

CAT-II  
LAO

110111  
001001  
Date

Q2)

Table for modified Booth's algorithm:

$i+1$	$i$	$i-1$	operation
0	0	0	$0 \times M$
0	0	1	$+1 \times M$
0	1	0	$+1 \times M$
0	1	1	$+2 \times M$
1	0	0	$-2 \times M$
1	0	1	$-1 \times M$
1	1	0	$-1 \times M$
1	1	1	$0 \times M$

110011  
110100

CAT-II

CAO

Q2)

Booth (Modified)

$M = 12 \Rightarrow 001100 \Rightarrow -M = 110100$

$Q = 7 \Rightarrow 000111 \rightarrow +2M = 011000$

comment	A	Q	Q-1	SC
initial	000000 110100	000111	0	3
Subtract M	110100	000111	0	
Ashr	111010	000011	1	2
Ashr	111101 011000	000001	1	
Add +2M	010101	000001	1	
Ashr	001010	100000	1	1
Ashr	000101	010000	0	
Add +0M				
Ashr	000010	101000	0	0
Ashr	000001	010100	0	

Answer  $\Rightarrow 1010100_2$

decimal  $\Rightarrow 64 \times 1 + 1 \times 16 + 1 \times 4$

$= 84$  Ans.



Q4) a)

$$\text{Main memory} = 8 \text{ MB} = 8 \times 2^{20} = 2^{23}$$

$$\text{cache memory} = 16 \text{ KB} = 16 \times 2^{10} = 2^{14}$$

$$\text{Byte per blocks} = 8 \text{ bytes} = 2^3$$

$$\therefore \text{no. of blocks} = \frac{\text{cache memory}}{\text{Bytes per block}}$$

$$= \frac{2^{14}}{2^3} = 2^9 = \boxed{512 \text{ blocks}}$$

$$\Rightarrow 2^9 = 512$$

$$\Rightarrow \boxed{x = 9}$$

$$\rightarrow \text{we know } 2^x + w = 2^{23} \quad (\text{pre-calculated})$$

$$\Rightarrow \boxed{x + w = 23}$$

$$\text{Bytes per block} = 2^3 \Rightarrow \boxed{w = 3}$$

$$\Rightarrow x = 20$$

$$x = 9$$

$$x - 9 = 11$$

$$w = 3$$

$\therefore$  memory address will look like!

$$s+w=23$$

$$s=9, w=3$$

tag	line	word
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for direct mapping.

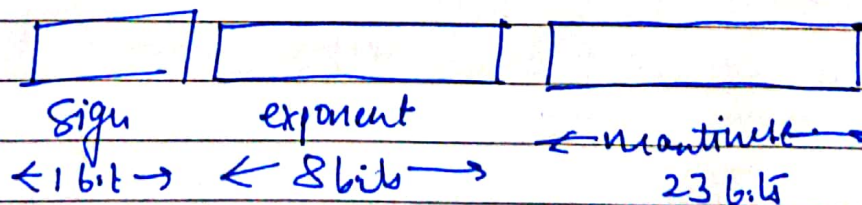
Q2) Table

4	2	0	1	2	6	1	4	6	2
		0	1	1	1	1	1	X	2
2	2	2	2	2	2	2	2	2	2
4	4	4	4	4	4	4	4	4	4
hit					hit				

page fault = 6

- Q1) b) (i) Direct  
(ii) Associative  
(iii) Key-associative.  
(iv) Set-associative.

Q3) 1.85 in IEEE 754



$$1.85 = (-1)^{\text{Sign}} \left( 2^{\text{Component} - 127} \right) (1. \text{mantissa})$$



as 18.5 is +ve then sign  $\rightarrow 0$

$$16 = 2^4 = 2^{\text{exponent}} = 128$$

then exponent =  $127 + 4 = 131$

$$(10000011)_2$$

$$\text{So } 18.5 = 16 \times (1.4)$$

$$18.5 = 16 \times (1 + 0.125)$$

$$= 1.1563 = 1 + 0.14$$

$$= 0.M = 0.15625 = (001010 \dots)_2$$

then :-

[illegible]

Q1) a) data word: 010100111011

$$C_1 = 1, C_2 = 1, C_4 = 1, C_8 = 0$$

$$\text{Data bits} = 01010110$$

calculating check bits :-

$$C_1 = D_1 \oplus D_2 \oplus D_4 \oplus D_5 \oplus D_7$$

$$C_2 = D_1 \oplus D_3 \oplus D_4 \oplus D_6 \oplus D_7$$

$$C_4 = D_2 \oplus D_3 \oplus D_4 \oplus D_8$$

$$C_8 = D_5 \oplus D_6 \oplus D_7 \oplus D_8$$

$$C_1 = 0 \oplus 1 \oplus 0 \oplus 1 \oplus 1 = 1$$

$$C_2 = 0 \oplus 1 \oplus 0 \oplus 0 \oplus 1 = 0$$

$$C_4 = 1 \oplus 1 \oplus 0 \oplus 0 = 0$$

$$C_8 = 1 \oplus 0 \oplus 1 \oplus 0 = 0$$

taking (XOR) of old & new check bits :-

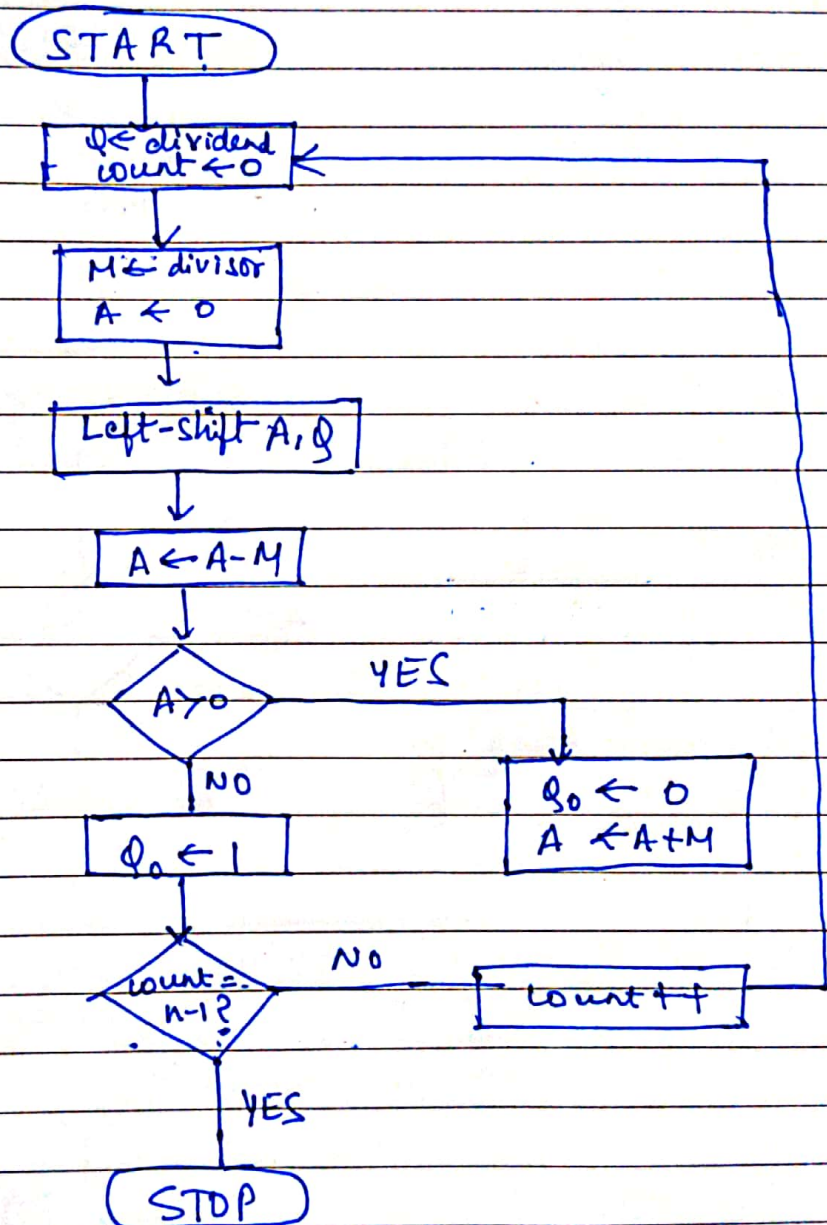
$$1000 \oplus 1110 = (0110)_2 = (6)_{10}$$

$\begin{smallmatrix} 0 \\ 0 \end{smallmatrix}$  we have error in 6<sup>th</sup> position which has  $D_3$



Q3) a)

flowchart for restoring division algorithm:



n m 11100 A 10 0

7 0001000 = 0000000 = 100110

0000001 001110?

01111001 001110.3

0000001 0011100

SLAQ

A = A - M

Q[0] = 0

ResA

6 0000010 011100? = 1

1111010 011100?

0000010 0111000 = 1

SLAQ

A = A - M

Q[0] = 0

ResA

5 0000100 111000?

111100 111000?

1 = 0000100 1110000

0 = 1 0 0 0 0 0 1 0 0 = 1

0 = 0 0 0 0 1 0 1 1 = 1

SLAQ

A = A - M

Q[0] = 0 ResA

4 0001001 0110000?

0000001 110000?

000000 1100001

SLAQ

A = A - M

Q[0] = 1

0110 = 0110 = 01100001

Answer = 1001.11