

Figure 3.39. A model \mathcal{M} .

2. Consider the system of Figure 3.39. For each of the formulas ϕ :

- (a) $G a$
- (b) $a U b$
- (c) $a U X (a \wedge \neg b)$
- (d) $X \neg b \wedge G (\neg a \vee \neg b)$
- (e) $X (a \wedge b) \wedge F (\neg a \wedge \neg b)$

LTL

- (i) Find a path from the initial state q_3 which satisfies ϕ .
- (ii) Determine whether $\mathcal{M}, q_3 \models \phi$.

2. (a) (i) $q_3 \rightarrow q_4 \rightarrow q_3 \rightarrow q_4 \rightarrow \dots$

(ii) $\mathcal{M}, q_3 \models G a$ does not hold
进入 q_1 或 q_2 则不满足

(b) (i) $q_3 \rightarrow q_2 \rightarrow q_2 \rightarrow q_2 \rightarrow \dots$

(ii) $\mathcal{M}, q_3 \models a U b$ does not hold

反例: $q_3 \rightarrow q_1 \rightarrow q_2 \rightarrow \dots$

(c) (i) $q_3 \rightarrow q_4 \rightarrow q_3 \rightarrow \dots$

(ii) $\mathcal{M}, q_3 \models a U X (a \wedge \neg b)$ does not hold

需至少隔一步后再次回到 q_3

若 $q_3 \rightarrow q_2 \rightarrow q_2 \rightarrow \dots$ 则不成立

(d) (i) $q_3 \rightarrow q_1 \rightarrow q_2 \rightarrow q_2 \rightarrow \dots$

(ii) $\mathcal{M}, q_3 \models X \neg b \wedge G(\neg a \vee \neg b)$ does not hold

反例: $q_3 \rightarrow q_4 \rightarrow \dots$

(e) (i) $q_3 \rightarrow q_4 \rightarrow q_3 \rightarrow q_1 \rightarrow \dots$

(ii) $\mathcal{M}, q_3 \models X(a \wedge b) \wedge F(\neg a \wedge \neg b)$ does not hold

反例: $q_3 \rightarrow q_1 \rightarrow \dots$

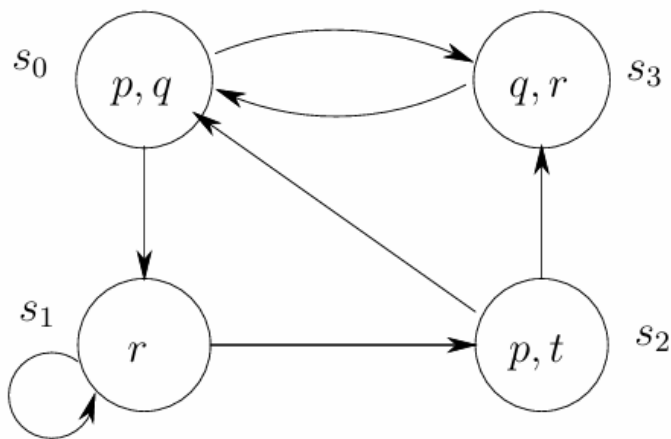


Figure 3.41. Another model with four states.

8. Consider the model \mathcal{M} in Figure 3.41. Check whether $\mathcal{M}, s_0 \models \phi$ and $\mathcal{M}, s_2 \models \phi$ hold for the CTL formulas ϕ :

CTL

- (a) $AF q$
- (b) $AG(EF(p \vee r))$
- (c) $EX(EX r)$
- (d) $AG(AF q)$.

8. (a) $\mathcal{M}, s_0 \models AF q$ holds

s_0 本身包含 q , all paths starting from s_0 hit a q .

$\mathcal{M}, s_2 \models AF q$ holds

s_2 进入 s_0 或 s_3 , 二者都含 q

(b) $\mathcal{M}, s_0 \models AG(EF(p \vee r))$ holds

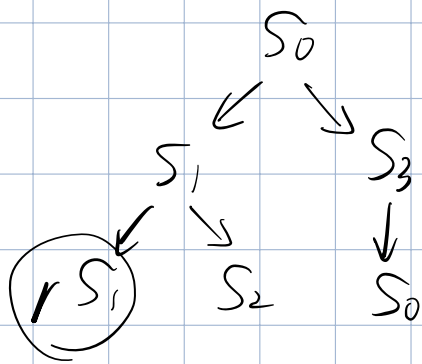
$p \vee r$ 对每个状态都成立

则上式必定成立

$\mathcal{M}, s_2 \models AG(EF(p \vee r))$ holds

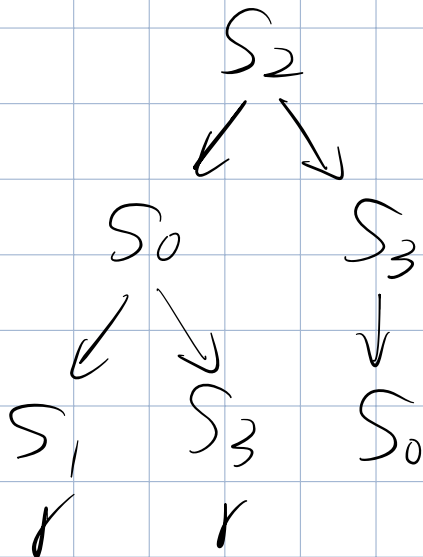
同上理

(c) $\mathcal{M}, s_0 \models EX(EX r)$ holds



存在 $s_0 \rightarrow s_1 \rightarrow s_1 \rightarrow \dots$ 这条 path 满足

$\mathcal{M}, s_2 \models EX(EX r)$ holds



同上理

(d) $\mathcal{M}, s_0 \models AG(AF g)$ does not hold

$\mathcal{M}, s_2 \models AG(AF g)$ does not hold

即每条 path 上的每个状态 s_i 都要满足 $M, s_i \models AF \varphi$

反例: $\dots \rightarrow s_i \rightarrow s_i \rightarrow s_i \rightarrow \dots$ (一直留在 s_i)

此条路径上的这几个 s_i 状态及后续的所有 s_i 状态都不满足条件.