

# Relations

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# What is a Relation?

In discrete mathematics, relation is a way of showing a relationship between any two sets.

- ❑ Relationship between any program and its variable.
- ❑ Relationship between pair of cities linked by railway in a network.

# Necessity for studying Relation

- Relational Database model is based on the concept of relation.

## Cartesian Product

- Given two sets  $A$  and  $B$ , their **cartesian product**  $A \times B$ , is defined as

$$A \times B = \{(a, b) \mid a \in A, b \in B\}$$

## Ordered Pairs

- The elements of  $A \times B$  are called **ordered pairs** with the elements of  $A$  as the first entry and elements of  $B$  as the second entry.
- Order matters

## Special Case:

$$A^2 = A \times A = \{(a_1, a_2) \mid a_1, a_2 \in A\}$$

Similarly,

$$A^n = A \times A \times \cdots \times A (n \text{ times}) = \{(a_1, a_2, \dots, a_n) \mid a_1, a_2, \dots, a_n \in A\}$$

Relation is the subset of the cartesian product of the sets.

## n-ary Relation

- Let  $\{A_1, A_2, \dots, A_n\}$  be  $n$  sets.
- An **n-ary relation**  $R$  on  $A_1 \times A_2 \times \dots \times A_n$  is a subset of  $A_1 \times A_2 \times \dots \times A_n$ .
- If  $A_i = A ; \forall i$  , then  $R$  is called the **n-ary relation on  $A$** .



## Empty and Universal Relation

- If  $R = \emptyset$ , then  $R$  is called the **empty** or **void relation**.
- If  $R = A_1 \times A_2 \times \dots \times A_n$ , then  $R$  is called the **universal relation**.

## Definition (Binary Relation)

- Given two sets  $A$  and  $B$ , a relation between  $A$  and  $B$  is a **subset of  $A \times B$** .
- If  $R$  is a relation on  $A \times B$  (i.e.,  $R \subseteq A \times B$ ) and  $(a, b) \in R$ , we say “ **$a$  is related to  $b$** ”.
- It can also be written as  **$aRb$** .

### Example:

Let  $A = \{a, b\}$  and  $B = \{2, 3, 4\}$

$R = \{(a, 3), (b, 2), (b, 4)\}$  is a relation from  $A$  to  $B$ .

# Binary Relation on a set

- A binary relation  $R$  on a set  $A$  is a **subset of  $A \times A$** .

## Examples:

1. “Taller -than ” is a relation on people.  
 $(a, b) \in$  “Taller -than” if person  $a$  is taller than person  $b$ .
2. “ $\geq$ ” is a relation on real set  $\mathbf{R}$ .  
“ $\geq$ ” =  $\{(x, y) \in \mathbf{R} \mid x, y \in \mathbf{R}, x \geq y\}$

## Examples (Cont..)

3. Let  $A = \{1, 2, 3, 4, 5, 6\}$ .

If  $R = \{(a, b) \mid a \text{ divides } b\}$  is a relation from  $A$  to  $B$  then ordered pairs in the relation  $R$  are

$\{(1,1), (1,2), (1,3), (1,4), (1,5), (1,6), (2,4), (2,6), (3,6)\}$

## Examples (Cont..)

Let  $A = \{1, 2, 3\}$

$A \times A$

$= \{(1,1), (1,2), (1,3), (2,1), (2,2), (2,3), (3,1), (3,2), (3,3)\}$

- Here,  $A \times A$  is an universal relation on  $A$ .
- $\emptyset$  is an empty relation on  $A$ .

## Examples (Cont..)

$$“=” = \Delta = \{(1,1), (2,2), (3,3)\}$$

$$“<” = \{(1,2), (1,3), (2,3)\}$$

$$“>” = \{(2,1), (3,1), (3,2)\}$$

$$“\leq” = \{(1,1), (1,2), (1,3), (2,2), (2,3), (3,3)\}$$

$$“\geq” = \{(1,1), (2,1), (2,2), (3,1), (3,2), (3,3)\}$$

$$“|” = \{(1,1), (1,2), (1,3), (2,2), (3,3)\}$$

$$“\text{multiple of}” = \{(1,1), (2,1), (2,2), (3,1), (3,3)\}$$

# Representing Relations

Relations can be represented in two ways:



Matrix

Graph

# Representation of Relations as Matrix

- If  $R$  is a relation on set  $A = \{a_1, a_2, \dots, a_n\}$  and  $|A| = n$ , then it can be represented as  $n \times n$  Boolean Matrix  $M_R$ .

$M_R$  can be defined as:

$$M_R = [m_{ij}]_{n \times n}$$

where,  $m_{ij} = \begin{cases} 0 & ; \text{if } (a_i, a_j) \notin R \\ 1 & ; \text{if } (a_i, a_j) \in R \end{cases}$



## Examples

- Let  $A = \{1, 2, 3\}$
- Let  $R = \{(1,1), (1,2), (2,1), (2,3), (3,2), (3,3)\}$  be a relation on  $A$ .

$$M_R = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

## Examples (Cont..)

- Let  $A = \{1, 2, 3\}$  and  $B = \{1, 2, 3, 4, 5\}$ . Which ordered pairs are in the relation  $R$  represented by the matrix

$$M_R = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 \end{bmatrix} ?$$

- Because  $R$  consists of those ordered pairs with  $a_{ij} = 1$ , it follows that:

$$R = \{(1, 2), (2, 1), (2, 3), (2, 4), (2, 5), (3, 1), (3, 3), (3, 5)\}.$$

# Representation of Relations as a Digraph (Directed Graph)

- The graph of a relation  $R$  over  $A$  is a directed graph with nodes corresponding to the elements of  $A$ . There is an edge from node  $x$  to  $y$  if and only if  $(x, y) \in R$ .
- An edge of the form  $(x, x)$  is called a self-loop.

## Examples

- Let  $A = \{1, 2, 3\}$
- Let  $R_1 = \{(1, 2), (1, 3), (2, 3)\}$  be a  $<$  relation on  $A$ .

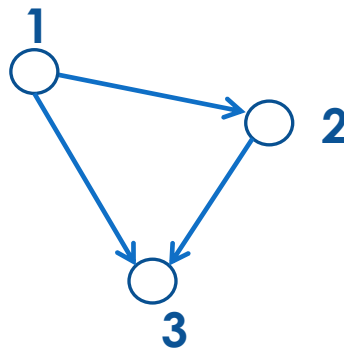
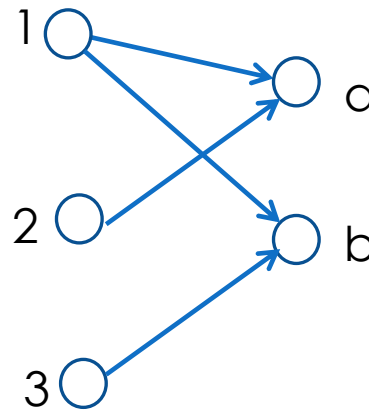


Figure 1

## Examples (Cont..)

- Let  $A = \{1,2,3\}$  and  $B = \{a,b\}$
- Let  $R_2 = \{(1, a), (1, b), (2, a), (3, b)\}$  be a relation from  $A$  to  $B$ .



**Figure 2**

# Domain and Range

**Domain** of Relation  $R$  = set of all first co-ordinates

**Range** of Relation  $R$  = set of all second co-ordinates

## Example

$< = \{(1,2), (1,3), (2,3)\}$  on  $A = \{1,2,3\}$

□ Domain of  $< = \{1,2\}$

□ Range of  $< = \{2,3\}$

## Equality of Two relations

- Let  $R_1$  be an  $n$ -ary relation on  $A_1 \times A_2 \times \dots \times A_n$ .
- Let  $R_2$  be an  $m$ -ary relation on  $B_1 \times B_2 \times \dots \times B_m$ .
- Then,  $R_1 = R_2$   
If and only if
  - ❖  **$n=m$**
  - ❖  **$A_i = B_i; \forall i, 1 \leq i \leq n$**
  - ❖ **and,  $R_1$  &  $R_2$  are equal set of ordered pairs.**

## Example

- Let  $A = \{a, b\}, B = \{1, 2\}, C = \{1, 2, 3\}$
- Let  $R_1 = \{(a, 1), (b, 2)\}$  is a relation on  $A \times B$
- Let  $R_2 = \{(a, 1), (b, 2)\}$  is a relation on  $A \times C$

$$R_1 = R_2?$$

**No**

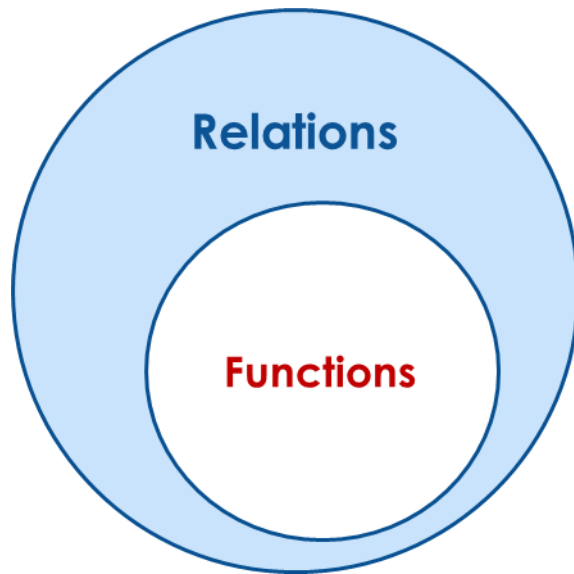
How many number of relations  
are there on a set  $A$  having  $n$   
elements?

$$2^{n^2}$$



Thank  
you!!!





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# Properties of Relations

- Reflexive
- Symmetric
- Transitive
- Irreflexive
- Asymmetric
- Antisymmetric



# Reflexive Relations

- $R$  is **reflexive** iff  $(x, x) \in R$  for every element  $x \in A$ .

## Examples

1. Let  $A = \{1, 2, 3\}$

Suppose  $R_1 = \{(1,1), (2,2), (2,3)\}$  be a relation on  $A$ .

Is  $R_1$  reflexive?

No

2.  $=, A \times A, \leq, \geq, |, \text{multiple of}$

Reflexive? Yes

3.  $\emptyset, <, >$

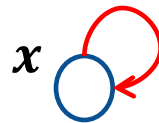
Reflexive? No

# Reflexive Relation in Matrix and Graph

- If  $R$  is a **reflexive** relation, all the elements on the main diagonal of  $M_R$  are equal to 1.

$$M_R = \begin{bmatrix} 1 & \cdots & \\ \vdots & \ddots & \vdots \\ \cdots & & 1 \end{bmatrix}$$

- A loop must be present at all vertices in the graph.



# Symmetric Relations

- $R$  is **symmetric** iff  $(y, x) \in R$  whenever  $(x, y) \in R$  for all  $x, y \in A$ .

## Examples

1. Let  $A = \{1, 2, 3\}$

Suppose  $R_2 = \{(1, 2), (2, 1), (2, 3)\}$  be a relation on  $A$ .

Is  $R_1$  Symmetric?

No

2. "sibling-of" is **symmetric**, but "sister-of" is **not**.

3.  $A \times A, \emptyset, =$  **Symmetric?** **Yes**

4.  $<, >, \leq, \geq, |, \text{multiple of}$  **Symmetric?** **No**

# Symmetric Relation in Matrix and Graph

- $R$  is a symmetric relation if and only if  $m_{ji} = 1$ , whenever  $m_{ij} = 1$ .

$$M_R = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$

- If  $(x, y)$  is an edge in the graph, then there must be an edge  $(y, x)$  also.





## Transitive Relations

- A relation  $R$  on a set  $A$  is called **transitive** if whenever  $(x, y) \in R$  and  $(y, z) \in R$ , then  $(x, z) \in R$ , for all  $x, y, z \in A$ .

### Examples

1. Let  $A = \{1, 2, 3\}$

Suppose  $R_3 = \{(1, 3), (3, 1)\}$  be a relation on  $A$ .

Is  $R_3$  Transitive?

No

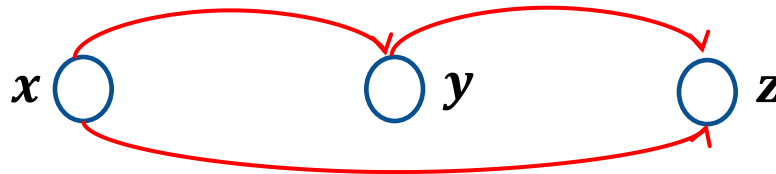
2.  $A \times A$ ,  $\emptyset$ ,  $=$ ,  $<$ ,  $>$ ,  $\leq$ ,  $\geq$ ,  $|$ , multiple of

Transitive?

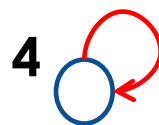
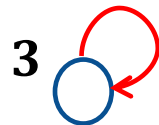
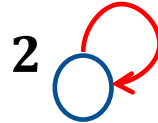
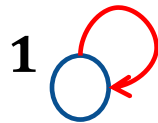
Yes

## Transitive Relations in Graph

- $R$  is transitive iff in its graph, for any three nodes  $x, y$  and  $z$  such that there is an edge  $(x, y)$  and  $(y, z)$ , there exists an edge  $(x, z)$ .



# Examples

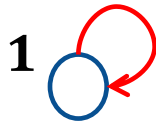


Equality Relation on  
 $A = \{1, 2, 3, 4\}$

$$M_R = \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1 \end{bmatrix}$$

- Reflexive?  
Yes
- Symmetric?  
Yes
- Transitive?  
Yes

## Examples (Cont..)



- Reflexive?

**No**

- Symmetric?

**Yes**

- Transitive?

**Yes**

## Examples (Cont..)

- Suppose that the relation  $R$  on a set is represented by the matrix

$$M_R = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

- Reflexive?

Yes

- Symmetric?

Yes

How many number of **Reflexive Relations** are there on set A having n elements?

$$2^{n(n-1)}$$

How many number of **Symmetric Relations** are there on set A having n elements?

$$2^{n(n+1)/2}$$

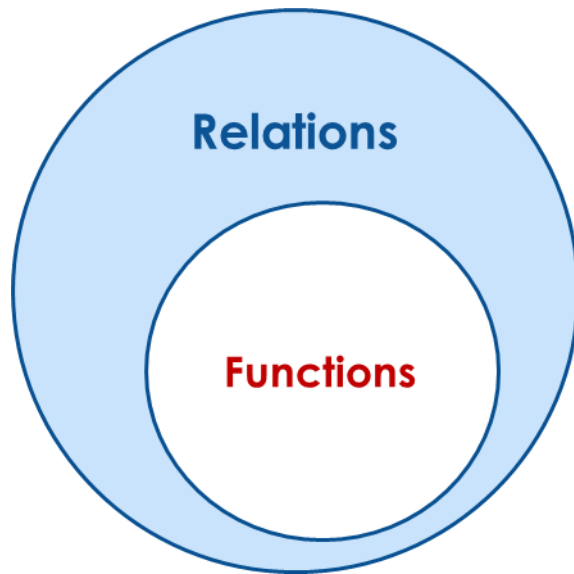
How many number of **Transitive Relations** are there on set  $A$  having  $n$  elements?

*No closed form found*



Thank  
you!!!





# Relations

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# Properties of Relations

- Reflexive
- Symmetric
- Transitive
- Irreflexive
- Asymmetric
- Antisymmetric

# Irreflexive Relations

- $R$  is **irreflexive** iff  $(x, x) \notin R$  for every element  $x \in A$ .
- No Reflexive ordered pair should belong to the relation.

## Examples

1. Let  $A = \{1, 2, 3\}$

Suppose  $R_1 = \{(1,1), (2,2), (2,3)\}$  be a relation on  $A$ .

Is  $R_1$  Irreflexive?

No

2.  $\emptyset, <, >$

Irreflexive? Yes

3.  $\Delta, A \times A, \leq, \geq, |, \text{multiple of}$

Irreflexive? No



# Irreflexive Relation in Matrix and Graph

- If  $R$  is an irreflexive relation, all the elements on the main diagonal of  $M_R$  are equal to 0.

$$M_R = \begin{bmatrix} 0 & \cdots & \\ \vdots & \ddots & \vdots \\ \cdots & & 0 \end{bmatrix}$$

- No vertex should contain self-loop in the graph.



# Asymmetric Relations

- A relation  $R$  on a set  $A$  such that for all  $x, y \in A$ , if  $(x, y) \in R$  then  $(y, x) \notin R$ , is called **asymmetric**.

## Examples

Let  $A = \{1, 2, 3\}$

1. Suppose  $R_2 = \{(1, 2)\}$  be a relation on  $A$ .

Is  $R_2$  Asymmetric?

Yes

2. Suppose  $R_3 = \{(1, 3), (3, 1), (2, 3)\}$  be another relation on  $A$ .

Is  $R_3$  Asymmetric?

No

3.  $\emptyset, <, >$

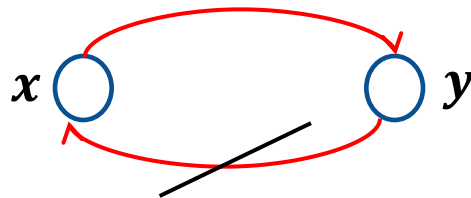
Asymmetric? Yes

4.  $A \times A, \leq, \geq, |$ , multiple of

Asymmetric? No

# Asymmetric Relations in Graph

- If  $(x, y)$  with  $x \neq y$  is an edge, then  $(y, x)$  is not an edge.
- There must also be no self loop.





# Antisymmetric Relations

- A relation  $R$  on a set  $A$  such that for all  $x, y \in A$ , if  $(x, y) \in R$  and if  $(y, x) \in R$ , then  $x = y$ , is called **antisymmetric**.

If  $x \neq y$  and if  $(x, y)$  is present, then  $(y, x)$  should not be present there.

## Examples

1. Let  $A = \{1, 2, 3\}$

Suppose  $R_1 = \{(1, 2), (2, 1), (2, 3)\}$  be a relation on  $A$ .

Is  $R_1$  Antisymmetric?

No

2.  $\emptyset, \Delta, <, >, \leq, \geq, |$ , multiple of

Antisymmetric? Yes

3.  $A \times A$

Antisymmetric? No

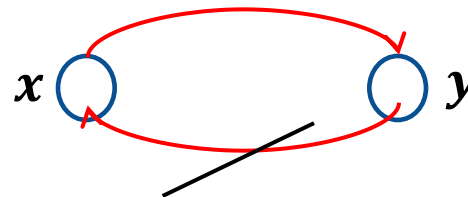
# Antisymmetric Relation in Matrix and Graph

- $R$  is an antisymmetric relation if and only if  $m_{ji} = 0$ , or  $m_{ij} = 0$ , when  $i \neq j$ .

$$M_R = \begin{bmatrix} & 1 \\ 0 & \end{bmatrix}$$

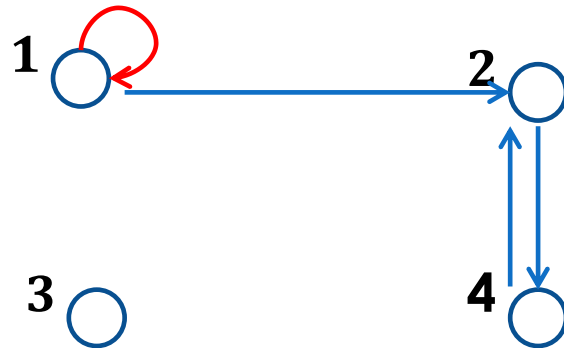
The matrix is a 2x2 matrix with red numbers. The top-right element is 1, the bottom-left element is 0, and the diagonal elements are 0. Blue lines cross the diagonal from top-left to bottom-right and bottom-left to top-right.

- If  $(x, y)$  with  $x \neq y$  is an edge, then  $(y, x)$  is not an edge.



Self-loops can be there.

# Example



- Reflexive?  
No
- Symmetric?  
No
- Transitive?  
No
- Irreflexive?  
No
- Antisymmetric?  
No
- Asymmetric ?  
No

## Some Points to remember

- There can be a relation which is neither reflexive nor irreflexive.

### Example

1. Let  $A = \{1, 2, 3\}$

Suppose  $R_3 = \{(1,1), (2,2), (2,3)\}$  be a relation on  $A$ .

**Neither Reflexive nor Irreflexive**

## Some Points to remember (Cont..)

- There can be a relation which is neither symmetric nor antisymmetric.

### Example

Let  $A = \{1, 2, 3\}$

Suppose  $R_5 = \{(1,2), (2,3), (3,2)\}$  be a relation on  $A$ .

**Neither Symmetric nor Antisymmetric**

How many number of **Irreflexive Relations** are there on set A having n elements?

$$2^{n(n-1)}$$

How many number of **Asymmetric Relations** are there on set A having n elements?

$$3^{n(n-1)/2}$$

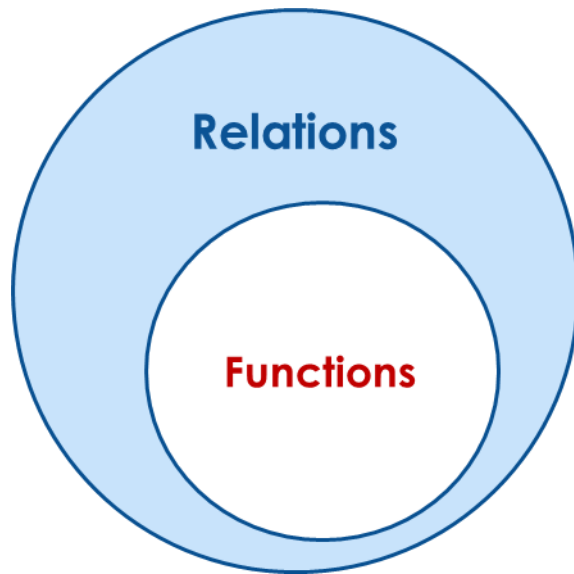


How many number of  
**Antisymmetric Relations** are there  
on set A having n elements?

$$2^n 3^{n(n-1)/2}$$

Thank  
you!!!





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## Inverse Relation

- If  $R \subseteq A \times B$  then  $R^{-1} \subseteq B \times A$ , and is defined as:

$$R^{-1} = \{(b, a) | (a, b) \in R\}$$

$R$	$R^{-1}$
$<$	$>$
$\leq$	$\geq$
<i>divides</i>	<i>multiple of</i>
<i>subset</i>	<i>superset</i>

## Example

- Let  $A = \{1, 2, 3\}$  and  $B = \{3, 4, 5\}$   
Suppose  $R = \{(1,3), (1,5), (2,4), (3,5)\}$

$$R^{-1} = \{(3,1), (5,1), (4,2), (5,3)\}$$

# Complementary Relations

- Let  $R$  be a relation from  $A$  to  $B$ , then complementary relation  $R^C$  is defined as:

$$R^C = \{(a, b) | (a, b) \notin R \text{ and } (a, b) \in A \times B\}$$

## Example

- Let  $A = \{1, 2, 3\}$  and  $B = \{3, 4, 5\}$

Suppose  $R = \{(1,3), (1,5), (2,4), (3,5)\}$

$$A \times B = \{(1,3), (1,4), (1,5), (2,3), (2,4), (2,5), (3,3), (3,4), (3,5)\}$$

$$R^C = \{(1,4), (2,3), (3,3), (3,4), (2,5)\}$$



## Combining Relation

- Given two relations  $R_1$  and  $R_2$ , these can be combined by using basic set operations to form new relations such as  $R_1 \cup R_2$ ,  $R_1 \cap R_2$ ,  $R_1 - R_2$ , and  $R_2 - R_1$ .

- $R_1 \cup R_2 = \{(a, b) \mid (a, b) \in R_1 \text{ or } (a, b) \in R_2 \text{ or both}\}$
- $R_1 \cap R_2 = \{(a, b) \mid (a, b) \in R_1 \text{ and } (a, b) \in R_2\}$
- $R_1 - R_2 = \{(a, b) \mid (a, b) \in R_1 \text{ and } (a, b) \notin R_2\}$
- $R_2 - R_1 = \{(a, b) \mid (a, b) \in R_2 \text{ and } (a, b) \notin R_1\}$

## Example

- Let  $A = \{1, 2, 3\}$  and  $B = \{3, 4, 5\}$

Suppose  $R = \{(1,3), (1,5), (2,4), (3,5)\}$ ,

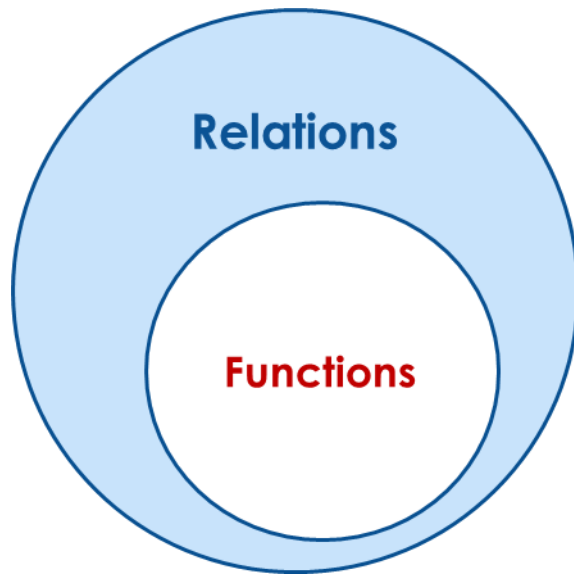
$R_1 = \{(1,4), (2,3), (2,5), (3,3), (3,5)\}$  and

$R_2 = \{(1,3), (1,4), (2,3), (3,4), (3,5)\}$

- $R_1 \cup R_2 = \{(1,3), (1,4), (2,3), (2,5), (3,3), (3,4), (3,5)\}$
- $R_1 \cap R_2 = \{(1,4), (2,3), (3,5)\}$
- $R_1 - R_2 = \{(2,5), (3,3)\}$
- $R_2 - R_1 = \{(1,3), (3,4)\}$

Thank  
you!!!





# Relations

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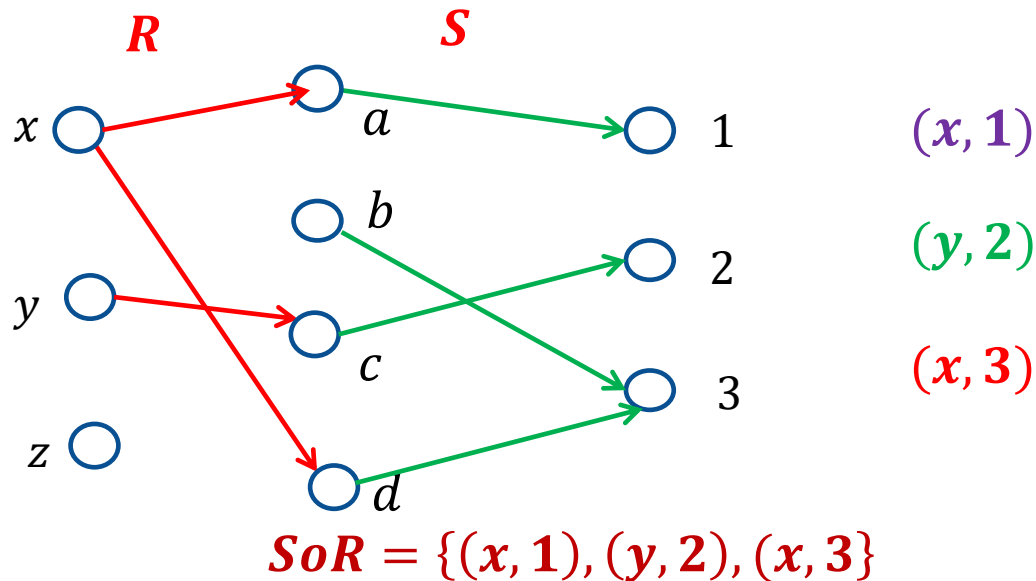
## Composition of Relations

- If  $R \subseteq A \times B$  and  $S \subseteq B \times C$  are two relations, then the composition (or composite) of  $S$  with  $R$  is a relation from  $A$  to  $C$  and is defined as:

$$SoR = \{(a, c) \mid \exists b \in B \text{ such that } (a, b) \in R \text{ and } (b, c) \in S\}$$

# Representing the Composition of Relations

- Let  $A = \{x, y, z\}$ ,  $B = \{a, b, c, d\}$  and  $C = \{1, 2, 3\}$ .
- Suppose  $R = \{(x, a), (x, d), (y, c)\}$  be a relation from  $A$  to  $B$ .
- Suppose  $S = \{(a, 1), (b, 3), (c, 2), (d, 3)\}$  be a relation from  $B$  to  $C$ .





## Power of Relations

- If  $R \subseteq A \times A$ , then

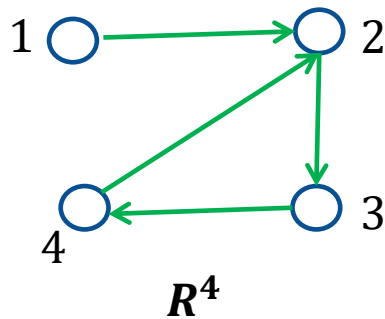
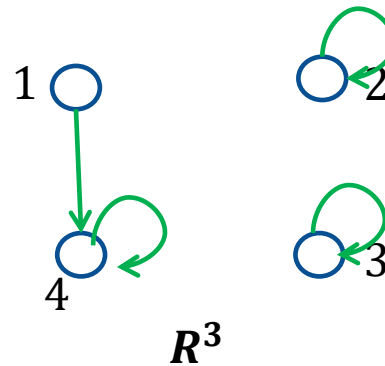
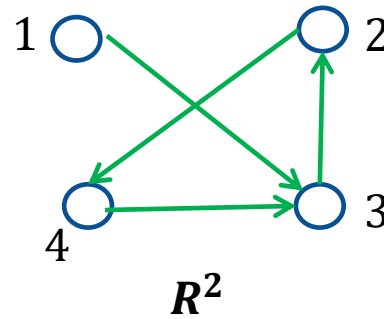
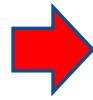
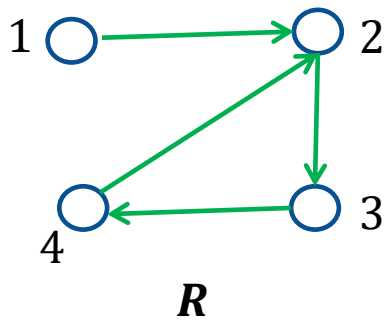
$$R^2 = R \circ R$$

$$R^3 = R^2 \circ R$$

$$\vdots$$

$$R^n = R^{n-1} \circ R$$

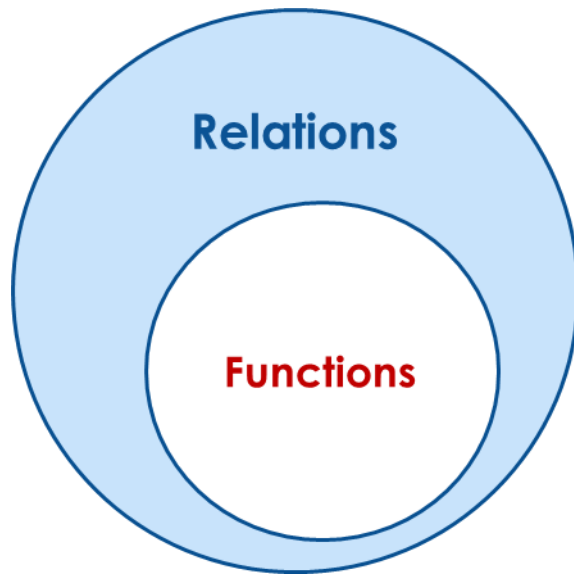
# Example



The pair  $(a, b)$  is in  $R^n$  if there is a path of length  $n$  from  $a$  to  $b$  in  $R$ .

Thank  
you!!!





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# Equivalence Relation

- Let  $R$  be a relation on set  $A$ , then  $R$  is called equivalence relation if it is:
  - Reflexive
  - Symmetric
  - Transitive

# Examples

• Let  $A = \{1, 2, 3\}$

1.  $\emptyset$  i.e. Empty Relation on  $A$

Reflexive?

Symmetric?

Transitive?

**Not an Equivalence Relation**

2.  $\Delta = \{(1, 1), (2, 2), (3, 3)\}$

Reflexive?

Symmetric?

Transitive?

**Equivalence Relation on  $A$**

**Smallest Equivalence Relation on  $A$**



## Examples (Cont..)

- Let  $A = \{1, 2, 3\}$

### 3. Universal Relation on $A$ i.e. $A \times A$

Reflexive?

Symmetric?

Transitive?

**Equivalence Relation**

**Largest Equivalence Relation on  $A$**

- 4. Let  $A = \{1, 2, 3, 4\}$

$$R = \{(1, 1), (1, 2), (2, 1), (2, 2), (3, 3), (3, 4), (4, 3), (4, 4)\}$$

Reflexive?

Symmetric?

Transitive?

**Equivalence Relation on  $A$**

If  $R_1$  and  $R_2$  are two equivalence relations on  $A$ , then which of the following is always true?

- I.*  $R_1 \cap R_2$  is an Equivalence Relation.
- II.*  $R_1 \cup R_2$  is an Equivalence Relation.

- (a) Only *I*
- (b) Only *II*
- (c) Both are true
- (d) Both are false

## Exercise

- Let  $R$  be a relation defined on *set of integers* as:

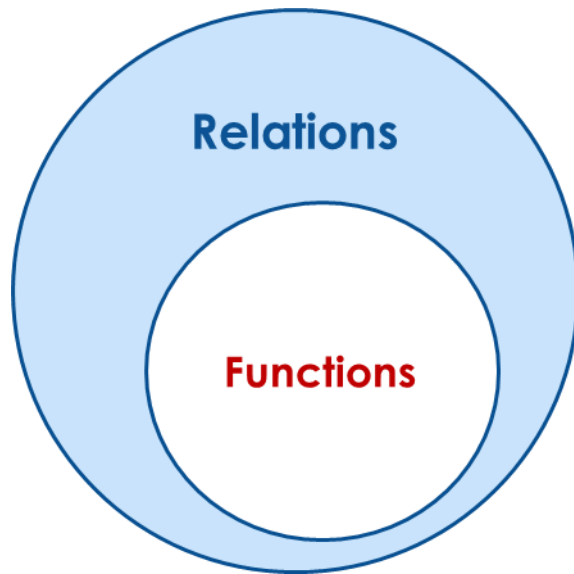
$$xRy \text{ iff } x + y \text{ is even}$$

Is  $R$  an equivalence relation?

Thank  
you!!!







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## Equivalence Class

- Let  $R$  be an equivalence relation on  $A$ , and  $a \in A$ .
- The equivalence class of  $a$ , denoted as  $[a]$  or  $\bar{a}$ , is defined as:

$$\bar{a} = [a] = \{b \in A \mid (a, b) \in R\}$$



## Examples

- Let  $R = \{(1,1), (1,2), (2,1), (2,2), (3,3), (3,4), (4,3), (4,4)\}$  on  $A = \{1, 2, 3, 4\}$

First Check whether  $R$  is an equivalence relation on  $A$  or not.

Reflexive?

Symmetric?

Transitive?

**Equivalence Classes:-**

$$[1] = \{1, 2\}$$

$$[2] = \{1, 2\}$$

$$[3] = \{3, 4\}$$

$$[4] = \{3, 4\}$$

## Examples

- Let  $R = \{(1,1), (2,2), (3,3), (4,4)\}$  on  $A = \{1, 2, 3, 4\}$

**Equivalence Classes:-**

$$[1] = \{1\}$$

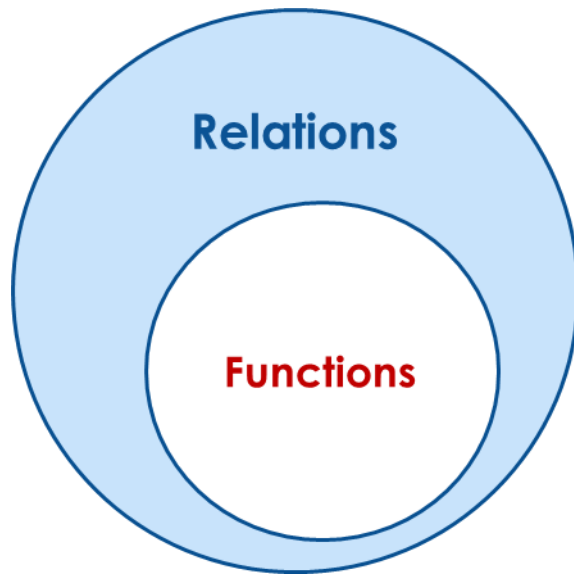
$$[2] = \{2\}$$

$$[3] = \{3\}$$

$$[4] = \{4\}$$

Thank  
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## Equivalence Relation to Partition

- Let  $R = \{(1,1), (1,2), (2,1), (2,2), (3,3), (3,4), (4,3), (4,4)\}$  on  $A = \{1, 2, 3, 4\}$ .
- $R$  is an equivalence relation on  $A$ .

### Equivalence Classes:-

$$[1] = \{1, 2\} = [2]$$

$$[3] = \{3, 4\} = [4]$$

$$\square \text{ Partition } P = \{\{1, 2\}, \{3, 4\}\}$$

## Partition to Equivalence Relation

- Let  $A = \{1, 2, 3, 4\}$  be a set and  $P = \{\{1, 3\}, \{2, 4\}\}$  be a partition on  $A$ .
- Find Equivalence relation on  $A$ .

$$\{1, 3\} \rightarrow \{(1, 1), (1, 3), (3, 1), (3, 3)\}$$

$$\{2, 4\} \rightarrow \{(2, 2), (2, 4), (4, 2), (4, 4)\}$$

Therefore, the equivalence relation on  $A$  is:

$$\square R = \{(1, 1), (1, 3), (3, 1), (3, 3), (2, 2), (2, 4), (4, 2), (4, 4)\}$$



## Result

- There is a one-to-one correspondence between partitions of  $A$  and Equivalence Relation on  $A$ .
- Therefore, if  $|A| = n$ , then

**Number of Partitions of  $A$  = Number of Equivalence Relations on  $A$   
=  $B_n$  (Bell Number)**

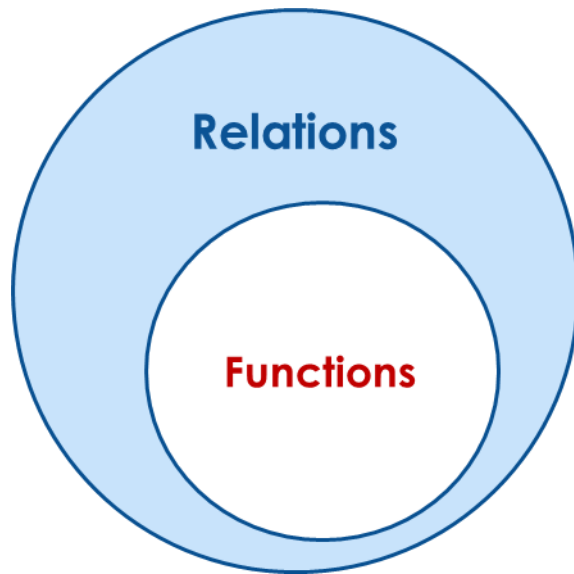
**Bell Number:**

$$B_n = \sum_{k=0}^{n-1} n - 1 C_k B_k$$

**where,  $B_0 = 1$**

Thank  
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# Closure of Relations

- Reflexive Closure
- Symmetric Closure
- Transitive Closure

# Reflexive Closure

- A relation is called **reflexive closure**  $R_r$  of relation  $R$  if:
  - 1) It is reflexive.
  - 2) It contains  $R$ .
  - 3) It is the minimal relation satisfying conditions (1) and (2).

## Examples

1. Let  $A = \{1, 2, 3\}$

Suppose  $R_1 = \{(1,1), (2,2), (2,3)\}$  be a relation on  $A$ .

$$R_r = \{(1,1), (2,2), (2,3), (3,3)\}$$



## Examples (Cont..)

2.  $R$  is a relation defined on set of positive integers such that  $aRb$  if  $a < b$ .

Reflexive Closure?

Result:

- $R_r = R \cup \Delta$
- $R_r = R$  iff  $R$  is Reflexive.

# Symmetric Closure

- A relation is called **symmetric closure**  $R_s$  of relation  $R$  if:
  - 1) It is symmetric.
  - 2) It contains  $R$ .
  - 3) It is the minimal relation satisfying conditions (1) and (2).

## Examples

- Let  $A = \{1, 2, 3\}$

Suppose  $R_1 = \{(1,1), (2,2), (2,3)\}$  be a relation on  $A$ .

$$R_s = \{(1,1), (2,2), (2,3), (3,2)\}$$

## Result

- $R_s = R \cup R^{-1}$
- $R_r = R$  iff  $R$  is Symmetric.

# Transitive Closure

- A relation is called Transitive closure  $R^*$  of relation  $R$  if:
  - 1) It is transitive.
  - 2) It contains  $R$ .
  - 3) It is the minimal relation satisfying conditions (1) and (2).

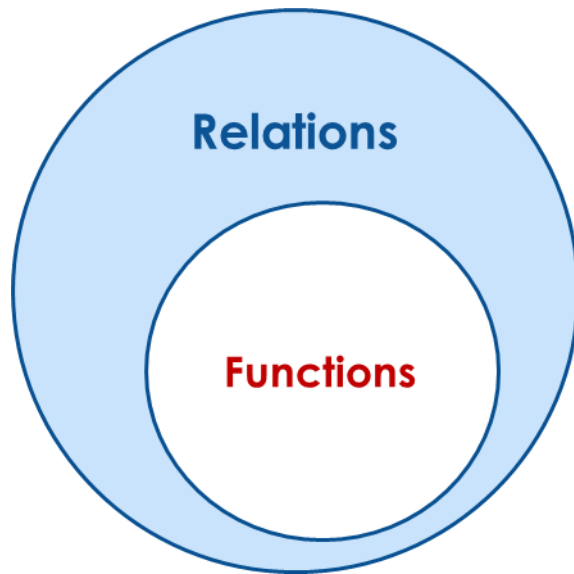
## Result

1. Let  $|A| = n$ ,  
then,  $R^* = R^1 \cup R^2 \cup \dots \cup R^n$
2.  $R$  is transitive iff  $R^* = R$



Thank  
you!!!





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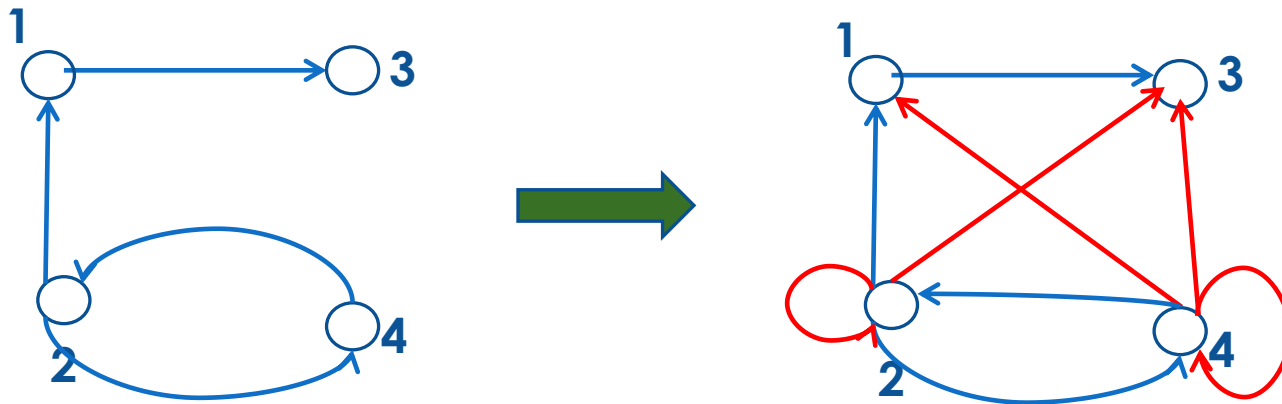
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# Warshall's Algorithm

- Computes the transitive closure of a relation

Example of transitive closure:



## Warshall's Algorithm (Cont..)

- Main concept: a path exists between two vertices  $i, j$ , iff
  - there is an edge from  $i$  to  $j$ ; or
  - there is a path from  $i$  to  $j$  going through vertex 1; or
  - there is a path from  $i$  to  $j$  going through vertex 1 and/or 2; or
  - there is a path from  $i$  to  $j$  going through vertex 1, 2, and/or 3; or
  - ...
  - there is a path from  $i$  to  $j$  going through any of the other vertices

## Warshall's Algorithm (Cont..)

- On the  $k^{th}$  iteration, the algorithm determine if a path exists between two vertices  $i, j$  using vertices among  $1, \dots, k$  allowed as intermediate

$$W^{(k)}[i, j] = \begin{cases} W^{(k-1)}[i, j] \\ \text{or} \\ (W^{(k-1)}[i, k]) \text{ and } (W^{(k-1)}[k, j]) \end{cases}$$

## Warshall's Algorithm (Cont..)

- Recurrence relating elements  $W^{(k)}$  to elements of  $W^{(k-1)}$  is:

$$W^{(k)}[i,j] = W^{(k-1)}[i,j] \text{ or } (W^{(k-1)}[i,k] \text{ and } W^{(k-1)}[k,j])$$

- It implies the following rules for generating  $W^{(k)}$  from  $W^{(k-1)}$  is:

1. If an element in row  $i$  and column  $j$  is 1 in  $W^{(k-1)}$ , it remains 1 in  $W^{(k)}$ .
2. If an element in row  $i$  and column  $j$  is 0 in  $W^{(k-1)}$ , it has to be changed to 1 in  $W^{(k)}$  if and only if the element in its row  $i$  and column  $k$  and the element in its row  $k$  and column  $j$  are both 1's in  $W^{(k-1)}$ .

## Warshall's Algorithm (Cont..)

- The procedure for computing  $W^{(k)}$  from  $W^{(k-1)}$  is as follows:
  1. First transfer all 1's in  $W^{(k-1)}$  to  $W^{(k)}$ .
  2. List the locations  $p_1, p_2, \dots$ , in column  $k$  of  $W^{(k-1)}$ , where the entry is 1, and the locations  $q_1, q_2, \dots$ , in row  $k$  of  $W^{(k-1)}$ , where the entry is 1.
  3. Put 1's in all the positions  $(p_i, q_i)$  of  $W^{(k)}$  (if they are not already there).



# Warshall's Algorithm (Cont..)

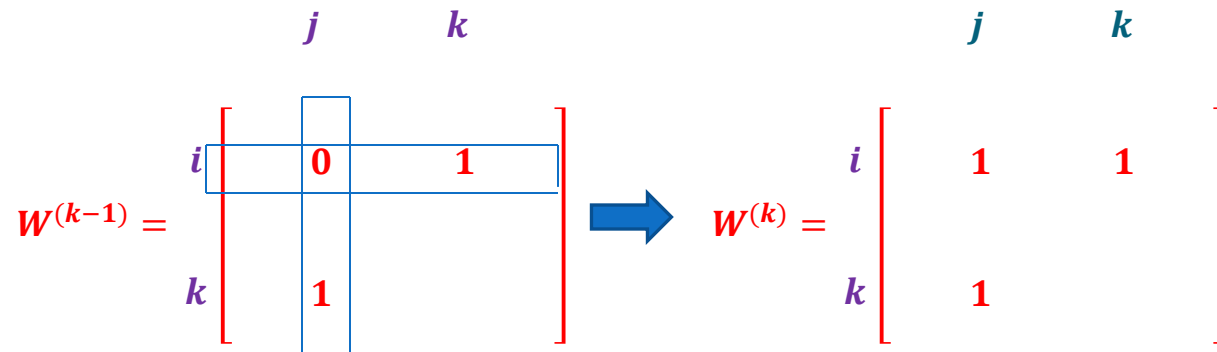


Figure 1: Step for Changing zeros in Warshall's Algorithm

## Example

1. Find transitive closure of relation  $R$  represented by following matrix (using Warshall's algorithm):

$$M_R = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

**Solution 1:**

$$W^{(0)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

$$W^{(1)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

$$W^{(2)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

$$W^{(3)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

$$W^{(4)} = W^{(3)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

Answer is :

$$M_{R^*} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

$$R^* = \{(1, 1), (1, 2), (1, 3), (1, 4), (2, 1), (2, 2), (2, 3), (2, 4), (3, 4)\}$$



Thank  
you!!!

