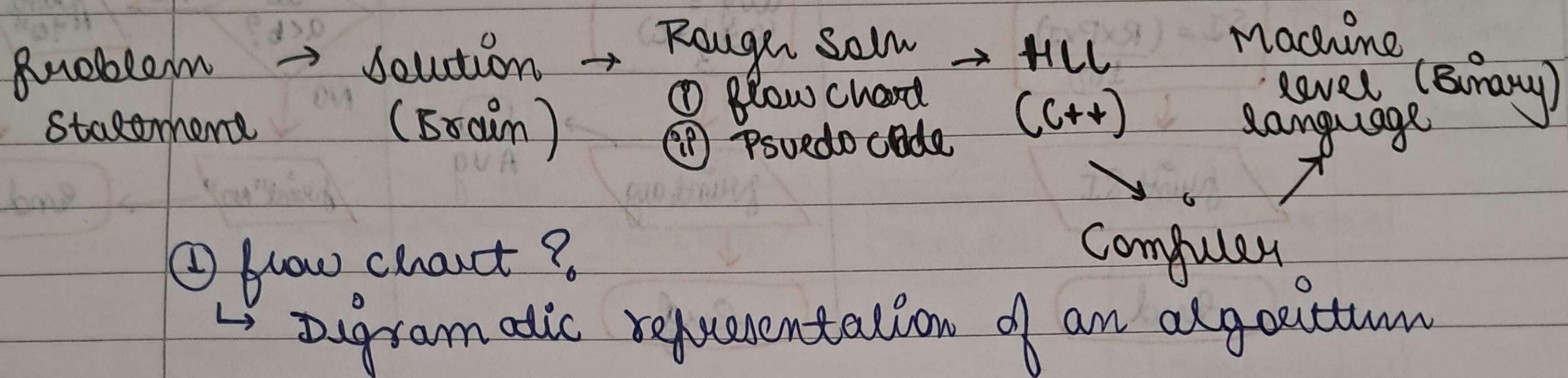


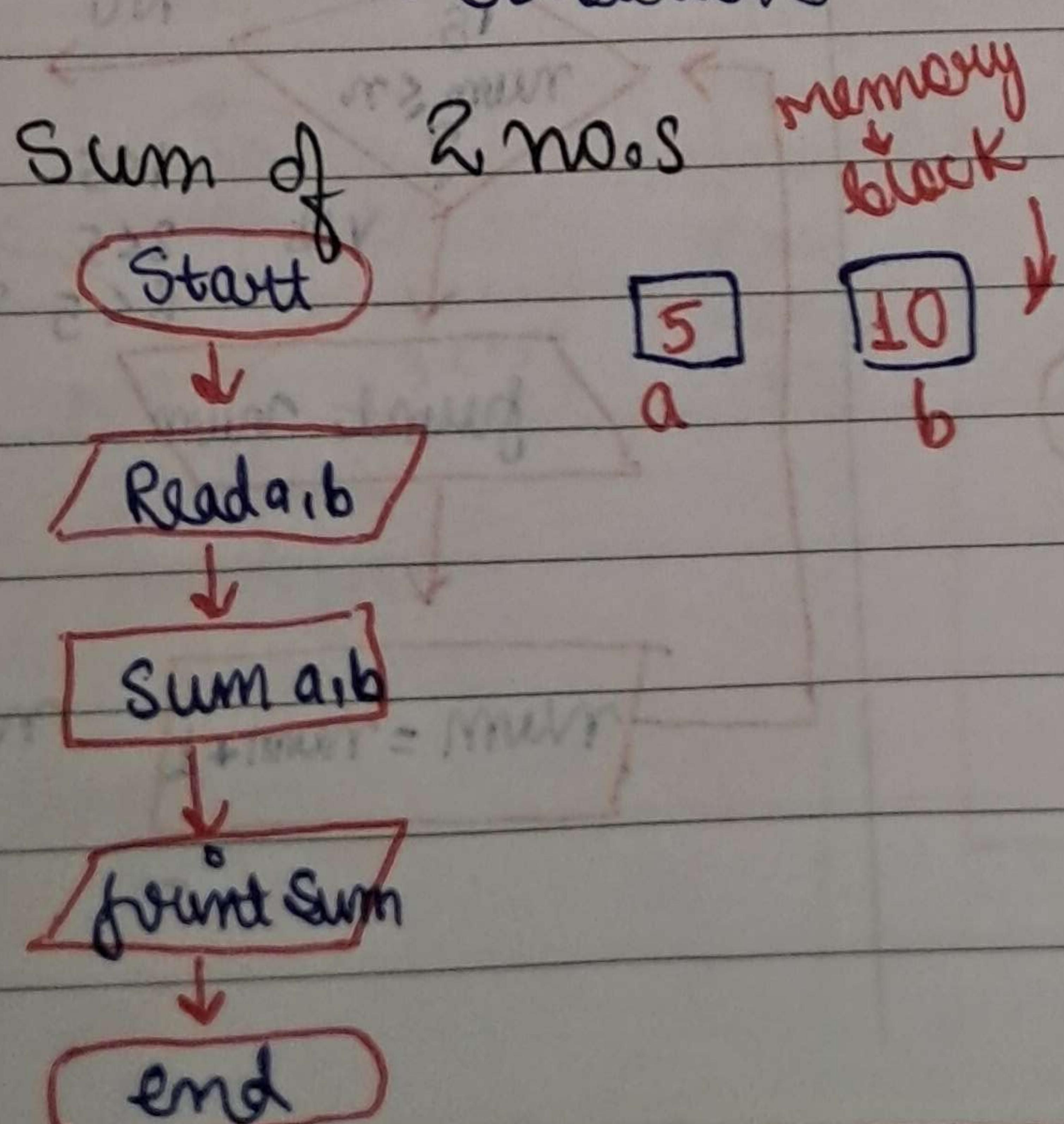
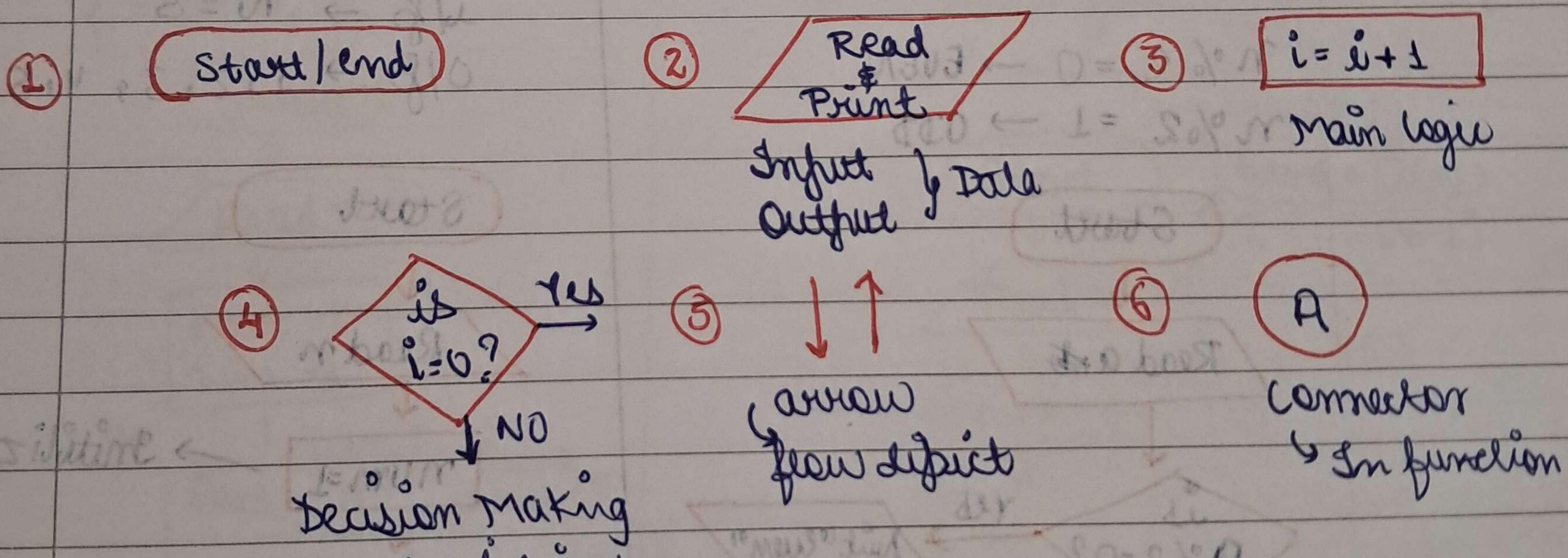
## LEC-1

### How to approach problem

1. Understand the problem
2. Enter Appropriate input values
3. Create Logics / Algorithm

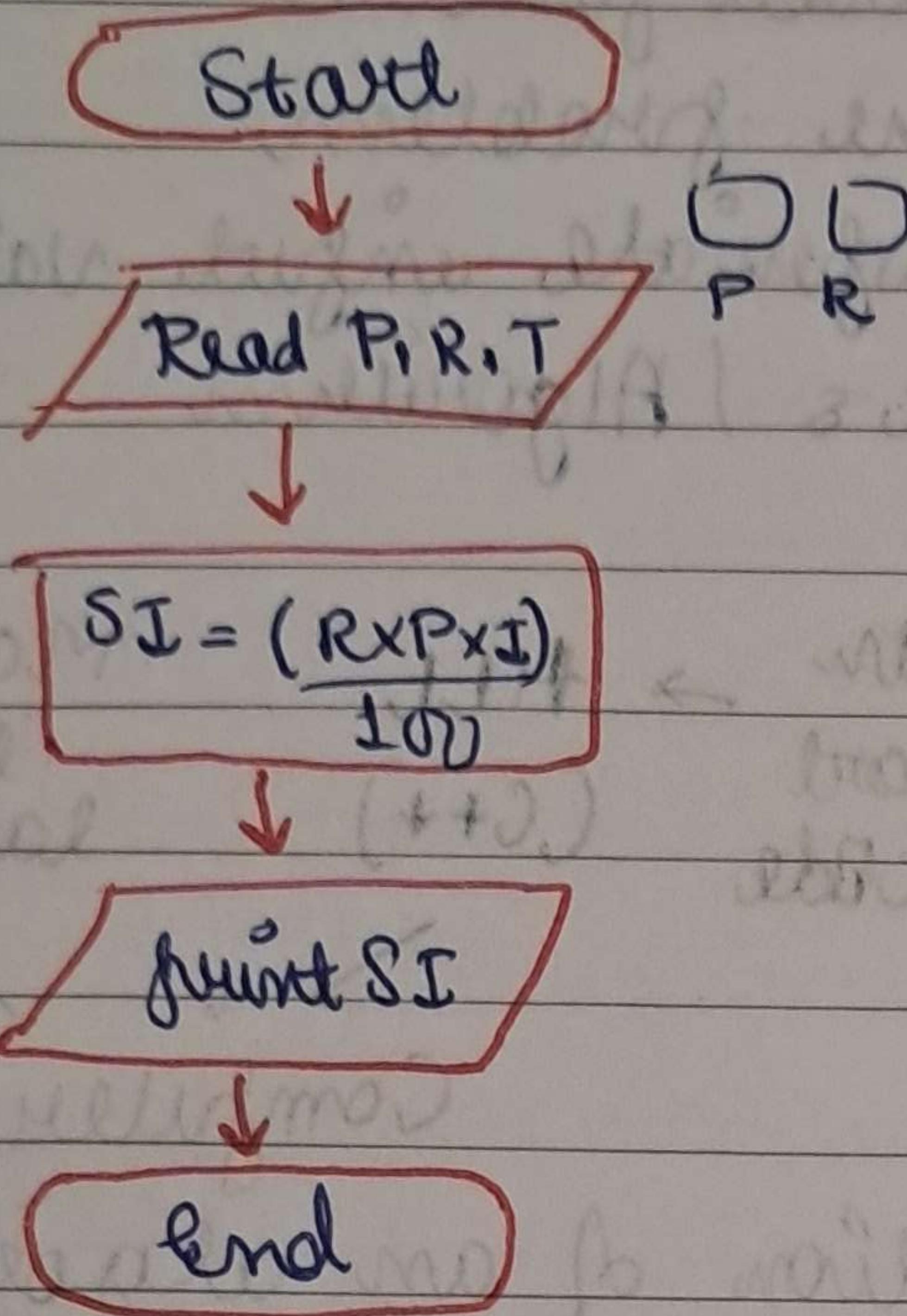


### Component of flow chart

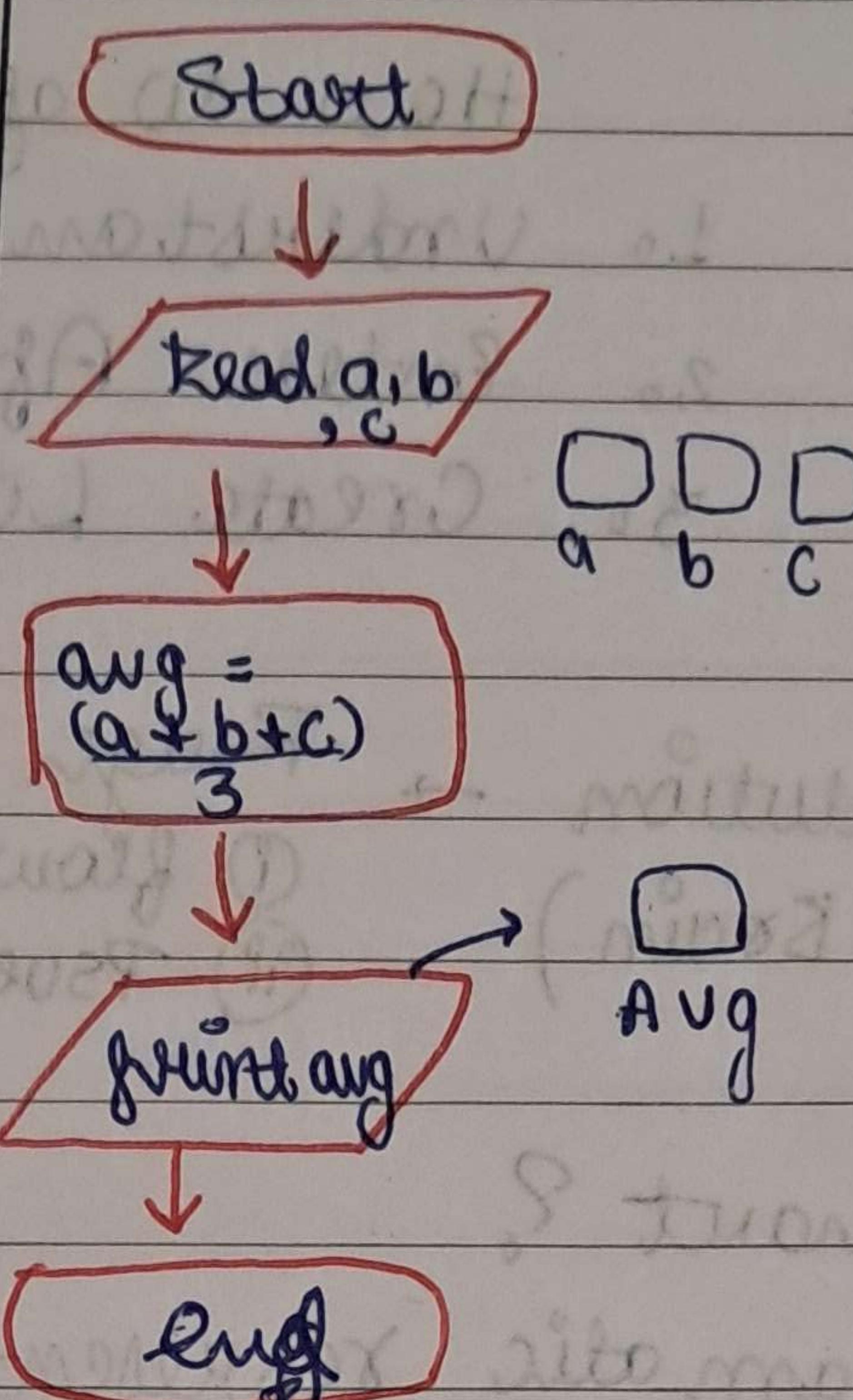


- Pseudo code → generic way of representing code
1. Read 2 no.s  $a$  and  $b$
  2. Sum  $a = a + b$
  3. Print sum
  4. Read  $a$
  5. Read  $b$
  6. Let  $sum = 0$
  7.  $sum = a + b$
  8. Print sum
- various person to person

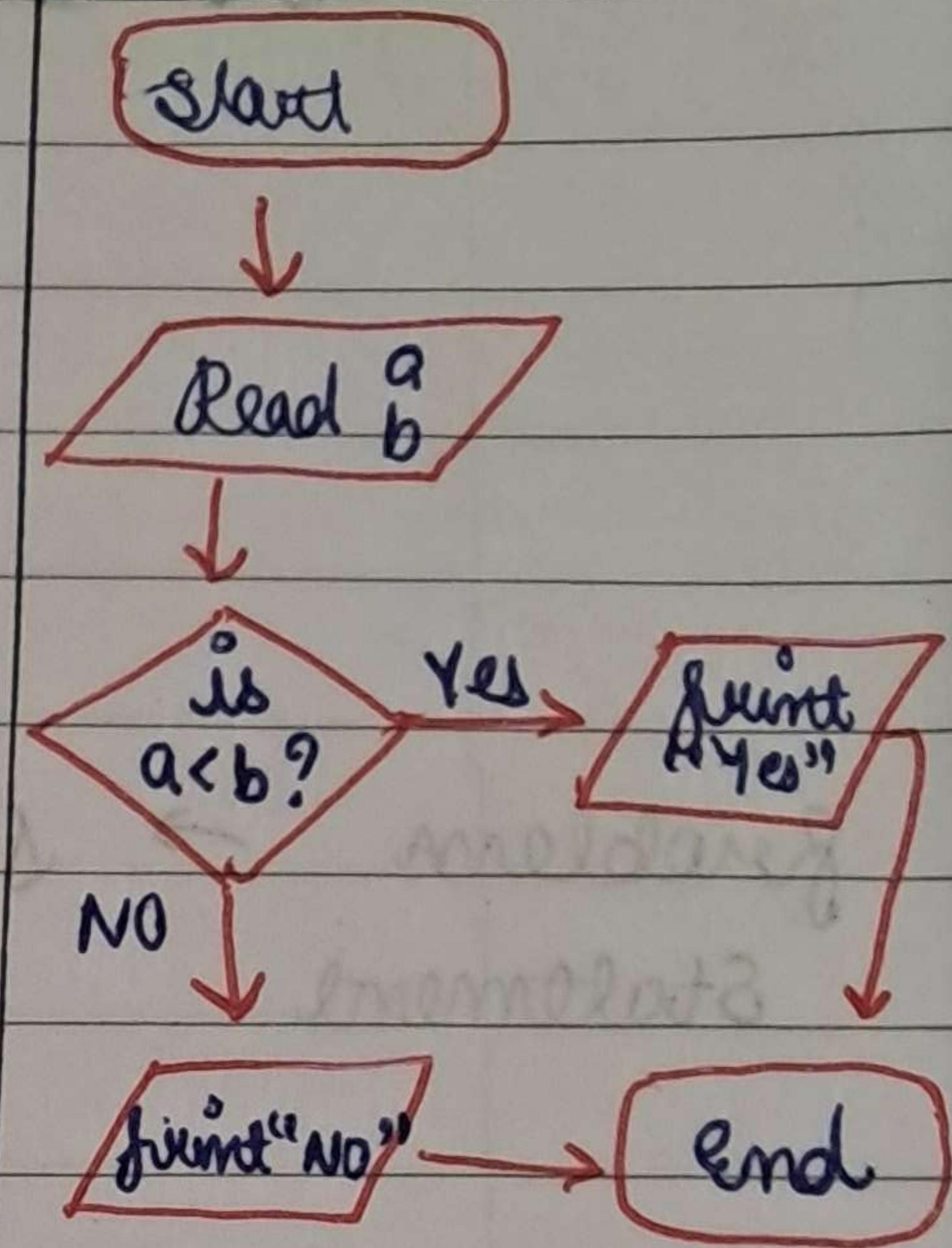
② Simple Interest  
 $S.I. = \frac{R \times P \times T}{100}$



③ Avg of 3 nos



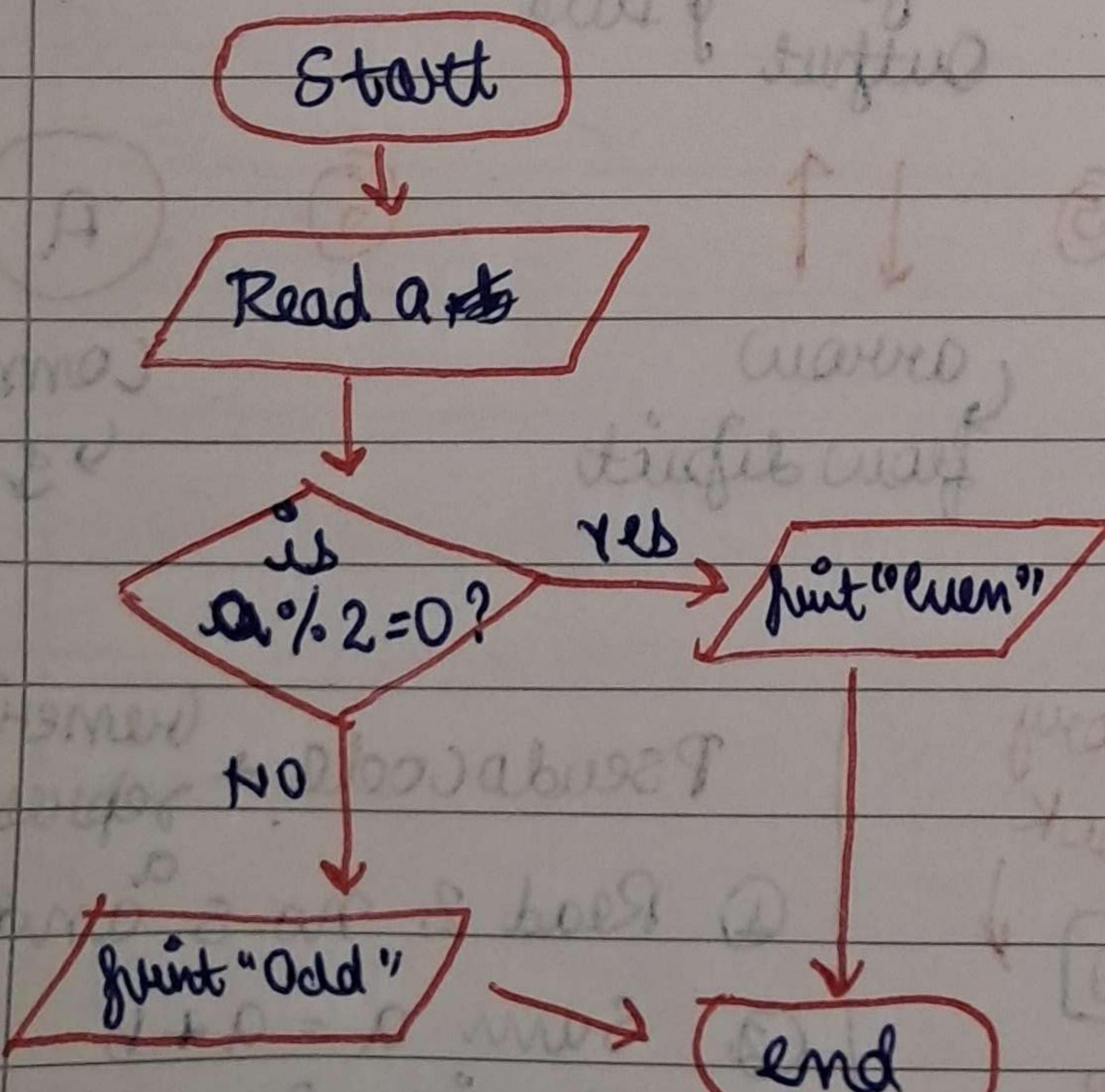
④ a < b? "Yes" / "No"



⑤ Odd or even ??

$$n \% 2 = 0 \rightarrow \text{EVEN}$$

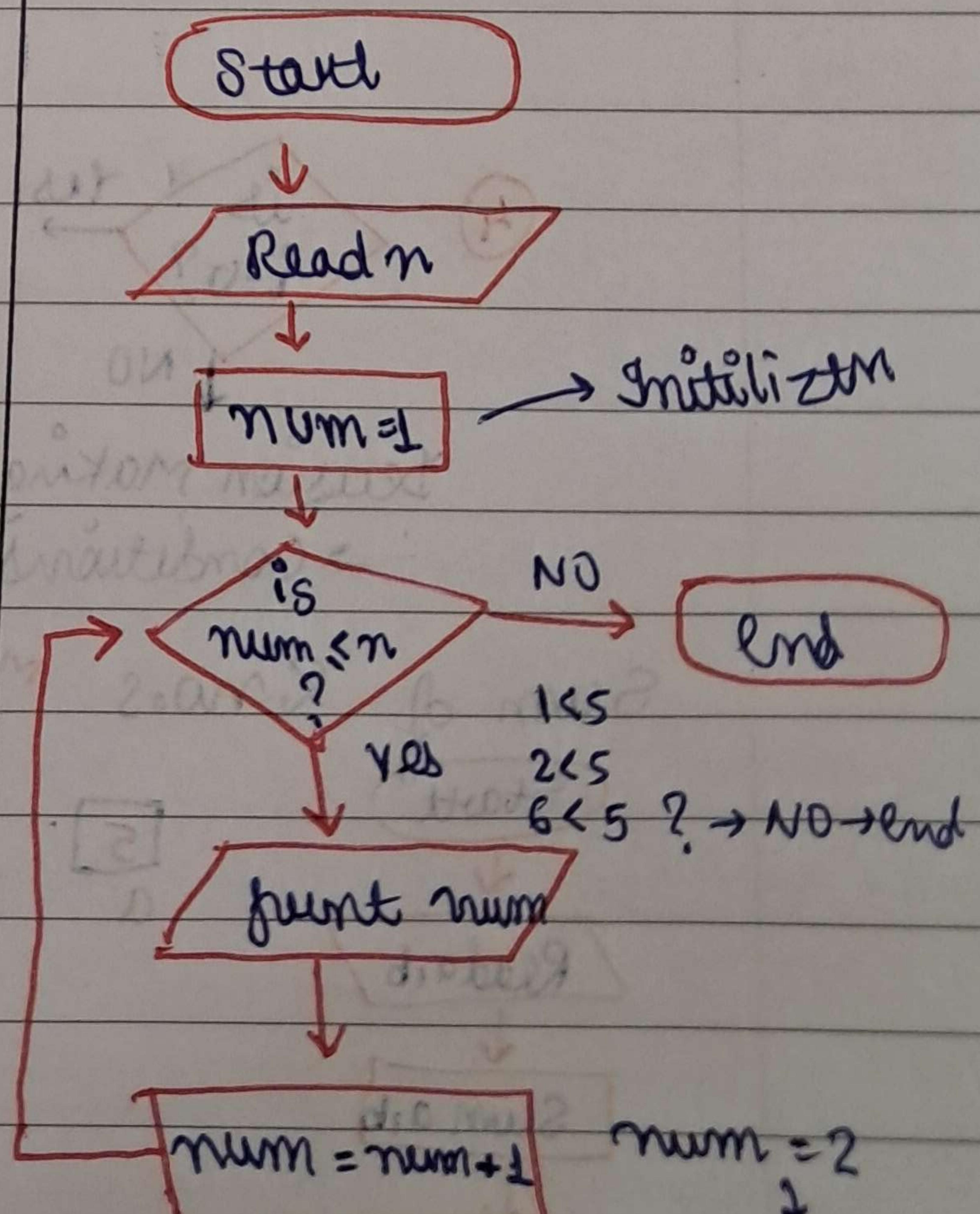
$$n \% 2 = 1 \rightarrow \text{ODD}$$



⑥ Print 1 to N

$$i/f \rightarrow N = 5$$

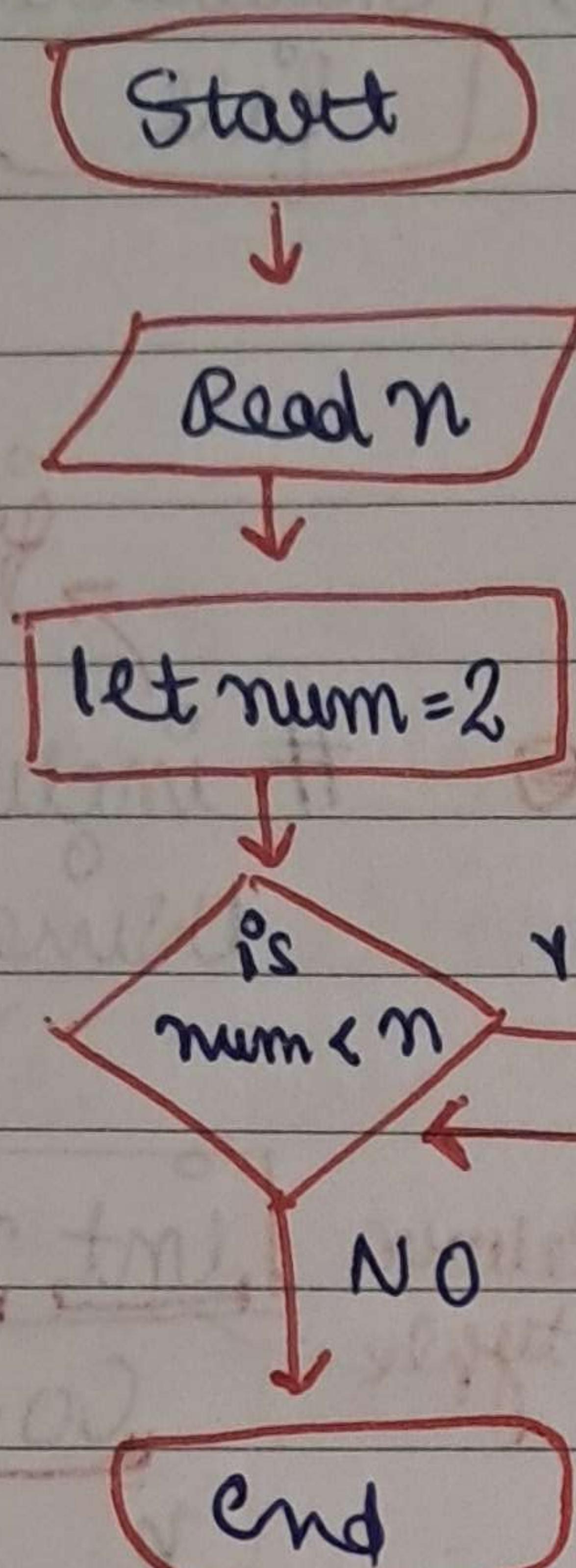
o/f  $\rightarrow 1, 2, 3, 4, 5$



6 1 to n (even print)

if  $n = 5$

0/fp  $\rightarrow 2, 4$



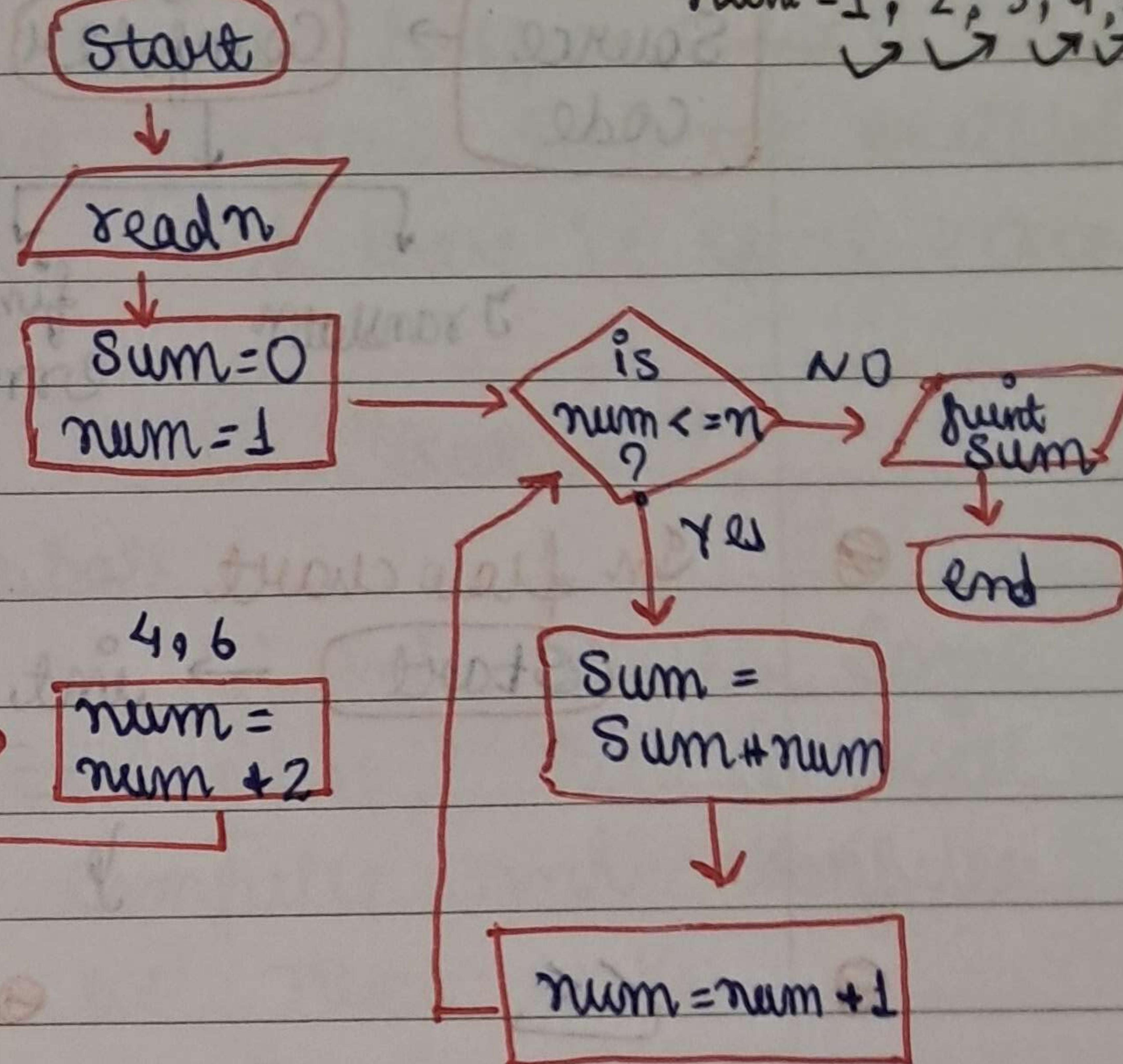
7 find sum (1 to n)

$n = 5$

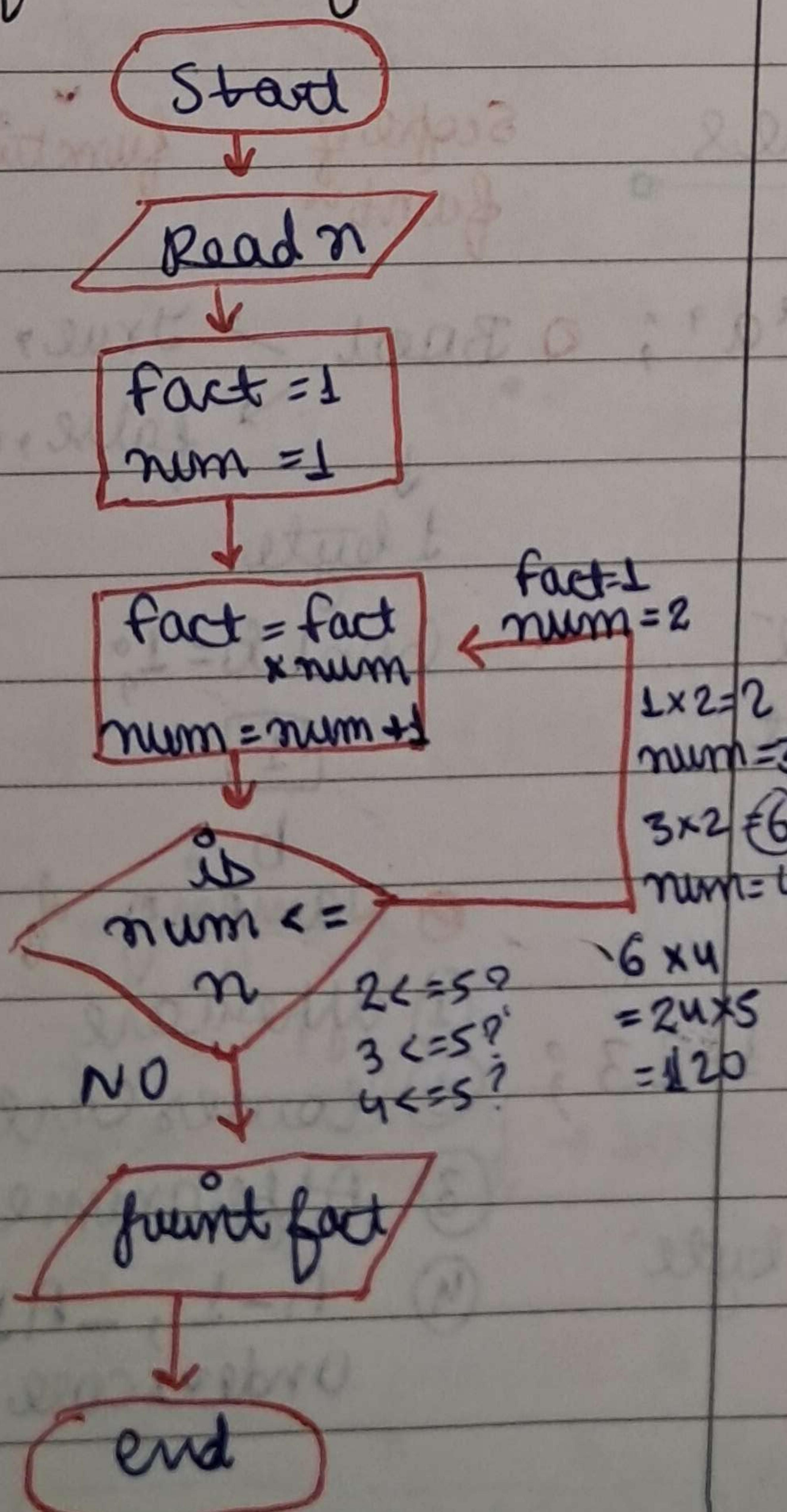
0/fp  $\rightarrow 1+2+3+4+5 = 15$

$n = 5$

Sum = 0, 3 6 10  
num = 1, 2, 3, 4, 5  
→ → → → →



8 factorial of n

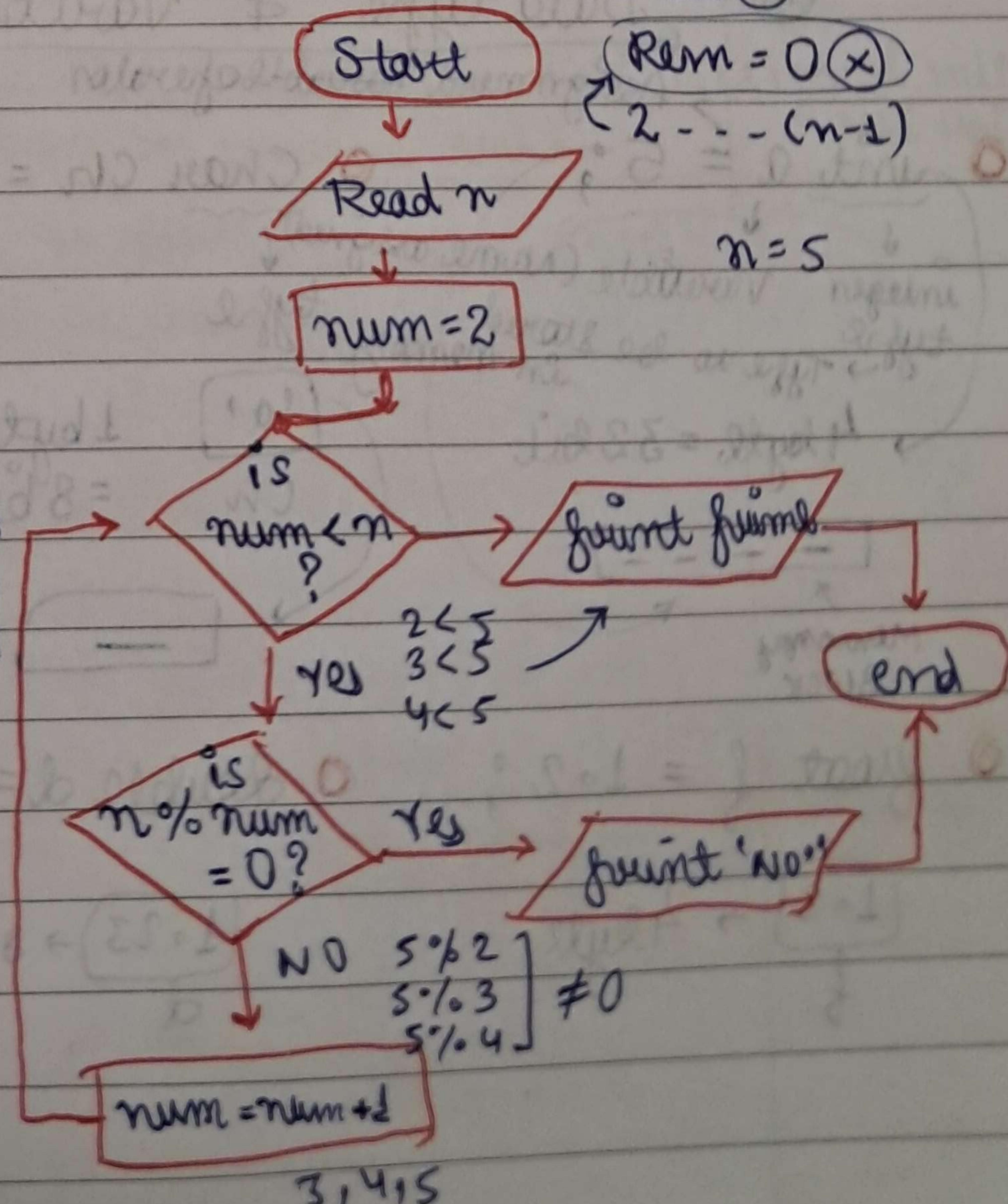


9 Check prime or not?

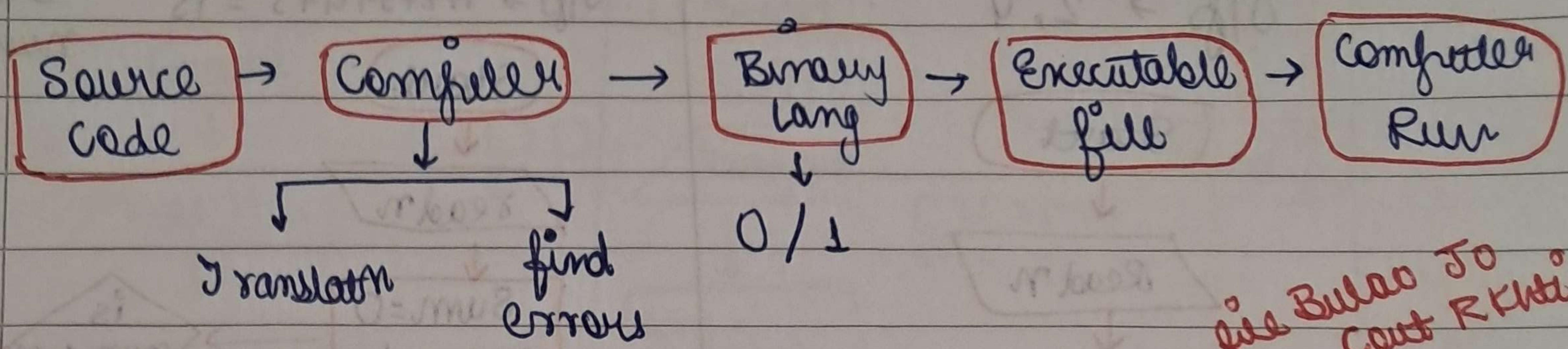
$1, \dots, (n-1), n$

table  $\times$   
 $2, \dots, (n-1)$   
 $n = 5$

Rem = 0  $\times$



## LEC-2



② In flow chart

Start → int main() {

② <<  
↓ to write / display  
Something

② endl / \n  
↓ newline

Return type int main() {  
cout << " ";

② ; → line Khatu Kne  
Ke liye

② Data type & Variables

Assignment Variable operators

Scope of function name

point to  
shree,,  
mai ho

② int a = 5;

integer variable (Name assigned)  
type → Type to be stored in memory type

4 byte = 32 bit



Memory block

'a'  
ch

1 byte  
= 8 bit

② float f = 1.2;

1.2

→ 4 byte

② double d = 1.23;

1.23

→ 8 byte

bool b = 1;

1

b

② Naming of Variable

- ① Uppercase
- ② Lowercase
- ③ Alphanumeric → abc1
- ④ A-1, -A1  
Underscore → 1abc X

→ // or Ctrl + / :- used to Comment

line of code ignore by compiler

→ return 0; → This mean that the program executed ~~successfully~~

→ return 1; → This mean the program does not executed successfully and there is some error.

→ cout << "Babbar" ; cout << "Babbar" ;

Output :- Babbar Babbar

Explanation → If we use ';' and write some code just after the Semicolon, then it will get executed by the compiler and compiler will treat that code as next task

→ Declaration :-

Ex - int age;

→ Initialisation :-

Ex - int age = 21 ;

If we print a variable → Jiske koi assigned value gives → random unique value

Ex → int random ;

cout << random ;

Output  
43426387

↳ unique  
Always

# how data is stored

①

int a = 8

↓  
4 byte

Binary  
Convert → 1000

-----

4 bit

8 8 8 8

0 0000 0000 0000 1000

32 bit

Last 4 bit = 1000  
Rest 28 bit = 0000-----

② Char ch = 'a'

↓ 1 bit

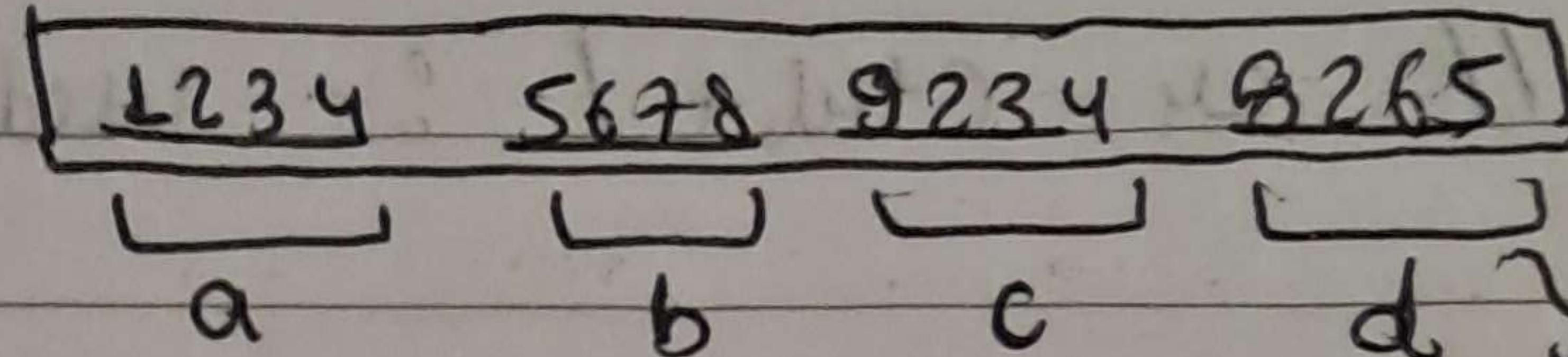
-----

8 bit

→ ASCII Table

↓  
97 → Binary

Suppose if



any integers

4 char

confusion  
between by  
datatype

## # Type Casting

int a = 'a';  
cout << a << endl;  
O/p → 97

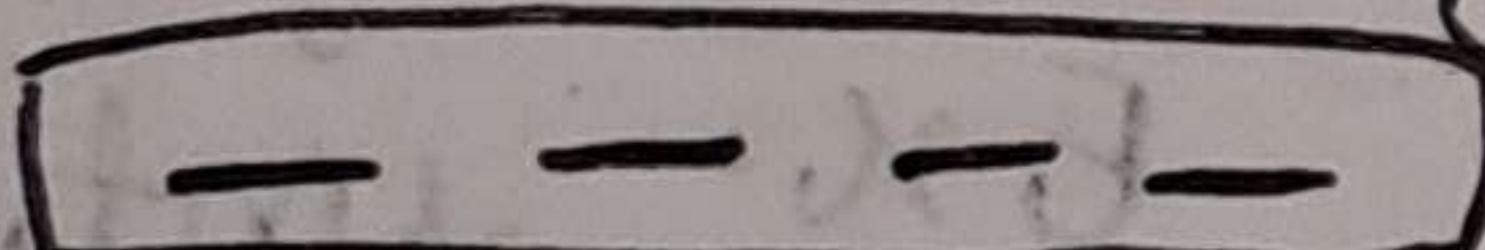
Since a is mapped to 97  
in ASCII table

char ch = 98;

O/p → b

int typecasted into char

If int → 4 byte = 32bit



↳ Max →  $2^{32}-1$

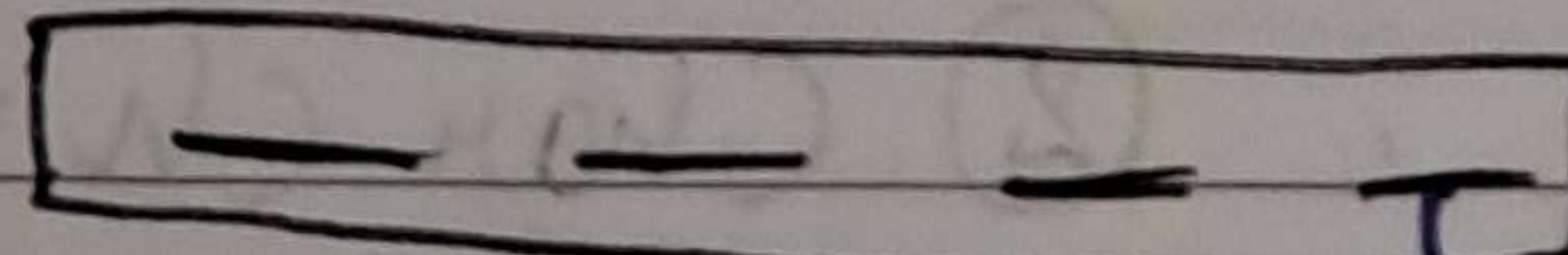
Store ↳ Min → 0

Similarly Char → 1 byte = 8 bit

↳ Max →  $2^8-1$

Min → 0

If we have integer of  $2^{16}$  → less than  $2^{32}-1$   
then use store kro Mai But greater  $2^8-1$



Char ch = 123456;

cout << ch << endl;

O/p → @

→ 64 bit jayega last value check kro

use small  $2^8-1$   
form Mai Bna

lga jayega aur  
output milega

1-125 = std::max

float a = 5.466;

int b = int(a);

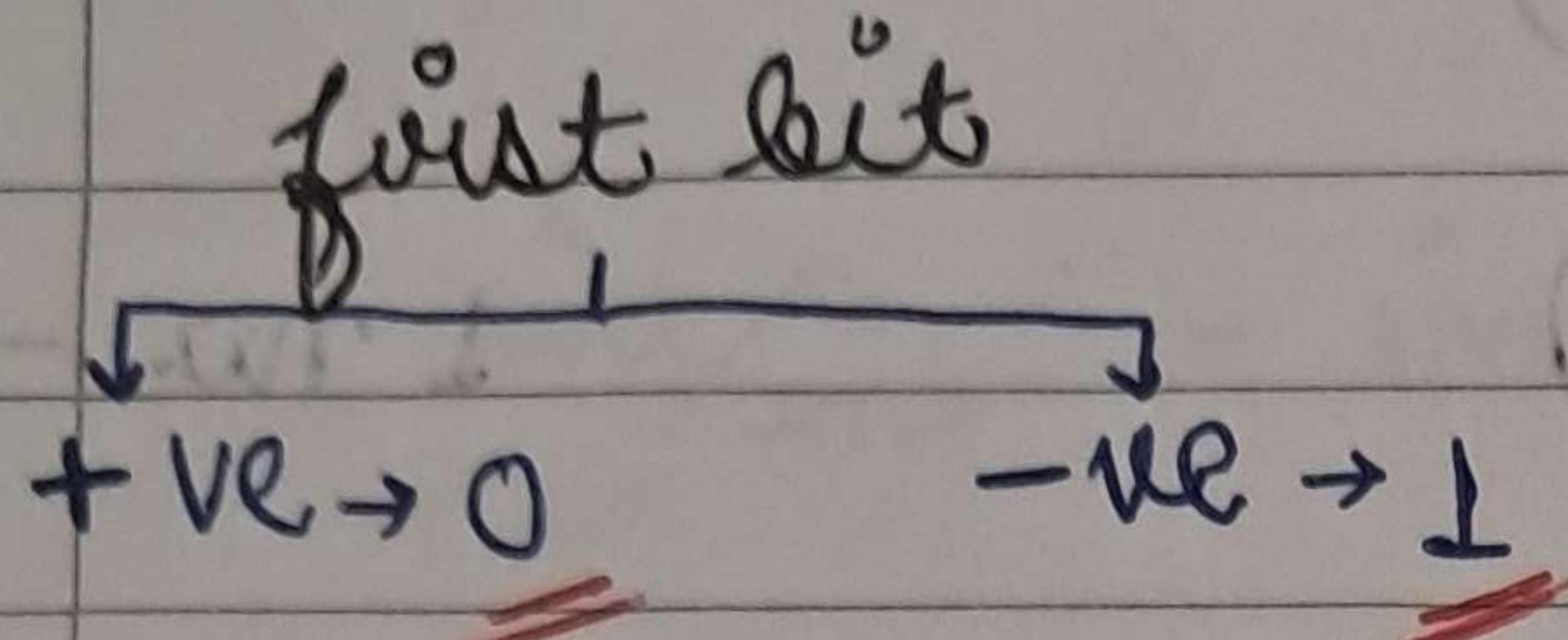
cout << b;

↳ o/p → 5

## Explicit Typecasting

#

How -ve numbers are stored?



-ve no. → (-5) → store

- ① Ignore the -ve sign → 5
- ② Convert into binary representation

00 00 00 101

- ③ Take 2's Complement & store

00 000	0000	0000	0000 101
11111	1111	1111	11010
+ 1			

2's Complement  
Bit exchange

11111 1111 1111 11011

11111 1111 1111 11011

11111 1111 1111 11011

11111 1111 1111 11011

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11111 1111 1111 11011

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11111 1111 1111 11011

11111 1111 1111 1101

Oct 31 = 31 days

$$\text{Mem bytes} = 2^{31} - 1$$

0 | - - - - - - - -

→ Rest letters will be considered **+ve**

1 - - - 1 - - -

→ Rest is it - will be considered

$$0|0 \rightarrow -1 \sim -1 \rightarrow 0$$

1 m. → " 0 , ,

$$\boxed{110 \dots} \rightarrow -0 \rightsquigarrow 0$$

2 Representation  
use to display

Do Jane that one reg  $\rightarrow$  2's campument

$$-(2^{31}-1) \quad - \quad - \quad - \quad - \quad - \quad 2^{31}-1$$

↓ to ↓

$$= 2^{31} - \dots - 2^{81} - 1$$

1 Rep same

Signed data

unsigned data

unsigned int a1 = -112;

front a 1

→ Huge Value

# Detail in vs coda

# # OPERATORS

## ① Arithmetic

↓      ↓      ↓      ↓  
+      -      \*      /      %

## ② Relational

↓      ↓      ↓      ↓      ↓  
=      >      <      >=      <=      !=

int/int = int

## ③ Logical

↓      ↓  
& &      ||      !  
and      or      not

double/int = Double

Cin >> " " space 32  
" \t " tab 10 don't Read X  
" \n " enter 9 Read ✓  
a = Cin.get()

ASCII table

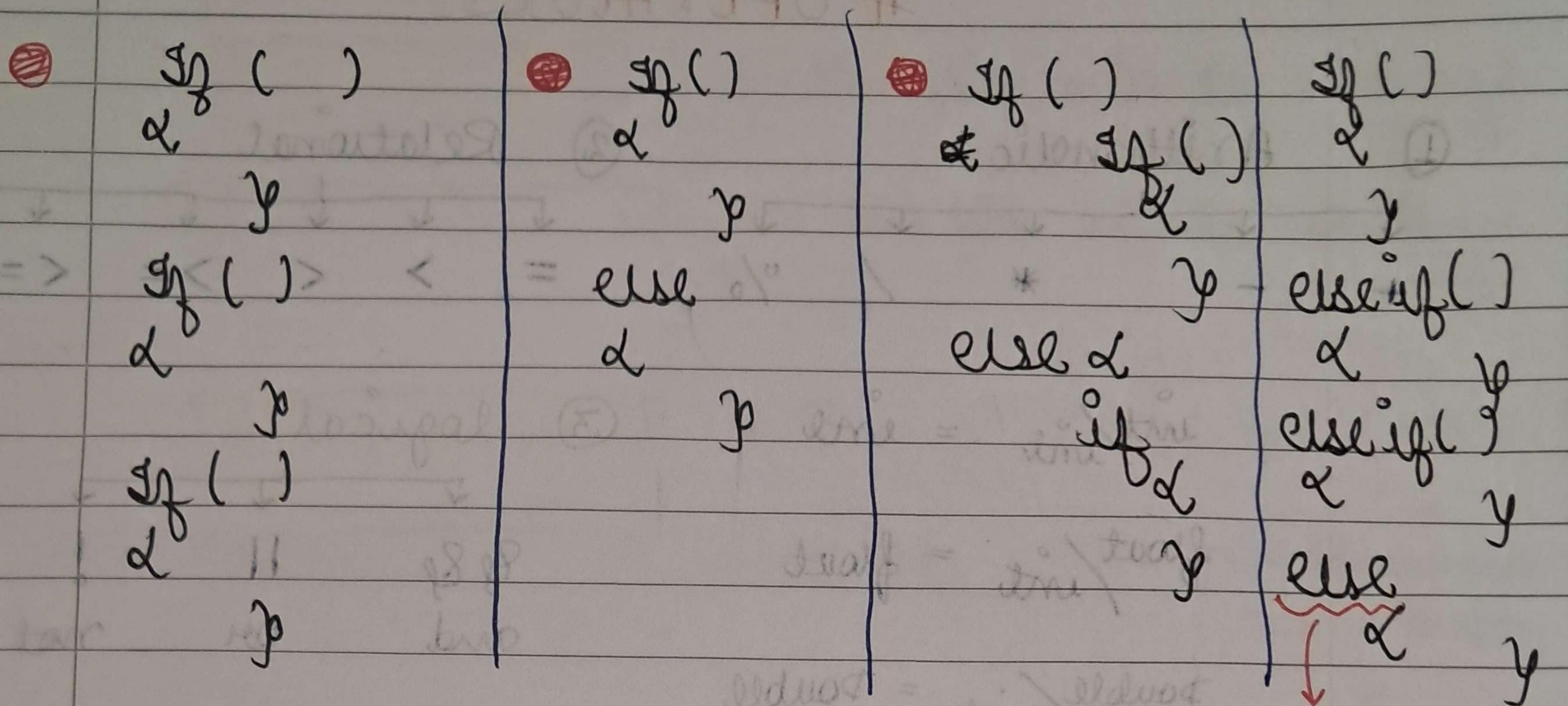
# # Conditionals

If (✓) &  
y —

else &

If (✓) &  
y —  
else &  
y —

else if () &  
y —



## # While loop.

while ( ) WTF nha True han

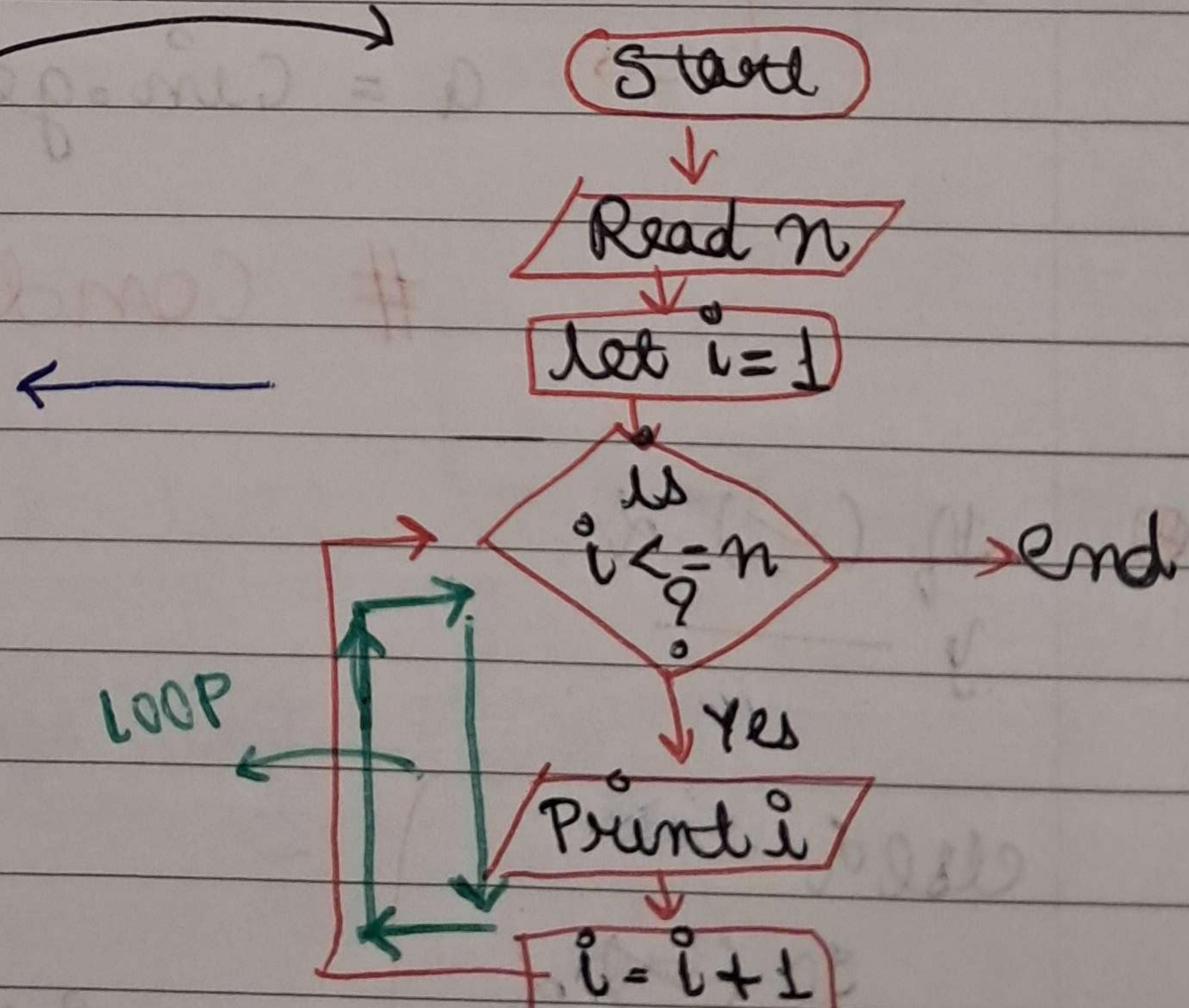
int n;  
cin >> n;  
int i=1

while (i <= n)

    cout << i;

    i = i + 1;

y



## Pattern Ques (Grif on loop)

★ ★ ★ ★  
★ ★ ★ ★  
★ ★ ★ ★  
★ ★ ★ ★

int  $i = 1$ ;

while ( $i < n$ ) {

    y  $i = 1, 2, 3$

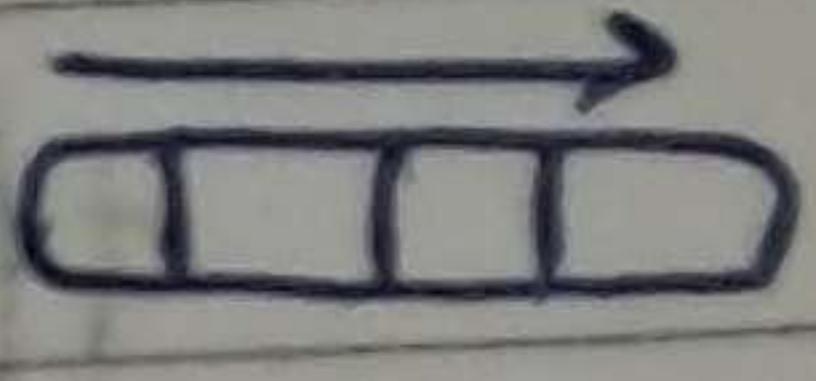
    84

    int  $i = 0$ ;

    while ( $i < n$ ) {

        y  $i = 0, 1, 2$

        LEC-4

Observation 1  $\rightarrow$  row = 4 

Observation 2  $\rightarrow$  first row mai kitne column hai

$\downarrow$  1st row mai 4 column hai

    2nd ----- 4

    3rd ----- 4

    4th ----- 4

Now find rem 3/w Row & column

$\hookrightarrow$  here;  $\rightarrow$  Title Row = column

$\hookrightarrow$  Lec 4 have 20 (appuon) pattern find problem

    I have write the approach  
    to solve them in  
    the code itself

$\hookrightarrow$  using //Comment

## LES - 5

### # Bitwise Operators

AND      OR      NOT      XOR  
 &      |      ~      ^  
 &      |      ~      ^

① AND  $\rightarrow 8f$

$$a = 2, b = 3$$

$$a \& b \rightarrow 2 \rightarrow 10$$

$$\begin{array}{r} 1 \\ 0 \\ \hline 10 \end{array}$$

$$0/f \rightarrow 2$$

i/p		0/f
x	y	z
0	0	0
0	1	0
1	0	0
1	1	1

② OR  $\rightarrow 1$

$$a = 2$$

$$a \mid b$$

$$b = 4$$

$$\hookrightarrow 10$$

$$\begin{array}{r} 100 \\ 110 \end{array}$$

$$\rightarrow 0/f \rightarrow 6$$

i/p		0/f
x	y	z
0	0	0
0	1	1
1	0	1
1	1	1

③ NOT  $\sim$

$$a = 2$$

$$\sim a = -3$$

x	z
0	1
1	0

$0000000-0010$

$\sim a$

$\sim a = 111111111101$

$-3$

$$1^{\text{st}} \text{ Comp} \rightarrow 000000010$$

$$2^{\text{nd}} \text{ Comp} \rightarrow \begin{array}{r} +1 \\ \hline 000001011 \end{array}$$

$$\rightarrow 3$$

④  $\text{XOR} \rightarrow \wedge$

0/p		0/p
x	y	z
0	0	0
0	1	1
1	0	1
1	1	0

Both bits → Similar → 0  
Both bits → Disimilar → 1

eg of  $\text{XOR} (\wedge)$

$$a = 2 \rightarrow 010$$

$$b = 4 \rightarrow 100$$

$$\begin{array}{r} 110 \\ \hline 0/p \rightarrow 6 \end{array}$$

⑤ Left Shift ( $<<$ )

$$5 << 1$$

$$\begin{array}{r} 00000 \dots 00101 \\ \hline 00000 \dots 001010 \end{array}$$

padding  
0/p → 10

left shift  
by 1 bit

⑥ Right Shift ( $>>$ )

$$5 >> 2$$

$$\begin{array}{r} 000000 \dots 00101 \\ \hline 000000 \dots 001 \end{array}$$

padding  
0/p → 1

Right shift  
fitne Basg shift utne  
Bar dide 2 se

$$3 << 2$$

$$\begin{array}{r} 0000 \dots 0011 \\ \hline 0000 \dots 001100 \end{array}$$

0/p → 12

⑦ Big no → -ve  
small → Double

$$5 \times 2 \rightarrow 10$$

$$3 \times 2 * 2 \rightarrow 12$$

} Summary

$$5 >> 1 \rightarrow 5/2 = 2$$

$$5 >> 2 \rightarrow 5/2 * 2 = 5/4 = 1$$

⑧ left shift / Right Shift

padding with 0 (+ve)

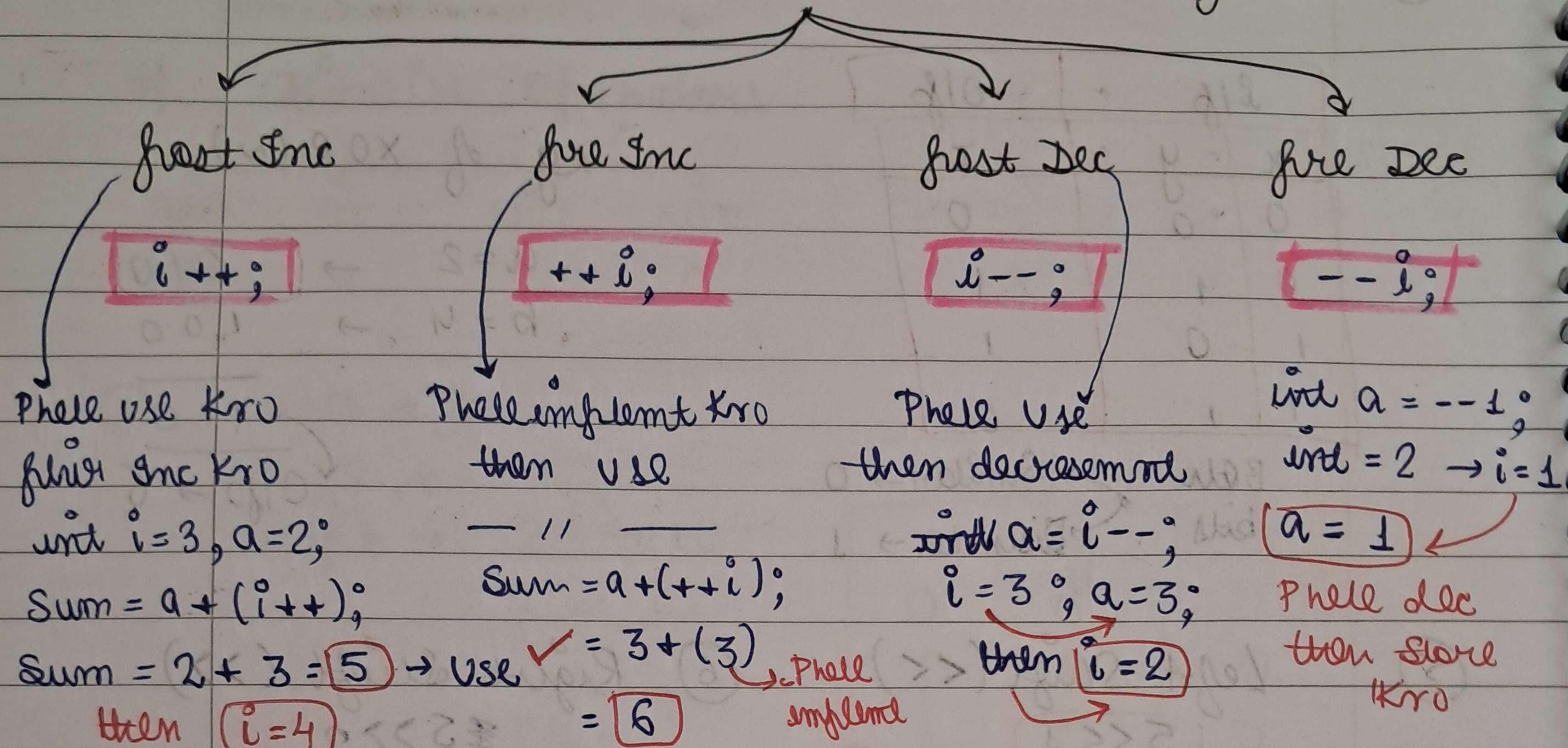
(-ve) no → padding

computer dependent

①  $i = i + 1$   
 $i + = 1;$

②  $\text{Sum} = \text{Sum} + b$   
 $\text{Sum} + = b;$

## # Increment and Decrement Operators



## # for loop

for (int  $i = 0$  ;  $i \leq n$  ;  $i++$ )  $\alpha$   
 initialisation      CondN      Updating  
 \_\_\_\_\_      \_\_\_\_\_      inc/dec

for (   ;   )  $\alpha$   
 \_\_\_\_\_      infinite loop

## # Multiple CondN check

for (int  $a = 0, b = 1$  ;  $a = 0, b \geq 1$  ;  $a--, b--$ )  $\alpha$   
 cout << a << " " << b;

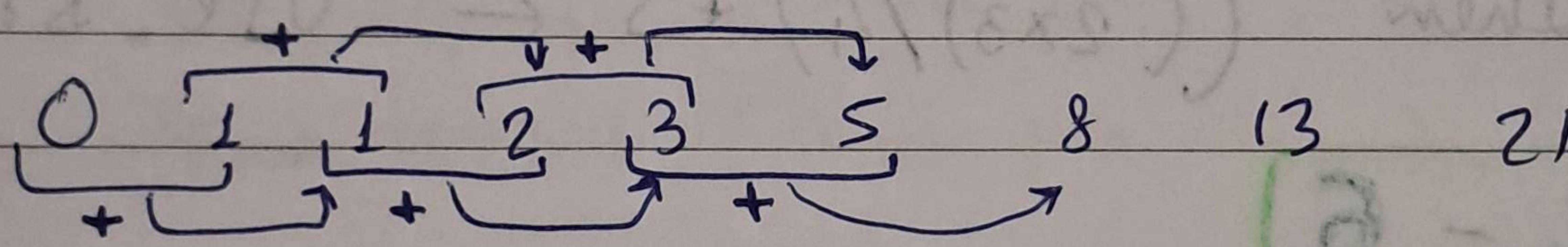
```

    with i = 1;
    for (; ; ) {
        if (i <= n) {
            cout << i << endl;
        }
        else {
            break;
        }
        i++;
    }

```

→ Mention ~~else~~ - then  
do loop  
get out of current loop

### fibonacci Series



$$n = (n-1) + (n-2)$$

# Continue → Iteration Skip  
→ Continue ke niche wala part unreachable hogaya

# Scope → The lifetime of a variable, where does the variable exist and after what line of code will get destroyed.

a is accessible  
throughout the  
main function  
after it declared

```

int main() {
    int a = 3;
    cout << a;
    if (+ve) {
        cout << a;
    }
}

```

```

int main() {
    int a = 3;
    cout << a; → local to
}
if (true) {
    int a = 5;
    cout << a; → local to if block, Only
}

```

accessible within the if block

a

b } Two same  
b } variable under  
Same scope (X)

y

a

b  
d  
b

y

Want this to be used

Do not create  
this one

strictly say  
cont < b,

u will get  
upper b.

#

Operator Precedence → like BodMas

Brackets, we can also performe calcult using  
Brackets

$$\textcircled{1} \quad 2 \times 3/4 + 5 \rightarrow 2 \times 0 + 5$$

= 5

$$\textcircled{2} \quad 2 \times \frac{3}{4} + 5 = \frac{6}{4} + 5 = 6 \quad \Rightarrow \neq$$

Then  $((2 \times 3)/4) + 5 \leftarrow$  use Brackets

## Lec - 5

### Decimal to Binary

① Divide by 2

2 → 10

② Store remainder in answer

5 → 101

③ Repeat above 2 step  
until  $n \neq 0$

10 → 1010

④ Read Ans in Reverse  
order

8 → 1000

$n = 10$

Divide

$$10/2 = 5$$

Rem

0

$$5/2 = 2$$

1

Remainder

$$2/2 = 1$$

0

$$1/2 = 0$$

1

→ 1010