

# 16. Random Walks

## Markov Chain

A generalized random walk on a directed graph.

$$G = (V, E)$$

$$|V| = n, |E| = m$$

We have a *transition probability matrix*  $P_{n \times n}$

$P_{ij}$  : Probability that  $j$  is the next state given  $i$  is the current state.

$X_t$  : State of the markov chain at time  $t$ .

$p_i(t)$  : Probability of being at state  $s_i$  at time  $t$

Using this, we can define a probability vector

$$p_t = (p_1(t), p_2(t), \dots, p_n(t))$$

### Memory-Less

$$\mathbb{P}[X_t = a_t | X_{t-1} = a_{t-1}, \dots, X_0 = a_0] = \mathbb{P}[X_t = a_t | X_{t-1} = a_{t-1}] = P_{a_{t-1}, a_t}$$

Hence, the future behaviour of the Markov Chain is only dependent upon the current state.

we know that

$$\begin{aligned} p_{t+1} &= p_t P \\ \implies p_{t+m} &= p_t \cdot P^m \end{aligned}$$

### Irreducible Markov Chain

A markov chain is irreducible if  $G$  is strongly connected,  $\forall s_i \neq s_j, \exists$  a directed path from  $s_i \rightarrow s_j$ .

$$\forall i \neq j \quad \mathbb{P}[\exists m > 0, X_{t+m} = s_j | X_t = s_i] > 0$$

### Period of a State

We define the period of a state  $s_i$  in a markov chain as

$$d(s_i) = \gcd(\{t \in \mathbb{N} | (P^t)_{ii} > 0, t > 0\})$$

Essentially, we find the time periods at which the probability of being at position  $i$  is non-zero, and find the gcd of the time periods to get the period of the state.

Say we get the period as  $d(s_i) = T$  for some state, and at time  $t$ , we get  $p_t(i) > 0$ , then  $t$  belongs to the AP

$$\{a + i \cdot T | i \geq 0\}$$

### Aperiodic -

A markov chain is aperiodic if  $d(s_i) = 1 \forall s_i \in \{s_1, s_2, \dots, s_n\}$

### Other Properties

$p_t(i, j) = \mathbb{P}[X_t = s_j | X_0 = s_i]$  (Probability of being at  $s_j$  at time  $t$  given you start at  $s_i$ )

$r_t(i, j) = \mathbb{P}[X_t = s_j \text{ and } \forall i \leq u \leq t-1, X_u \neq s_j | X_0 = s_i]$  (First visit to  $s_i$  from  $s_j$ )

Starting from  $s_i$  probability of a transition into state  $s_j$  is given by

$$f_{ij} = \sum_{t>0} r_t(i, j)$$

Expected number of time steps to reach  $s_j$  starting from  $s_i$ .

$$h_{ij} = \sum_{t>0} t \cdot r_t(i, j)$$

**Theorem** - If a markov chain is finite, irreducible and aperiodic, then

$\exists T < \infty$  such that  $(P^t)_{ij} > 0, \forall t \geq T, \forall i, j \in V, P_{t+1} = P_t$ .

**Stationary Distribution** - A probability distribution  $\pi$  for a markov chain such that

$$\pi = \pi P$$

## Fundamental Theorem of Markov Chains

For any finite, irreducible and aperiodic markov chain,

1. There exists a unique stationary distribution  $\pi$  such that  $\forall 1 \leq i \leq n, \pi_i > 0$
2.  $\forall 1 \leq i \leq n, f_{ii} = 1, h_{ii} = \frac{1}{\pi_i}$
3. Let  $N(i, t)$  be the number of time the markov chain visits state  $i$  in  $t$  steps, then  

$$\lim_{t \rightarrow \infty} \frac{N(i, t)}{t} = \pi_i$$

## Random Walks

We define the transition probability matrix as

$$P_{u,v} = \begin{cases} \frac{1}{deg(u)} & \text{if } (u,v) \in E \\ 0 & \text{otherwise} \end{cases}$$

Given the graph  $G = (V, E)$ , we work with a non-bipartite, undirected & connected graph.

**Claim** -  $G$  is aperiodic

**Proof** -

We know cycles of length 2 exist already for every state. To ensure the period of every state is 1, we just need to prove every state has an odd cycle as well.

What non-bipartiteness ensures is that there exists a cycle in the graph with an odd number of nodes. Any state, whether inside the cycle or outside, can use this to create a cycle which is of odd length.

Going and coming back to the cycle for a state have the same length, hence making it even as we do it twice, and taking the cycle once gives us an odd length path.

**Lemma** -  $\forall v \in V, \pi_v = \frac{deg(v)}{2m}$

**Proof** -

$\pi$  is stationary if  $\pi = \pi P$ . Therefore,  $\forall v \in V$

$$\pi_v = (\pi P)_v = \sum_u \pi_u P_{uv}$$

We are claiming that  $\pi_u = \frac{deg(u)}{2m}$  is a solution.

$$\sum_u \pi_u P_{uv} = \sum_{u, (u,v) \in E} \frac{deg(u)}{2m} \cdot \frac{1}{deg(u)} = \frac{deg(v)}{2m} = \pi_v$$

Hence proved.

And from the fundamental theorem, we know this stationary distribution is unique.

## Pagerank Algorithm

1. Each node in the graph is assigned a pagerank value of  $\frac{1}{n}$ .
2. Each node divides their page rank value by their out-degree, and send it to each neighbour.

3. Pagerank of each node is updated to sum of all values received in the previous iteration.

Can be modelled as a markov chain, with the transition probability matrix

$$P_{ij} = \begin{cases} \frac{1}{d_{\text{out}}(i)} & \text{if page } i \text{ links to } j \\ 0 & \text{otherwise} \end{cases}$$

$$\text{pagerank}^t(j) = \sum_{i, (i,j) \in E} \frac{\text{pagerank}^{t-1}(i)}{d_{\text{out}}(i)}$$

If the web is finite, irreducible and aperiodic, we can apply the fundamental theorem to reach a unique stationary distribution.