Using Dip With OS

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Abstract

In this document we describe how to use the Decomposition in Integer Programming (Dip) package with the Optimization Services (OS) package. The code for this example is contained in the folder `OS/examples/osDip`.

1 Using the OS-Dip Example

Currently, the Decomposition in Integer Programming (**Dip** package is not a dependency of the Optimization Services (**OS**) package. In order to run the **osDip** example it is necessary to download both the **OS** and **Dip** package. Download order is not relevant. In the discussion that follows we assume that for both **OS** and **Dip** the user has successfully completed a `configure`, `make`, and `make install`.

I have tested this on several simple plant location problems and generalized assignment problems. I have tested on both the Mac and Linux and it seems to be working. The configure step should generate a working Makefile for your platform. The OS project still does not include Dip as an external so you need to point to the Dip root. There is a variable in the generated Makefile,

DIPPATH =

that need to be filled in after configure. This "should" be the only line in the Makefile you need to edit. So on my machine, I have

DIPPATH = /Users/kmartin/coin/dip-trunk/vpath/

and there exists a directory

/Users/kmartin/coin/dip-trunk/vpath/lib

with the necessary libs and a directory

DIPPATH = /Users/kmartin/coin/dip-trunk/vpath/include

with the necessary headers.

After building the executable run

./osdip –param osdip.parm

Look at the osdip.parm file. You can see by commenting and uncommenting you can run one of three problems that will also get downloaded.

spl1.osil – a simple plant location problem spl2.osil – a second simple plant location problem genAssign.osil – a generalize assignment problem

The osol files (the option files) determine behavior. For example, if you use

osolFiles/spl1-b.osol

then the assingment constraints are the block constraints. If you use

osolFiles/spl1.osol

then the setup forcing constraints are the block constraints. This new example also exhibits the problems I filed ticked on.

2 Simple Plant/Lockbox Location Example

The problem minimizing the sum of the cost of capital due to float and the cost of operating the lock boxes is the problem.

Parameters:

$m-$ number of customers to be assigned a lock box

$n-$ number of potential lock box sites

$c_{ij}-$ annual cost of capital associated with serving customer j from lock box i

f_i- annual fixed cost of operating a lock box at location i

Variables:

$x_{ij}-$ a binary variable which is equal to 1 if customer j is assigned to lock box i and 0 if not

y_i- a binary variable which is equal to 1 if the lock box at location i is opened and 0 if not

The integer linear program for the lock box location problem is

$$
\begin{aligned}
& & \min \quad & \sum_{i=1}^{n}\sum_{j=1}^{m} c_{ij}x_{ij} + \sum_{i=1}^{n} f_i y_i & & (1)\\
&(LB) & \text{s.t.} \quad & \sum_{i=1}^{n} x_{ij} = 1, \qquad j = 1,\ldots,m & & (2)\\
& & & x_{ij} - y_i \leq 0, \qquad i = 1,\ldots,n,\ j = 1,\ldots,m & & (3)\\
& & & x_{ij},\ y_i \in \{0,1\}, \quad i = 1,\ldots,n,\ j = 1,\ldots,m. & & (4)
\end{aligned}
$$

The objective (1) is to minimize the sum of the cost of capital plus the fixed cost of operating the lock boxes. The requirement that every customer be assigned a lock box is modeled by constraint (2). Constraints (3) are forcing constraints and play the same role as constraint set (**??**) in the dynamic lot size model.

Location Example 1: A three-by-five example.

		CUSTOMER					
		1	2	3	4	5	FIXED COSTS
	1	2	3	4	5	7	2
PLANT	2	4	3	1	2	6	3
	3	5	4	2	1	3	3

Location Example 2: A three-by-three example.

$$\min 2x_{11} + x_{12} + x_{13} + x_{21} + 2x_{22} + x_{23} + x_{31} + x_{32} + 2x_{33} + y_1 + y_2 + y_3$$

$$
\begin{aligned}
\text{s.t.} \quad & x_{11} + x_{21} + x_{31} = 1 \\
& x_{12} + x_{22} + x_{32} = 1 \qquad Ax \geq b \text{ constraints} \\
& x_{13} + x_{23} + x_{33} = 1
\end{aligned}
$$

$$
\begin{aligned}
& x_{11} \leq y_1 \leq 1 \\
& x_{12} \leq y_1 \leq 1 \\
& x_{13} \leq y_1 \leq 1 \\
& x_{21} \leq y_2 \leq 1 \\
& x_{22} \leq y_2 \leq 1 \qquad Bx \geq b \text{ constraints} \\
& x_{23} \leq y_2 \leq 1 \\
& x_{31} \leq y_3 \leq 1 \\
& x_{32} \leq y_3 \leq 1 \\
& x_{33} \leq y_3 \leq 1 \\
& x_{ij}, y_i \geq 0,\ i = 1,\ldots,n,\ j = 1,\ldots,m.
\end{aligned}
$$

3 Generalized Assignment Problem Example

A problem that plays a prominent role in vehicle routing is the *generalized assignment problem.* The problem is to assign each of n tasks to m servers without exceeding the resource capacity of the servers.

Parameters:

$n-$ number of required tasks

$m-$ number of servers

$f_{ij}-$ cost of assigning task i to server j

b_j- units of resource available to server j

$a_{ij}-$ units of server j resource required to perform task i

Variables:

$x_{ij}-$ a binary variable which is equal to 1 if task i is assigned to server j and 0 if not

The integer linear program for the generalized assignment problem is

$$\min \sum_{i=1}^{n}\sum_{j=1}^{m} f_{ij}x_{ij} \tag{5}$$

$$(GAP) \qquad \text{s.t.} \quad \sum_{j=1}^{m} x_{ij} = 1, \qquad i = 1,\ldots,n \tag{6}$$

$$\sum_{i=1}^{n} a_{ij}x_{ij} \le b_j, \qquad j = 1,\ldots,m \tag{7}$$

$$x_{ij} \in –0, 1'', \; i = 1,\ldots,n, \; j = 1,\ldots,m. \tag{8}$$

The objective function (5) is to minimize the total assignment cost. Constraint (6) requires that each task is assigned a server. The requirement that the server capacity not be exceeded is given in (7).

The test problem

```
 min     2 X11 + 11 X12 + 7 X21 + 7 X22 + 20 X31 +
2 X32 + 5 X41 + 5 X42
s.t.
X11 + X12 =    1
X21 + X22 =    1
X31 + X32 =    1
X41 + X42 =    1

3 X11 + 6 X21 + 5 X31 + 7 X41 <=   13
2 X12 + 4 X22 + 10 X32 + 4 X42 <=   10
```