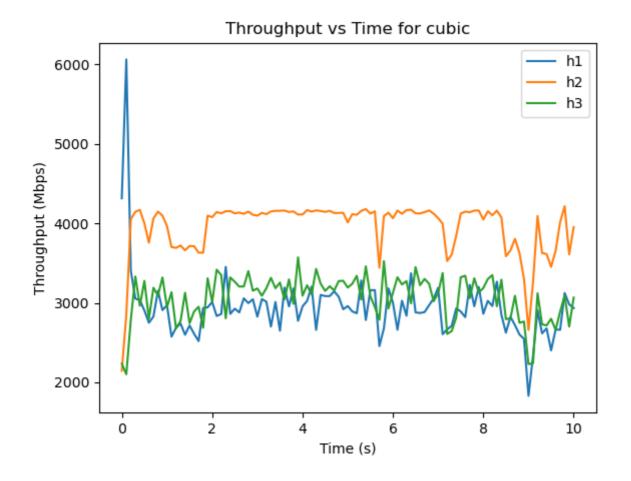
**(c)** The following commands were used to get the throughput plots for different hosts using the mentioned congestion control algorithm.

**Note:** File when run at different times might produce different plots depending on which host first starts sending packets.

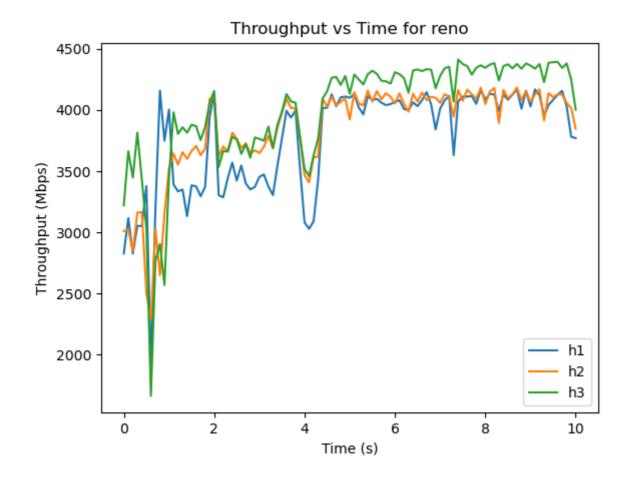
sudo python q2.py --config c -cc BBR



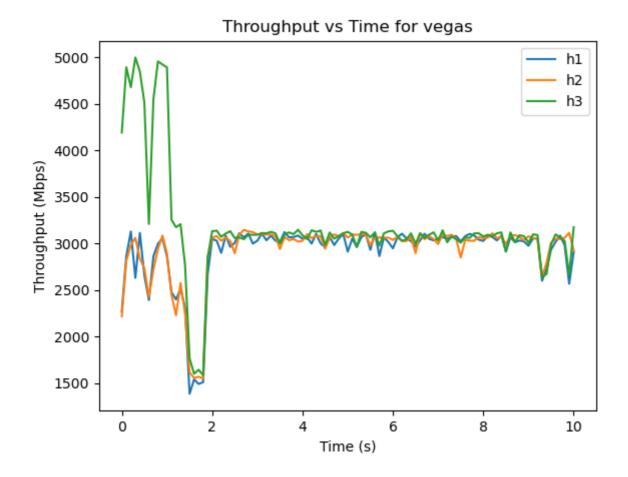
sudo python q2.py --config c -cc Cubic



sudo python q2.py --config c -cc Reno



sudo python q2.py --config c -cc Vegas

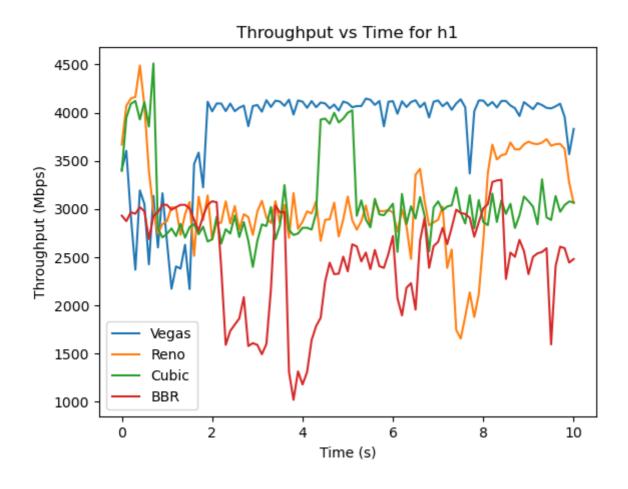


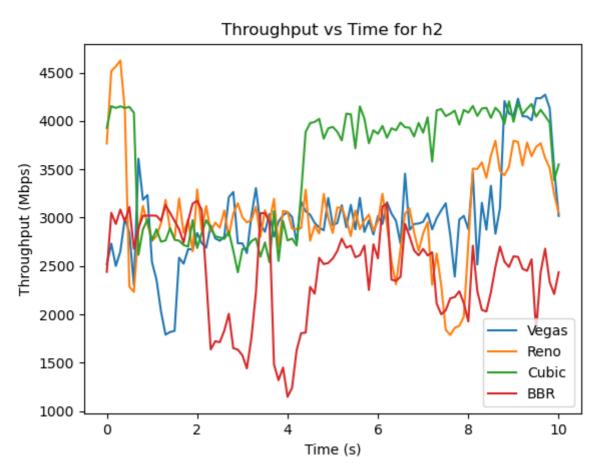
It can be observed that BBR congestion control gives less throughput, vegas appears to be fair for different hosts. BBR, and Reno congestion control show frequent changes. In case of BBR, it appears that all the hosts sharing the bandwidth respond similarly to the congestion control algorithm. Vegas appears to be the first to start fair sharing of bandwidth among the hosts, and it does not to show frequent changes in throughput if the network conditions are not changed. This may be because of the fact that Vegas uses RTT to throttle bandwith usage, and it does not use packet loss as a metric. Cubic and Reno use packet loss as a metric to throttle the andwidth usage, and hence they show frequent changes in throughput.

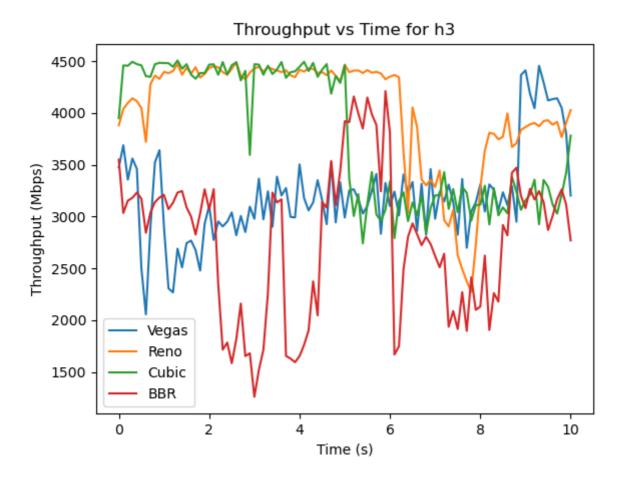
From questions (b) and (c) it can be observed that both network conditions and congestion control algorithm affect the throughput.

The following command was used to get the throughput plots for different congestion control algorithm for a given host.

sudo python q2.py --config c

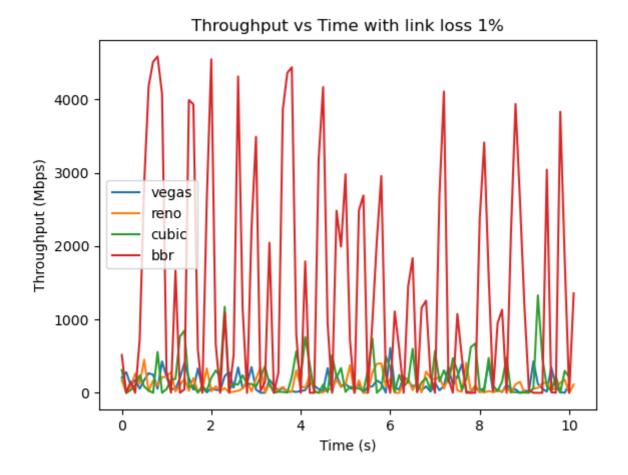




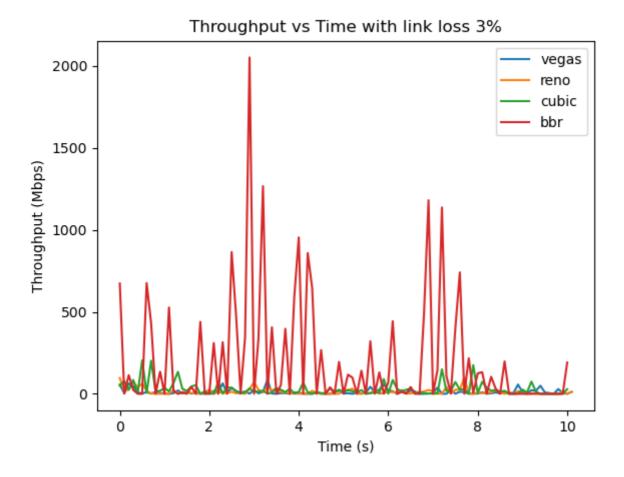


(d) The following command was used to get the throughput plots on client side for different congestion control algorithm for a given link loss $\{1, 3\}$ .

```
sudo python q2.py --config b -ll 1
```



sudo python q2.py --config b -11 3



It appears that BBR shows the highest throughput in case of packets losses and very high fluctuations in throughput during loss as compared to others. Vegas shows very less throughput change. This might be becuase of the fact that vegas uses RTT to throttle bandwith usage, and it does not use packet loss as a metric. From this we can say that BBR causes more congestion in the network as compared to others (which are sensitive to congestion in the network).

## References:

- 1. Iperf
- 2. Mininet
- 3. Mininet Python API
- 4. Mininet Router Implementation
  - 5. Congestion Control Algorithms