

Geometric algorithms for conformational analysis of protein loops

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Agenda

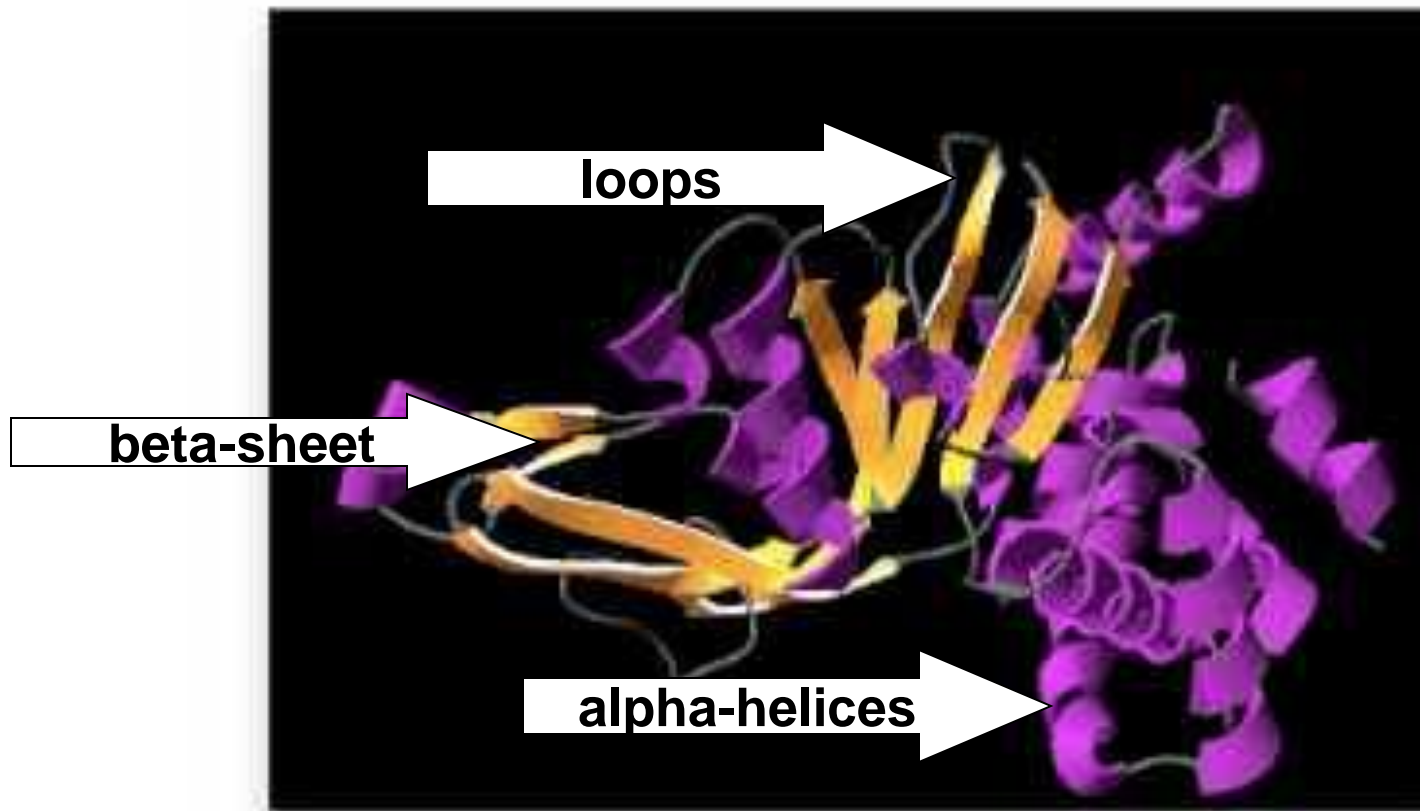
- Introduction and Motivation
- Problem Formulation
- Conformational sampling
- Conformational Space Exploration
- Results
- Conclusions
- Prospects



Introduction-outline

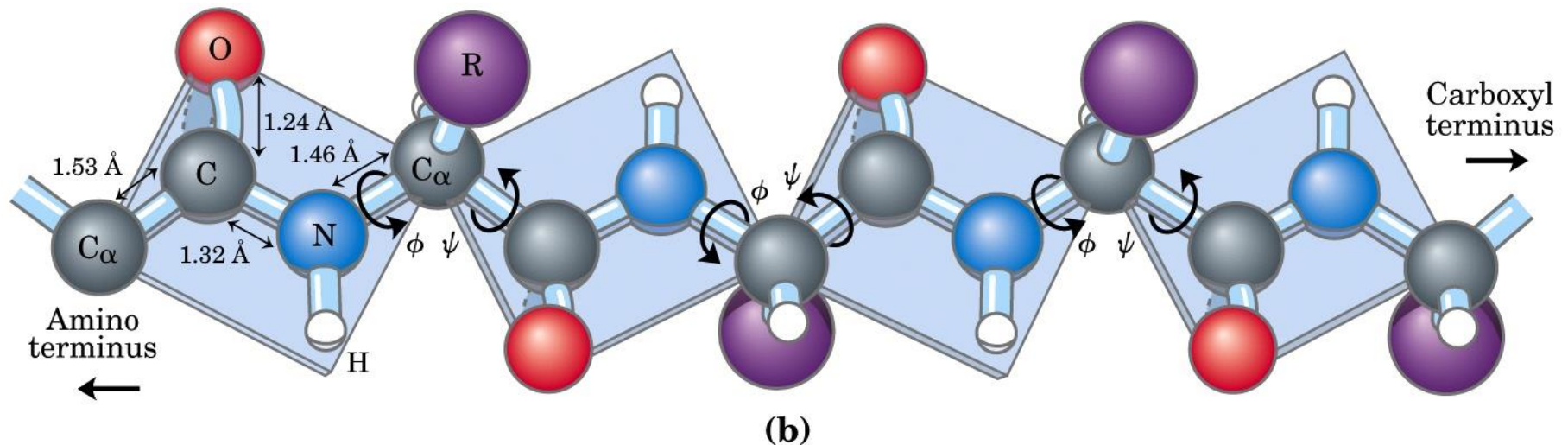
- Protein secondary structure
- Protein loop structure
- Loop closure and steric clashes
- Loop closure/clash filtering approaches
- Analogy to Robotics

Protein secondary structure



Protein loop structure

- Backbone chain
- Side chains (R)

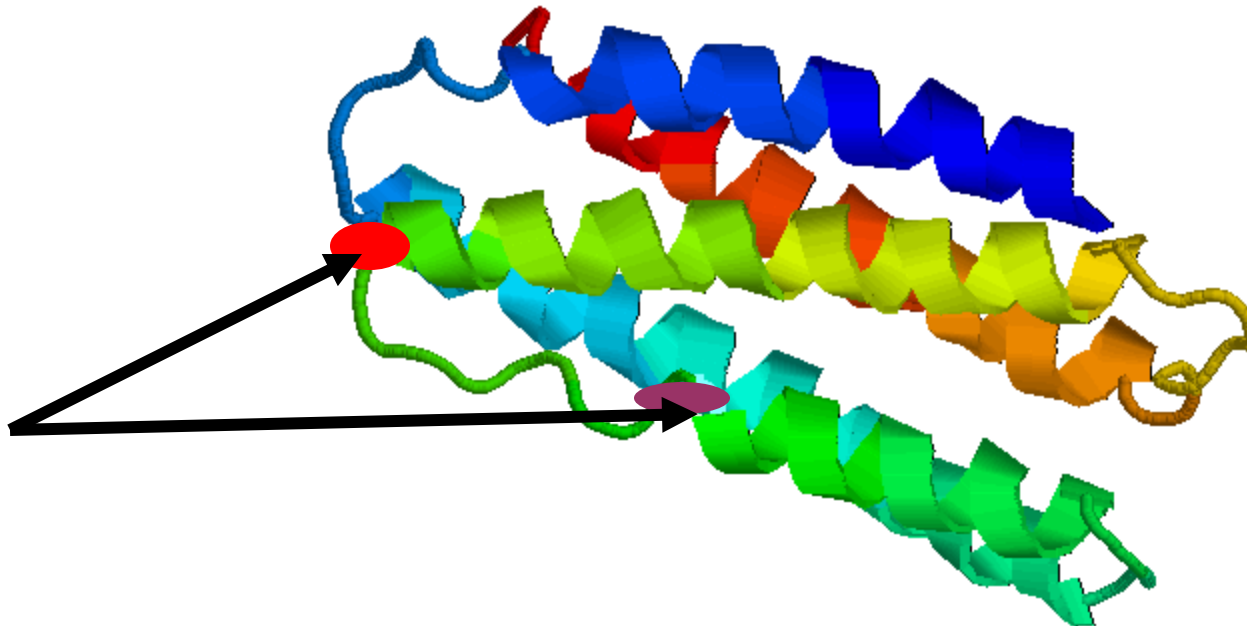


Constants : bond lengths, bond angles

Variables : dihedral angle

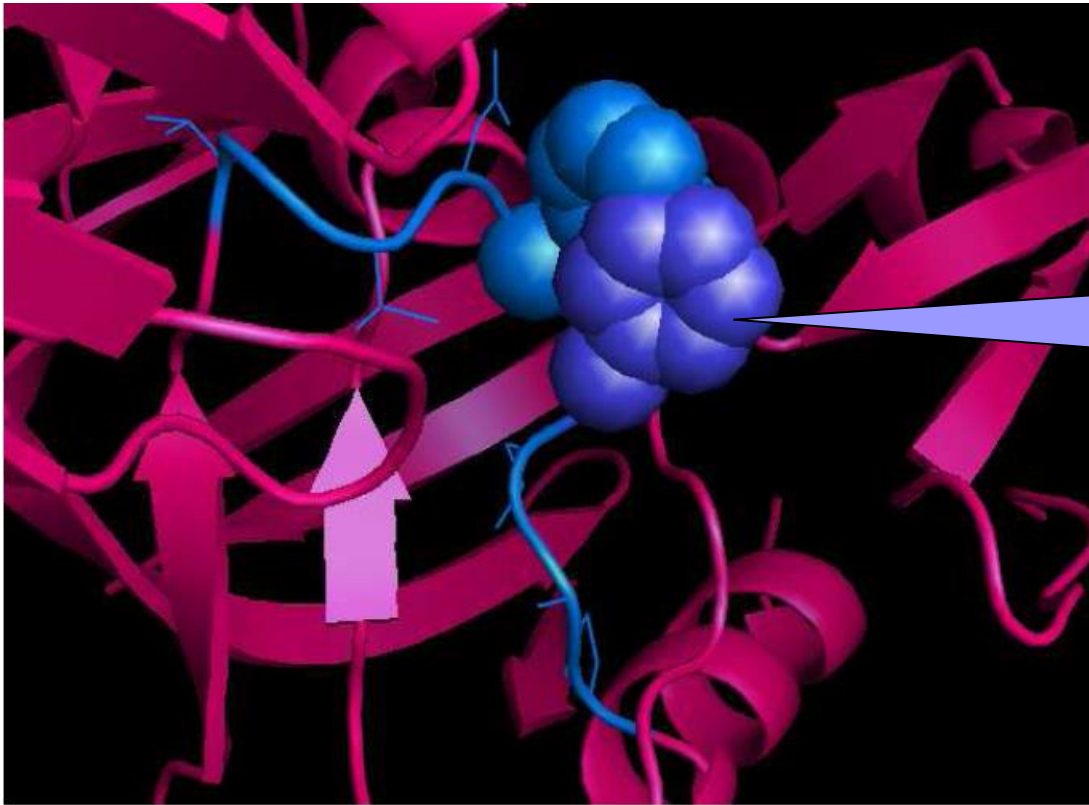
Loop closure

- The first and last atoms of the treated segment of a molecular chain must remain bonded with their neighbor atoms



Steric clash

- Distance(Atom1, Atom2) > Sum **van der Waals** radii of (Atom1, Atom2)



Example of
clashes between
2 side-chains in
1MPP

(a)

Self-clashes and Internal clashes are forbidden

Loop modelling approaches

- *Ab initio* loop prediction
 - Conformational search for minimum energy structures (simulating annealing, gradient decent)
- Database approach to loop prediction

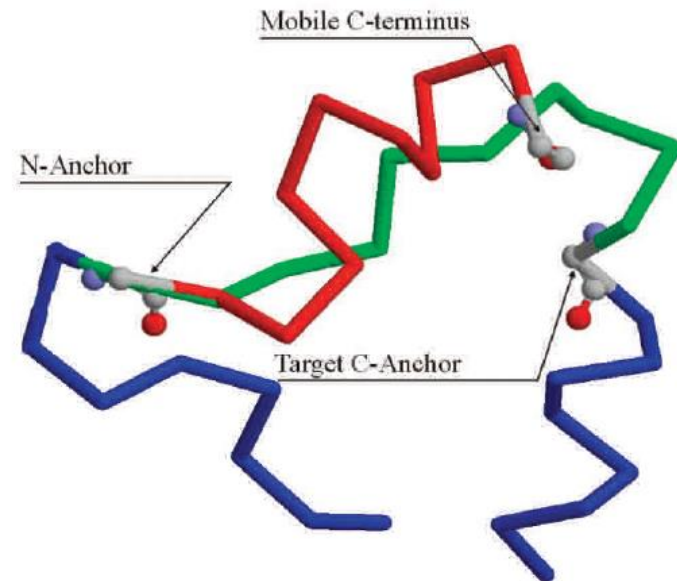
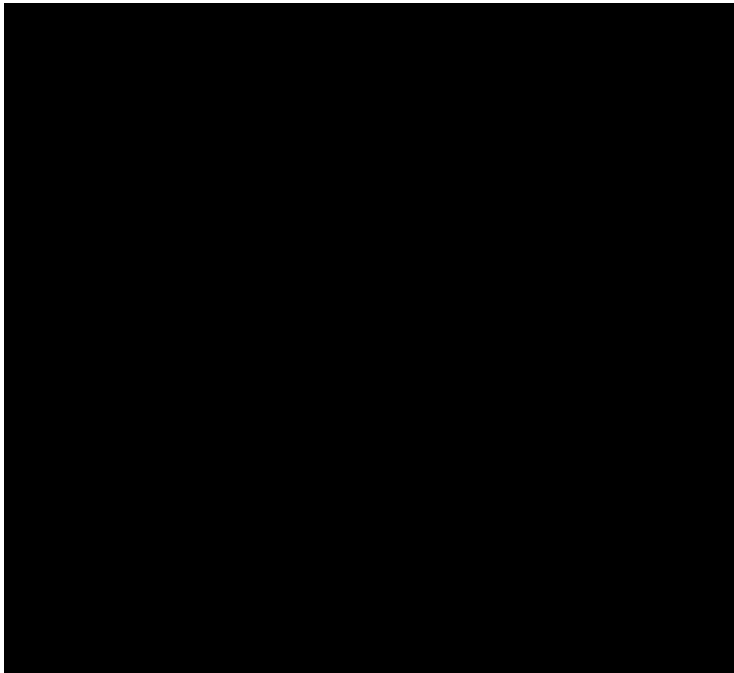
Loop Closure Approaches(1)

- Analytical – inverse kinematics
 - Wedemeyer and Scheraga(1999) solution for tripeptides with 6 degree of freedom, etc.
- Database based methods

Difficulty increases with the length of the molecular chain

Loop Closure Approaches(2)

- Optimization – e.g. CCD



Difficulty increases with the length of the molecular chain

Clash Filtering Approaches

- Energetic – accepting/rejecting a conformation according to some energetic (repulsive VdW energy) cutoff
- Geometric – “clash grids”

The higher the number of atoms → more critical efficiency



In Robotics – motion planning

- Collision with obstacles/self-collisions are forbidden
- For kinematic loops - closure constraints exist
- Exploration of the configurations space, searching for feasible configurations
- Sampling-based techniques capture the topology of the feasible space within a data-structure (graph or a tree) by performing random exploration



Motivation

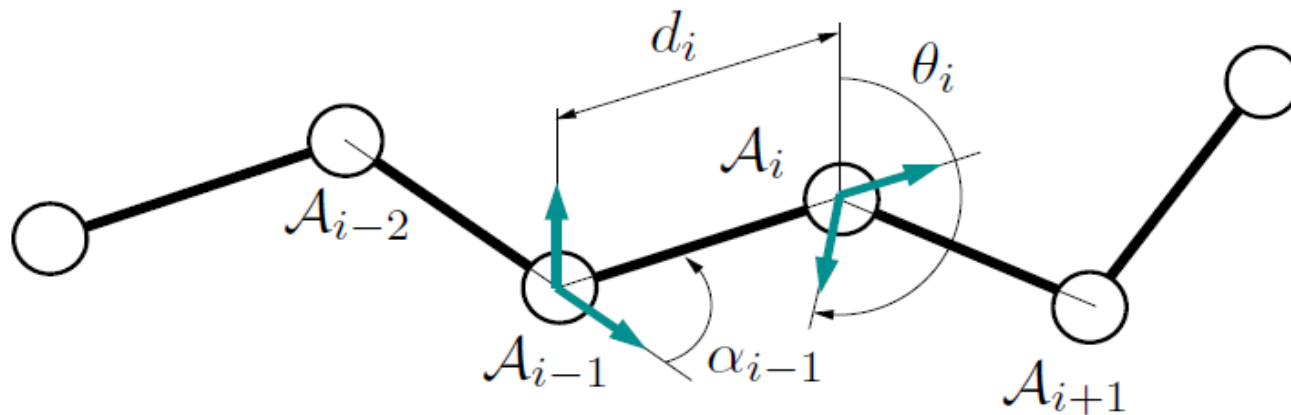
- Filter unfeasible loop conformations to aid searching conformational space for various application:
 - Protein loop modeling
 - Molecular simulations: conformational changes under environmental conditions



Problem formulation - outline

- Geometric Model of the molecule
- Geometric Constraints

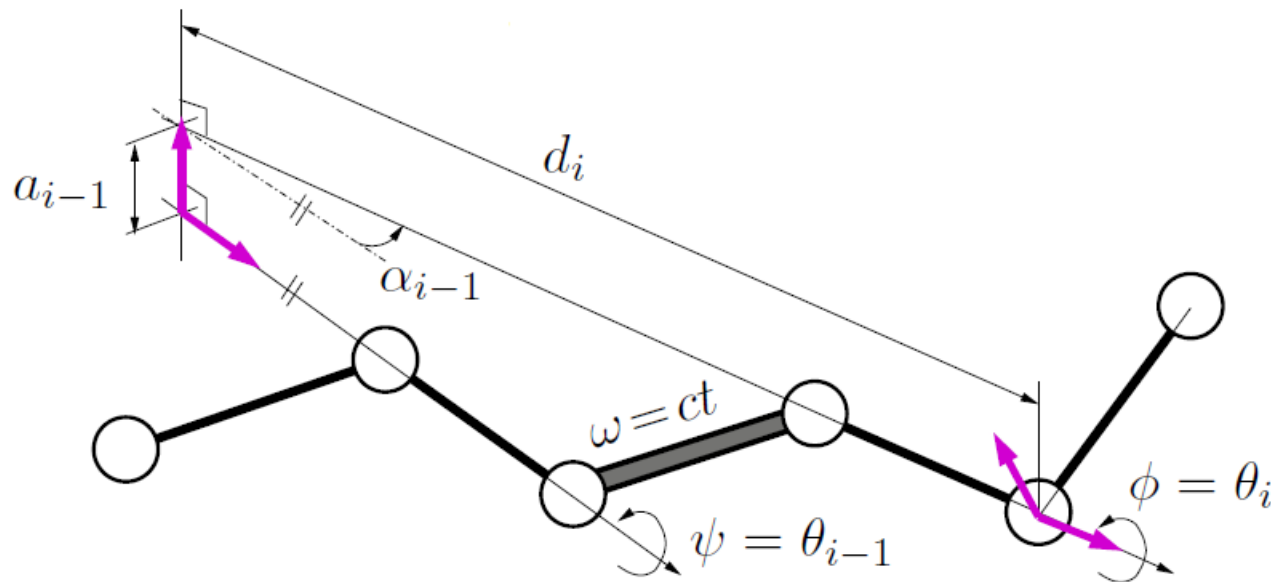
Geometric Model: Molecular chain model



- Molecules – A_{i-2}, \dots, A_{i+1}
- Bond length, d_i , between A_{i-1}, A_i
- Bond angle, α_{i-1} , between A_{i-2}, A_{i-1}, A_i
- Dihedral angle, θ_i , formed by $A_{i-2}, A_{i-1}, A_i, A_{i+1}$

Geometric Model: Articulated mechanism

- We assume that only dihedral angles are variable parameters (rigid geometry)



The Homogeneous Transformation Matrix

- We defines the relative location of consecutive frames as follows:

$${}^{i-1}T_i = \begin{pmatrix} C\theta_i & -S\theta_i & 0 & 0 \\ S\theta_i C\alpha_{i-1} & C\theta_i C\alpha_{i-1} & -S\alpha_{i-1} & -S\alpha_{i-1}d_i \\ S\theta_i S\alpha_{i-1} & C\theta_i S\alpha_{i-1} & C\alpha_{i-1} & C\alpha_{i-1}d_i \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

- All parameters as in Molecular chain model that was defined previously
- S - sinus
- C - cosines

Geometric Constraints

- Loop Closure

$${}^0T_n = {}^0T_1 {}^1T_2 \dots {}^{n-1}T_n$$

- Clash avoidance(70% of VdW radii)

(Condition must be satisfied between atoms of the articulated segment and between atoms of the rest of the molecule)



Conformational Sampling - outline

- Random Loop Conformation
- Backbone Conformation with closure
- Loop decomposition
- RLG algorithm
- RLG Algorithm: example
- Clashes and Side-Chain Conformation

Random Loop Conformation

Compute random conformation achieving loop-closure and clash avoidance constraints in 3D.

Algorithm 1: RANDOMLOOPCONF

→ **input** : the *loop*, the rest of the *protein*
→ **output** : the conformation q
begin
→ $q_b \leftarrow \text{RANDOMBACKBONECONF}(\text{loop.bkb});$
→ **if not** CLASHCHECK(q_b , *loop.bkb*, *protein*) **then**
→ **if** $q_s \leftarrow \text{GENERATESIDECHAINS}(q_b, \text{loop}, \text{protein})$ **then**
→ $q \leftarrow \text{COMPOUNDCONF}(\text{loop}, q_b, q_s);$
→ **else return** *Failure*;
 else return *Failure*;
end

Array of dihedral angles: $\theta_1, \theta_2, \dots, \theta_n$

A generic 3D collision detection algorithm (T. Siméon, C. van Geem, 2001)

Sample angles randomly at random side-chain order.
Check for clashes

Solving multivariable non-linear algebraic equations ${}^0T_n = {}^0T_1 {}^1T_2 \dots {}^{n-1}T_n$

- No efficient general solution exists
- Known – 6 variables are dependent
- How to obtain independent variables satisfy closure equations?

Random Loop Generator (RLG)

Closure interval of J_k - values of J_k that make possible to close the mechanism are the ones for which P_{k+1}^k is reachable (in position and orientation) by the chain

$$J_n \rightarrow J_{k+1}$$

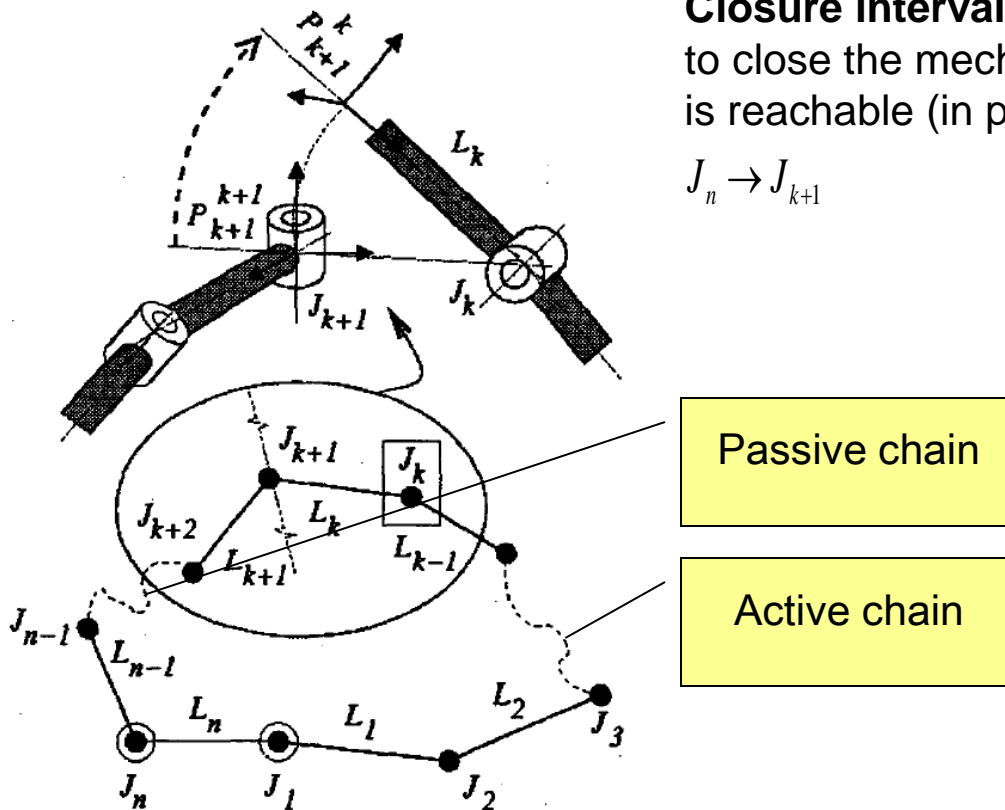
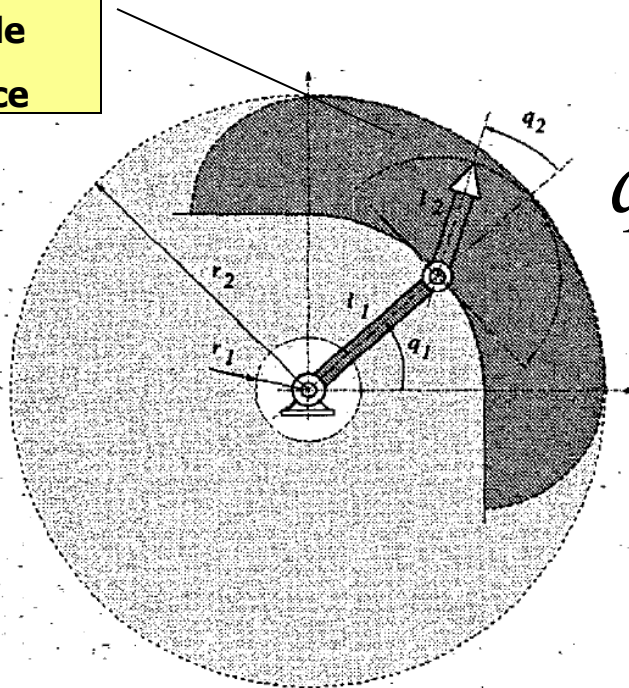


Figure 3: Illustration of the principle of RLG.

RLG: Approximated reachable Workspace

Reachable
workspace



$$q_1 \in [0, \pi/2]; q_2 \in [-\pi/2, \pi/2]$$

Figure 4: Conservative approximation of the reachable workspace of a RR planar chain.

RLG: Closure Interval Computation

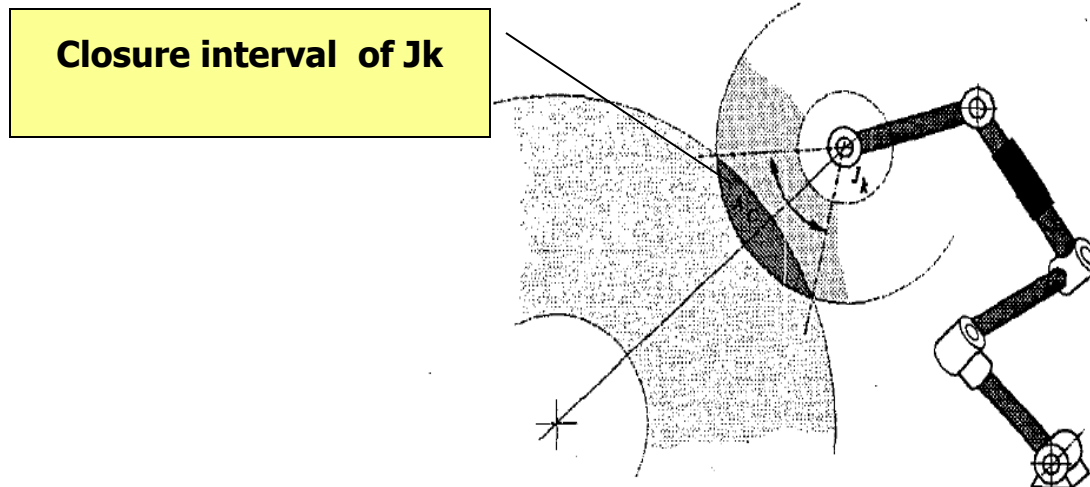


Figure 5: Illustration of the intersections computed for the closure interval estimation in the case of a revolute joint.

RLG Algorithm

Algorithm 2: RANDOMBACKBONECONF

input : the *backbone*

output : the conformation q_b

begin

$q^a \leftarrow \text{SAMPLE_}q^a(\text{backbone});$

if $q^p \leftarrow \text{COMPUTE_}q^p(\text{backbone}, q^a)$ then

└ $q_b \leftarrow \text{COMPOUNDCONF}(\text{backbone}, q^a, q^p);$

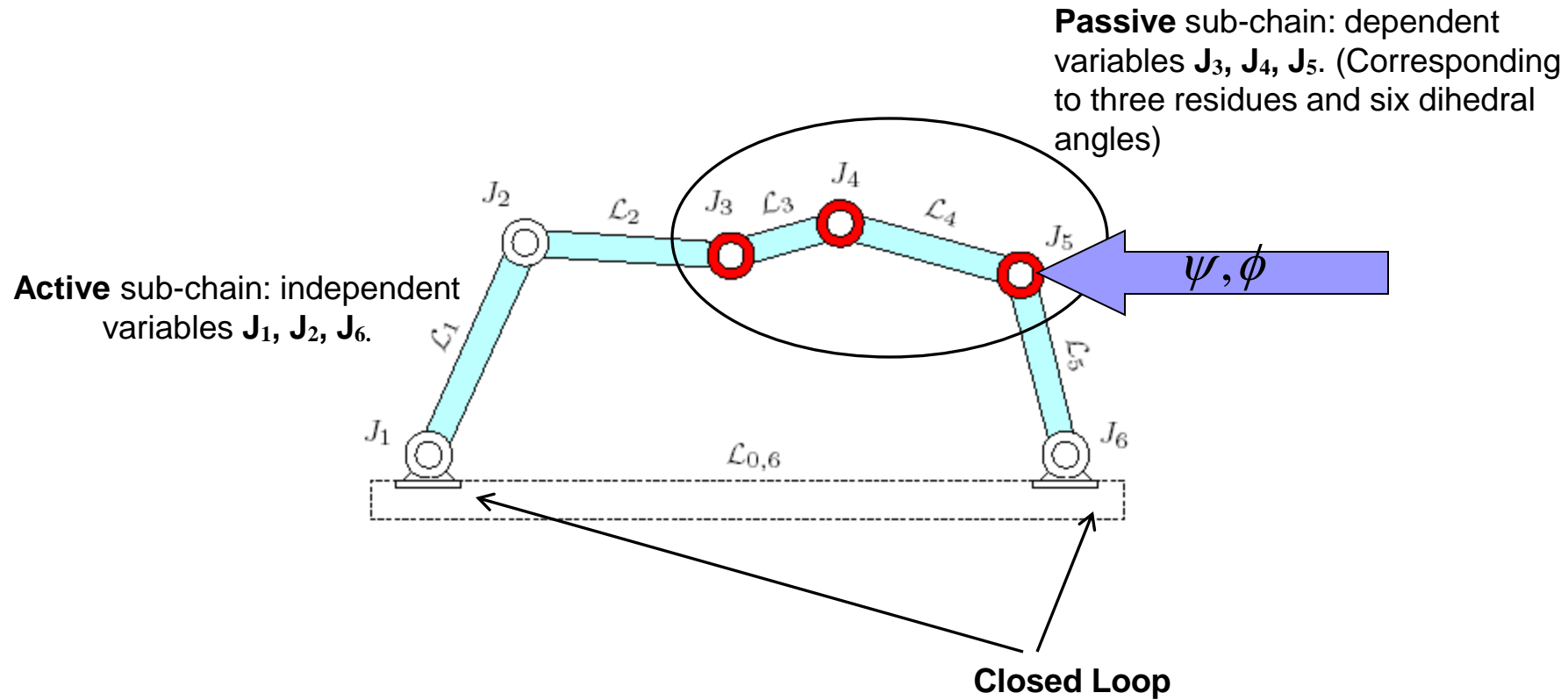
else return *Failure*;

end

A standard
inverse
kinematics
problem

Compute
configuration
parameters
of active sub-chain

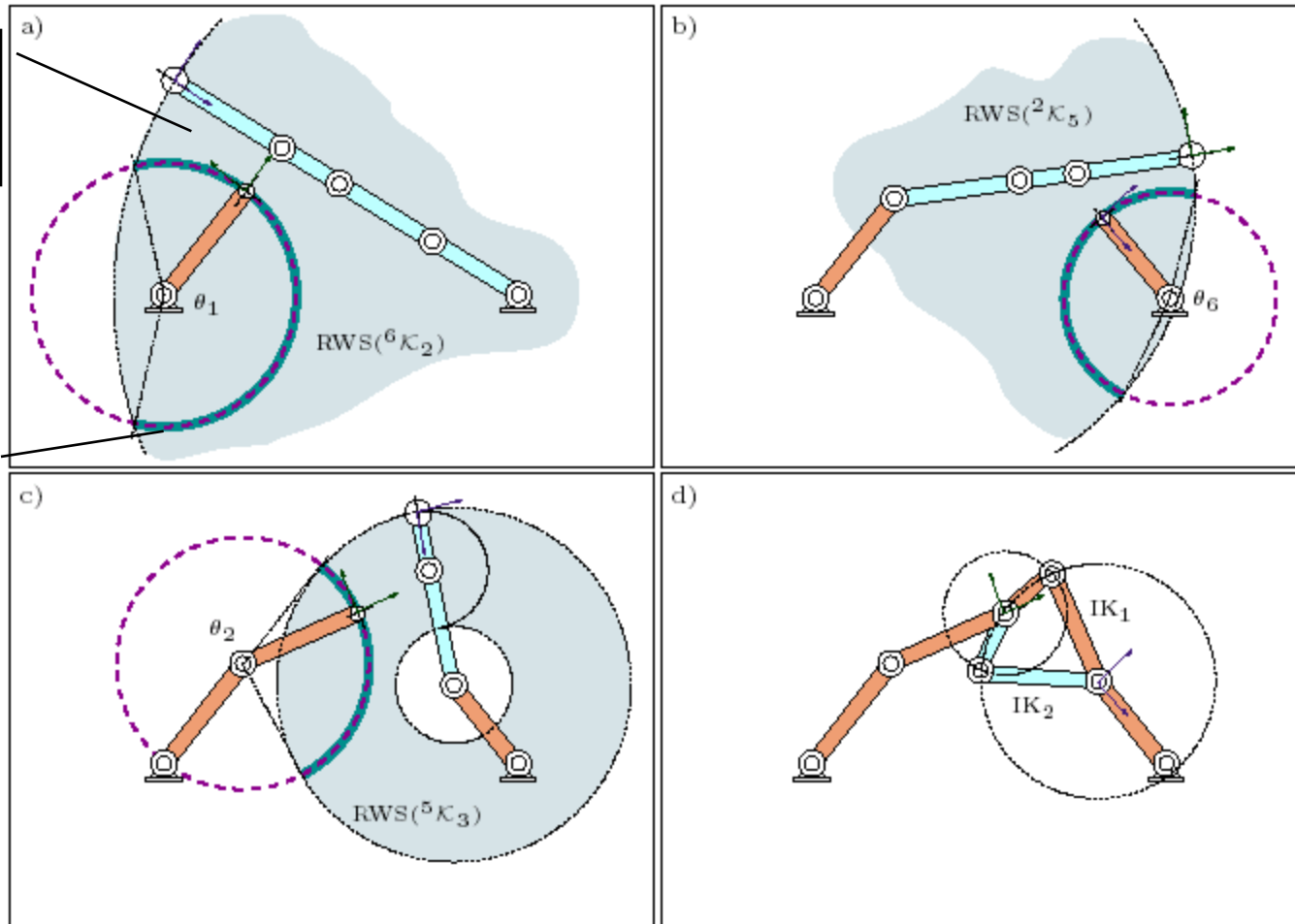
Loop decomposition-6R planar linkage



RLG Algorithm: example

Reachable
WorkSpace of
Chain6-2

Closure
Interval of θ_1



→ Solving the positional-reachable problem is simple and fast approximation to the exact closure range

RLG Algorithm: Backbone Generation

Algorithm 3: $\text{SAMPLE-}q^a$

```
input    : the backbone
output   : the active variables  $q^a$ 
begin
1   $(J_b, J_e) \leftarrow \text{INITSAMPLER}(\textit{backbone});$ 
   while not  $\text{ENDACTIVECHAIN}(\textit{backbone}, J_b)$  do
        $I_c \leftarrow \text{COMPUTECLOSURERANGE}(\textit{backbone}, J_b, J_e);$ 
       if  $I_c = \emptyset$  then goto line 1;
        $\text{SETJOINTVALUE}(J_b, \text{RANDOM}(I_c));$ 
        $J_b \leftarrow \text{NEXTJOINT}(\textit{backbone}, J_b);$ 
       if not  $\text{ENDACTIVECHAIN}(\textit{backbone}, J_e)$  then
           SWITCH( $J_b, J_e$ );
   end
end
```



Sampling side-chain conformation

- Built upon a feasible backbone conformation
- Generated by randomly sampled dihedral angles
- Tested for collisions
- Only the conformation of clashing side-chains is resampled



Conformational Space Exploration - outline

- Sampling-based Motion Planning Techniques
- RTT Algorithm

Probabilistic RoadMap(PRM) principle

- Graph (roadmap) captures the topology of collision-free configuration
- Nodes are randomly sampled
- Edges are short feasible paths linking “nearby” nodes

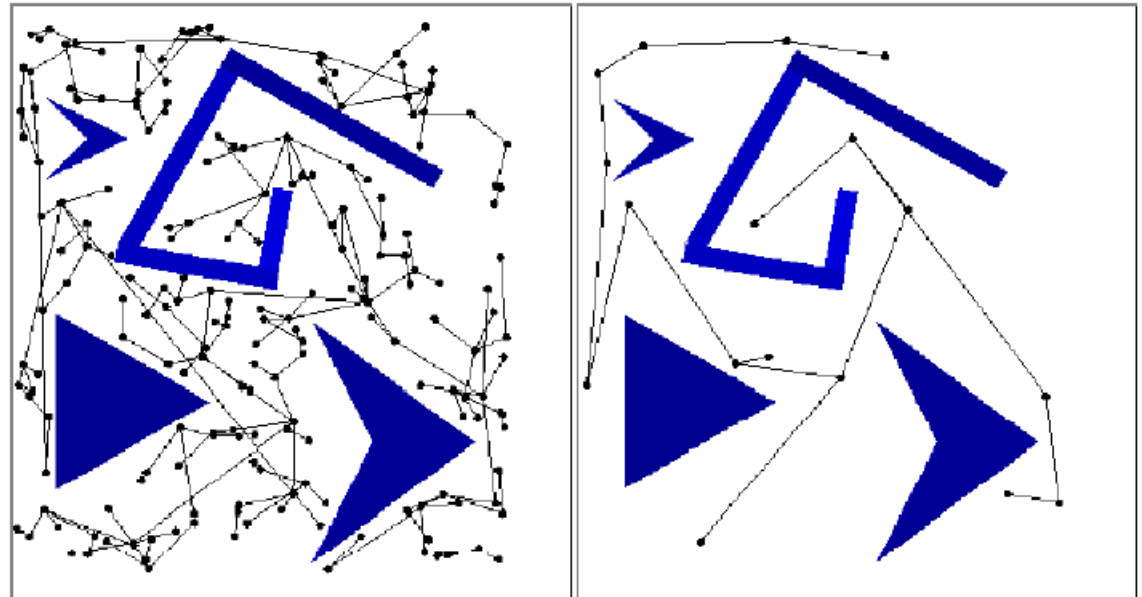
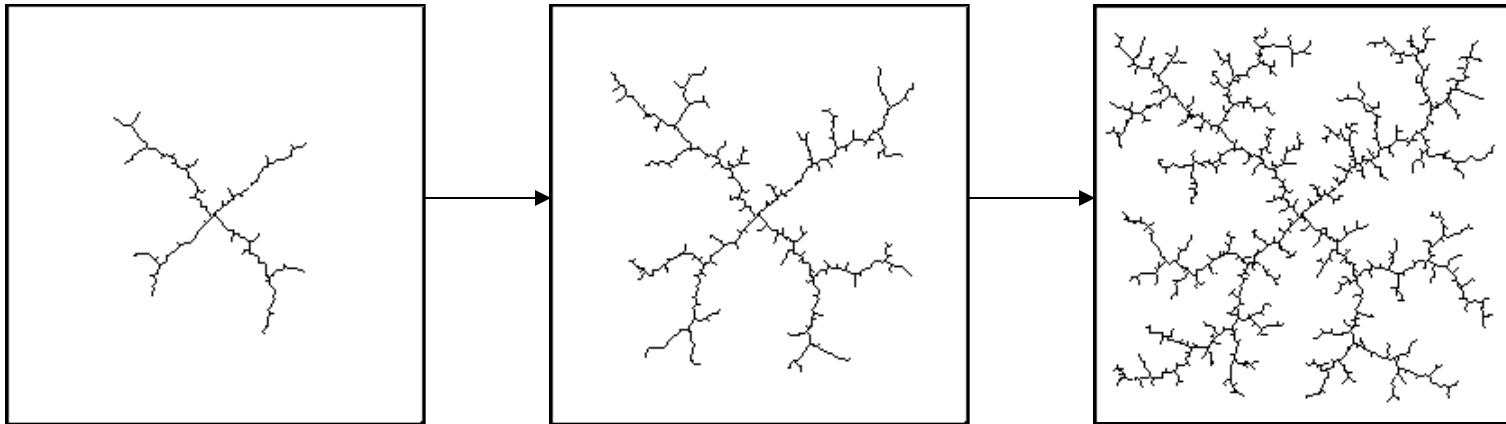


Figure 2: Basic-PRM and visibility-PRM in the same 2D environment.

Rapidly Exploring Random Tree (RRT)

developed by M.LaValle and J.Kuffner



Properties:

- Expands quickly
- Unbiased relative to random walk.
- Vertices are uniformly distributed
- Short paths
- Quickly searches high-dimensional spaces

We search for....

$$C_{lowE} \subset C_{feas} \subset C$$
$$\uparrow$$
$$C_{feas} = C_{clos} \cap C_{free}$$

Assumption: All energetically feasible configurations are considered

RTT Algorithm

Algorithm 4: ExploreByRRT

```

input   : the loop, the rest of the protein,  $q_{init}$ 
output  : the tree  $\mathcal{T}$ 
begin
   $G \leftarrow \text{INITTREE}(q_{init});$ 
   $n_{fail} \leftarrow 0;$ 
  while not STOPCONDITION( $\mathcal{T}$ ) do
     $q_{rand} \leftarrow \text{GUIDEDRANDOMCONF}(\text{loop});$ 
     $q_{near} \leftarrow \text{NEARESTNEIGHBOR}(q_{rand}, \mathcal{T});$ 
     $q_{feas} \leftarrow q_{near};$ 
     $state \leftarrow OK;$ 
    while  $state = OK$  do
       $q_{step} \leftarrow \text{MAKESTEP}(q_{feas}, q_{rand});$ 
      if FEASIBLECONF( $q_{step}$ ) then  $q_{feas} \leftarrow q_{step};$ 
      else  $state \leftarrow FAIL;$ 
    if not TOOSIMILARCONF( $q_{near}, q_{feas}$ ) then
       $q_{new} \leftarrow \text{INTERMEDIATECONF}(q_{near}, q_{feas});$ 
      GROWTREE( $q_{new}, q_{near}, \mathcal{T}$ );
       $n_{fail} \leftarrow 0;$ 
    else  $n_{fail} \leftarrow n_{fail} + 1;$ 
  end

```

Random
conf. or
from PDB

Sample
 q^a

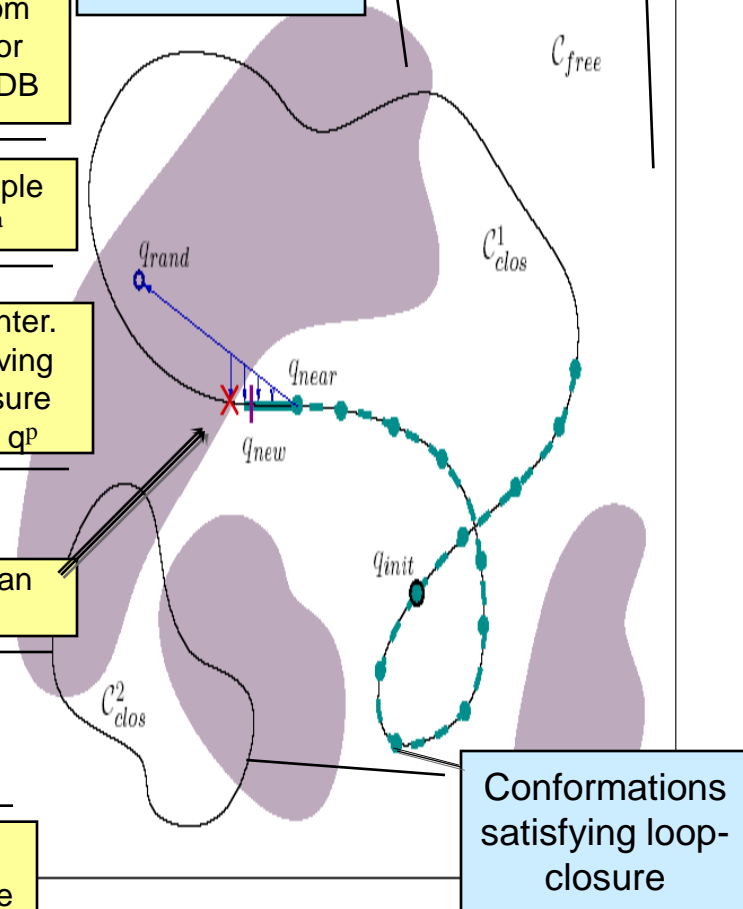
Linear Inter.
and solving
the closure
eq. for q^p

Gaussian
smpl

Believed to be an
estimate to coverage

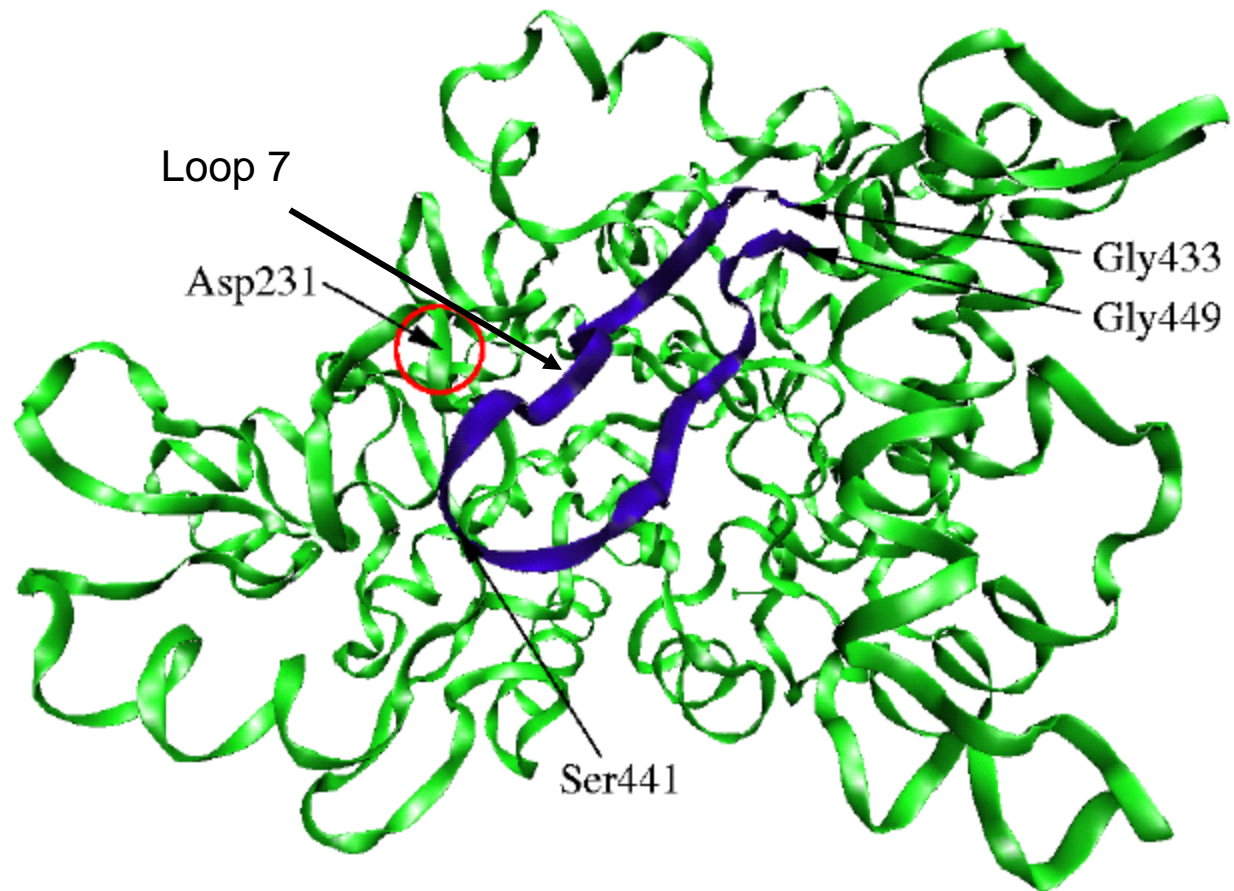
Conformations
with clashes

Clash-Free
conformation
subspace

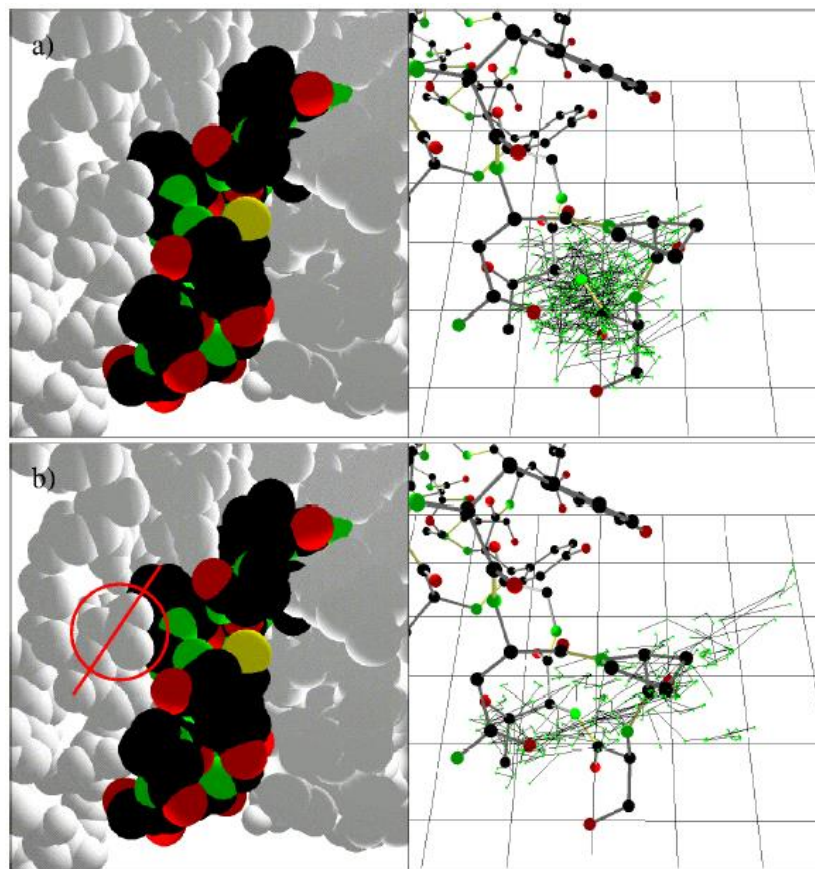


Results-structure of Amylosucrase

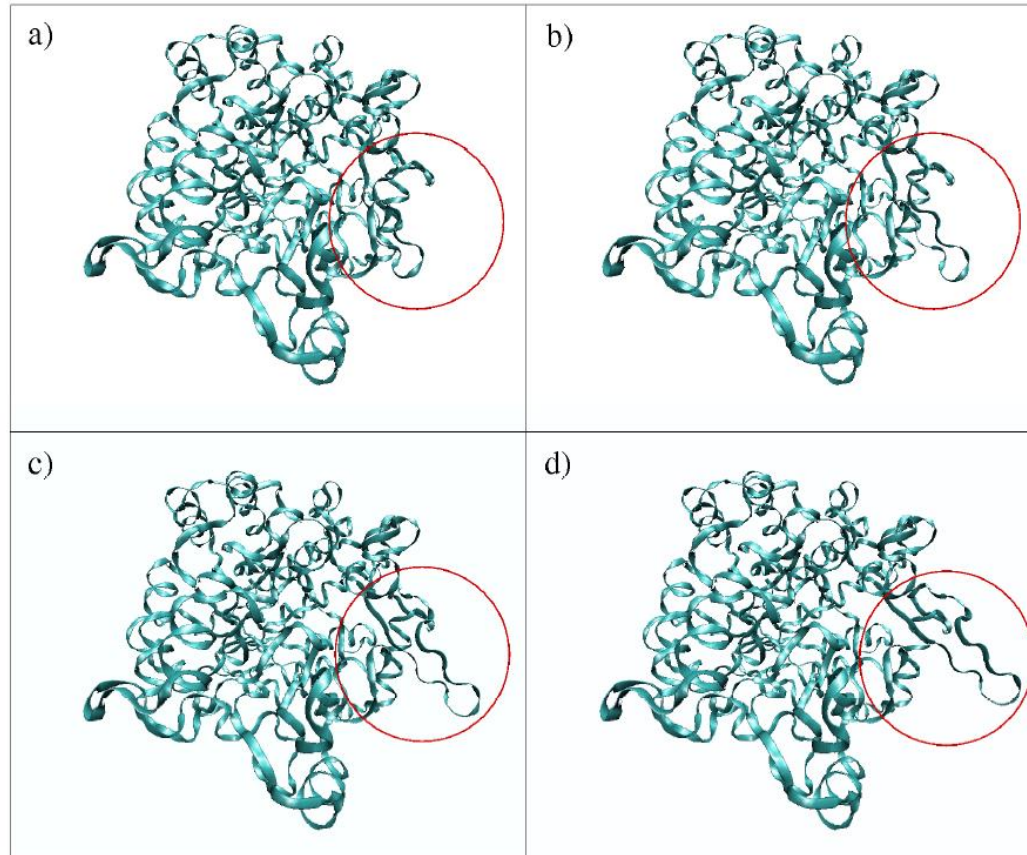
Motion of Loop 7 may have a pivotal rule in facilitating molecules interactions.



Results-exploration with/without side chain of Asp



Results-simulated conformational change in loop 7



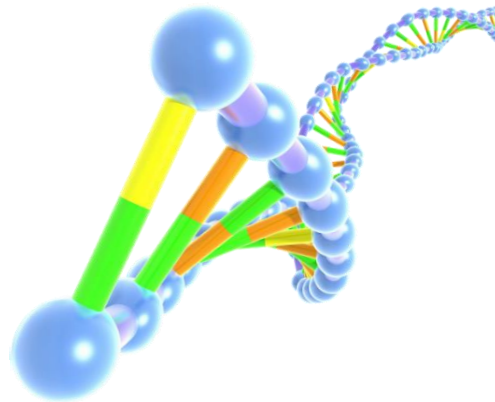
Conclusions

- Filter technique for conformational sampling and search methods was proposed
- Solution to the loop closing problem is computationally efficient
- RLG computes exact closure solutions while CCD outputs approximated solutions

Prospects

- Tailor a collision detection algorithm for the molecular application (Collision detection is by far the most computation expensive task)
- Check clashes after each stage of RLG between backbone & static env.
- Avoid steric clashes by choosing statistically preferred values for dihedral angles
- Incorporate energetic analysis (constraints) into the incremental search technique

Questions ?





Thank You

Polypeptide Extension (approximation)

l_{π} – length of polypeptide chain when all the dihedral angles at π .

\tilde{l} – upper bound on the chain's length. It is the sum of the distances between consecutive C_{α} atoms.

The extension of a chain is randomly sampled from a distribution between l_{π} and \tilde{l} .

