

NAME:	Shubham Vishwakarma
UID No.	2021700071
BRANCH:	S.Y CSE-DS
BATCH:	D
SUBJECT	Design and Analysis of Algorithms
EXPERIMENT No.	1B
Date of Performance	30/01/2023
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AIM:	Experiment on finding the running time of an algorithm.
Program 1	
PROBLEM STATEMENT :	<p>For this experiment, you need to implement two sorting algorithms namely Insertion and Selection sort methods. Compare these algorithms based on time and space complexity. Time required to sorting algorithms can be performed using <code>high_resolution_clock::now()</code> under namespace <code>std::chrono</code>.</p> <p>You have to generate 1,00,000 integer numbers using C/C++ <code>Rand</code> function and save them in a text file. Both the sorting algorithms use these 1,00,000 integer numbers as input as follows. Each sorting algorithm sorts a block of 100 integers numbers with array indexes numbers <code>A[0..99]</code>, <code>A[0..199]</code>, <code>A[0..299]</code>, ..., <code>A[0..99999]</code>. You need to use <code>high_resolution_clock::now()</code> function to find the time required for 100, 200, 300, ..., 100000 integer numbers. Finally, compare two algorithms namely Insertion and Selection by plotting the time required to sort 100000 integers using LibreOffice Calc/MS Excel. The x-axis of 2-D plot represents the block no. of 1000 blocks. The y-axis of 2-D plot represents the running time to sort 1000 blocks of 100, 200, 300, ..., 100000 integer numbers.</p> <p>Note – You have to use C/C++ file processing functions for reading and writing randomly generated 100000 integer numbers.</p>
ALGORITHM/ THEORY:	<p>Sorting Algorithms</p> <p>Sorting Algorithms are a class of algorithms which are used to arrange the elements of a list or an array in a particular order.</p> <p>There are a lot of sorting algorithms including bubble sort, selection sort, quick sort, insertion sort, heap sort, etc. This experiment focuses on insertion and selection sort. For</p>

our examples, we will consider sorting in ascending order.

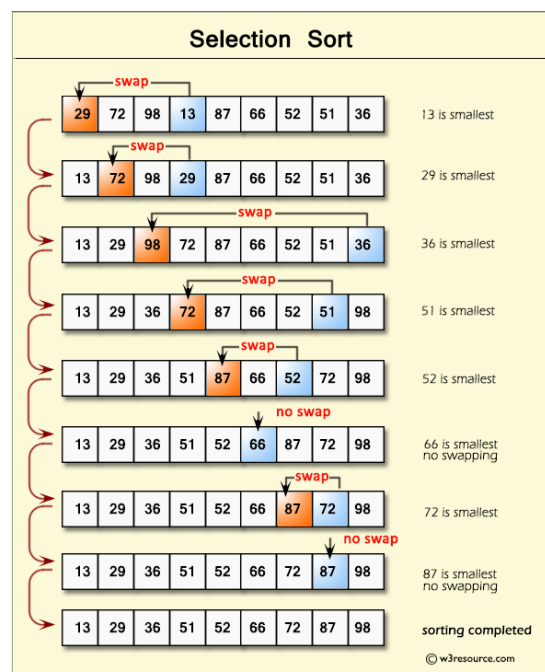
Selection Sort

The selection sort algorithm sorts an array by repeatedly finding the minimum element (considering ascending order) from unsorted part and putting it at the beginning. The algorithm maintains two subarrays in a given array.

- 1) The subarray which is already sorted.
- 2) Remaining subarray which is unsorted.

In every iteration of selection sort, the minimum element (considering ascending order) from the unsorted subarray is picked and moved to the sorted subarray.

Following example explains the above steps:



The time complexities for selection sort are given below:

Best	Average	Worst
$\Omega(n^2)$	$\theta(n^2)$	$O(n^2)$

Since it always has an n^2 time complexity, it is inefficient for large lists. We will discuss its algorithm later.

Insertion Sort

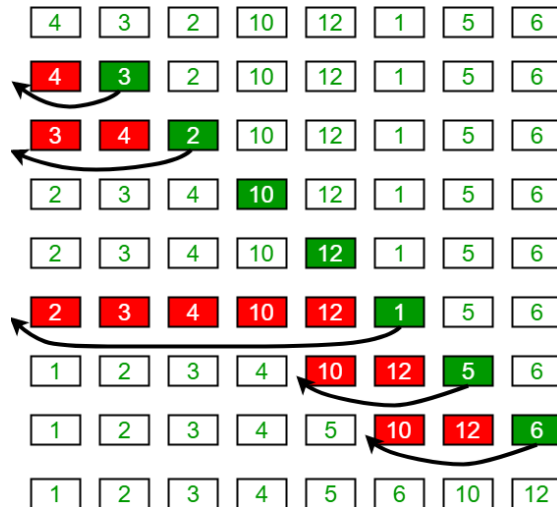
Insertion sort is a simple sorting algorithm that works similar to the way you sort playing cards in your hands. The array is virtually split into a sorted and an unsorted part. Values from the unsorted part are picked and placed at the correct position in the

sorted part.

To sort an array of size n in ascending order:

- a. Iterate from $\text{arr}[1]$ to $\text{arr}[n]$ over the array.
- b. Compare the current element (key) to its predecessor.
- c. If the key element is smaller than its predecessor, compare it to the elements before. Move the greater elements one position up to make space for the swapped element.

Insertion Sort Execution Example



The time complexities for insertion sort are given below:

Best	Average	Worst
$\Omega(n)$	$\theta(n^2)$	$O(n^2)$

Since it always has a best-case time complexity of n , it is more efficient compared to selection sort.

Algorithms:

Selection Sort:

- 1) START
- 2) For $i = 0$ to $n-2$:
 - a. $\text{minIndex} = i$
 - b. For $j = i+1$ to $n-1$
 - i. If $\text{arr}[\text{minIndex}] > \text{arr}[j]$
 1. $\text{minIndex} = j$
 - c. Swap elements of array at positions i and minIndex
- 3) END

Insertion Sort:

- 1) START
- 2) For $i = 1$ to $n-1$
 - a. $\text{Key} = \text{arr}[i]$
 - b. $j = i - 1$
 - c. While $j \geq 0$ and $\text{arr}[j] > \text{Key}$:
 - i. $\text{arr}[j+1] = \text{arr}[j]$
 - ii. $j = j - 1$
 - d. $\text{arr}[j+1] = \text{Key}$
- 3) STOP

PROGRAM:

```
#include <stdio.h>
#include <time.h>
#include <stdlib.h>

void print(int * A, int n)
{
    for(int i = 0; i<n; i++)
    {
        printf("%d ", A[i]);
    }
}

double numberGenertor(int n)
{
    FILE *fp = fopen("./num.txt", "w");
    clock_t start, end;
    double cpu_time_used;
    int p;
    start = clock();
    for(int i = 0; i < n; i++)
    {
        p = (rand() + (rand()*133));
        fprintf(fp, " %d\n", p);
    }
    end = clock();
    cpu_time_used = ((double) (end - start)) / CLOCKS_PER_SEC;
    fclose(fp);
    return cpu_time_used;
}
```

```

void InsertionSort(int arr[], int n)
{
    int i, key, j;
    for (i = 1; i < n; i++)
    {
        key = arr[i];
        j = i - 1;

        while (j >= 0 && arr[j] > key)
        {
            arr[j + 1] = arr[j];
            j = j - 1;
        }
        arr[j + 1] = key;
    }
}

void SelectionSort(int * A, int n)
{
    for(int i = 0; i<n-1; i++)
    {
        int min = A[i];
        int pos = i;
        int temp;
        for(int j = i; j<n; j++)
        {
            if(A[j]<min)
            {
                min = A[j];
                pos = j;
            }
        }
        temp = A[i];
        A[i] = min;
        A[pos] = temp;
    }
}

int main()
{
    double createTime = numberGenertor(100000);
    printf("\n%lf\n", createTime);
    FILE * read, * fpsel, * fpins;
    fpsel = fopen("./selection.csv", "w");
    fpins = fopen("./insertion.csv", "w");
    if(!fpsel)

```

```

        return 0;
    if(!fpins)
        return 0;
    fprintf(fpse1, "Blocks, time\n");
    fprintf(fpins, "Blocks, time\n");

    for(long int x = 100; x<=100000; x+=100)
    {
        read = fopen("num.txt", "r");
        int A[x], B[x];
        clock_t start1, end1, start2, end2;
        // if(x % 1000 == 0)
        // {
        //     printf("x = %ld.\n", x);
        // }

        for(long int i = 0; i<x; i++)
        {
            fscanf(read, "%d\n", &A[i]);
            B[i] = A[i];
        }

        start1 = clock();
        SelectionSort(A, x);
        end1 = clock();

        start2 = clock();
        InsertionSort(B, x);
        end2 = clock();

        double t1 = (double) (end1 - start1) / CLOCKS_PER_SEC;
        double t2 = (double) (end2 - start2) / CLOCKS_PER_SEC;
        printf("%6d | %1f | %1f\n", x, t1, t2);

        fprintf(fpse1, "%ld, %f\n",x,t1);
        fprintf(fpins, "%ld, %f\n",x,t2);

        fclose(read);
    }
    fclose(fpse1);
    fclose(fpins);
}

```

RESULT:

● PS C:\Users\smsa\Desktop\SEM 4\DAA\Practicals\Exp1\Exp1b\new1b> & .\"exp1b.exe"

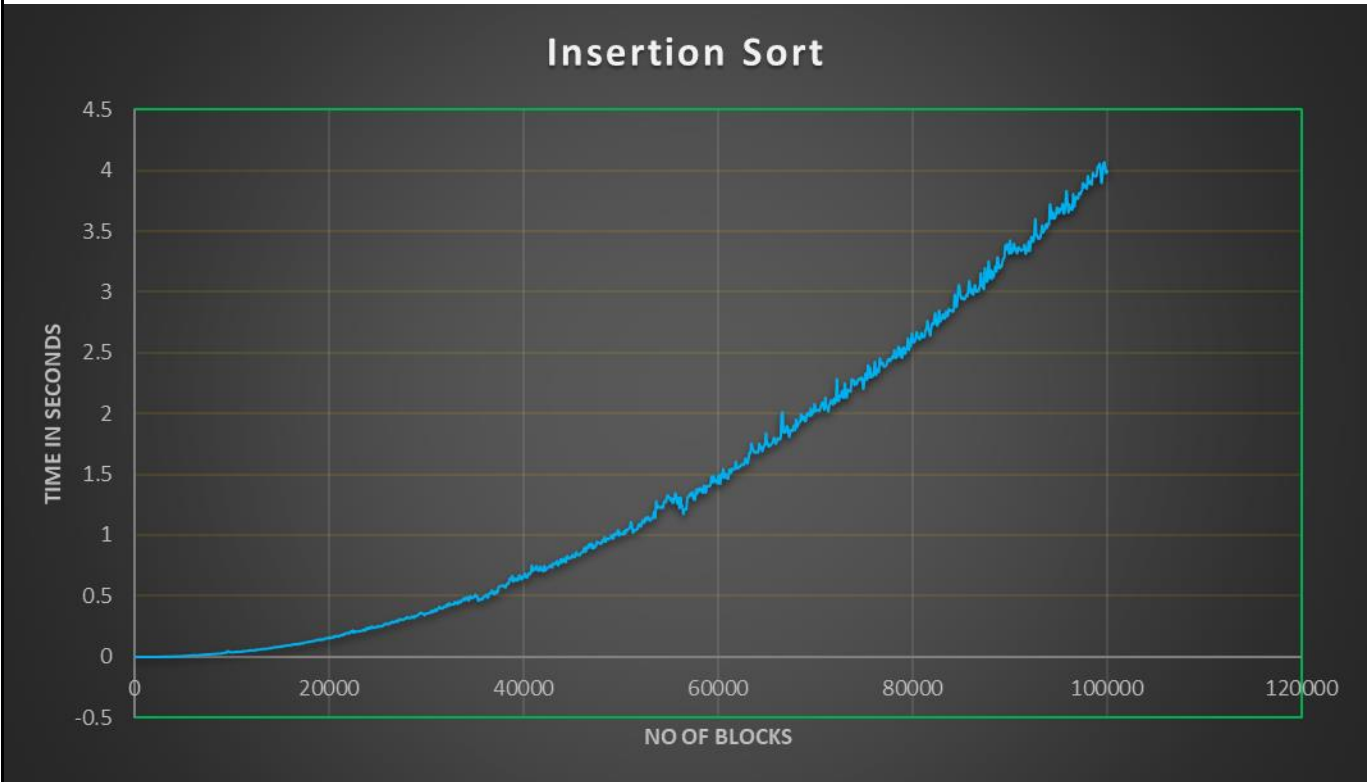
0.020000

100	0.000000	0.000000
200	0.000000	0.000000
300	0.000000	0.000000
400	0.000000	0.000000
500	0.000000	0.001000
600	0.000000	0.000000
700	0.000000	0.000000
800	0.000000	0.001000
900	0.000000	0.001000
1000	0.001000	0.000000
1100	0.000000	0.001000
1200	0.001000	0.001000
1300	0.001000	0.001000
1400	0.001000	0.001000
1500	0.001000	0.001000
1600	0.002000	0.001000
1700	0.002000	0.001000
1800	0.002000	0.001000

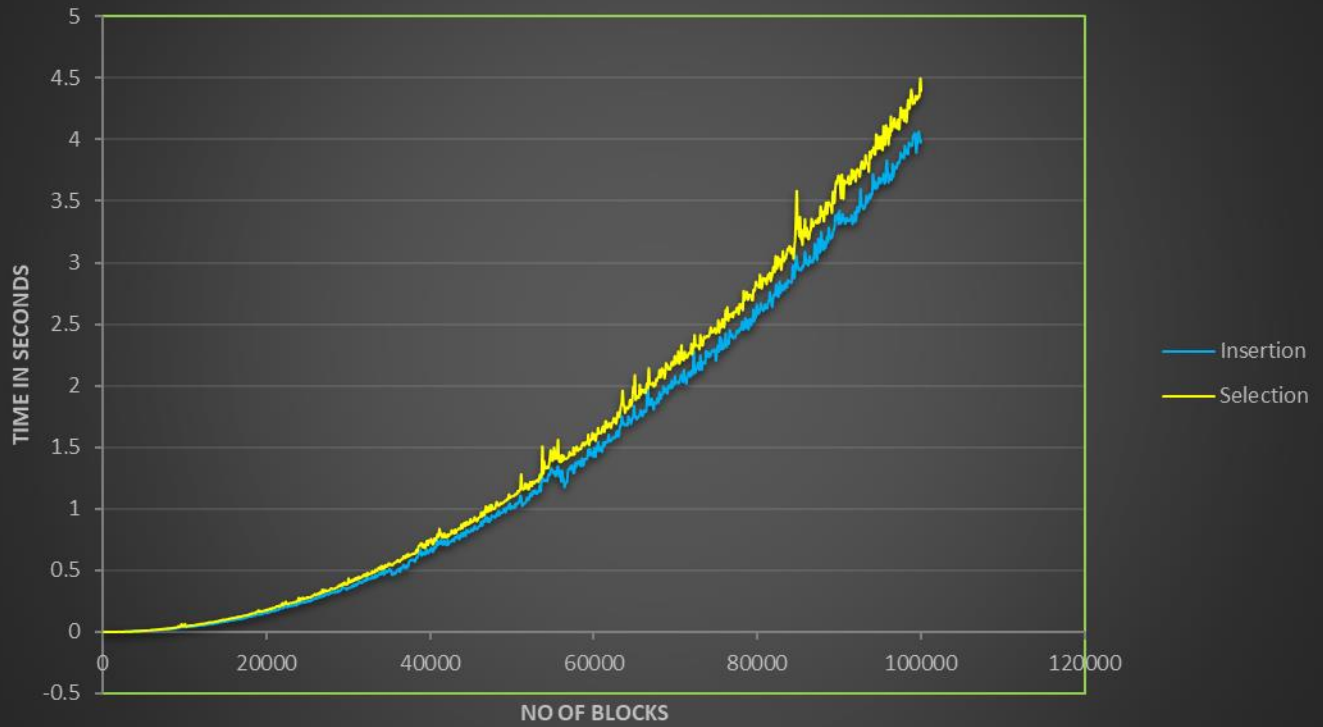
95000	3.930000	3.644000
95100	3.960000	3.663000
95200	4.031000	3.686000
95300	3.916000	3.675000
95400	4.112000	3.721000
95500	3.976000	3.643000
95600	3.962000	3.656000
98400	4.326000	3.882000
98500	4.252000	3.976000
98600	4.258000	3.955000
98700	4.352000	3.952000
98800	4.409000	3.947000
98900	4.352000	3.953000
99000	4.293000	4.030000
99100	4.313000	4.020000
99200	4.298000	4.053000
99300	4.358000	4.014000
99400	4.345000	3.893000
99500	4.327000	3.957000
99600	4.358000	4.050000
99700	4.344000	4.064000
99800	4.365000	4.020000
99900	4.501000	3.975000
100000	4.398000	3.991000

PS C:\Users\smsa\Desktop\SEM 4\DAA\Practicals\Exp1\Exp1b\new1b> █

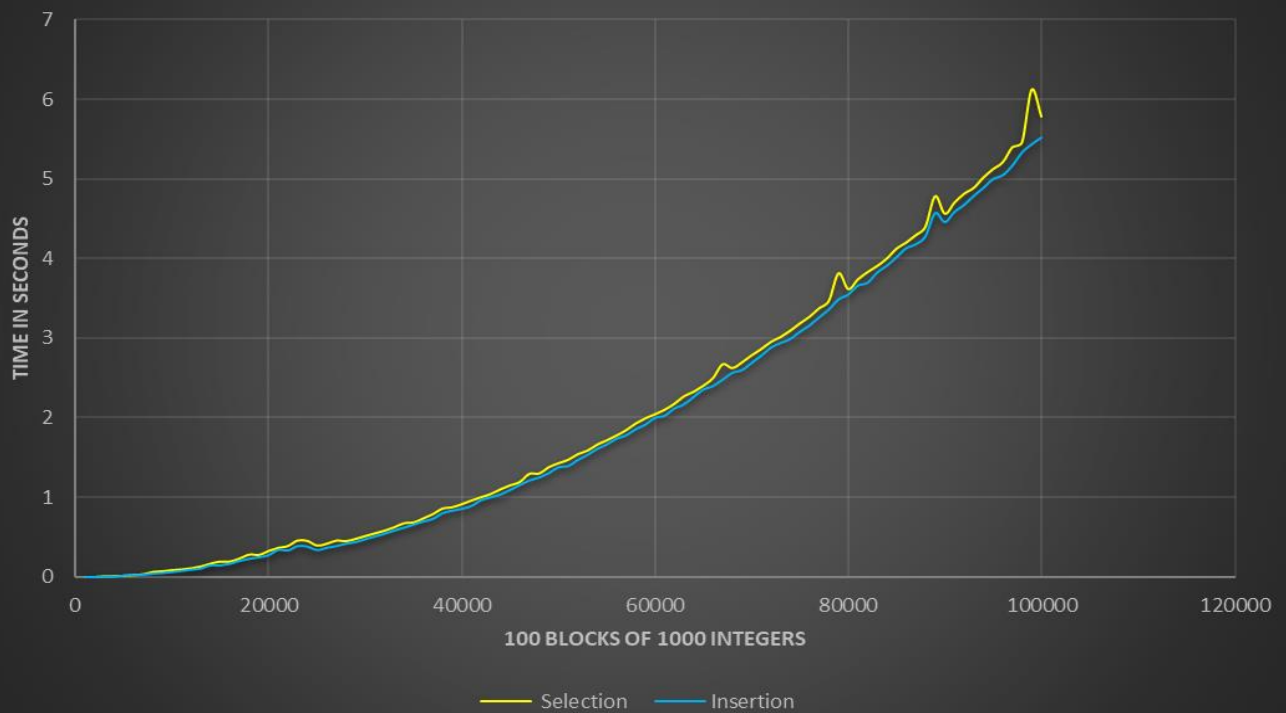
GRAPH:



Selection Sort VS Insertion Sort



Selection Vs Insertion



From the graph we can see that insertion sort have better running time than selection sort
The spikes in the graph due to following reasons

1. Due to a greater number of swaps at that given range
2. Due to CPU over heating (Machine Dependent)
3. Due to nature of random number (large difference between consecutive numbers)
4. Due to other process running in the background during the execution of program

	Best	Average	Worst
Selection Sort	$\Omega(n^2)$	$\theta(n^2)$	$O(n^2)$
Insertion Sort	$\Omega(n)$	$\theta(n^2)$	$O(n^2)$

The time complexity for all 3 cases of selection sort is n^2 . Meanwhile in insertion sort, the average and worst-case time complexity is the same – n^2 but the best-case time complexity is linear (in the case where the input is already sorted).

Hence, we can say that in general, insertion sort is more efficient compared to selection sort.

CONCLUSION :

In this experiment, we wrote a program which given 100000 integers in 1000 blocks of 100 numbers, uses the selection and insertion sort algorithms to sort them in ascending order.

In selection sort, during every iteration, we move the minimum element of the unsorted subarray to the end of the sorted subarray.

In insertion sort, during every iteration, we move the first element of the unsorted subarray to the appropriate position in the sorted subarray.

We compared the running time of insertion and selection sort by comparing the running time graph of 2 algorithms. We also analyzed the 2 algorithms in which we found that while the time complexity of selection sort is n^2 , the best-case time complexity of insertion sort is linear (in the case where the input is already sorted). Hence, we concluded that in general, insertion sort is

	<p>more efficient compared to selection sort.</p>
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	<p>Thus, we successfully accomplished the aim of this experiment.</p>
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