

Relations

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Contents

- Relations and Its Introduction
- Representation of Relations:
 - Using Matrices
 - Using Diagram
- Properties of Relations
- Inverse and Complementry Relations
- Combining Relations and Composite Relations
- Equivalence Relations
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- Closure of Relations
- Warshall's Algorithm
- Partial Ordering and Partially Ordered Set
- Lexicographic Ordering
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What is a Relation?

In discrete mathematics, relation is a way of showing a relationship between any two sets.

- Relationship between any program and its variable.
- Relationship between pair of cities linked by railway in a network.

Necessity for studying Relation

- Relational Database model is based on the concept of relation.

Cartesian Product

- Given two sets A and B, their **cartesian product** $A \times B$, is defined as

$$A \times B = \{(a, b) \mid a \in A, b \in B\}$$

Ordered Pairs

- The elements of $A \times B$ are called **ordered pairs** with the elements of A as the first entry and elements of B as the second entry.
- Order matters

Special Case:

$$A^2 = A \times A = \{(a_1, a_2) \mid a_1, a_2 \in A\}$$

Similarly,

$$A^n = A \times A \times \cdots \times A(n \text{ times}) = \{(a_1, a_2, \dots, a_n) \mid a_1, a_2, \dots, a_n \in A\}$$

Relation is the subset of the cartesian product of the sets.

n-ary Relation

- Let $\{A_1, A_2, \dots, A_n\}$ be n sets.
- An *n-ary relation* R on $A_1 \times A_2 \times \dots \times A_n$ is a subset of $A_1 \times A_2 \times \dots \times A_n$.
- If $A_i = A ; \forall i$, then R is called the *n-ary relation on A*.

Empty and Universal Relation

- If $R = \emptyset$, then R is called the **empty** or **void relation**.
- If $R = A_1 \times A_2 \times \dots \times A_n$, then R is called the **universal relation**.

Definition (Binary Relation)

- Given two sets A and B , a relation between A and B is a subset of $A \times B$.
- If R is a relation on $A \times B$ (i.e., $R \subseteq A \times B$) and $(a, b) \in R$, we say “ a is related to b ”.
- It can also be written as aRb .

Example:

Let $A = \{a, b\}$ and $B = \{2, 3, 4\}$

$R = \{(a, 3), (b, 2), (b, 4)\}$ is a relation from A to B .

Binary Relation on a set

- A binary relation R on a set A is a **subset of $A \times A$** .

Examples:

1. “Taller -than ” is a relation on people.
 $(a, b) \in$ “Taller -than” if person a is taller than person b.
2. “ \geq ” is a relation on real set \mathbf{R} .
$$\geq = \{(x, y) \in \mathbf{R} \mid x, y \in \mathbf{R}, x \geq y\}$$

Examples (Cont..)

3. Let $A = \{1, 2, 3, 4, 5, 6\}$.

If $R = \{(a, b) | a \text{ divides } b\}$ is a relation from A to B then ordered pairs in the relation R are

$$\{(1,1), (1,2), (1,3), (1,4), (1,5), (1,6), (2,4), (2,6), (3,6)\}$$

Examples (Cont..)

Let $A=\{1, 2, 3\}$

$$\begin{aligned}A \times A \\= \{(1,1), (1,2), (1,3), (2,1), (2,2), (2,3), (3,1), (3,2), (3,3)\}\end{aligned}$$

- Here, $A \times A$ is an universal relation on A.
- \emptyset is an empty relation on A.

Examples (Cont..)

$\text{"="} = \Delta = \{(1,1), (2,2), (3,3)\}$

$\text{"<} = \{(1,2), (1,3), (2,3)\}$

$\text{">} = \{(2,1), (3,1), (3,2)\}$

$\text{"\leq"} = \{(1,1), (1,2), (1,3), (2,2), (2,3), (3,3)\}$

$\text{"\geq"} = \{(1,1), (2,1), (2,2), (3,1), (3,2), (3,3)\}$

$\text{"|"} = \{(1,1), (1,2), (1,3), (2,2), (3,3)\}$

$\text{"multiple of"} = \{(1,1), (2,1), (2,2), (3,1), (3,3)\}$

Representing Relations

Relations can be represented in two ways:

Matrix

Graph

Representation of Relations as Matrix

- If R is a relation on set $A = \{a_1, a_2, \dots, a_n\}$ and $|A| = n$, then it can be represented as $n \times n$ Boolean Matrix M_R .

M_R can be defined as:

$$M_R = [m_{ij}]_{n \times n}$$

where, $m_{ij} = \begin{cases} 0 & ; \text{if } (a_i, a_j) \notin R \\ 1 & ; \text{if } (a_i, a_j) \in R \end{cases}$

Examples

- Let $A = \{1, 2, 3\}$
- Let $R = \{(1,1), (1,2), (2,1), (2,3), (3,2), (3,3)\}$ be a relation on A.

$$M_R = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

Examples (Cont..)

- Let $A = \{1, 2, 3\}$ and $B = \{1, 2, 3, 4, 5\}$. Which ordered pairs are in the relation R represented by the matrix

$$M_R = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 \end{bmatrix} ?$$

- Because R consists of those ordered pairs with $a_{ij} = 1$, it follows that:

$$R = \{(1, 2), (2, 1), (2, 3), (2, 4), (2, 5), (3, 1), (3, 3), (3, 5)\}.$$

Representation of Relations as a Digraph (Directed Graph)

- The graph of a relation R over A is a directed graph with nodes corresponding to the elements of A . There is an edge from node x to y if and only if $(x, y) \in R$.
- An edge of the form (x, x) is called a self-loop.

Examples

- Let $A = \{1, 2, 3\}$
- Let $R_1 = \{(1, 2), (1, 3), (2, 3)\}$ be a $<$ relation on A .

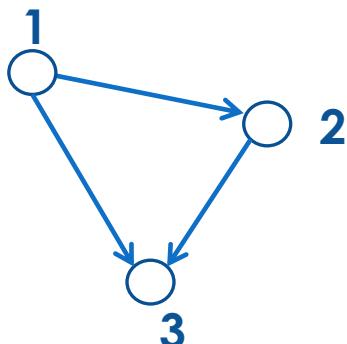


Figure 1

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Examples (Cont..)

- Let $A = \{1,2,3\}$ and $B = \{a,b\}$
- Let $R_2 = \{(1, a), (1, b), (2, a), (3, b)\}$ be a relation from A to B .

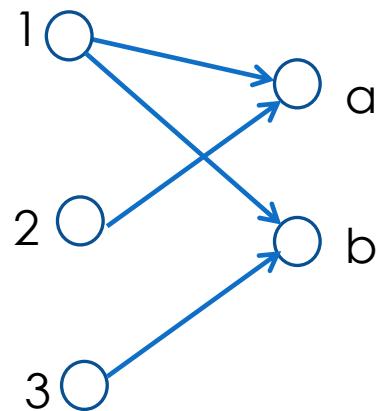


Figure 2

Domain and Range

Domain of Relation R= set of all first co-ordinates

Range of Relation R= set of all second co-ordinates

Example

“<”= $\{(1,2), (1,3), (2,3)\}$ on $A = \{1,2,3\}$

- Domain of “<”= $\{1,2\}$
- Range of “<”= $\{2,3\}$

Equality of Two relations

- Let R_1 be an n-ary relation on $A_1 \times A_2 \times \dots \times A_n$.
- Let R_2 be an m-ary relation on $B_1 \times B_2 \times \dots \times B_m$.
- Then, $R_1 = R_2$
If and only if
 - ❖ **n=m**
 - ❖ **$A_i = B_i; \forall i, 1 \leq i \leq n$**
 - ❖ **and, R_1 & R_2 are equal set of ordered pairs.**

Example

- Let $A = \{a, b\}, B = \{1, 2\}, C = \{1, 2, 3\}$

- Let $R_1 = \{(a, 1), (b, 2)\}$ is a relation on $A \times B$
- Let $R_2 = \{(a, 1), (b, 2)\}$ is a relation on $A \times C$

$R_1 = R_2?$

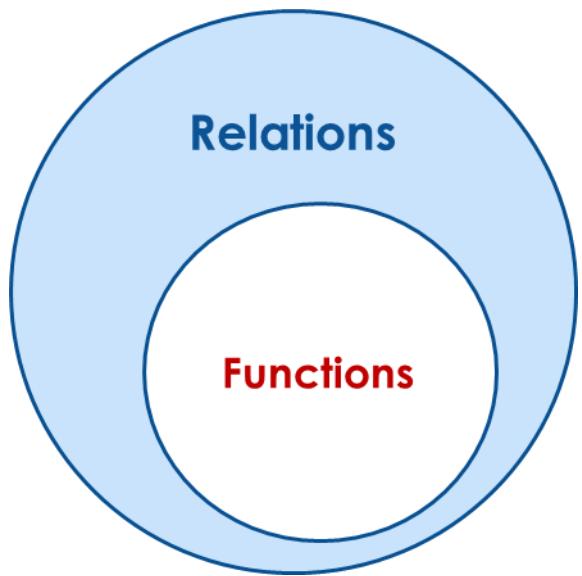
No

How many number of relations
are there on a set A having n
elements?

$$2^{n^2}$$

Thank
you!!!





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Properties of Relations

- Reflexive
- Symmetric
- Transitive
- Irreflexive
- Asymmetric
- Antisymmetric

Reflexive Relations

- R is **reflexive** iff $(x, x) \in R$ for every element $x \in A$.

Examples

1. Let $A = \{1, 2, 3\}$

Suppose $R_1 = \{(1,1), (2,2), (2,3)\}$ be a relation on A .

Is R_1 reflexive?

No

2. $=, A \times A, \leq, \geq, |, \text{multiple of}$ Reflexive? Yes

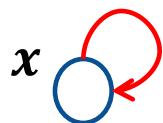
3. $\emptyset, <, >$ Reflexive? No

Reflexive Relation in Matrix and Graph

- If R is a **reflexive** relation, all the elements on the main diagonal of M_R are equal to 1.

$$M_R = \begin{bmatrix} 1 & \cdots & & \\ \vdots & \ddots & \vdots & \\ & \cdots & 1 & \end{bmatrix}$$

- A loop must be present at all vertices in the graph.



Symmetric Relations

- R is **symmetric** iff $(y, x) \in R$ whenever $(x, y) \in R$ for all $x, y \in A$.

Examples

1. Let $A = \{1, 2, 3\}$

Suppose $R_2 = \{(1,2), (2,1), (2,3)\}$ be a relation on A .

Is R_1 Symmetric?

No

2. "sibling-of" is **symmetric**, but "sister-of" is **not**.

3. $A \times A, \emptyset, =$ Symmetric? Yes

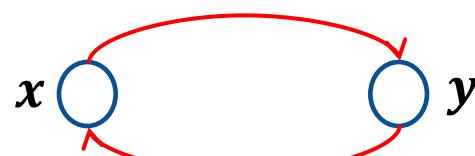
4. $<, >, \leq, \geq, |$, *multiple of* Symmetric? No

Symmetric Relation in Matrix and Graph

- R is a symmetric relation if and only if $m_{ji} = 1$, whenever $m_{ij} = 1$.

$$M_R = \begin{bmatrix} & & 1 \\ 1 & & \\ & 0 & \end{bmatrix}$$

- If (x, y) is an edge in the graph, then there must be an edge (y, x) also.



Transitive Relations

- A relation R on a set A is called **transitive** if whenever $(x, y) \in R$ and $(y, z) \in R$, then $(x, z) \in R$, for all $x, y, z \in A$.

Examples

1. Let $A=\{1, 2, 3\}$

Suppose $R_3 = \{(1,3), (3,1)\}$ be a relation on A .

Is R_3 Transitive?

No

2. $A \times A$, $\emptyset, =, <, >, \leq, \geq, |$, multiple of

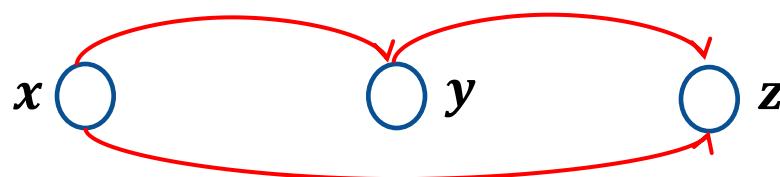
Transitive?

Yes

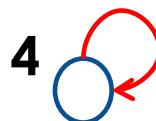
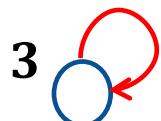
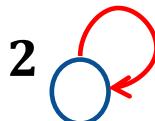
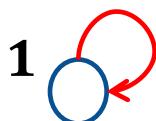
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Transitive Relations in Graph

- R is transitive iff in its graph, for any three nodes x, y and z such that there is an edge (x, y) and (y, z) , there exists an edge (x, z) .



Examples

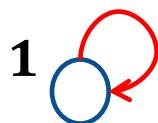


- Reflexive?
Yes
- Symmetric?
Yes
- Transitive?
Yes

Equality Relation on
 $A = \{1, 2, 3, 4\}$

$$M_R = \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1 \end{bmatrix}$$

Examples (Cont..)



- Reflexive?
No
- Symmetric?
Yes
- Transitive?
Yes

Examples (Cont..)

- Suppose that the relation R on a set is represented by the matrix

$$M_R = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

- Reflexive?
Yes
- Symmetric?
Yes

How many number of **Reflexive Relations** are there on set A having n elements?

$$2^{n(n-1)}$$

How many number of Symmetric Relations are there on set A having n elements?

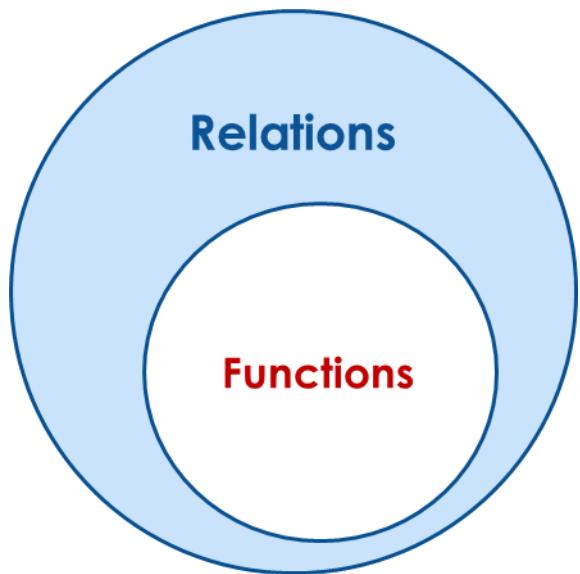
$$2^{n(n+1)/2}$$

How many number of Transitive Relations are there on set A having n elements?

No closed form found

Thank
you!!!





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Properties of Relations

- **Reflexive**
- **Symmetric**
- **Transitive**
- Irreflexive
- Asymmetric
- Antisymmetric

Irreflexive Relations

- R is **irreflexive** iff $(x, x) \notin R$ for every element $x \in A$.
- No Reflexive ordered pair should belong to the relation.

Examples

1. Let $A=\{1, 2, 3\}$

Suppose $R_1 = \{(1,1), (2,2), (2,3)\}$ be a relation on A .

Is R_1 Irreflexive?

No

2. $\emptyset, <, >$

Irreflexive? Yes

3. $\Delta, A \times A, \leq, \geq, |, \text{ multiple of }$

Irreflexive? No

Irreflexive Relation in Matrix and Graph

- If R is an irreflexive relation, all the elements on the main diagonal of M_R are equal to 0.

$$M_R = \begin{bmatrix} 0 & \cdots & \\ \vdots & \ddots & \vdots \\ & \cdots & 0 \end{bmatrix}$$

- No vertex should contain self-loop in the graph.

x 

Asymmetric Relations

- A relation R on a set A such that for all $x, y \in A$, if $(x, y) \in R$ then $(y, x) \notin R$, is called **asymmetric**.

Examples

Let $A = \{1, 2, 3\}$

1. Suppose $R_2 = \{(1, 2)\}$ be a relation on A .

Is R_2 Asymmetric?

Yes

2. Suppose $R_3 = \{(1, 3), (3, 1), (2, 3)\}$ be another relation on A .

Is R_3 Asymmetric?

No

3. $\emptyset, <, >$

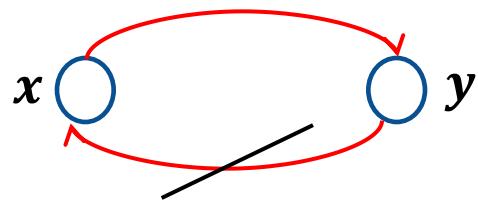
Asymmetric? Yes

4. $A \times A, \leq, \geq, |$, multiple of

Asymmetric? No

Asymmetric Relations in Graph

- If (x,y) with $x \neq y$ is an edge, then (y,x) is not an edge.
- There must also be no self loop.



Antisymmetric Relations

- A relation R on a set A such that for all $x, y \in A$, if $(x, y) \in R$ and if $(y, x) \in R$, then $x = y$, is called **antisymmetric**.

If $x \neq y$ and if (x, y) is present, then (y, x) should not be present there.

Examples

1. Let $A = \{1, 2, 3\}$

Suppose $R_1 = \{(1, 2), (2, 1), (2, 3)\}$ be a relation on A .

Is R_1 Antisymmetric?

No

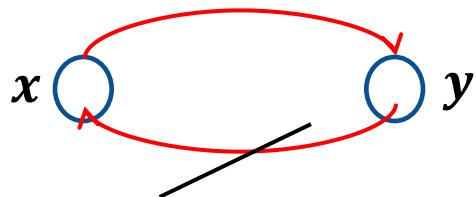
- | | | |
|--|----------------|-----|
| 2. $\emptyset, \Delta, <, >, \leq, \geq, ,$ multiple of | Antisymmetric? | Yes |
| 3. $A \times A$ | Antisymmetric? | No |

Antisymmetric Relation in Matrix and Graph

- R is a antisymmetric relation if and only if $m_{ji} = 0$, or $m_{ij} = 0$, when $i \neq j$.

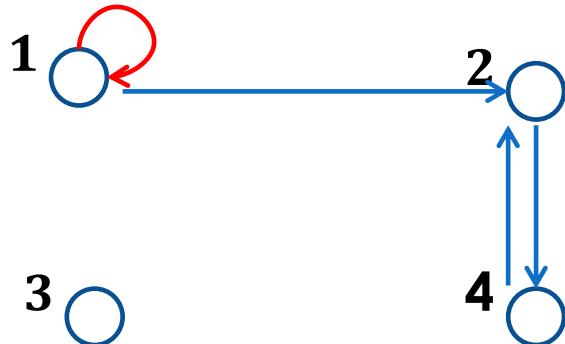
$$M_R = \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix}$$

- If (x, y) with $x \neq y$ is an edge, then (y, x) is not an edge.



Self-loops can be there.

Example



- **Reflexive?**
No
- **Symmetric?**
No
- **Transitive?**
No
- **Irreflexive?**
No
- **Antisymmetric?**
No
- **Asymmetric ?**
No

Some Points to remember

- There can be a relation which is neither reflexive nor irreflexive.

Example

1. Let $A=\{1, 2, 3\}$

Suppose $R_3 = \{(1,1), (2,2), (2,3)\}$ be a relation on A.

Neither Reflexive nor Irreflexive

Some Points to remember (Cont..)

- There can be a relation which is neither symmetric nor antisymmetric.

Example

Let $A=\{1, 2, 3\}$

Suppose $R_5 = \{(1,2), (2,3), (3,2)\}$ be a relation on A.

Neither Symmetric nor Antisymmetric

How many number of Irreflexive Relations are there on set A having n elements?

$$2^{n(n-1)}$$

How many number of **Asymmetric Relations** are there on set A having n elements?

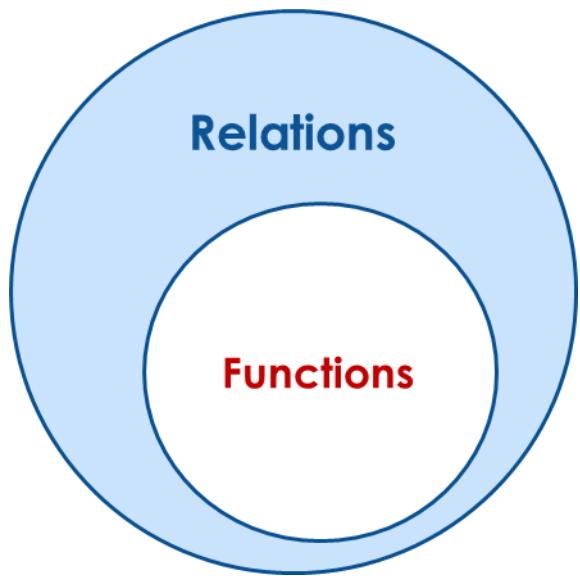
$$3^{n(n-1)/2}$$

How many number of
Antisymmetric Relations are there
on set A having n elements?

$$2^n 3^{n(n-1)/2}$$

Thank
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Inverse Relation

- If $R \subseteq A \times B$ then $R^{-1} \subseteq B \times A$, and is defined as:

$$R^{-1} = \{(b, a) | (a, b) \in R\}$$

R	R^{-1}
<	>
\leq	\geq
<i>divides</i>	<i>multiple of</i>
<i>subset</i>	<i>superset</i>

Example

- Let $A = \{1, 2, 3\}$ and $B = \{3, 4, 5\}$

Suppose $R = \{(1,3), (1,5), (2,4), (3,5)\}$

$$R^{-1} = \{(3,1), (5,1), (4,2), (5,3)\}$$

Complementary Relations

- Let R be a relation from A to B , then complementary relation R^C is defined as:

$$R^C = \{(a, b) | (a, b) \notin R \text{ and } (a, b) \in A \times B\}$$

Example

- Let $A = \{1, 2, 3\}$ and $B = \{3, 4, 5\}$

Suppose $R = \{(1,3), (1,5), (2,4), (3,5)\}$

$$A \times B = \{(1,3), (1,4), (1,5), (2,3), (2,4), (2,5), (3,3), (3,4), (3,5)\}$$

$$R^C = \{(1,4), (2,3), (3,3), (3,4), (2,5)\}$$

Combining Relation

- Given two relations R_1 and R_2 , these can be combined by using basic set operations to form new relations such as $R_1 \cup R_2$, $R_1 \cap R_2$, $R_1 - R_2$, and $R_2 - R_1$.
 - $R_1 \cup R_2 = \{(a, b) | (a, b) \in R_1 \text{ or } (a, b) \in R_2 \text{ or both}\}$
 - $R_1 \cap R_2 = \{(a, b) | (a, b) \in R_1 \text{ and } (a, b) \in R_2\}$
 - $R_1 - R_2 = \{(a, b) | (a, b) \in R_1 \text{ and } (a, b) \notin R_2\}$
 - $R_2 - R_1 = \{(a, b) | (a, b) \in R_2 \text{ and } (a, b) \notin R_1\}$

Example

- Let $A = \{1, 2, 3\}$ and $B = \{3, 4, 5\}$

Suppose $R = \{(1,3), (1,5), (2,4), (3,5)\}$,

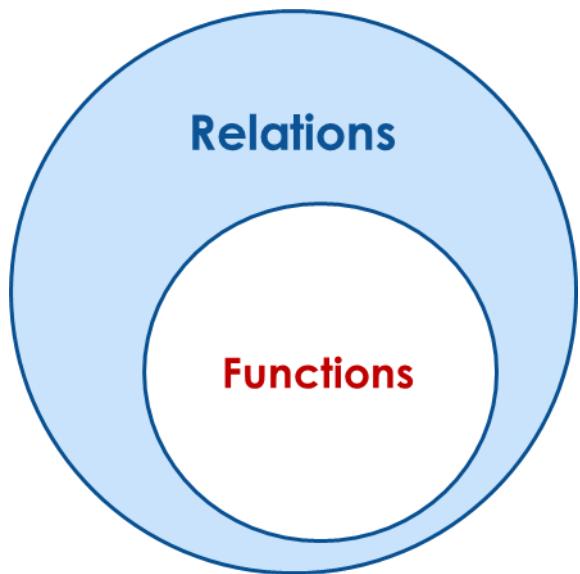
$R_1 = \{(1,4), (2,3), (2,5), (3,3), (3,5)\}$ and

$R_2 = \{(1,3), (1,4), (2,3), (3,4), (3,5)\}$

- $\square R_1 \cup R_2 = \{(1,3), (1,4), (2,3), (2,5), (3,3), (3,4), (3,5)\}$
- $\square R_1 \cap R_2 = \{(1,4), (2,3), (3,5)\}$
- $\square R_1 - R_2 = \{(2,5), (3,3)\}$
- $\square R_2 - R_1 = \{(1,3), (3,4)\}$

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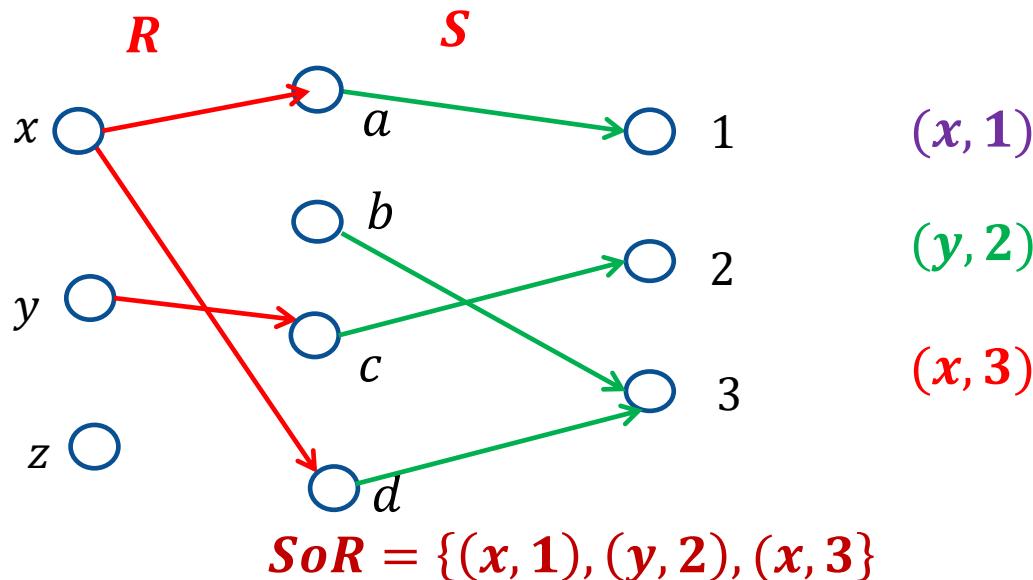
Composition of Relations

- If $R \subseteq A \times B$ and $S \subseteq B \times C$ are two relations, then the composition (or composite) of S with R is a relation from A to C and is defined as:

$$SoR = \{ \{a, c\} \mid \exists b \in B \text{ such that } (a, b) \in R \text{ and } (b, c) \in S \}$$

Representing the Composition of Relations

- Let $A = \{x, y, z\}$, $B = \{a, b, c, d\}$ and $C = \{1, 2, 3\}$.
- Suppose $R = \{(x, a), (x, d), (y, c)\}$ be a relation from A to B .
- Suppose $S = \{(a, 1), (b, 3), (c, 2), (d, 3)\}$ be a relation from B to C .



Power of Relations

- If $R \subseteq A \times A$, then

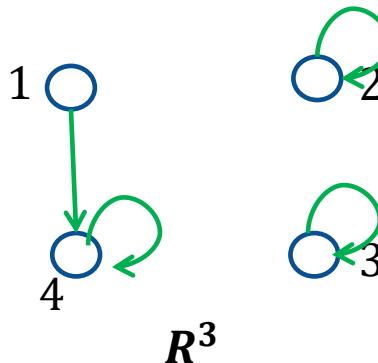
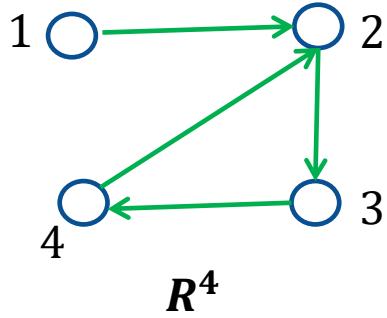
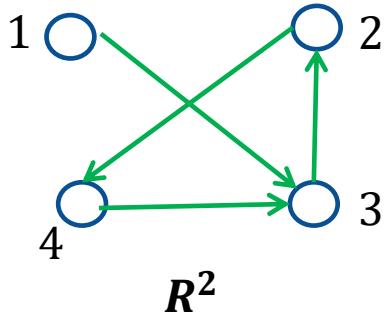
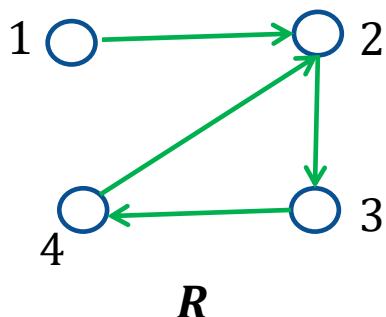
$$R^2 = R \circ R$$

$$R^3 = R^2 \circ R$$

⋮

$$R^n = R^{n-1} \circ R$$

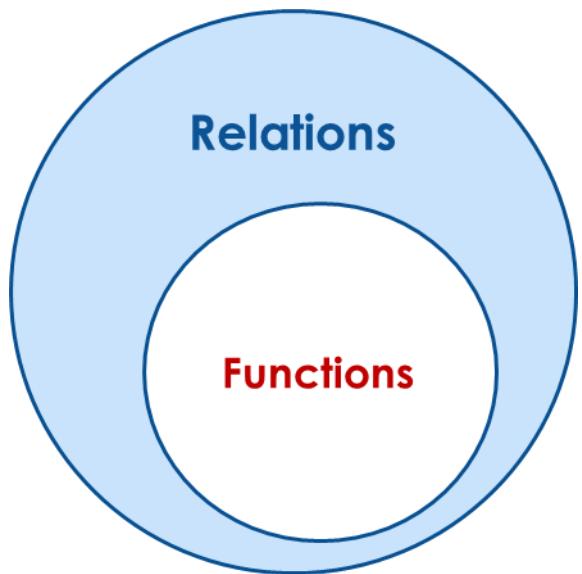
Example



The pair (a, b) is in R^n if there is a path of length n from a to b in R .

Thank
you!!!





Relations

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Equivalence Relation

- Let R be a relation on set A , then R is called equivalence relation if it is:
 - Reflexive
 - Symmetric
 - Transitive

Examples

- Let $A = \{1, 2, 3\}$

1. \emptyset i.e. Empty Relation on A

Reflexive?

Symmetric?

Transitive?

Not an Equivalence Relation

2. $\Delta = \{(1, 1), (2, 2), (3, 3)\}$

Reflexive?

Symmetric?

Transitive?

Equivalence Relation on A

Smallest Equivalence Relation on A

Examples (Cont..)

- Let $A = \{1, 2, 3\}$

3. Universal Relation on A i.e. $A \times A$

Reflexive?

Symmetric?

Transitive?

Equivalence Relation

Largest Equivalence Relation on A

- 4. Let $A = \{1, 2, 3, 4\}$

$R = \{(1, 1), (1, 2), (2, 1), (2, 2), (3, 3), (3, 4), (4, 3), (4, 4)\}$

Reflexive?

Symmetric?

Transitive?

Equivalence Relation on A

If R_1 and R_2 are two equivalence relations on A , then which of the following is always true?

- I.* $R_1 \cap R_2$ is an Equivalence Relation.
 - II.* $R_1 \cup R_2$ is an Equivalence Relation.
-
- (a) Only *I*
 - (b) Only *II*
 - (c) Both are true
 - (d) Both are false

Exercise

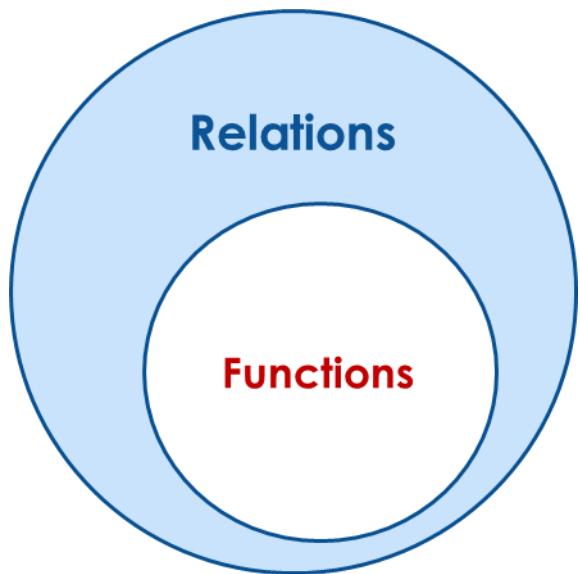
- Let R be a relation defined on *set of integers* as:

xRy iff $x + y$ is even

Is R an equivalence relation?

Thank
you!!!





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Equivalence Class

- Let R be an equivalence relation on A , and $a \in A$.
- The equivalence class of a , denoted as $[a]$ or \bar{a} , is defined as:

$$\bar{a} = [a] = \{b \in A | (a, b) \in R\}$$

Examples

- Let $R = \{(1,1), (1,2), (2,1), (2,2), (3,3), (3,4), (4,3), (4,4)\}$ on $A = \{1, 2, 3, 4\}$

First Check whether R is an equivalence relation on A or not.

Reflexive?

Symmetric?

Transitive?

Equivalence Classes:-

[1] = {1, 2}

[2] = {1, 2}

[3] = {3, 4}

[4] = {3, 4}

Examples

- Let $R = \{(1,1), (2,2), (3,3), (4,4)\}$ on $A = \{1, 2, 3, 4\}$

Equivalence Classes:-

[1] = {1}

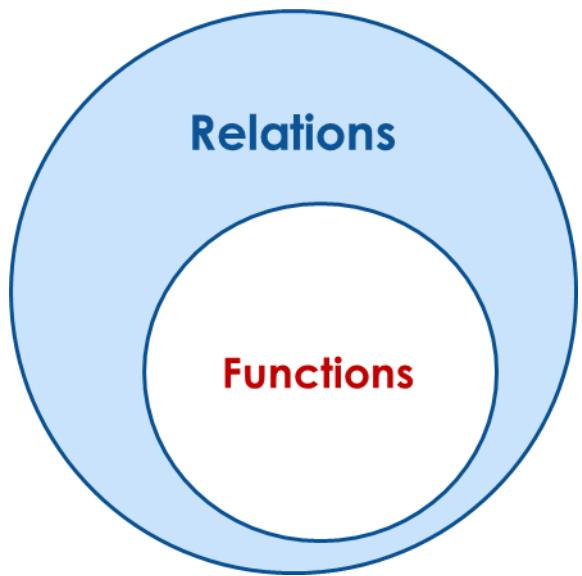
[2] = {2}

[3] = {3}

[4] = {4}

Thank
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Equivalence Relation to Partition

- Let $R = \{(1,1), (1,2), (2,1), (2,2), (3,3), (3,4), (4,3), (4,4)\}$ on $A = \{1, 2, 3, 4\}$.
- R is an equivalence relation on A .

Equivalence Classes:-

$$\begin{aligned}[1] &= \{1, 2\} = [2] \\ [3] &= \{3, 4\} = [4]\end{aligned}$$

- Partition $P = \{\{1, 2\}, \{3, 4\}\}$

Partition to Equivalence Relation

- Let $A = \{1, 2, 3, 4\}$ be a set and $P = \{\{1, 3\}, \{2, 4\}\}$ be a partition on A .
- Find Equivalence relation on A .

$$\{1, 3\} \rightarrow \{(1, 1), (1, 3), (3, 1), (3, 3)\}$$

$$\{2, 4\} \rightarrow \{(2, 2), (2, 4), (4, 2), (4, 4)\}$$

Therefore, the equivalence relation on A is:

$$\square R = \{(1, 1), (1, 3), (3, 1), (3, 3), (2, 2), (2, 4), (4, 2), (4, 4)\}$$

Result

- There is a one-to-one correspondence between partitions of A and Equivalence Relation on A .
- Therefore, if $|A| = n$, then

Number of Partitions of A = Number of Equivalence Relations on A
 $= B_n$ (Bell Number)

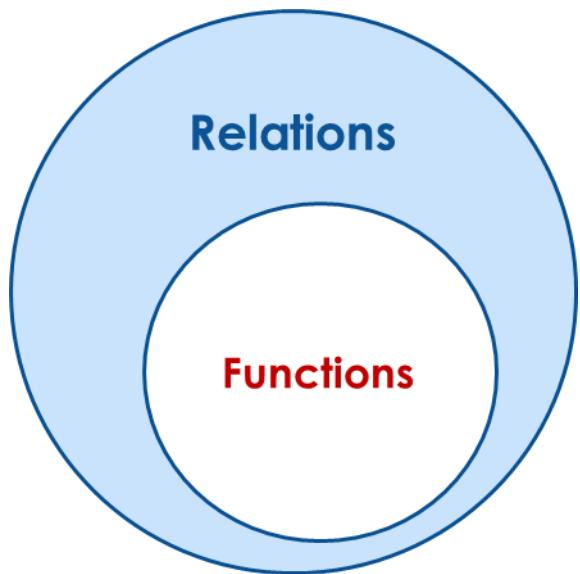
Bell Number:

$$B_n = \sum_{k=0}^{n-1} n - 1 c_k B_k$$

where, $B_0 = 1$

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Closure of Relations

- Reflexive Closure
- Symmetric Closure
- Transitive Closure

Reflexive Closure

- A relation is called **reflexive closure R_r** of relation R if:
 - 1) It is reflexive.
 - 2) It contains R .
 - 3) It is the minimal relation satisfying conditions (1) and (2).

Examples

1. Let $A=\{1, 2, 3\}$

Suppose $R_1 = \{(1,1), (2,2), (2,3)\}$ be a relation on A .

$$R_r = \{(1, 1), (2, 2), (2, 3), (3, 3)\}$$

Examples (Cont..)

2. R is a relation defined on set of positive integers such that aRb if $a < b$.

Reflexive Closure?

Result:

- $R_r = R \cup \Delta$
- $R_r = R$ iff R is Reflexive.

Symmetric Closure

- A relation is called **symmetric closure** R_s of relation R if:
 - 1) It is symmetric.
 - 2) It contains R .
 - 3) It is the minimal relation satisfying conditions (1) and (2).

Examples

- Let $A=\{1, 2, 3\}$

Suppose $R_1 = \{(1,1), (2,2), (2,3)\}$ be a relation on A .

$$R_s = \{(1,1), (2,2), (2,3), (3,2)\}$$

Result

- $R_s = R \cup R^{-1}$
- $R_r = R$ iff **R is Symmetric.**

Transitive Closure

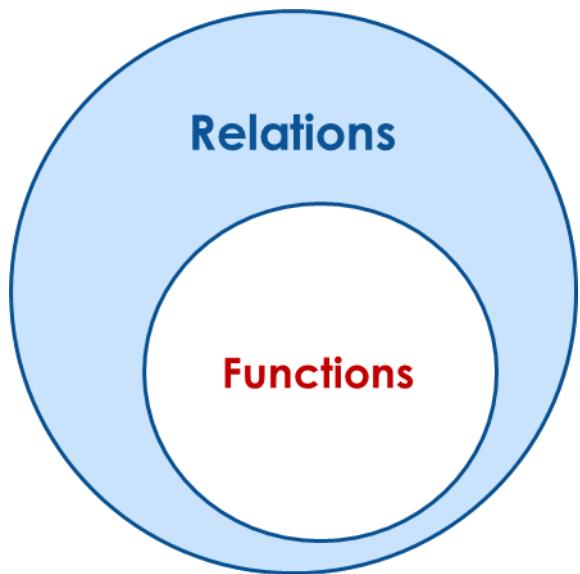
- A relation is called Transitive closure R^* of relation R if:
 - 1) It is transitive.
 - 2) It contains R .
 - 3) It is the minimal relation satisfying conditions (1) and (2).

Result

1. Let $|A| = n$,
then, $R^* = R^1 \cup R^2 \cup \dots \cup R^n$
2. R is transitive iff $R^* = R$

Thank
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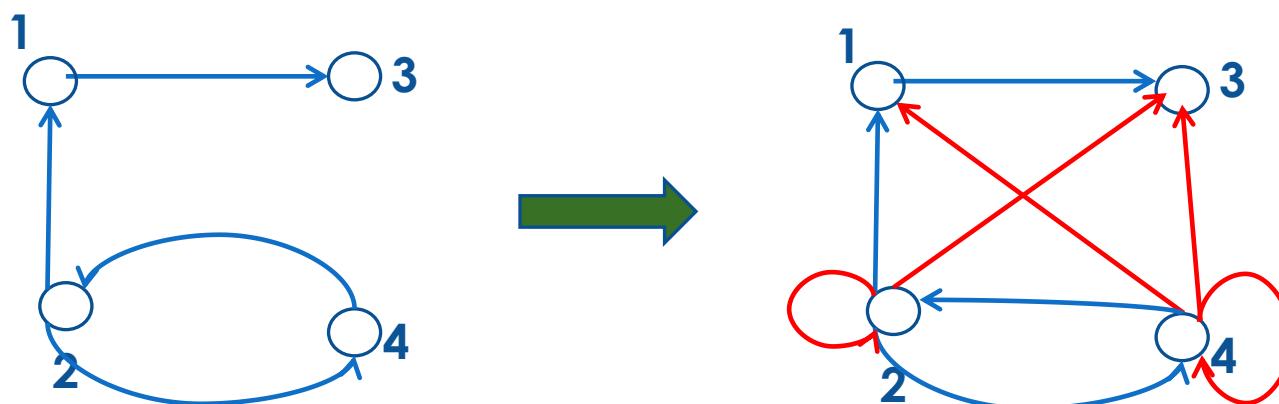
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Warshall's Algorithm

- Computes the transitive closure of a relation

Example of transitive closure:



Warshall's Algorithm (Cont..)

- Main concept: a path exists between two vertices i, j , iff
 - there is an edge from i to j ; or
 - there is a path from i to j going through vertex 1; or
 - there is a path from i to j going through vertex 1 and/or 2; or
 - there is a path from i to j going through vertex 1, 2, and/or 3; or
 - ...
 - there is a path from i to j going through any of the other vertices

Warshall's Algorithm (Cont..)

- On the k^{th} iteration, the algorithm determine if a path exists between two vertices i,j using vertices among $1,\dots,k$ allowed as intermediate

$$W^{(k)}[i,j] = \begin{cases} W^{(k-1)}[i,j] \\ \text{or} \\ (W^{(k-1)}[i,k]) \text{ and } (W^{(k-1)}[k,j]) \end{cases}$$

Warshall's Algorithm (Cont..)

- Recurrence relating elements $W^{(k)}$ to elements of $W^{(k-1)}$ is:

$$W^{(k)}[i, j] = W^{(k-1)}[i, j] \text{ or } (W^{(k-1)}[i, k] \text{ and } W^{(k-1)}[k, j])$$

- It implies the following rules for generating $W^{(k)}$ from $W^{(k-1)}$ is:

1. If an element in row i and column j is 1 in $W^{(k-1)}$, it remains 1 in $W^{(k)}$.
2. If an element in row i and column j is 0 in $W^{(k-1)}$, it has to be changed to 1 in $W^{(k)}$ if and only if the element in its row i and column k and the element in its row k and column j are both 1's in $W^{(k-1)}$.

Warshall's Algorithm (Cont..)

- The procedure for computing $W^{(k)}$ from $W^{(k-1)}$ is as follows:
 1. First transfer all 1's in $W^{(k-1)}$ to $W^{(k)}$.
 2. List the locations p_1, p_2, \dots , in column k of $W^{(k-1)}$, where the entry is 1, and the locations q_1, q_2, \dots , in row k of $W^{(k-1)}$, where the entry is 1.
 3. Put 1's in all the positions (p_i, q_i) of $W^{(k)}$ (if they are not already there).

Warshall's Algorithm (Cont..)

$$W^{(k-1)} = \begin{bmatrix} & j & k \\ i & \left[\begin{array}{cc} 0 & 1 \\ 1 & \end{array} \right] \\ k & \end{bmatrix} \rightarrow W^{(k)} = \begin{bmatrix} & j & k \\ i & \left[\begin{array}{cc} 1 & 1 \\ 1 & \end{array} \right] \\ k & \end{bmatrix}$$

Figure 1: Step for Changing zeros in Warshall's Algorithm

Example

1. Find transitive closure of relation R represented by following matrix (using Warshall's algorithm):

$$M_R = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \left[\begin{matrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{matrix} \right] \end{matrix}$$

Solution 1:

$$W^{(0)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \left[\begin{matrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{matrix} \right] \end{matrix}$$

$$W^{(1)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \left[\begin{matrix} 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{matrix} \right] \end{matrix}$$

$$W^{(2)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \left[\begin{matrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{matrix} \right] \end{matrix}$$

$$W^{(3)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \left[\begin{matrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{matrix} \right] \end{matrix}$$

$$W^{(4)} = W^{(3)} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \left[\begin{matrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{matrix} \right] \end{matrix}$$

Answer is :

$$M_{R^*} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \left[\begin{matrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{matrix} \right] \end{matrix}$$

$$R^* = \{(1, 1), (1, 2), (1, 3), (1, 4), (2, 1), (2, 2), (2, 3), (2, 4), (3, 4)\}$$

Thank
you!!!

