Amaya v2.0

System description for SMT-COMP 2024

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This is a brief overview of Amaya's submission to SMT-COMP 2024. Amaya is a standalone solver for quantified linear integer arithmetic (LIA) based on finite automata. It implements the standard automata-based decision procedure with several crucial optimizations that prune vast fractions of the underlying state space. Contrary to standard approaches, Amaya constructs a compact representation of all models of the LIA formula.

1 OVERVIEW

Amaya[1] is a standalone solver for quantified *linear integer arithmetic* (LIA)[2] based on finite automata (FAs). The foundation of the decision procedure implemented in Amaya is the work of Büchi [3] (in the context of showing automata-based decidability of the *weak monadic second-order theory of one successor*—WS1S—into which LIA can be encoded), which was later tailored for LIA (including the linear congruence operator) in the works [1, 2, 4]. The basic procedure works by constructing an FA $\mathcal{A}_{\varphi_{atom}}$ for each atomic formula φ_{atom}, which is of the form $a_1 x_1 + \ldots a_n x_n \sim c$ where $\sim \in \{=, \leq, \equiv_m\}$ (with $\equiv_m$ being the linear congruence modulo $m \in \mathbb{N} \setminus \{0\}$) for integers $a_1, \ldots, a_n, c$ and a natural number m. The FA $\mathcal{A}_{\varphi_{atom}}$ represents all solutions of the formula φ_{atom}, encoded as tuples of binary numbers in the two's complement encoding. FAs for complex formulae are then obtained inductively by performing standard automata operations: *intersection* for conjunction, *union* for disjunction, *complementation* for negation. Existential quantification is performed by the operation of *projection* on the alphabet of the corresponding FA (followed by saturation of the FAs accepting states).

Amaya implements the procedure with several crucial optimizations described in [5]. In particular, it fights the issue of constructing FAs with too many states (which emerge, e.g., when translating modulo congruence predicates) by directly creating FAs for non-atomic formulae using a derivative based construction with implicit pruning and subsumption. Moreover, the procedure also uses certain techniques of quantifier instantiation (based on formula monotonicity, finite range of models, and modulo linearization) as well as standard simplification techniques for on-the-fly rewriting of the formula for which the FA is being constructed.

[OL: preprocessing?]

2 IMPLEMENTATION

The high-level part of the decision procedure of Amaya is implemented in Python. Since the size of the underlying FA alphabet grows exponentially with the number of variables in the formula, needing to deal with alphabets having thousands up to millions of symbols, causing an explicit representation of the transition relation infeasible. To deal with this issue, Amaya encodes the transition relation using *multi-terminal binary decision diagrams* (MTBDDs) implemented within the Sylvan library [6], which often yields a much more compact representation. As Sylvan is a C library, Amaya includes a C++ backend that connects to the Python frontend via a `ctypes`

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[1]https://github.com/MichalHe/amaya

[2]also sometimes referred to as *Presburger arithmetic*

interface. The C++ backend also implements procedures for constructing FA $\mathcal{A}_{\varphi_{atom}}$ accepting the solutions of an atomic formula φ_{atom}, and the derivative-based construction.

3 AMAYA IN SMT-COMP 2024

We are submitting version v2.0 of Amaya to compete in the single query track of the division Arith, in logics LIA and NIA (nonlinear integer arithmetic; we participate in NIA because it contains many formulae that are linear in nature, but include the $\equiv_m$ predicate, which is not considered linear by SMT-LIB).

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