

# Loop Bound Analysis

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## 1 Motivating Examples

### 1.1 Example 1

However, we want to highlight the major weakness of existing approaches with the help of the following examples of C code snippet

$P_1$	$P_2$
<pre>int x , C; while(x &lt; C) {     x = x + 1; }</pre>	<pre>int x ,y,C; while(x + y &lt; C) {     x = x + 1;     y = y + 1; }</pre>

The potential-based techniques and amortized analysis base tool  $C^4B$  can easily find the bound for the program  $P_1$  where  $C^4B$  return the bound  $1.00| [0, C]|$  for the program  $P_1$ . But all the tools including  $C^4B$  failed to find the bound for the program  $P_2$ . VIAP can find the bound of both programs  $| [0, C - x]|$  and  $| [0, (C - x - y)/2]|$  for  $P_1$  and  $P_2$  respectively. To illustrate how VIAP works, consider the program  $P_2$ . Lin's [1] translation of the above program  $P$  generates the following a set of axioms  $\Pi_P^{\vec{X}}$  after some

simple simplifications where  $\vec{X} = \{x, y, C\}$

$$\begin{aligned} C_1 &= C, x_1 = x_3(n), y_1 = y_3(n), \\ x_3(0) &= x, y_3(0) = y, \\ x_3(n+1) &= x(n) + 1, y_3(n+1) = y_3(n) + 1, \\ \neg(x_3(N) + y_3(N) < C), \\ \forall n. n < N &\rightarrow (x_3(n) + y_3(n) < C) \end{aligned}$$

where  $x_1$ ,  $y_1$  and  $C_1$  denote the output values of  $a$ ,  $b$ ,  $z$ ,  $x$  and  $y$ , respectively,  $x_6(n)$  and  $y_6(n)$  the values of  $x$  and  $y$  during the  $n$ th iteration of the loop, respectively. Also  $N$  is a natural number constant, and the last two axioms say that it is exactly the number of iterations the loop executes before exiting. Our recurrence solving tool `recSolve(RS)` can find the closed-form solutions of  $x_3(n)$  and  $y_3(n)$ , which yields the closed-form solution  $x_3(n) = n - x$  and  $y_3(n) = n - y$  respectively. Those can be used to simplify the recurrence for  $y_3(n)$  into following set of axioms after eliminates recurrence relation for  $x_3()$ , and  $y_3()$

$$\begin{aligned} x_1 &= N + x, y_1 = N + y, C_1 = C, \\ (2 * N + x + y &\geq C) \\ \forall n. n < N &\rightarrow (2 * n + x + y < C) \end{aligned}$$

The system tried to derive using our algorithm derived  $N = (C - x - y)/2$ , then it tried to prove using SMT solver. If it is able to successfully prove that, then it successfully derived bound  $[0, (C - x - y)/2]$ .

## References

- [1] F. Lin, “A formalization of programs in first-order logic with a discrete linear order,” *Artificial Intelligence*, vol. 235, pp. 1 – 25, 2016. [Online]. Available: <http://www.sciencedirect.com/science/article/pii/S000437021630011X>