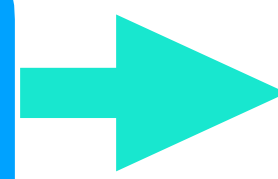
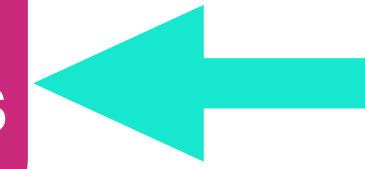


We have Items, every Item has weight and profit.



Unbounded
Knapsack Problems



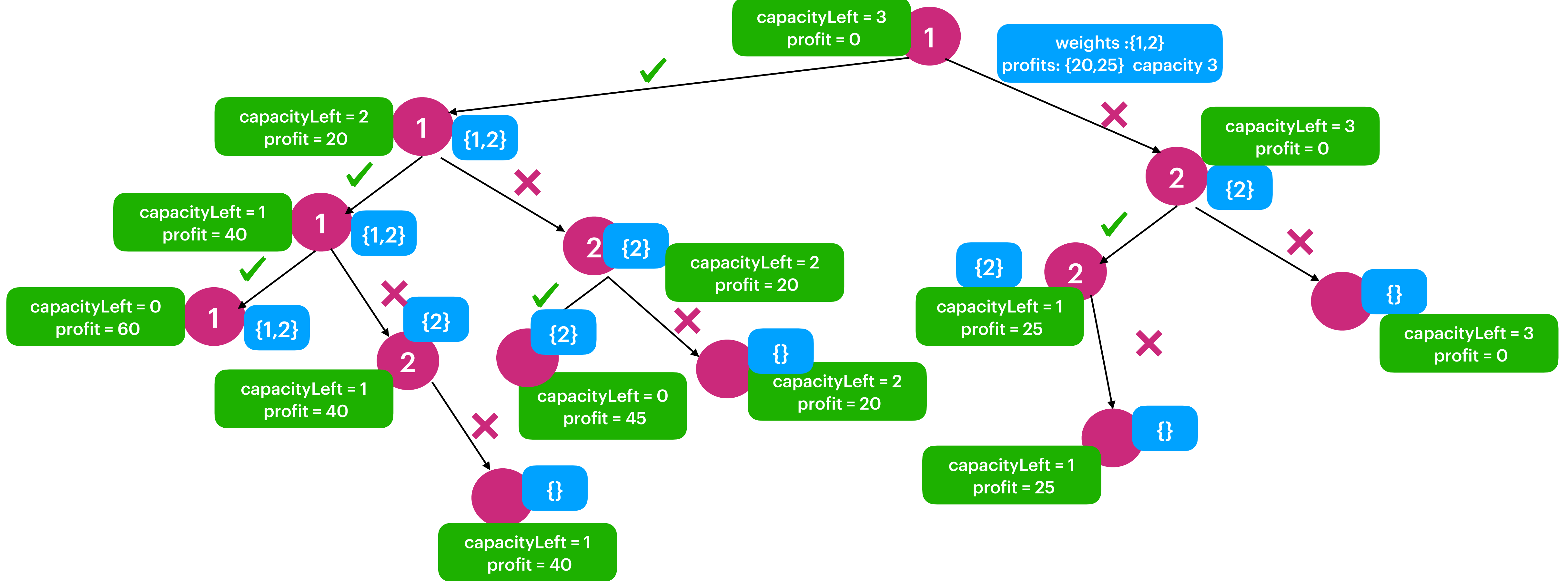
Problem Statement : Given item weights and their respective profits, itemWeight = {1,2,3} ItemsProfit = {5,2,4} , Get the max profit. We can have duplicate Items.

Constraints :

The knapsack capacity : 5

Expected Output :

Max Profit we can gain by choosing itemWeights {1,1,1,1,1} & respective profits {5,5,5,5,5} = 25



Time Complexity : 2^{n+c}
Space Complexity : $O(n+c)$

Only difference, we can move to next element only when we exclude item.

UnBounded Knapsack Problem

weights {1,2,3}
profits {5,2,4}

capacity : 2

Element

$\text{weights}[\text{element}] \leq \text{capacity}$



$\text{weights}[\text{element}] > \text{capacity}$



Element adds into the bag

Element can't be included into the bag

calculate profit with including current item

calculate profit without including current item

Just calculate profit without including current element.

$\text{dp}[i-1][c]$

$\max(\text{profits}[i] + \text{dp}[i][c - \text{weight}[i]], \text{dp}[i-1][c])$

Get max profit then return

itemWeight = {1,2,3}
ItemsProfit = {5,2,4},
capacity : 5

For zero capacity all the
items gives 0 profit.

When we are solving for the 1st
element.
i.e at index 0.

We don't have option of
comparison so

If weight of the 1st element is
≤ capacity
then value will be,
profits[0] +
dp[0][c-weight[0]].

f weight of the 1st element is
> capacity then value is 0.

Profits

5
2
4

Weight

1
2
3

Index

0
1
2

Target Sum					
0	1	2	3	4	5
0	5	10	15	20	25
0	5	10	15	20	25
0	5	10	15	20	25



Element
weight

≤ capacity

> capacity

Max(Including element , excluding element)
max(profits[i] + dp[i][c-weight[i]] , dp[i-1][c])

Exclude Element
dp[i-1][c]