

INT404 ARTIFICIAL INTELLIGENCE

Constraint Satisfaction

Lecture 10

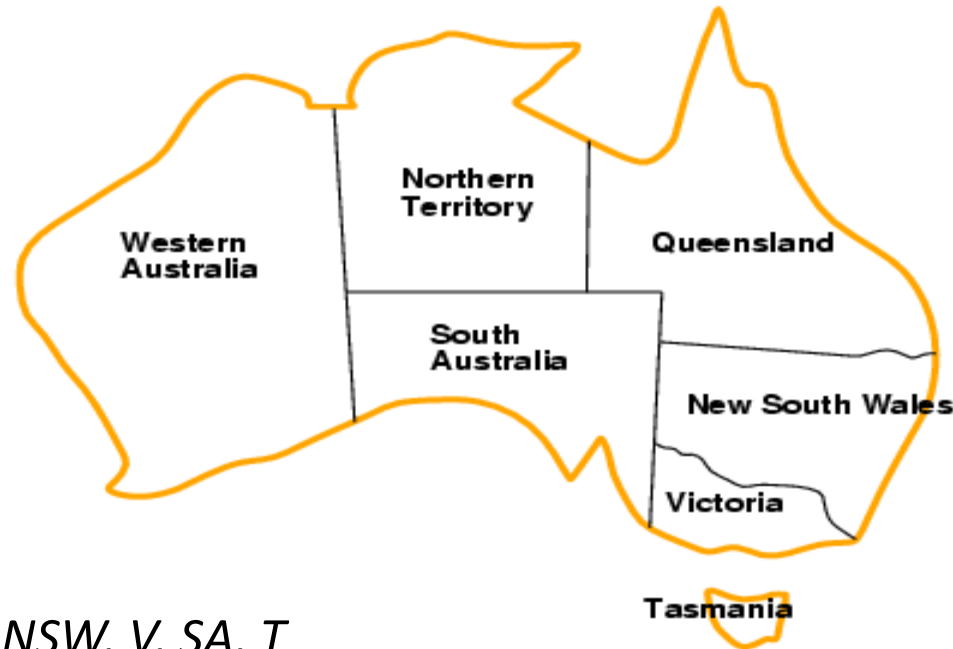
Formal Definition of CSP

- A constraint satisfaction problem (CSP) is a triple (V, D, C) where
 - V is a set of variables X_1, \dots, X_n .
 - D is the union of a set of domain sets D_1, \dots, D_n , where D_i is the domain of possible values for variable X_i .
 - C is a set of constraints on the values of the variables, which can be pairwise (simplest and most common) or k at a time.

CSPs vs. Standard Search Problems

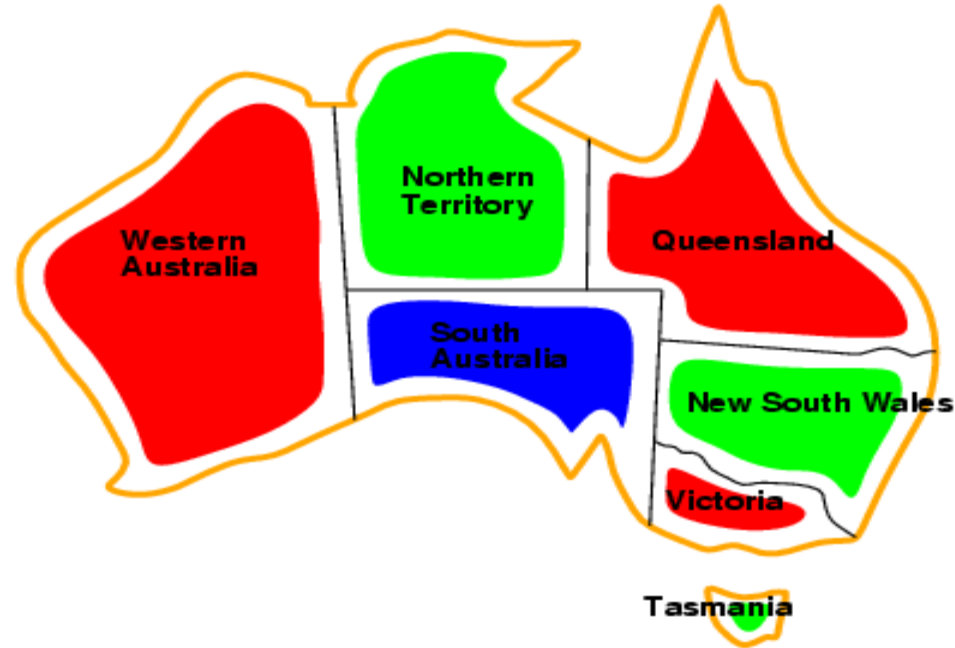
- Standard search problem:
 - state is a "black box" – any data structure that supports successor function, heuristic function, and goal test
- CSP:
 - state is defined by variables X_i with values from domain D_i
 - goal test is a set of constraints specifying allowable combinations of values for subsets of variables
- Simple example of a formal representation language
- **Allows useful general-purpose algorithms with more power than standard search algorithms**

Example: Map-Coloring



- Variables WA, NT, Q, NSW, V, SA, T
- Domains $D_i = \{\text{red, green, blue}\}$
- Constraints: adjacent regions must have different colors
- e.g., $WA \neq NT$, or (WA, NT) in $\{(\text{red, green}), (\text{red, blue}), (\text{green, red}), (\text{green, blue}), (\text{blue, red}), (\text{blue, green})\}$

Example: Map-Coloring



- **Solutions** are **complete** and **consistent** assignments, e.g., WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green



Algorithm : Constraint Satisfaction

1. Propagate available constraints. To do this, first set *OPEN* to the set of all objects that must have values assigned to them in a complete solution. Then do until an inconsistency is detected or until *OPEN* is empty:
 - (a) Select an object *OB* from *OPEN*. Strengthen as much as possible the set of constraints that apply to *OB*.
 - (b) If this set is different from the set that was assigned the last time *OB* was examined or if this is the first time *OB* has been examined, then add to *OPEN* all objects that share any constraints with *OB*.
 - (c) Remove *OB* from *OPEN*.
2. If the union of the constraints discovered above defines a solution, then quit and report the solution.
3. If the union of the constraints discovered above defines a contradiction, then return failure.
4. If neither of the above occurs, then it is necessary to make a guess at something in order to proceed. To do this, loop until a solution is found or all possible solutions have been eliminated:
 - (a) Select an object whose value is not yet determined and select a way of strengthening the constraints on that object.
 - (b) Recursively invoke constraint satisfaction with the current set of constraints augmented by the strengthening constraint just selected.

A Cryptarithmic Problem

Problem:

$$\begin{array}{r} \text{SEND} \\ + \text{MORE} \\ \hline \text{MONEY} \end{array}$$

Initial State:

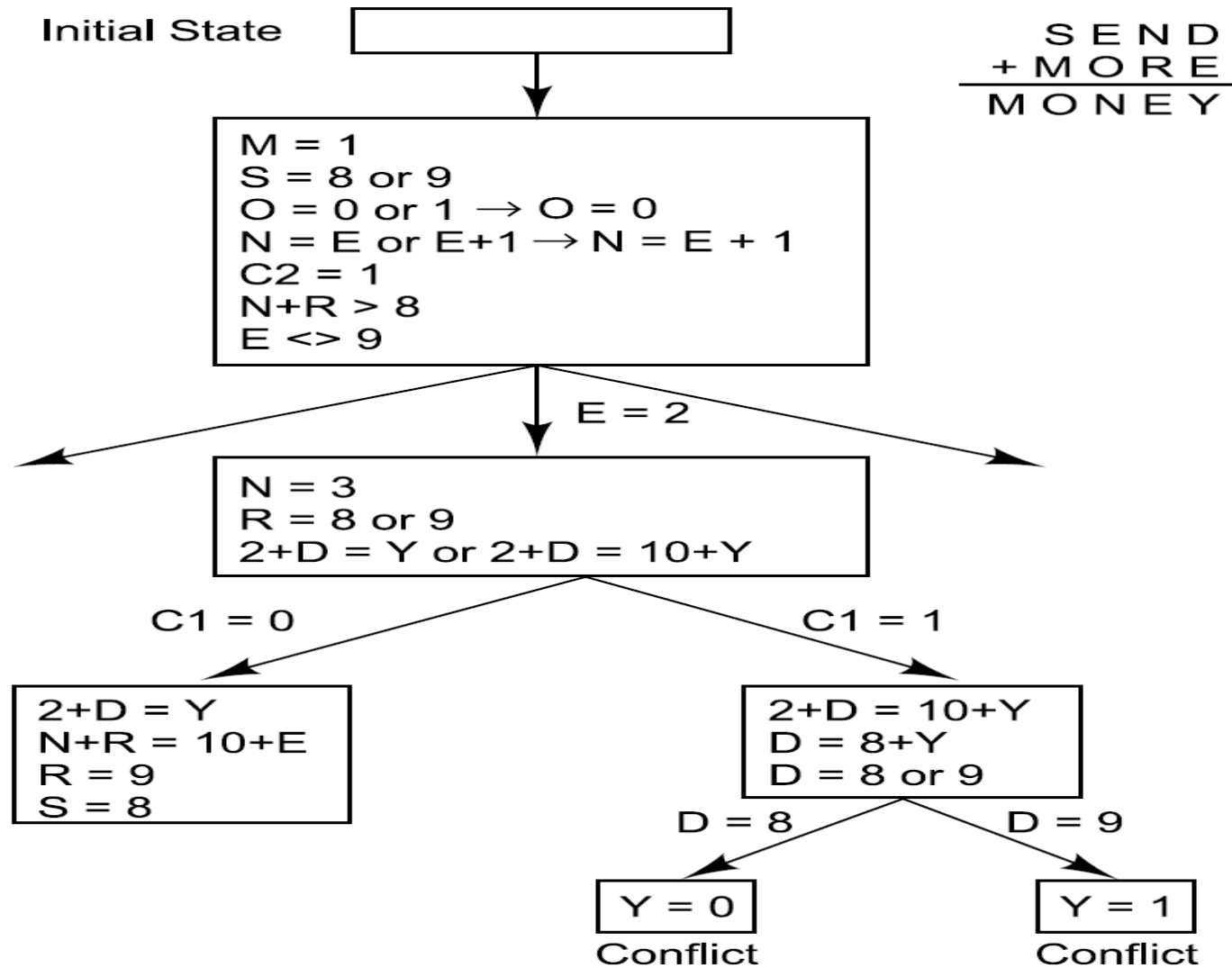
No two letters have the same value.
The sums of the digits must be as shown in the problem.

L
P
U

The diagram illustrates the addition of two words to form a third word. At the top, the word "SEND" is represented by four boxes containing the letters S, E, N, and D. Below it, a blue plus sign (+) is positioned to the left of the word "MORE", which is also represented by four boxes containing the letters M, O, R, and E. A horizontal blue line separates this from the bottom row, which contains a single word "MONKEY" represented by five boxes containing the letters M, O, N, E, and Y.

[illegible]

Solving a Cryptarithmic Problem





	S	E	N	D
+	M	O	R	E
<hr/>				
M	O	N	E	Y

From column 5, **M = 1** since it is the only carry-over possible from the sum of two single digit numbers in column 4.

Since there is a carry in column 5, and $M = 1$, then **O = 0**

Since $O = 0$ there cannot be a carry in column 4 (otherwise N would also be 0 in column 3) so **S = 9**.

If there were no carry in column 3 then $E = N$, which is impossible. Therefore there is a carry and $N = E + 1$.

If there were no carry in column 2, then $(N + R) \bmod 10 = E$, and $N = E + 1$, so $(E + 1 + R) \bmod 10 = E$ which means $(1 + R) \bmod 10 = 0$, so $R = 9$. But $S = 9$, so there must be a carry in column 2 so **R = 8**.

To produce a carry in column 2, we must have $D + E = 10 + Y$.

Y is at least 2 so $D + E$ is at least 12.

The only two pairs of available numbers that sum to at least 12 are (5,7) and (6,7) so either $E = 7$ or $D = 7$.

Since $N = E + 1$, E can't be 7 because then $N = 8 = R$ so **D = 7**.

E can't be 6 because then $N = 7 = D$ so **E = 5** and **N = 6**.

$D + E = 12$ so **Y = 2**.

L
P
U

N	O	O	N
M	O	O	N
S	O	O	N
J	U	N	E

[illegible]