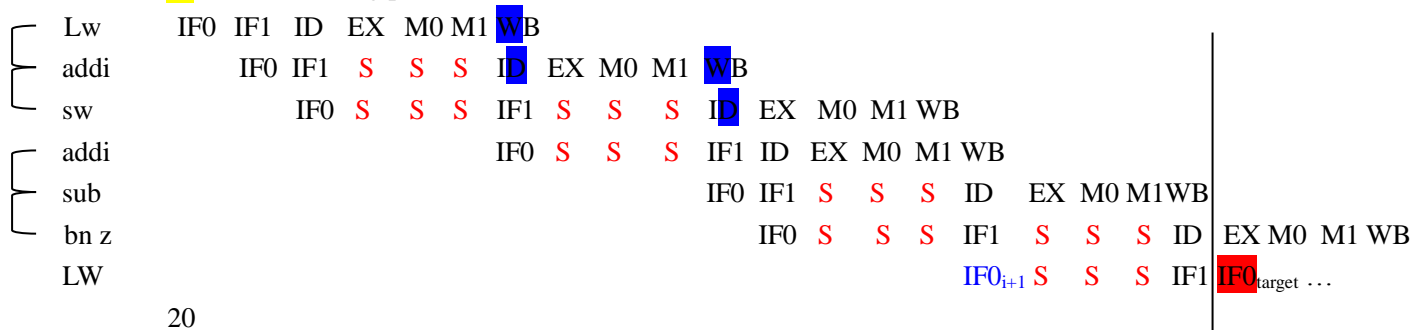


Answer for homework assignment 3

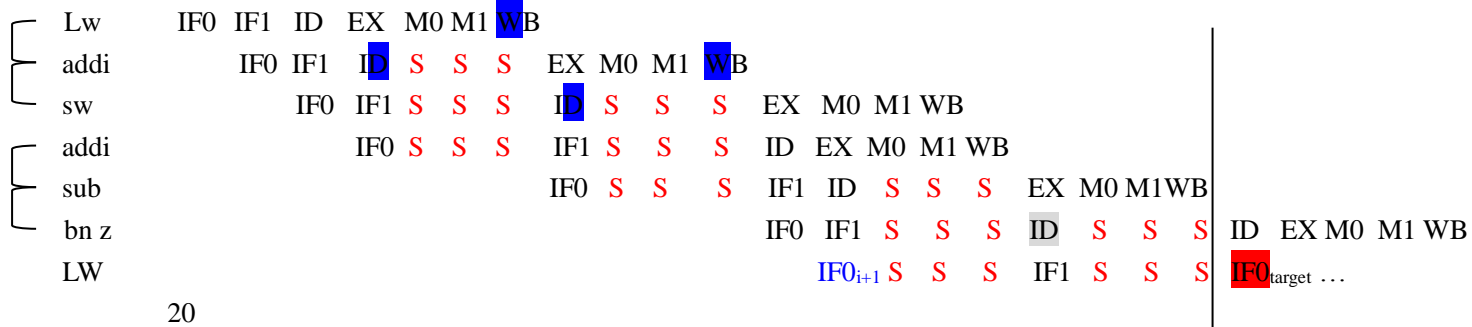
A.1

假设：转移目标地址计算和条件判断在 ID 级进行

(1) 已知：无任何 forwarding path，但在一个时钟周期内对寄存器堆可以进行一次写和一次读（前半周期写，后半周期读）。转移用 flushing 方法。

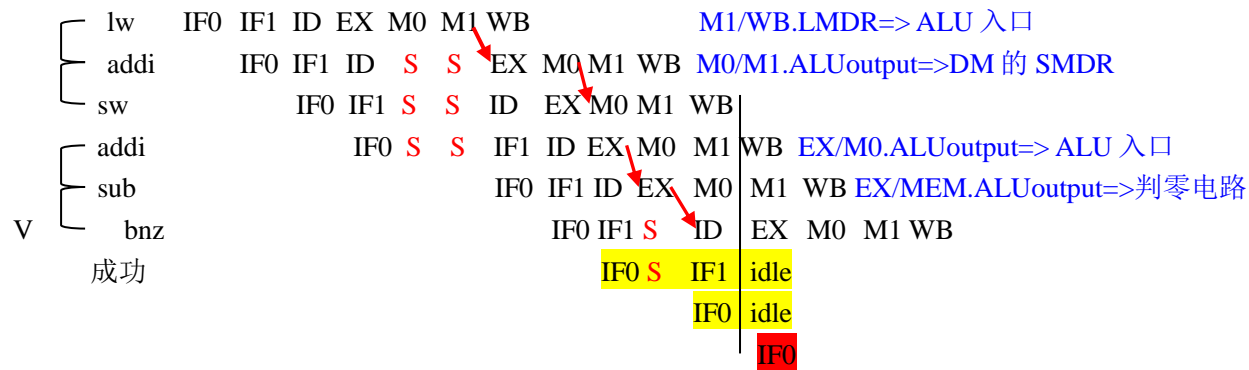


或者) (第二种画流水线的方法，这样也可以)



$$20 \times 98 + 24(20) = 1984(1980)$$

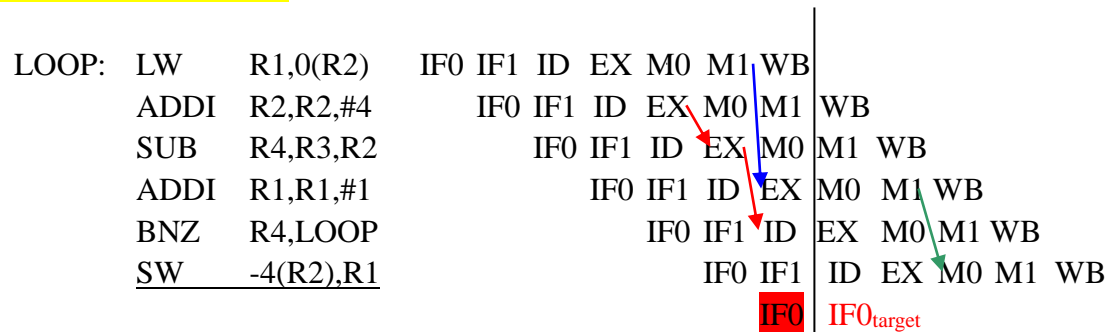
(2) 有 forwarding path, 硬件预测不成功,.



11

$$11 * 98 + 15(13) = 1093(1091)$$

(3) 有一个延时槽, 有 forwarding path, 利用延时转移技术, 可调整指令顺序或部分修改指令操作数, 但不允许减少循环内的指令数, 使得尽可能减少 Stall。



7

$$7 * 98 + 12 = 698$$

(1) 无 forwarding path, 转移目标地址计算和条件判断在 ID 级进行, flushing 方法

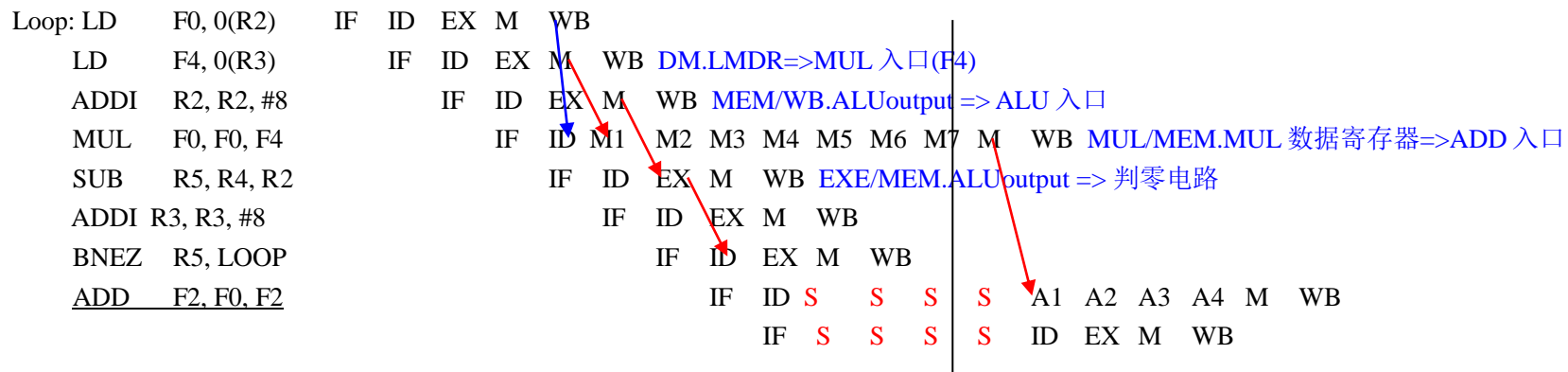
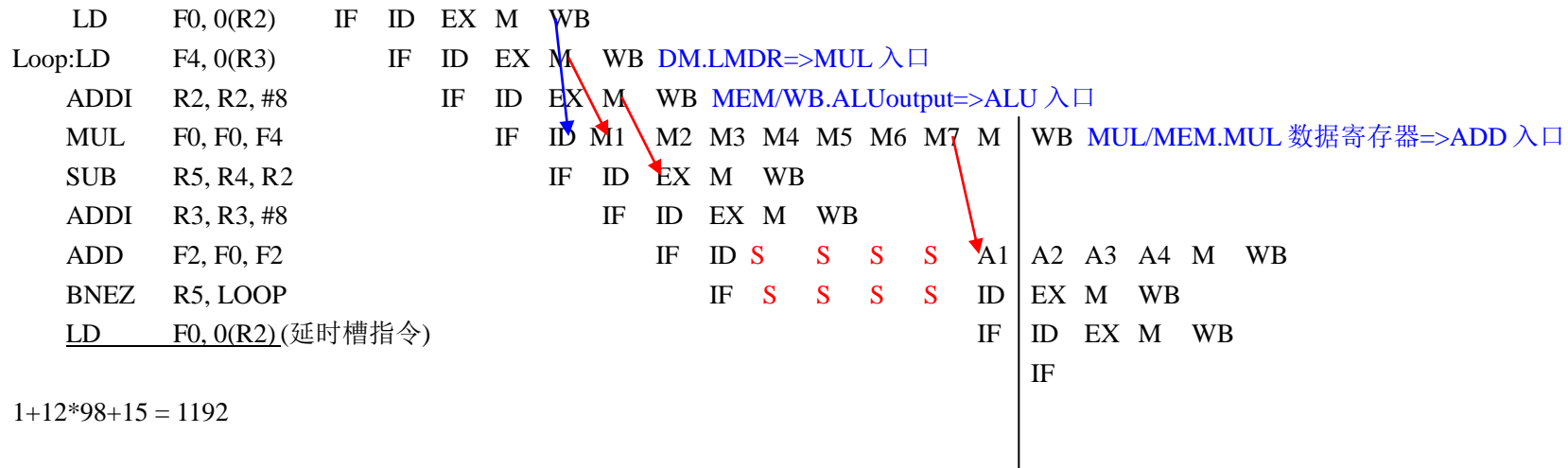
$$22 \cdot 98 + 25 = 2181$$

(2) 有 forwarding path, 转移目标地址计算和条件判断在 ID 级进行, 预测不成功

LD	F0, 0(R2)		IF	ID	EX	M	WB																
LD	F4, 0(R3)			IF	ID	EX	M	WB	MEM/WB.LMDR=>MUL入口														
MUL	F0, F0, F4				IF	ID	S	M1	M2	M3	M4	M5	M6	M7	M	A1	WB	MUL/MEM.MUL数据寄存器=>ADD入口					
ADD	F2, F0, F2					IF	S	ID	S	S	S	S	S	S	S		A2	A3	A4	M	WB写F2		
ADDI	R2, R2, #8							IF	S	S	S	S	S	S		S	ID	EX	M		WB		
ADDI	R3, R3, #8																	IF	ID	EX	M	WB	写R5
SUB	R5, R4, R2																		IF	ID	EX	M	WB EX/MEM.ALUoutput=>判零电路
BNEZ	R5, LOOP																			IF	S	ID	EX M WB
																					IF		idle
																							IF

17*98 + 19 = 1685

(3)



A.3

Given : unconditional resolved at the end of 2nd stage, conditional resolved at the end of 3rd stage.

Unconditional: L1 L2 L3 L4

L1 L1(target) 1 stall

conditional :

1. Flushing: L1 L2 L3 L4

L1 s L1 (target or pc+4) 2 stalls

2. predict-untaken:

untaken L1 L2 L3 L4

L1 L2 L3 L4 0 stall

Taken L1 L2 L3 L4

L1 s L1 (pc+4) 2 stall

3. Predict-taken:

taken L1 L2 L3 L4

L1 L1(t) L2 L3 1 stall

Taken L1 L2 L3 L4

L1 L1(t) L1 (pc+4) 2 stall

设 $CPI_{ideal} = 1$

$$CPI_{flushing} = CPI_{ideal} + 1\% * 1 + 15\% * 2 = 1.31 \quad speedup = 1.31$$

$$CPI_{predict-untaken} = CPI_{ideal} + 1\% * 1 + 15\% * (60\% * 2 + 40\% * 0) = 1.19 \quad speedup = 1.19$$

$$CPI_{predict-taken} = CPI_{ideal} + 1\% * 1 + 15\% * (60\% * 1 + 40\% * 2) = 1.22 \quad speedup = 1.22$$