Representation Learning on Graphs and Networks (L45) CST Part III / MPhil in ACS

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1 Primer on Graph Representations

1. Mathematical definition of graphs:

A graph $G = (V, \mathcal{E})$ is a collection of nodes V and edges $\mathcal{E} \subseteq V \times V$.

The edges can be represented by an *adjacency matrix*, $\mathbf{A} \in \mathbb{R}^{|\mathcal{V}| \times |\mathcal{V}|}$, such that

$$A_{uv} = \begin{cases} 1 & \text{if } (u, v) \in \mathcal{E} \\ 0 & \text{otherwise} \end{cases}$$

- 2. Some interesting graph types:
 - Undirected: $\forall u, v \in \mathcal{V}. (u, v) \in \mathcal{E} \iff (u, v) \in \mathcal{E} \text{ (i.e., } \mathbf{A}^T = \mathbf{A})$
 - Weighted: provided *edge weight* w_{uv} for every edge $(u, v) \in \mathcal{E}$
 - **Multirelational**: various *edge types*, i.e. $(u, t, v) \in \mathcal{E}$ if there exists an edge (u, v) linked by type t
 - Heterogeneous: various node types
- 3. Machine learning tasks on graphs by domain:
 - Transductive: training algorithm sees all observations, including the holdout observations
 - Task is to *propagate* labels from the training observations to the holdout observations
 - Also called semi-supervised learning
 - **Inductive**: training algorithm only sees the training observations during training, and only sees the holdout observations for prediction.
- 4. Node statistics:
 - Degree: amount of edges the node is incident to:

$$d_u = \sum_{v \in \mathcal{V}} A_{uv}$$

• **Centrality**: a measure of how "central" the node is in the graph: how often do infinite random walks visit the node?

$$d_u = \lambda^{-1} \sum_{v \in \mathcal{U}} A_{uv} e_v$$

where $\mathbf{e} \in \mathbb{R}^{|\mathcal{V}|}$ is the largest eigenvector of \mathbf{A} , with corresponding eigenvalue λ .

• Clustering coefficient: a measure of "clusteredness": are neighbours connected amongst each other?

$$c_{u} = \frac{\left| \{ (v_{1}, v_{2}) \in \mathcal{E} : v_{1}, v_{2} \in \mathcal{N}(u) \} \right|}{\binom{d_{u}}{2}}$$

5. Graph Laplacian:

Let **D** be the diagonal (out)-degree matrix of the graph, i.e., $D_{uu} = \sum_{v \in V} A_{ij}$. Then:

- The *unnormalised* graph Laplacian: L = D A
- The symmetric graph Laplacian: $\mathbf{L}_{\text{sym}} = \mathbf{D}^{-\frac{1}{2}} \mathbf{L} \mathbf{D}^{-\frac{1}{2}} = \mathbf{I} \mathbf{D}^{-\frac{1}{2}} \mathbf{A} \mathbf{D}^{-\frac{1}{2}}$
- The random walk graph Laplacian: $L_{RW} = D^{-1}L = I D^{-1}A$

Properties:

- For undirected graphs, **L** is *symmetric* ($\mathbf{L}^T = \mathbf{L}$) and *positive semi-definite* ($\forall \mathbf{x} \in \mathbb{R}^{|V|}$. $\mathbf{x}^T \mathbf{L} \mathbf{x} \geq 0$)
- For all $\mathbf{x} \in \mathbb{R}^{|\mathcal{V}|}$:

$$\mathbf{x}^T \mathbf{L} \mathbf{x} = \frac{1}{2} \sum_{u \in \mathcal{V}} \sum_{v \in \mathcal{V}} A_{uv} (x_u - x_v)^2 = \sum_{(u,v) \in \mathcal{E}} (x_u - x_v)^2$$

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• L has $|\mathcal{V}|$ nonnegative eigenvalues: $\lambda_1 \ge \cdots \ge \lambda_{|V|} = 0$

6. Spectral clustering:

• Two-way cut: partition the graph into $\mathcal{A} \subseteq \mathcal{V}$ and its complement $\mathcal{A}_c \subseteq \mathcal{V}$:

$$\mathrm{Cut}(\mathcal{A}) = \big| \{ (u, v) \in \mathcal{E} : u \in \mathcal{A} \land v \in \mathcal{A}_c \} \big|$$

Ratio cut metric:

$$RCut(\mathcal{A}) = Cut(\mathcal{A}) \left(\frac{1}{|\mathcal{A}|} + \frac{1}{|\mathcal{A}_c|} \right)$$

• Minimising $RCut(\mathcal{A})$:

Let $\mathbf{a} \in \mathbb{R}^{|\mathcal{V}|}$ be a vector representing the cut \mathcal{A} , defined as follows:

$$a_{u} = \begin{cases} \sqrt{\frac{\mathcal{A}_{c}}{\mathcal{A}}} & \text{if } u \in \mathcal{A} \\ -\sqrt{\frac{\mathcal{A}}{\mathcal{A}_{c}}} & \text{if } u \in \mathcal{A}_{c} \end{cases}$$

Then

$$\mathbf{a}^T \mathbf{L} \mathbf{a} = \sum_{(u,v) \in \mathcal{E}} (a_u - a_v)^2 = |\mathcal{V}| \mathrm{RCut}(\mathcal{A})$$

Minimising $\mathbf{a}^T \mathbf{L} \mathbf{a}$ corresponds to minimising $\mathrm{RCut}(\mathcal{A})$ (NP-hard as the condition is discrete)

- Relaxing: minimise $\mathbf{a}^T \mathbf{L} \mathbf{a}$ subject to $\mathbf{a} \perp \mathbf{1}$ and $||\mathbf{a}||^2 = |\mathcal{V}|$ Rayleigh–Ritz Theorem: The solution is exactly the second-smallest eigenvector of \mathbf{L} To obtain the cut, place u into \mathcal{A} or \mathcal{A}_c depending on the sign of a_u
- Can be generalised to *k*-clustering