Modern C++ for Parallelism in Scientific Computing

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Scientific computing parallelism

- Large amounts of data: often cartesian multi-dimensional arrays, sometimes unstructured d
- Large amounts of parallelism: each element of output array independent.
- No explicit threading parallelism created by some runtime
- Range algorithm notion: do some operation on each element of a dataset

Power method

I et A a matrix of interest I et x be a random vector For iterations until convergence compute the product $y \leftarrow Ax$ compute the norm $\gamma = ||y||$ normalize $x \leftarrow y/\gamma$

- ► Method for computing largest eigenvalue of a matrix
- ► Also Google Pagerank
- ► Stands for many scientific codes: Krylov methods, eigenvalue:

Stencil operations







- ▶ This rectangular m × n thing is the vector
 - ► The 4, -1,... stencil is / stands for the matrix.
 - Goes by: difference stencil, convolution, Toeplitz matrix

Array parallelism

Traditional C implementation:

- for (idxint i=0; i<m; i++) for (idxint 1-0: 1<n: 1++ out[IINDEX(i, j, m, n, b)] = in[IINDEX(i, j, m, n, b)
- ► Two / three-dimensional loop
- all dimensions large
- every output element independent

Reductions ℓ₂ reduction:

: for (idxint i=0; i<m; i++) for (idxint j=0; j<n; j++) { auto v = out[IINDEX(i, j, m, n, b)];</pre> sum_of_squares += v*v;

- return std::sqrt(sum_of_squares);
- Parallel except for the accumulation Obviously should not be done through at

Differential operator / image convolution

Stencil computation

Apply stencil to each (i,j) index:

► Structure can be more complicated in scientific codes

for (idxint i=0; i<m; i++) {
 for (idxint j=0; j<n; j++) {
 out [IINDEX(i,j,m,n,b)] - 4*in[IINDEX(i,j,m,n,b)]
 - in[IINDEX(i=1,j,m,n,b)] - in[IINDEX(i+1,j,m,n,b) - in[IINDEX(i=1,j,m,n,b)] - in[IINDEX(i=1,j,m,n,b]] - in[I

Tools: mdspan and cartesian_product

Data is declared as mdspan

```
private:
    real *_data{nullptr};
                                  //! pointer to the data as 2D array auto& data2d() {
                                    return cartesian_data; };
                                   const auto@ data2d() const
    md::dextents<idxint.2>
                                    return cartesian_data; };
             > cartesian_data;5
```

mdspan and cartesian_product

Index range is declared as range::views::cartesian product;

```
const auto@ s = data2d();
int b = this->border();
a idxint
     lo_m = static_cast<idxint>(b),
     hi_m = static_cast<idxint>(s.extent(0)-b),
     lo n - static cast<idxint>(b).
   hi_n = static_cast<idxint>(s.extent(1)-b);
range2d = rng::views::cartesian_product
o (rng::views::iota(lo_m,hi_m),rng::views::iota(lo_n,hi_n));
```

- ▶ Vector allocated with size $(m+2b) \times (n+2b)$ to include border
- ▶ for handling of boundary conditions / halo regions in PDEs.

Implementation 1: OpenMP parallelism

Annotate loops as parallel and/or reduction:

```
*pragma omp parallel for reduction(+:sum_of_squares)
for ( idxint i=0; i<m; i++ )
  for ( idxint j=0; j<n; j++ ) {
    auto v = out[ IINDEX(i, j, m, n, b) ];</pre>
      sum_of_squares += v*v;
return std::sqrt(sum_of_squares);
```

- Static assigment of iterations to threads by default
- Highly controlled affinity
- ▶ 'oned' as above, 'clps' for both loops collapsed
- Can be formulated as range algorithm.

Implementation 2: range over indices

Range-based for loop:

- auto array = this->data2d(); #pragma omp parallel for reduction(+:sum_of_squares)
 for (auto ij : this->inner()) {
 auto [i, j] = ij; auto v = array[i, j];
 sum_of_squares += v*v;
- return std::sqrt(sum_of_squares); ► Range over indices, not over data
- Indices are a subset of the full data!

Stencil operation

// span.cpp auto out = this->data2d(); auto out = this>data2d(); const auto in = other.data2d(); #pragma omp parallel for for (auto ij : this>inner()) { auto [i, j] = ij; out[i, j] = 4*in[i, j] - in[i-1, j] - in[i+1, j] - in[i, j-1] - in[i, j+1

Hard to formulate as range algorithm

Most complicated operation of the bunch:

Performance not necessarily determined by floating point op

Implementation 3: Sycl

Open standard, but mostly pushed by Intel

```
q.submit([%] (handler %h) {
  accessor D_a(Buf_a,h,write_only);
h.parallel_for
(range<2>(msize-2,nsize-2),
         [Fanget2](ms12e=2, ns12e=2),
[=](auto index)
auto row = index.get_id(0) + 1;
auto col = index.get_id(1) + 1;
D_a[row][col] = 1.;
11.weit():
```

- Heterogeneous CPU/GPU code. transparent data movement
- Range algorithm-like syntax. but explicit task queue

Implementation 4: Kokkos

Open Source heterogeneous execution laver

```
Kokkos::parallel_for
   ("Update x",
Kokkos::MDRangePolicy<Kokkos::Rank<2>>
   ({1, 1}, {msize-1, nsize-1}),

KOKKOS_LAMBDA(int i, int j) {

x(i, j) = Ax(i, j) / norm;
```

- Same code for CPU and GPU
- Implicit task queue
- Two-dimensional indexing
- Range algorithm-like philosophy

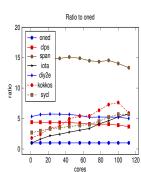
Comparing models (Intel)

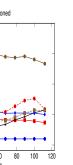
oned dps - span → iota 104 → diy2e - + - kokkos -#- sycl

60

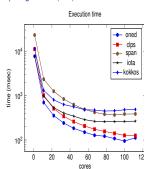
100

Ratio to fastest (Intel)

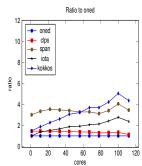




Comparing models (Gcc)



Ratio to fastest (Gcc)



Where do we lose performance?

Conclusion and Acknowledgement

Hint: perf output on the 'span' variant:

⇒std::ranges::cartesian product viewstd::ranges::iota viewlong ⇒long>, std::ranges::iota_view<long, long> ⇒>:: Iterator<true>::operator+= 18.73% [.] __divti3 11.33% [.] $\hookrightarrow \texttt{linalg::bordered_array_span} \land \texttt{float} \gt :: \texttt{central_difference_from}$ 5.37% [.] \hookrightarrow linalg::bordered_array_span<float>::scale_interior 5.01% [.] linalg::bordered_array_span<float>::12norm

2.69% [.] divti3@plt Index calculations take lots of time

▶ 'Fancy' schemes suffer from indexing overhead strongly implementation and compiler dependent

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Intel compiler. C-style variant fastest.

20 40

Gcc compiler, less variance between variants