

*Course: Computer and Communication Networks*

*Topic: An Introduction to Queues and Queueing Theory*

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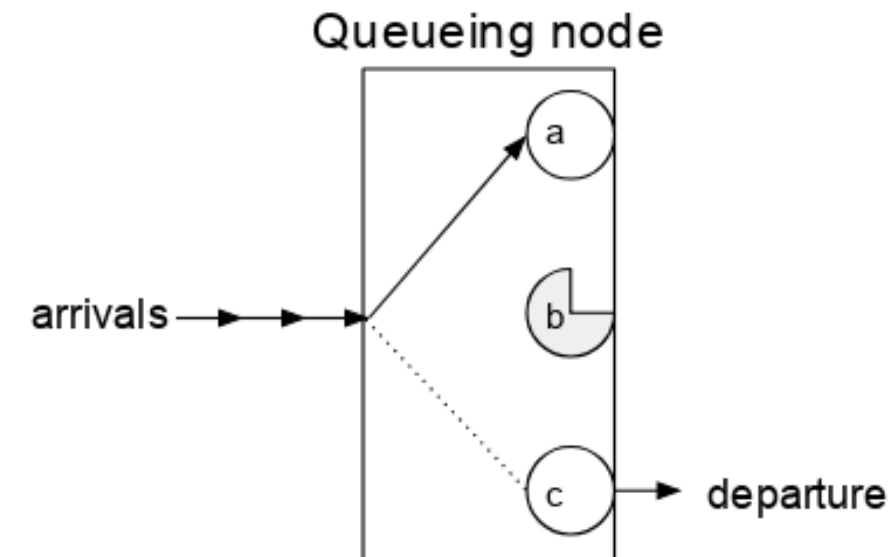
# An Introduction to Queues and Queueing Theory

## Queue Model

1. **First Come First Served (FCFS)**
2. **Last Come First Served (LCFS)**
3. **Service in Random Order (SIRO)**
4. **Priority Service (PS)**

Queueing theory is the mathematical study of waiting lines, or queues.

A queueing model is constructed so that queue lengths and waiting time can be predicted.



# An Introduction to Queues and Queueing Theory

## 1. First Come First Served

Process	Duration (sec)	Order	Arrival Time (sec)
P1	24	1	0
P2	3	2	0
P3	4	3	0

P1 waiting time is 0 second

P2 waiting time is 24second

P3 waiting time is 27second

Average waiting time is

=  $(0+24+27)/3=17$  seconds

E.g. checkout counter at super market

**Completion Time:** Time at which process completes its execution.

**Turn Around Time:** Time Difference between completion time and arrival time. Turn Around Time = Completion Time – Arrival Time

**Waiting Time(W.T):** Time Difference between turn around time and burst time. Waiting Time = Turn Around Time – Burst Time

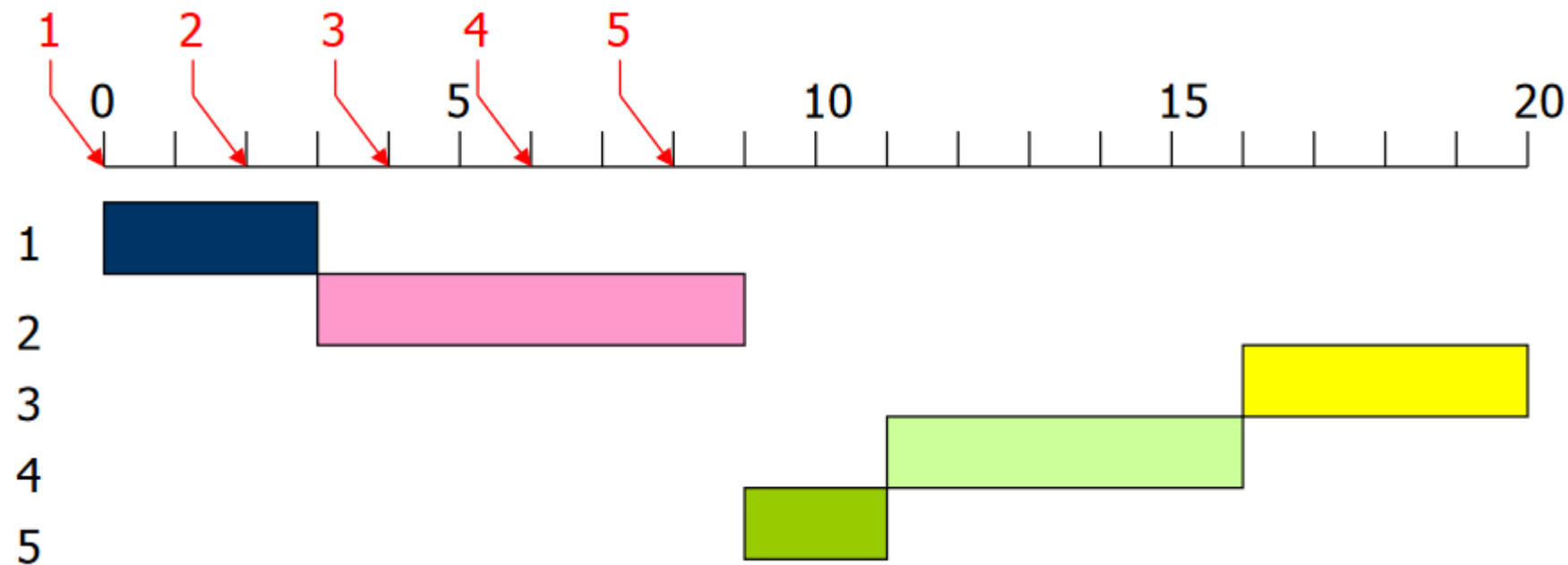
*If arrival times as 0, turn around and completion times are same.*



# An Introduction to Queues and Queueing Theory

## 2. Last Come First Served

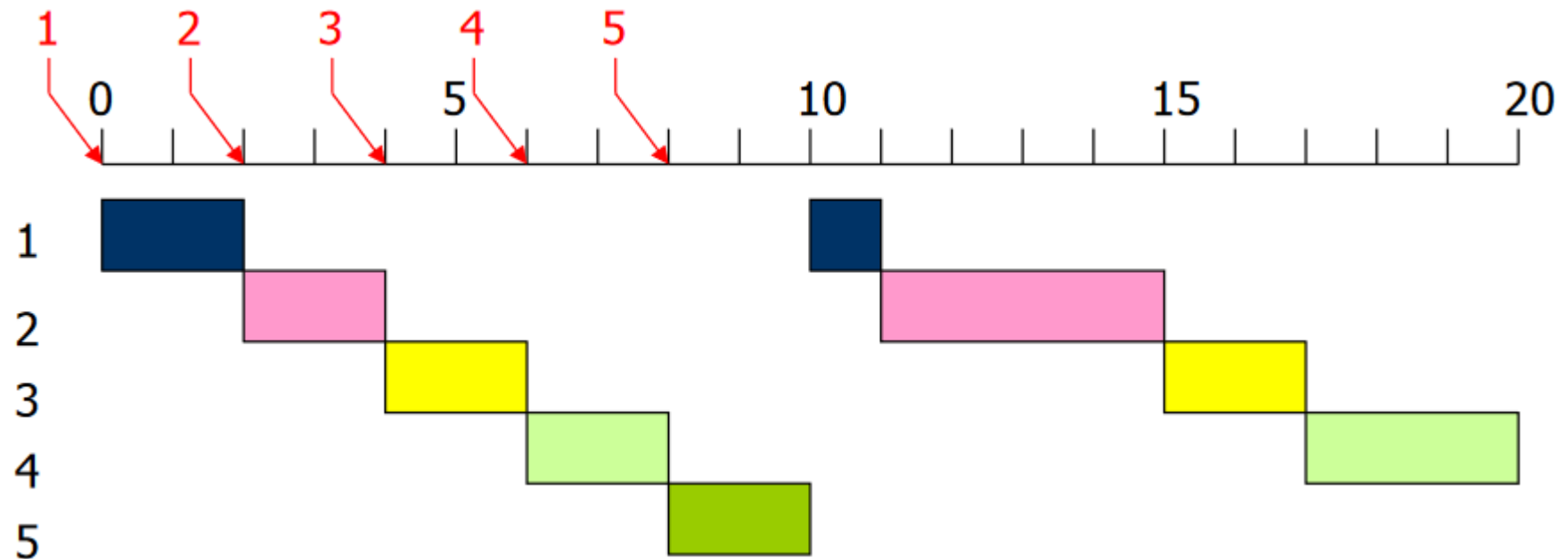
- Execution of threads in reversed order of arrival at the ready list.
- Occupation of processor until end or voluntary yield.
- Remark: Rarely used in the pure form



# An Introduction to Queues and Queueing Theory

## Last Come First Served - Preemptive Resume

- Newly arriving thread at ready list preempts the currently running thread.
- Preempted thread is appended to ready list.
- In case of no further arrivals, the ready list is processed without preemption.
- Goal: Preference to short threads.
- A short thread has a good chance to finish before another thread arrives.
- A long thread is likely to be preempted several times.

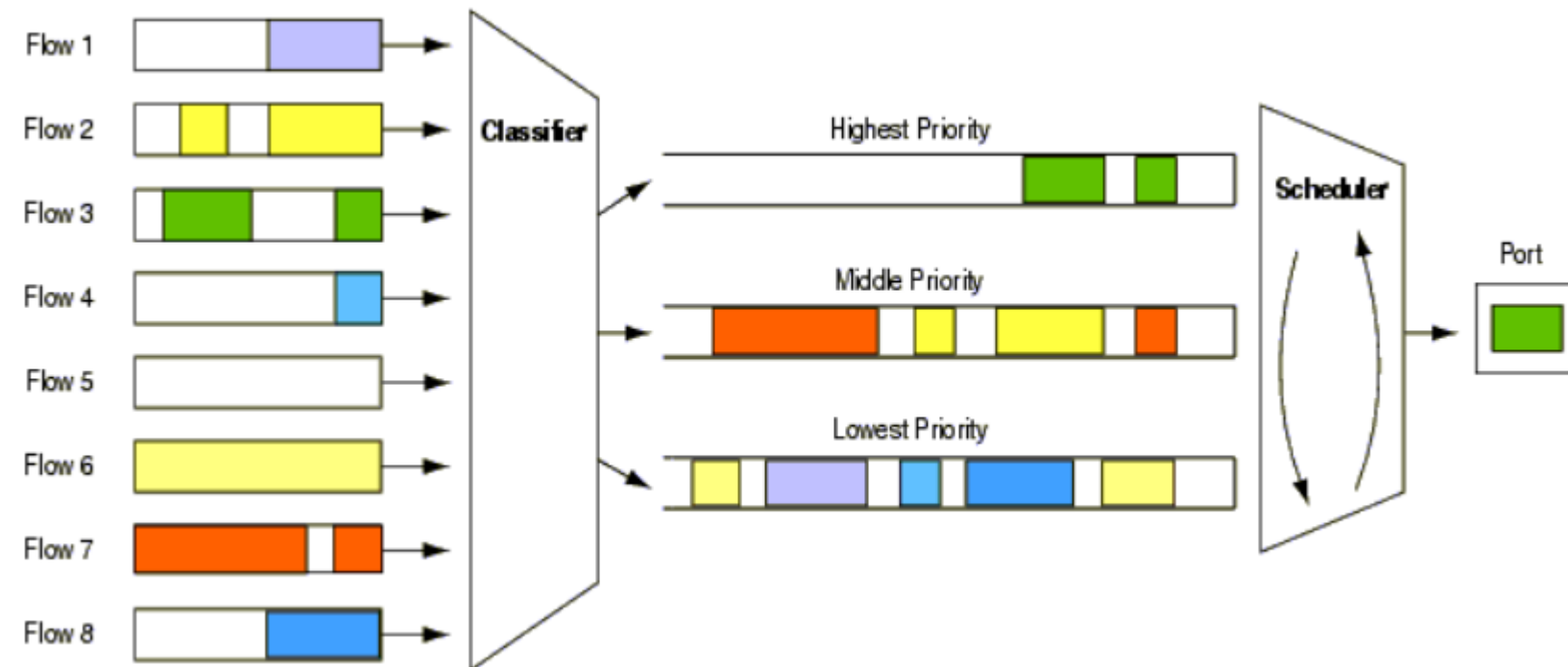


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## 3. Service in Random Order (SIRO)

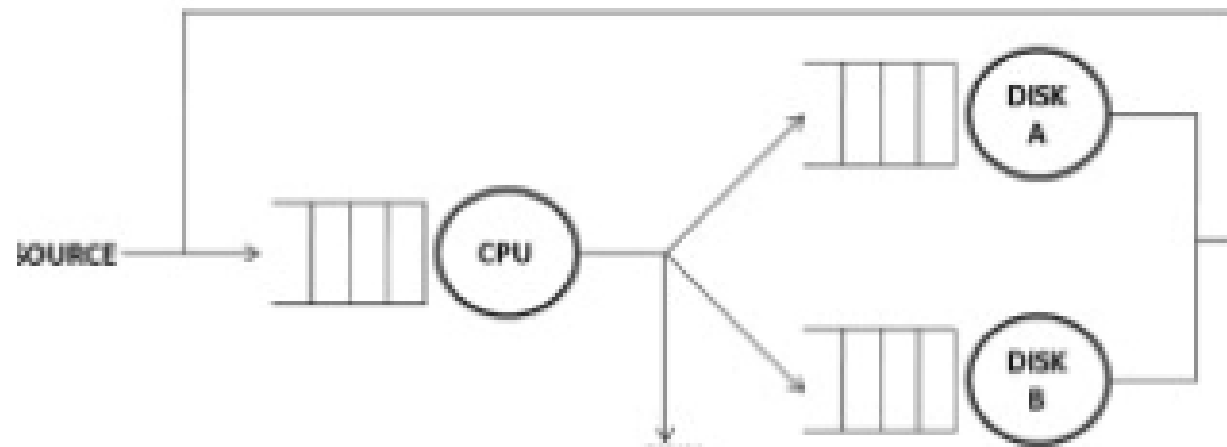
If the next customer to enter service is randomly chosen from those customers waiting for the service it is called SIRO

## 4. Priority Service:

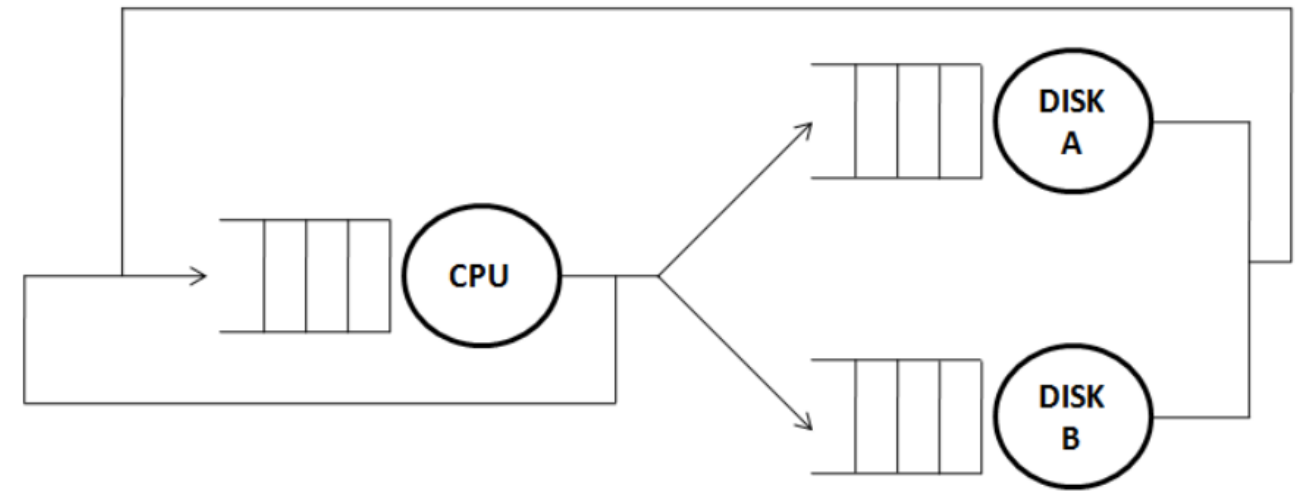


# An Introduction to Queues and Queueing Theory

An Open Queueing Network



A Closed Queueing Network



# Little's Theorem

Little's Theorem states:

$$N = \lambda T$$

Where  $N$  is the average number of customers in a queue,  $T$  is the average time a customer spends queuing and  $\lambda$  is the average rate of arrivals to the queue.

If a lot of people are in a queue ( $N$  is large) then they will have long delays ( $T$  is large);

if few people arrive in a queue ( $\lambda$  is small) then the average number of people in the queue is small ( $N$  is small).

Let us first make precise the definitions of  $N$ ,  $\lambda$  and  $T$  and then make clear the assumptions on which the theorem rests.



$N(\tau)$  is the number of customers in the system at time  $\tau$ .

$\alpha(\tau)$  is the number of customers who arrived in the interval  $[0,\tau]$ .

$\beta(\tau)$  is the number of customers who have departed in the interval  $[0,\tau]$ .

$t_i$  is the time at which the  $i$ th customer arrived.

$T(i)$  is the time spent queuing by the  $i$ th customer.

If  $N_t$  is the mean value of  $N(\tau)$  taken over the interval  $[0,t]$  then it is clear that:

$$N_t = \frac{1}{t} \int_0^t N(\tau) d\tau$$

$$N = \lim_{t \rightarrow \infty} N_t$$

(Note that this limit is not guaranteed to exist — imagine, for example, a queue which keeps growing.) If the limit exists,  $N$  is the steady state time average of  $N(\tau)$ .

We can next define the average arrival rate over the time period  $[0,t]$ .

$$\lambda_t = \frac{\alpha(t)}{t}$$

and, again, we assume that the following limit exists:

$$\lambda = \lim_{t \rightarrow \infty} \lambda_t$$

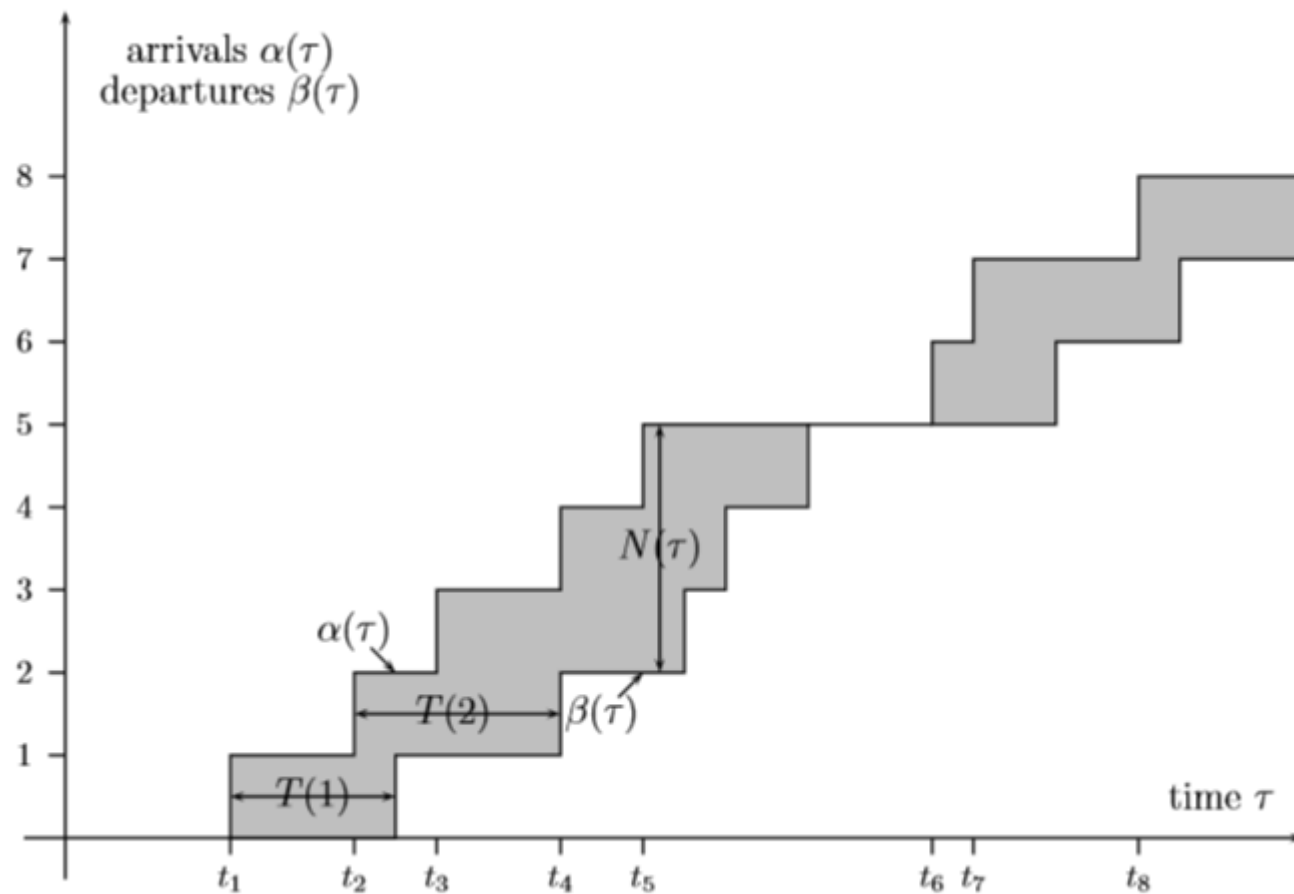
Finally, the average delay experienced by those customers who enter the system at times in  $[0,t]$  is given by:

$$T_t = \sum_{i=1}^{\alpha(t)} \frac{T(i)}{\alpha(t)}$$

And, for a third time, we assume that the following limit exists:

$$T = \lim_{t \rightarrow \infty} T_t$$

# Little's Theorem Proof assuming FIFO



$$N(\tau) = \alpha(\tau) - \beta(\tau)$$

It is clear that if we choose a time  $t$  when the system again becomes empty then we can calculate the area of the shaded area  $A(t)$ :

$$A(t) = \int_0^t N(\tau) d\tau$$

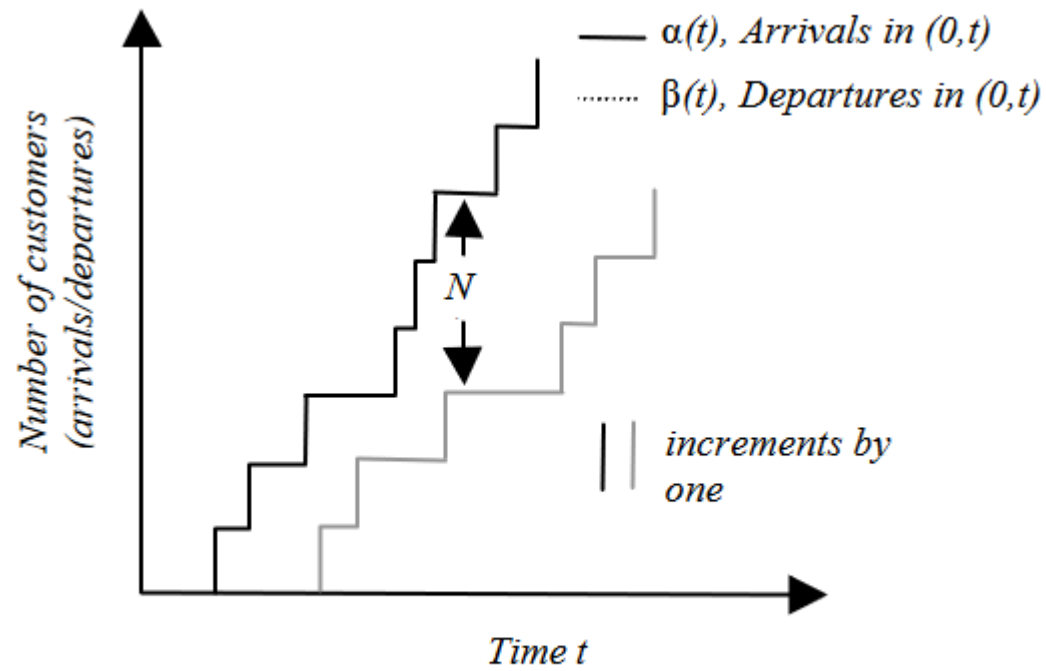
However, equally, we can consider the shaded area to be composed of horizontal strips of height 1 and width  $T(i)$  (for the  $i$ th customer). In this case, we have:

$$A(t) = \sum_{i=1}^{\alpha(t)} T(i)$$

$$\frac{1}{t} \int_0^t N(\tau) d\tau = \frac{1}{t} \sum_{i=1}^{\alpha(t)} T(i) = \frac{\alpha(t)}{t} \frac{\sum_{i=1}^{\alpha(t)} T(i)}{\alpha(t)}$$

$$N_t = \lambda_t T_t$$

# Graphical verification of Little's Theorem



Graphical Illustration/Verification of Little's Result

Consider the time interval  $(0,t)$  where  $t$  is large, i.e.  $t \rightarrow \infty$

$$Area(t) = \text{area between } \alpha(t) \text{ and } \beta(t) \text{ at time } t = \int_0^t [\alpha(t) - \beta(t)] dt$$

$$\text{Average Time } W \text{ spent in system} = \lim_{t \rightarrow \infty} \frac{Area(t)}{\alpha(t)}$$

$$\text{Average Number } N \text{ in system} = \lim_{t \rightarrow \infty} \frac{Area(t)}{t} = \lim_{t \rightarrow \infty} \frac{\alpha(t)}{t} \frac{Area(t)}{\alpha(t)}$$

$$\text{Since, } \lambda = \lim_{t \rightarrow \infty} \frac{\alpha(t)}{t}$$

$$\text{Therefore, } N = \lambda W$$