

# AI 1: lab session 3

## Computational Logic

### Instructions

1. This is a ***team assignment*** to be completed in a team of two or three students and consists of a Python coding element and a report.
2. Answer all questions and ***formulate your answers concisely and clearly***.
3. Remember ***it is not allowed to use code not developed by your team***. Doing so constitutes plagiarism. If you use external sources (videos, websites, discussion fora, etc.) to develop your code make sure to ***clearly refer to them in your report and in your code***. Failing to disclose these sources constitutes again an instance of plagiarism.
4. ***Reports need to be written in L<sup>A</sup>T<sub>E</sub>X*** and submitted as a pdf ***separately from the code***. A template is available on Brightspace (under ‘Course overview’). The text should be ***authored by your team and your team only***.
5. Comply with the deadline provided on Brightspace. ***Deadlines are strict*** (penalties apply for late submissions, see below).
6. The code should be submitted in a ***separate zip file***.

**Assessment** The grade of this lab assignment will counts for ***10% of your final grade***. We ***subtract***  $2^{n-1}$  grade points for a submission that is between  $n - 1$  and  $n$  days late ( $n \geq 1$ ).

**Questions?** Get in touch with us! You can do so during labs, tutorials, using the discussion boards on Brightspace . Make sure you read the instructions to contact the us that are provided on Brightspace.

# 1 Model checking in propositional logic [20 pts]

In this programming exercise, we will implement model checking for propositional logic. A major part of the code has already been prepared for you, and can be found in the file `model.py`. The script runs without error, but it will not result in a complete program.

The program reads from the input two sets of sentences in propositional logic. The first set is the KB, and the second is a set of propositional sentences that we want to infer (if possible) from the KB. The parsing of the input has been implemented completely: you need not change anything in this part of the code, nor are you requested to understand the parsing of the input. There are also routines to evaluate a set of propositional sentences given a model. An example routine `evaluate_random_model` is available to show you how to use these evaluation routines. The routine `evaluate_random_model` generates a random model (random assignment of boolean values) for all identifiers (atoms) and then evaluates the truth value of the sets KB and INFER. Study this routine carefully.

An example input for the program can be found in the file `model1`:

```
{'KB': [
    'p+(q*r) <=> (p+q)*(p+r)',
    'p=>q <=> !p+q',
    'p=>q',
    'q=>r',
    'p'
],
'INFER': [
    'q',
    'r',
    'q*r',
    'true'
]
}
```

This example input represents the following sets KB and INFER:

$$\begin{aligned}\text{KB} &= \{(p \vee (q \wedge r)) \Leftrightarrow ((p \vee q) \wedge (p \vee r)), (p \Rightarrow q) \Leftrightarrow (\neg p \vee q), p \Rightarrow q, q \Rightarrow r, p\} \\ \text{INFER} &= \{q, r, q \wedge r, \text{true}\}\end{aligned}$$

Compile the program and run it using the above input. Your session should look like:

```
$python model.py model1
Complete input: {'KB': ['p+(q*r) <=> (p+q)*(p+r)', 'p=>q <=> !p+q', 'p=>q',
'q=>r', 'p'],
'INFER': ['q', 'r', 'q*r', 'true']}
Identifiers: ['p', 'q', 'r']
=====
KB:
((p + (q * r)) <=> ((p + q) * (p + r)))
((p => q) <=> (!p + q))
(p => q)
(q => r)
```

```

p
INFER:
q
r
(q * r)
True
=====

Randomly chosen model: {'p': True, 'q': False, 'r': False}
    KB evaluates to: False
    INFER evaluates to: False

KB entails INFER

The function check_all_models is not implemented yet
The goal of this assignment is to implement this yourself
Currently, this function always returns True

```

The program chose the random model `[p=true,q=false,r=false]` and then evaluated KB and INFER. Both evaluations returned false.

Implement the routine `check_all_models()` that determines whether  $KB \models INFER$ . If  $KB \models INFER$  is not the case, then the program should print a counter example on the output. Built-in Python libraries like `itertools` can be useful! Note that only built-in libraries are allowed in these exercises.

### Questions:

Run your code on the files `model1` `model2` and `model3` and report the results

## 2 Resolution in propositional logic [30 pts]

On Brightspace, you can find the files `resolution.py`. The file `resolution.py` contains the source code of a complete program that performs resolution on a (propositional) KB in conjunctive normal form (CNF).

In the file `resolution.py` the following KB is hard coded in the routine `init`:

$$KB = \{a \vee \neg b \vee \neg c \vee \neg d, c \vee \neg d, \neg a \vee \neg b, b \vee \neg d\}$$

Note that this KB is the CNF equivalent of:

$$KB = \{\neg a \Rightarrow (\neg b \vee \neg c \vee \neg d), \neg c \Rightarrow \neg d, a \Rightarrow \neg b, \neg b \Rightarrow \neg d\}$$

From the latter KB, it is easy to see that  $KB \models \neg d$ :

- Assume  $a$ : From  $a \Rightarrow \neg b$  we conclude  $\neg b$ . Using  $\neg b \Rightarrow \neg d$  we conclude  $\neg d$ .

- Assume  $\neg a$ : From  $\neg a \Rightarrow (\neg b \vee \neg c \vee \neg d)$  we conclude  $\neg b \vee \neg c \vee \neg d$ . Since  $\neg b \Rightarrow \neg d$  and  $\neg c \Rightarrow \neg d$ , we can conclude  $\neg d \vee \neg d \vee \neg d = \neg d$ .

So, if we add  $\neg \neg d = d$  to the KB, we should be able to infer the empty clause (i.e. false) using resolution. The resolution proof tree is given in the following figure:

$$\begin{array}{c}
 \frac{a \vee \neg b \vee \neg c \vee \neg d \quad c \vee \neg d}{a \vee \neg b \vee \neg d} \quad \frac{\neg a \vee \neg b}{b \vee \neg d} \\
 \frac{\quad}{\neg b \vee \neg d} \quad \frac{\quad}{\neg d} \quad \frac{\quad}{d} \\
 \frac{\quad}{\emptyset}
 \end{array}$$

The negation of the conclusion ( $\neg \neg d = d$ ) is added to the KB in the last lines of the routine `init` of `resolution.py`.

Running the script `resolution.py` gives you the following output:

```

$ python resolution.py
KB={~a,~b], [a,~b,~c,~d], [b,~d], [c,~d], [d]}
KB after resolution={~a,~b], [a,~b,~c,~d], [b,~d], [c,~d], [d], [~b,~c,~d],
[~a,~d], [a,~c,~d], [a,~b,~d], [a,~b,~c], [b], [c], [~b,~d], [~b,~c], [~a],
[~c,~d], [a,~d], [a,~c], [a,~b], [~b], [~d], [~c], [a], []=FALSE}
Resolution proof completed.

Proof:
Implement the function recursive_print_proof() yourself!

```

As you can see, the program generates all possible clauses that can be inferred from the KB. However, it generates many clauses that are not needed at all in the proof of the goal  $\neg d$ . For example, the clause  $[a, b, c]$  (i.e.  $a \vee \neg b \vee \neg c$ ) is inferred, but is not used in the proof of  $\neg d$ .

Study the code in `resolution.py`. Extend the program such that after resolution, a proof is printed in the routine `recursive_print_proof`. For your solution, it should not be necessary to change anything in the `Clause` class or in the helper functions provided. The output should be as follows:

```

Proof:
[a,~b,~d] is inferred from [a,~b,~c,~d] and [c,~d].
[~b,~d] is inferred from [~a,~b] and [a,~b,~d].
[~d] is inferred from [b,~d] and [~b,~d].
[]=FALSE is inferred from [d] and [~d].

```

Change the routine `init` such that your program can read a KB from standard input. The input for the above KB would be:

```

KB=[[~a,~b],[a,~b,~c,~d],[b,~d],[c,~d],[d]]

```

**Question:** Consider the following KB:

```
KB=[[a,~b],[a,b,~c,~d],[b,c,e]]
```

Add clauses to the KB such that *false* can be inferred. All clauses should be part of the inference. Include the proof generated by your code in your submission.

### 3 Prolog assignments [40 pts]

For the following exercises we will be using `swipl`, which is a freely available prolog implementation. It is already installed on the lab computers. For documentation, visit <http://www.swi-prolog.org>.

You are only allowed to use the built-in functions of Prolog, as well as any predicates already supplied by the assignment.

#### 3.1 [4 pts] Biblical family

Load the file `biblical.pl`, study the rules in the KB and answer the following questions.

**Questions:**

- a) Which Prolog query determines who is the grandfather of Lot? What is the answer?
- b) Which Prolog query determines all grandsons of Terach? What is the answer?

#### 3.2 [21 pts] Arithmetic with natural numbers

Download the file `arith.pl`. In this file some operations on so-called *Peano integers* are defined. For the sake of readability, Arabic numerals are used in the questions below. Be sure to use Peano numbers in your exercises!

Study the rules in the KB and answer the following questions.

**Questions:**

- a) What is a suitable query to ask the system whether  $3+2=5$ ? What does the system answer?
- b) What is a suitable query to ask the system whether  $3+2=6$ ? What does the system answer?
- c) Add predicates `even(N)` and `odd(N)`, that determine whether *N* is even or odd.
- d) Add a predicate `div2(N,D)` that determines whether the integer division  $N/2$  is equal to *D*. Your solution should not use the predicate `times`. Test your predicate for  $4/2 = 2$ ,  $3/2 = 1$  and  $0/2 = 0$ .

- e) Add a predicate `divi2(N,D)`, that computes the same result as `div2(N,D)`, but now using the predicate `times`. Of course, you need to test this predicate as well using the same numbers as before.
- f) Find an  $n$  such that  $2^n = 8$  using a suitable query. Add a predicate `log(X,B,N)` that determines whether  $B^N = X$ .
- g) Extend the KB with a predicate `fib(X,Y)`, where `fib` denotes the Fibonacci function. The predicate returns true if and only if  $fib(X) = Y$ . [Note:  $fib(0) = 0, fib(1) = 1, fib(n) = fib(n-1) + fib(n-2)$ ]
- h) In the course *Imperative programming* you learned that  $B^N$  can be computed in  $O(\log N)$  steps using the rules:  $a^{2b} = (a^2)^b$  and  $a^{2b+1} = a \cdot a^{2b}$ . Extend the KB with the predicate `power(X,N,Y)`, based on these rules. The predicate returns true if  $X^N = Y$ . Is this predicate an improvement over a direct  $O(N)$  computation? Why (not)?

### 3.3 [8 pts] Lists

In Prolog it is possible to use lists. For example, the following snippet of code determines the length of a list:

```
len([],0).
len([_|T],N) :- len(T,N1), N is N1+1.
```

Note that lists are not types, i.e. the elements of a list can be anything. For example, the query `len([1,2,[artificial,intelligence]],X)` will be answered with `X=3`.

#### Questions:

Download the file `arith.pl`, and extend it with the following functionality:

- a) `member(X,L)` returns true if `X` is a member of the list `L`.
- b) `concat(L,X,Y)` returns true if `L` is the concatenation of the lists `X` and `Y`.
- c) `reverse(L,R)` returns true if `R` is the reversal of the list `L`.
- d) `palindrome(L)` returns true if `L` is a palindrome.

### 3.4 [7pts] Maze

Consider the following maze:

m	n	o	p
i	j	k	l
e	f	g	h
a	b	c	d

**Questions:**

- a) Write a Prolog KB that represents the maze.
- b) Extend the KB with a predicate `path(X,Y)` that returns true if there exists a path from X to Y. Test the query `path(a,p)`, it should succeed.
- c) Try also `path(a,m)`. What is the result?