

On ArrayList

ArrayList : internal using array :

ArrayList
size 10

5 elements to ArrayList : 5 to []

when size Of ArrayList > $n/2$, creates new array with double size then copy element in new array, removes the old one.

size 20 : [...5elements , 15 empty]

add 5 elements :

20 : [...10elements , 10 empty]

when you try to add 11 elements

$11 > 20/2$

40 : [...11elements , 29 empty]

Remove element from Array

int arr[] = {10 [0] , 11[1] , 12[2], 13[3] };

size : 4

index start 0 to n-1 : 0 to 3

Access element through index in constant time

arr[40] = ? returns value of index 40 in a constant time.

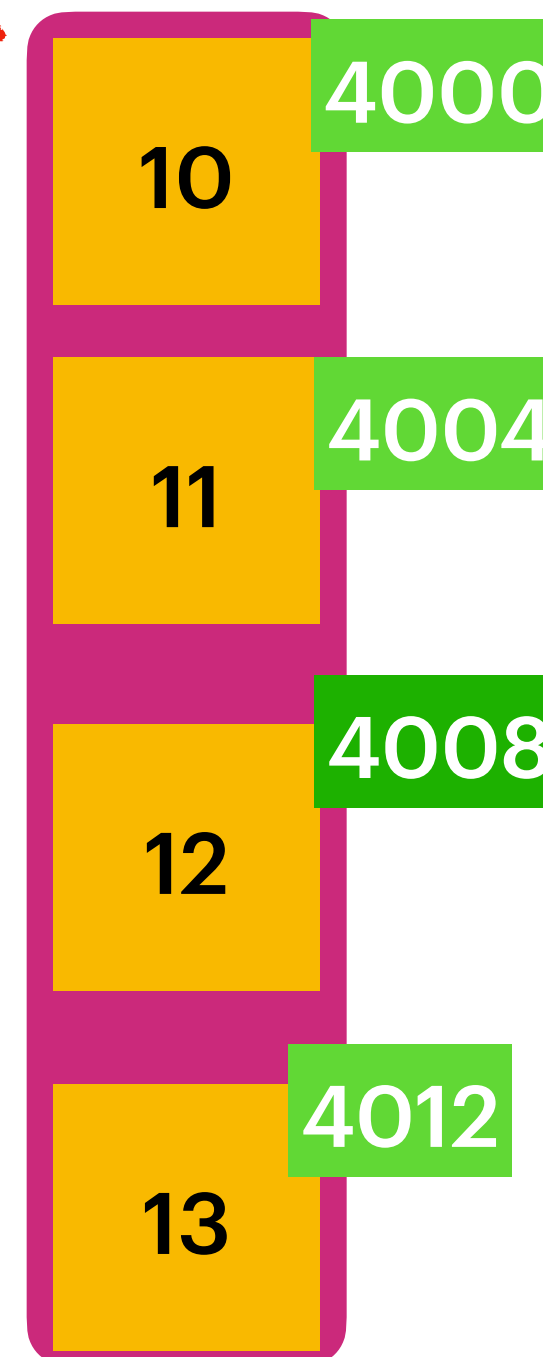
Array ::

```
int arr[] = {10 [0] , 11[1] , 12[2], 13[3] };
```

arr

$\text{arr}[3] = \text{arrAddr} + 3 * 4$
 $= 4000 + 12 = 4012 = 13$

$\text{arr}[0] = \text{arrAddr} + 0 * 4$
 $= 4000 + 0 = 4000$
 $= 10$



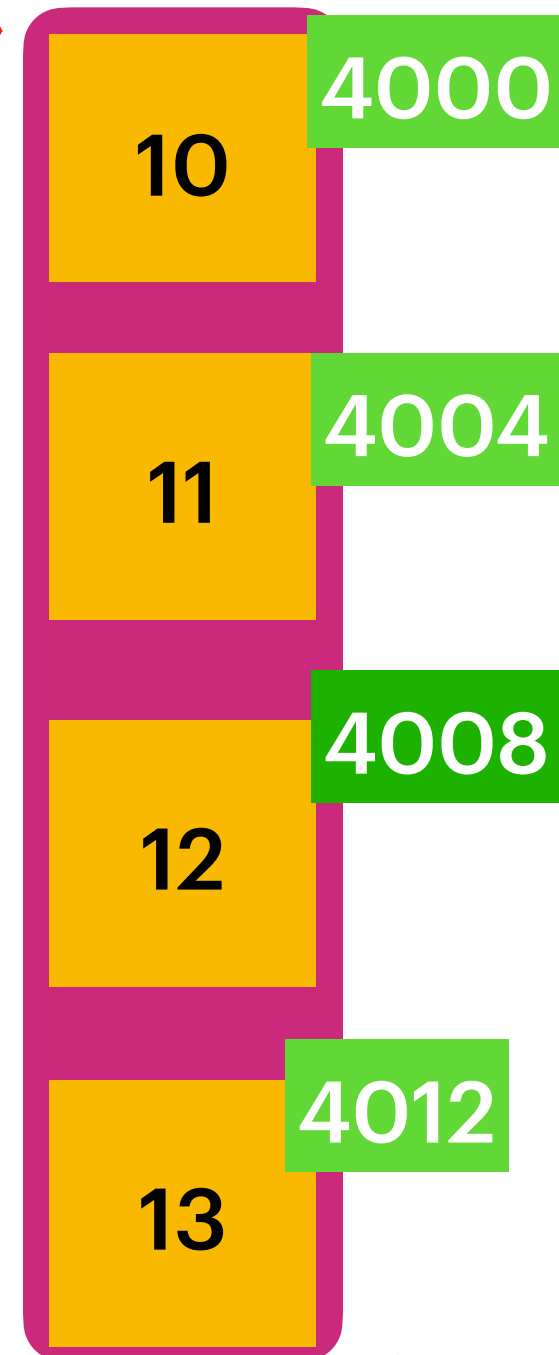
Start Address Hash : 4000

Start Address Hash : 4016

Array ::

```
int arr[] = {10 [0] , 11[1] , 12[2], 13[3] };
```

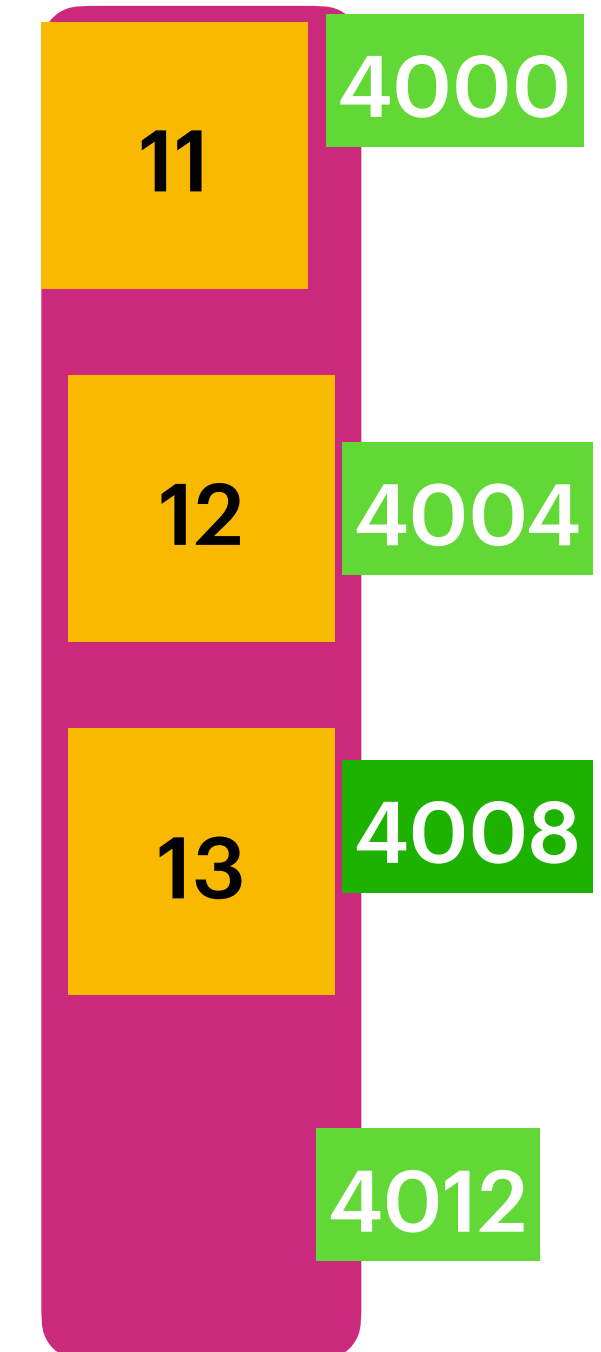
arr



Start Address Hash : 4000

Remove 10

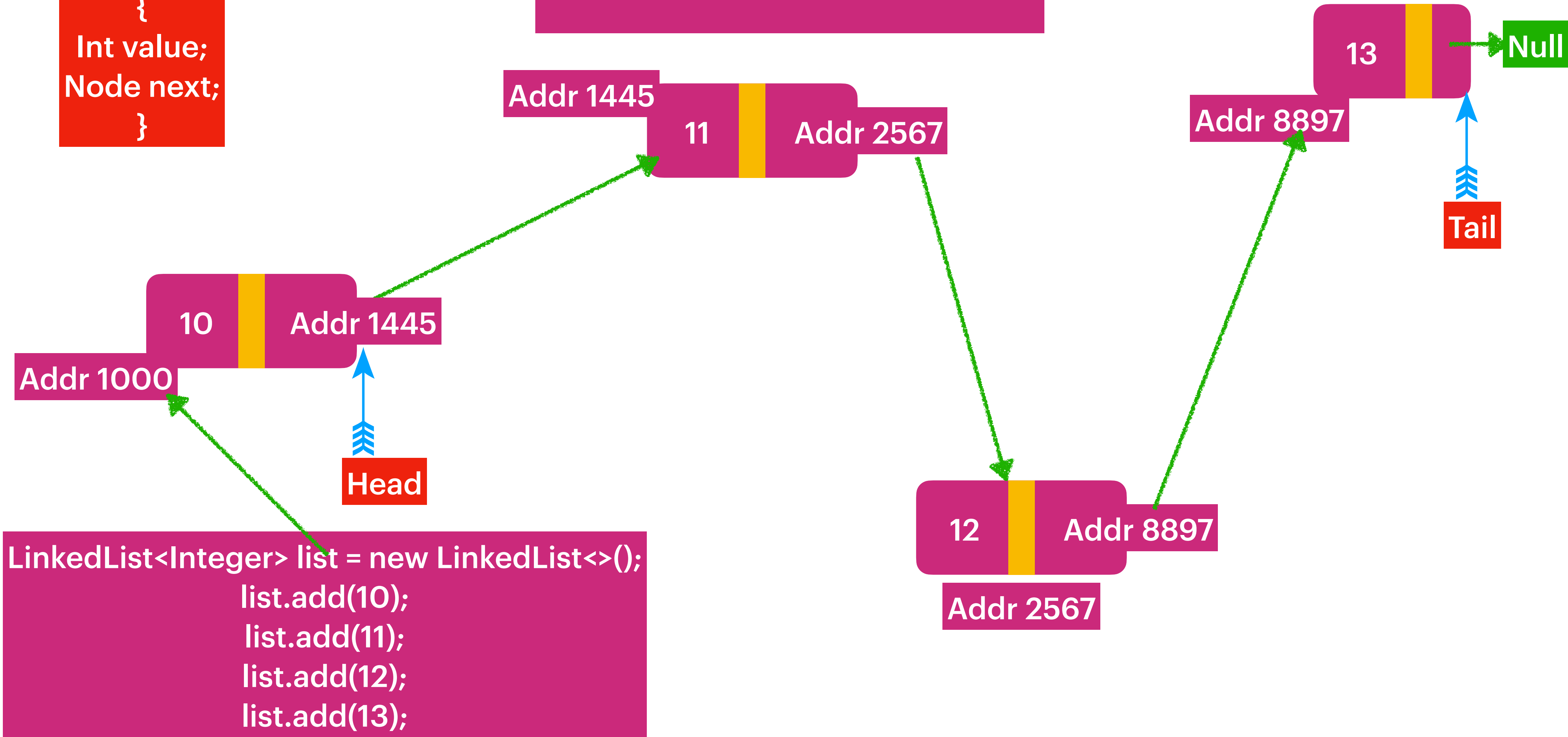
Start Address Hash : 4016



On ArrayList :
Insertion : $O(1)$
Deletion : $O(n)$
Search : $O(1)$ on index based
Seach : $O(n)$ on value based

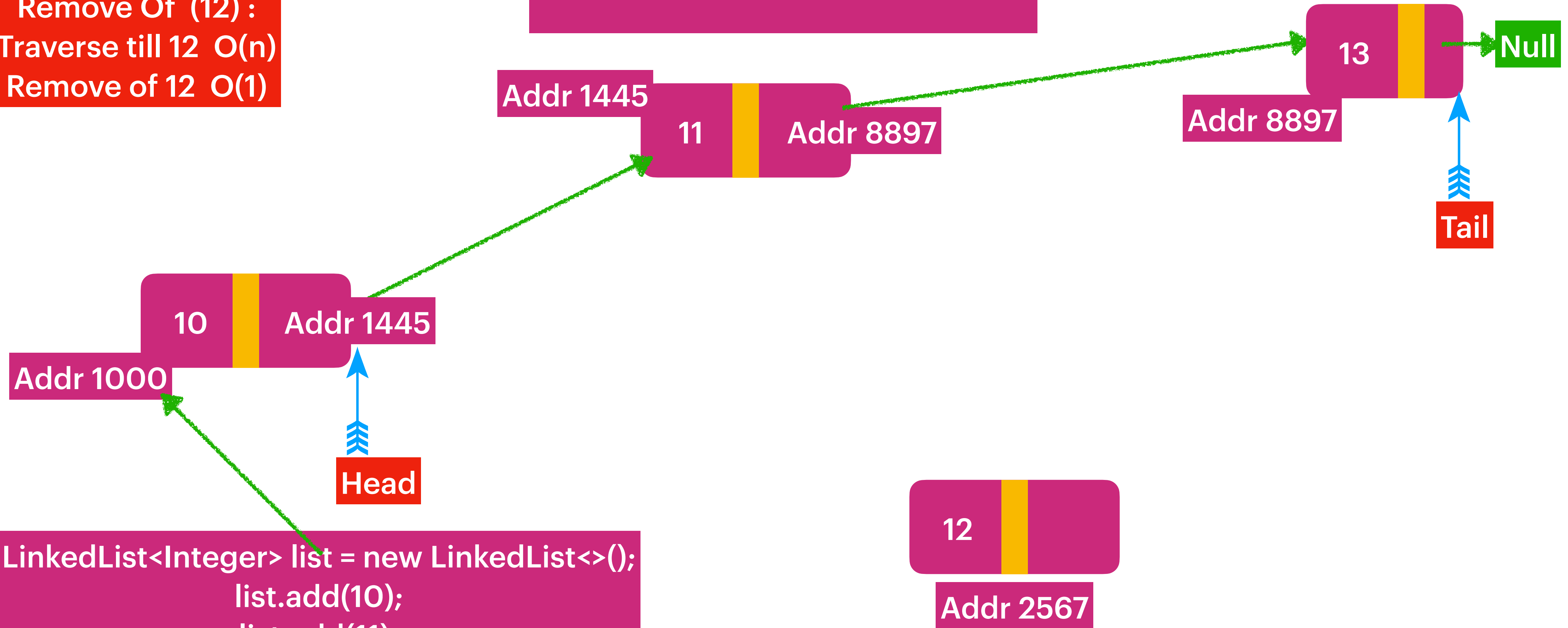
```
class Node
{
  Int value;
  Node next;
}
```

On Linked List :: (SingleLinkedList)
= {10 [0] , 11[1] , 12[2], 13[3] };



Remove Of (12) :
Traverse till 12 O(n)
Remove of 12 O(1)

On Linked List :: (SingleLinkedList)
= {10 [0] , 11[1] , 12[2], 13[3] };

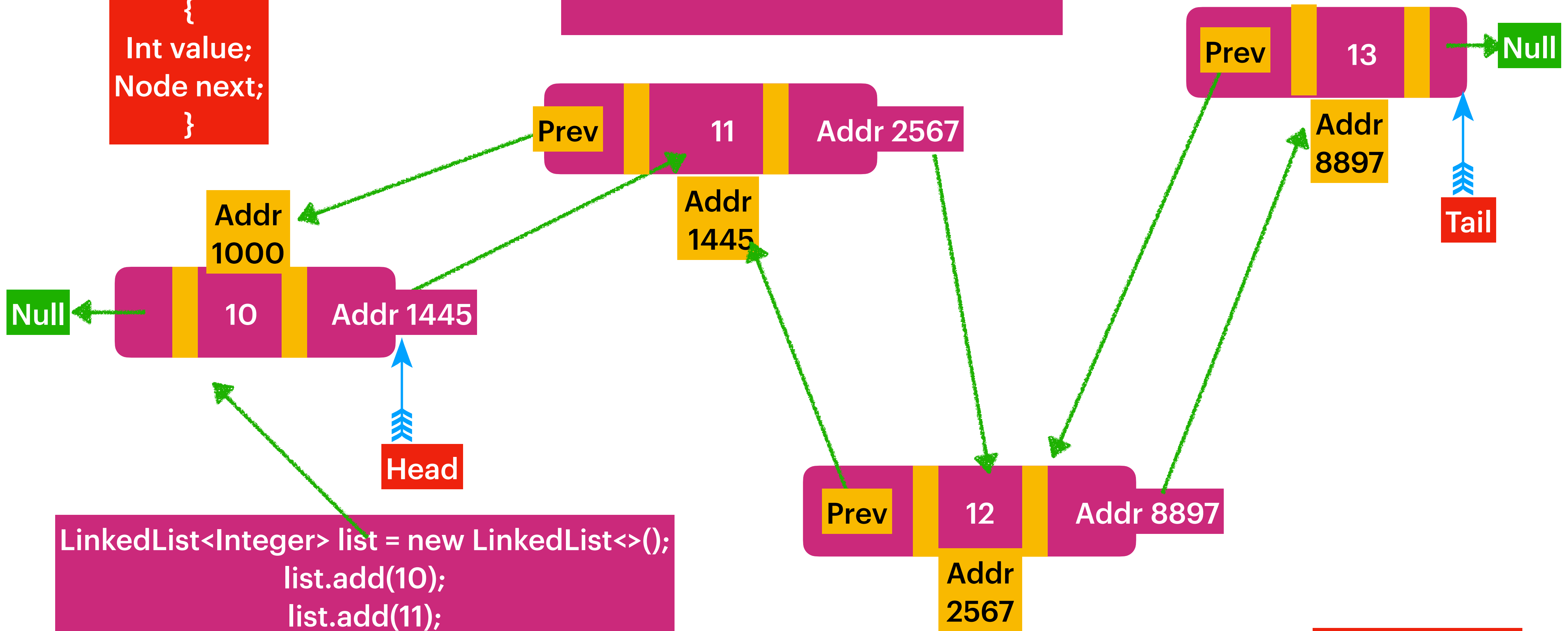


```
LinkedList<Integer> list = new LinkedList<>();  
list.add(10);  
list.add(11);  
list.add(12);  
list.add(13);
```

Deletion : O(n) Traverse + O(1) remove = O(n)

```
class Node
{
  Int value;
  Node next;
}
```

On Linked List :: (DoubleLinkedList)
= {10 [0] , 11[1] , 12[2], 13[3] };



```
LinkedList<Integer> list = new LinkedList<>();
list.add(10);
list.add(11);
list.add(12);
list.add(13);
```

Insertion : O(1)

On LinkedList :

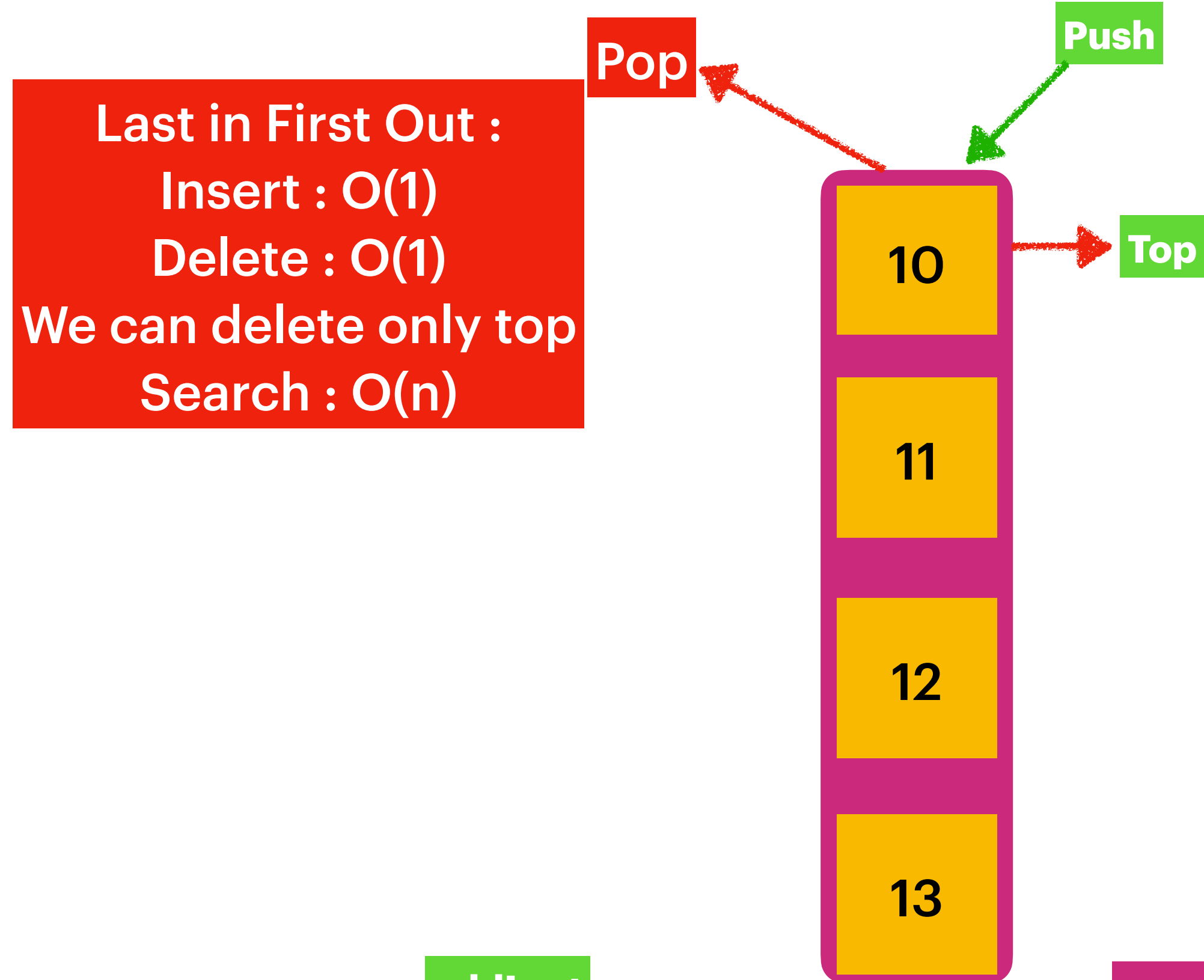
Insertion : $O(1)$

Deletion : $O(n)$ but $O(1)$ shifts

Search : $O(n)$ on index based

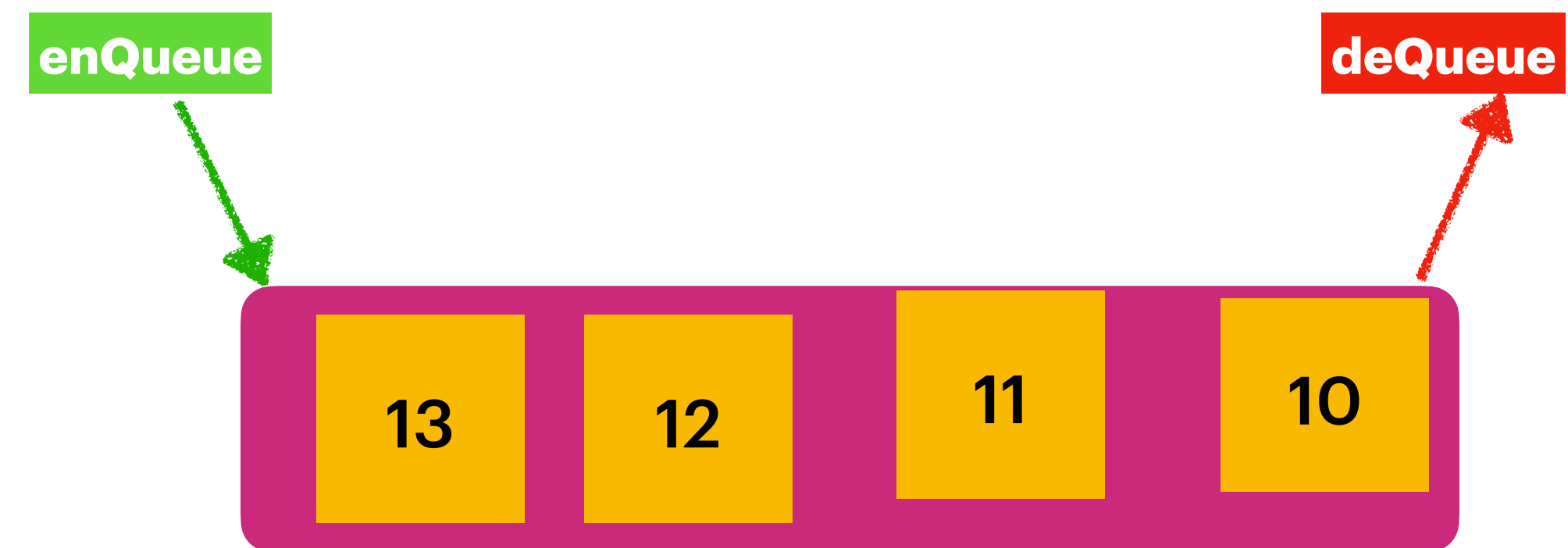
Seach : $O(n)$ on value based

Stack

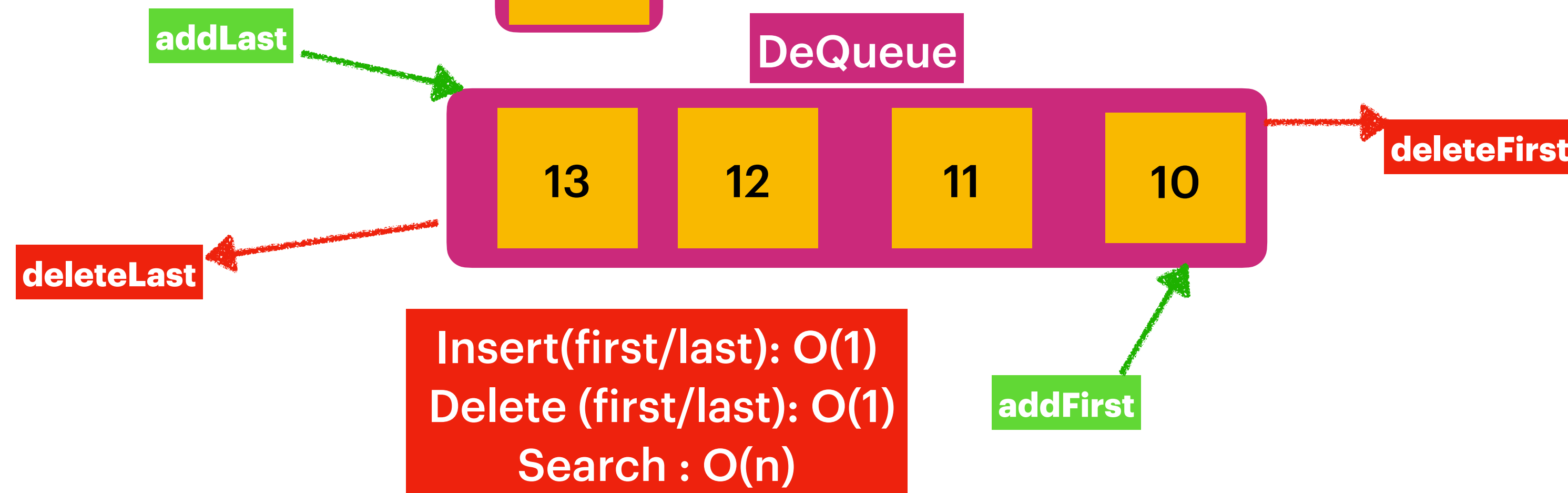


Last in First Out :
Insert : $O(1)$
Delete : $O(1)$
We can delete only top
Search : $O(n)$

Queue



First in First Out :
Insert : $O(1)$ adds to rear
Delete : $O(1)$ delete from front
We can delete only from front
Search : $O(n)$



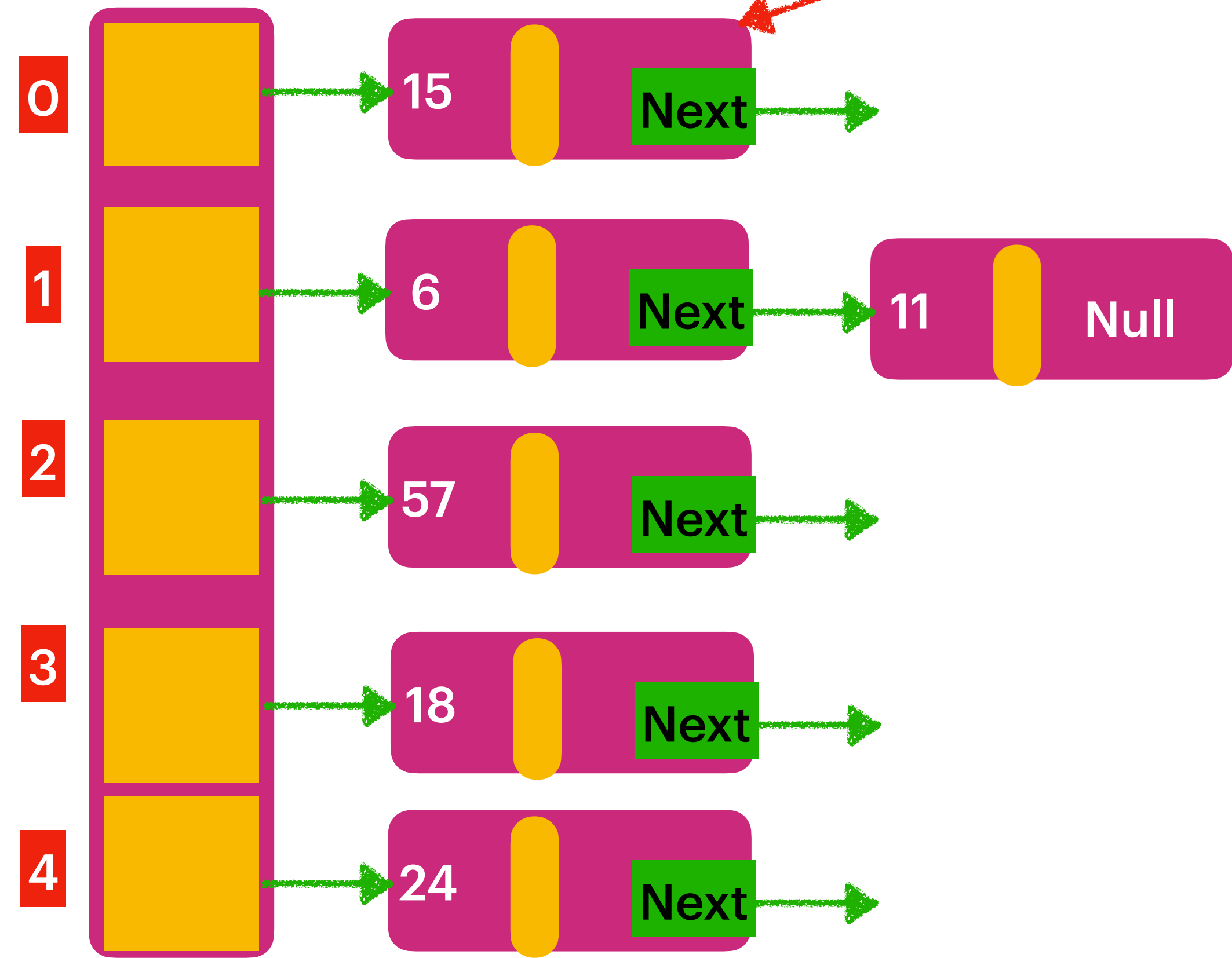
Insert(first/last): $O(1)$
Delete (first/last): $O(1)$
Search : $O(n)$

Hashing

6, 11, 15, 18, 57, 15, 24

$15\%5 = 0$

15%5 = 0 As key:15 is already exist so no insert.



6

Actual HashCode : 1156
hashCode()
method returns
its own value : 6

$6\%5 = 1$

$11\%5 = 1$

$18\%5 = 3$

$57\%5 = 2$

$24\%5 = 4$

search(11) :=> Get HashCode
of(11) = 11 => 11%5 = 1

delete(11) :=> Get HashCode
of(11) = 11 => 11%5 = 1

Insert : $O(1) / \log(n)$
Update : $O(1) / \log(n)$
Search : $O(1) / \log(n)$
Delete : $O(1) / \log(n)$

(6,12), (11,1), (15,22),
(18,6), (57,100), (15,27) (24,18)

6

Actual hashCode : 1156
hashCode()
method returns
its own value : 6

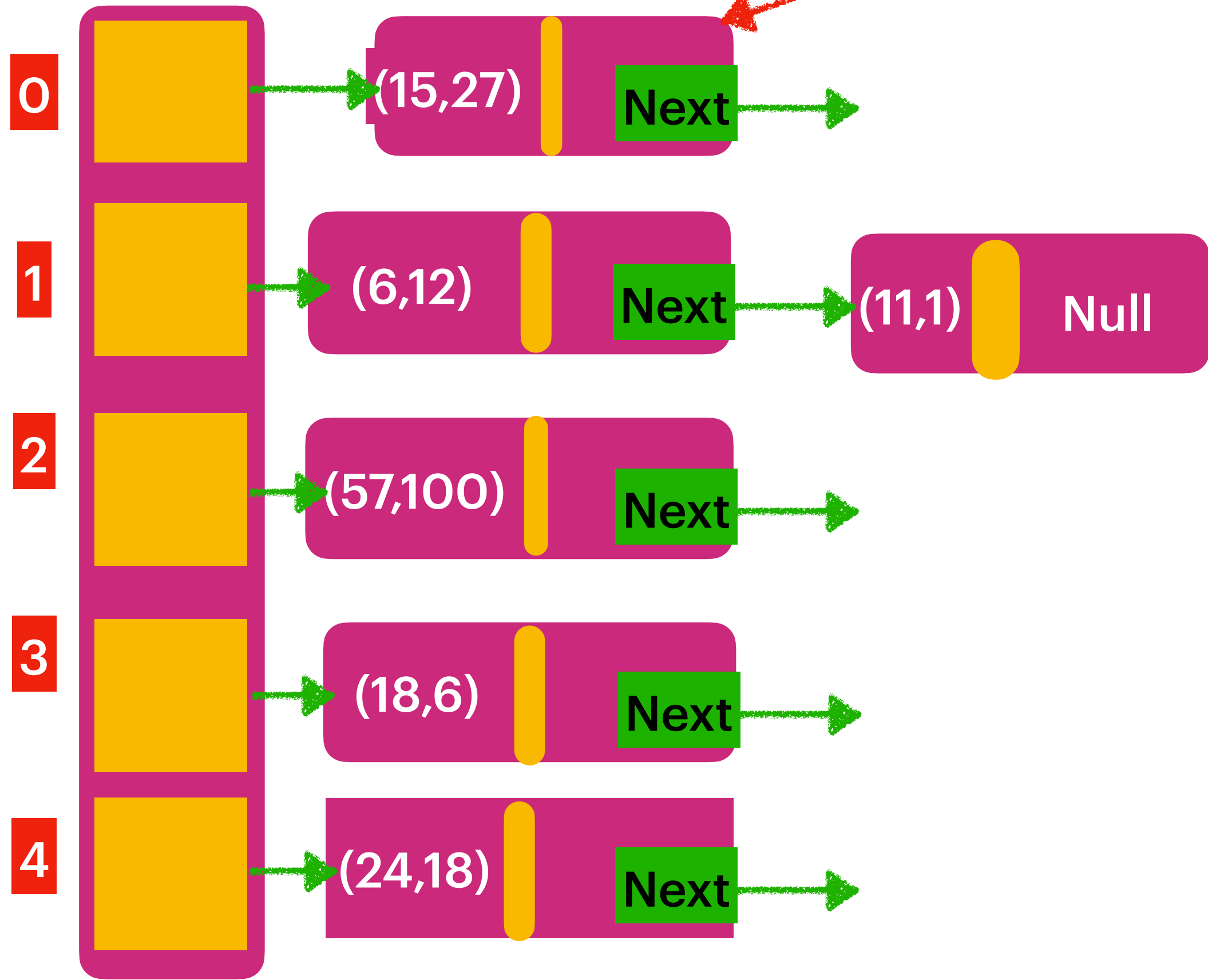
$6\%5 = 1$

$11\%5 = 1$

Hashing

$(15,22) = 15\%5 = 0$

$(15,27) \rightarrow 15\%5 = 0$ As key:15 is
already exist so no insert but
Replaces the Value.



search(11) :=> Get hashCode
of(11) = 11 => $11\%5 = 1$

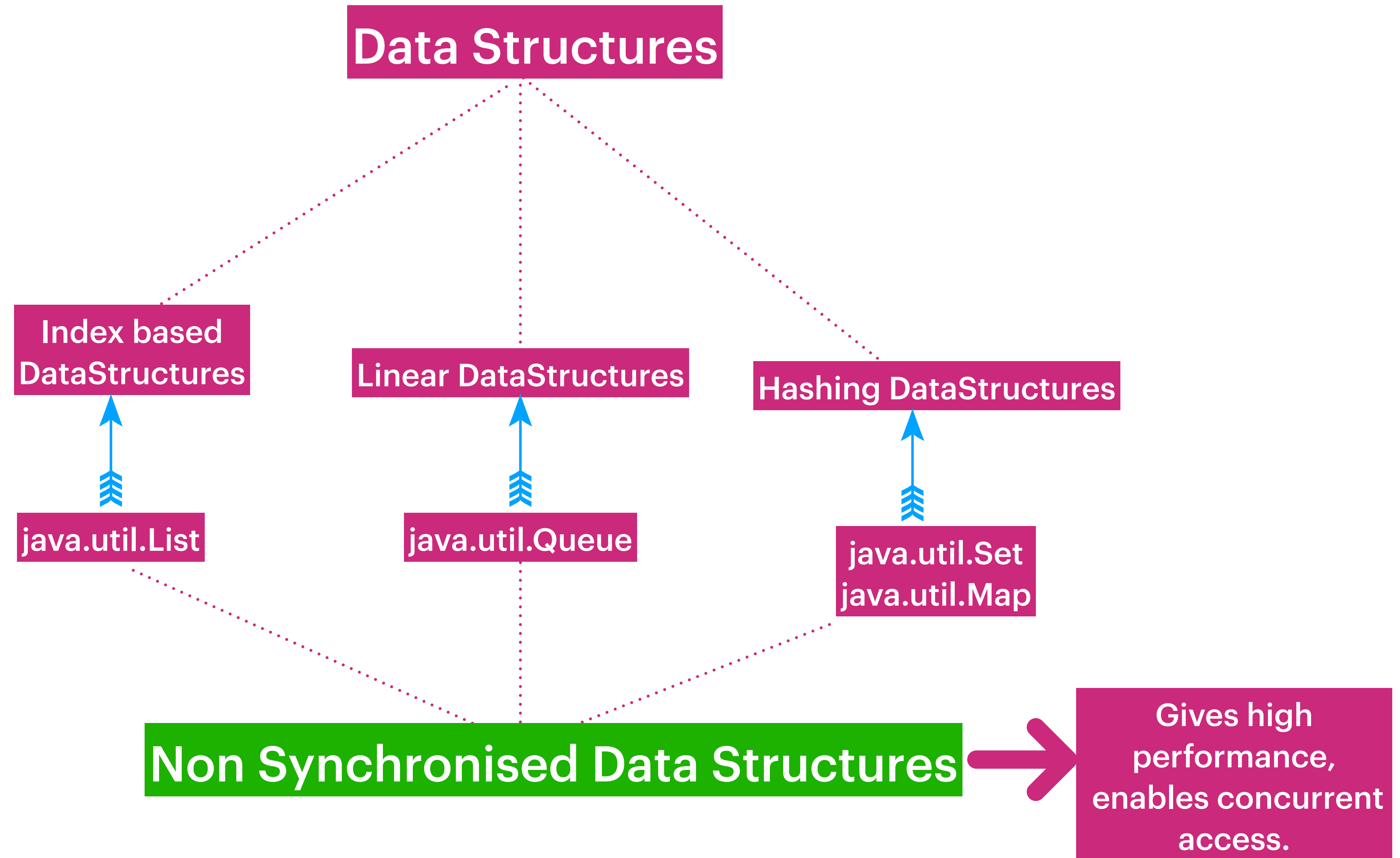
delete(11) :=> Get hashCode
of(11) = 11 => $11\%5 = 1$

Insert : $O(1) / \log(n)$
Update : $O(1) / \log(n)$
Search : $O(1) / \log(n)$
Delete : $O(1) / \log(n)$

$18\%5 = 3$

$57\%5 = 2$

$24\%5 = 4$



Synchronised Data Structures

Impacts the performance,
Use only in multithread
environment.

