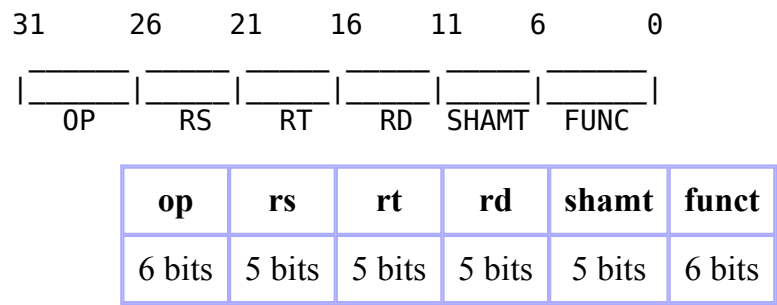


# A Minimalistic Introduction to MIPS Instruction

## General Format



**op**  
Operation code

**rs**  
First source register operand

**rt**  
Second source register operand

**rd**  
Destination register operand

**shamt**  
Shift amount - used in shift instructions

**funct**  
Select the variant of the operation in the op code field

## Specific Instruction Formats

| Format | 6 bits | 5 bits         | 5 bits | 5 bits            | 5 bits | 6 bits | Comments                   |
|--------|--------|----------------|--------|-------------------|--------|--------|----------------------------|
| R      | op     | rs             | rt     | rd                | shamt  | funct  | Arithmetic                 |
| I      | op     | rs             | rt     | address/immediate |        |        | Transfer, branch,immediate |
| J      | op     | target address |        |                   |        |        | Jump                       |

# MIPS Instruction Set

The MIPS instruction set illustrates four underlying principles of hardware design:

1. Simplicity favors regularity.
2. Smaller is faster.
3. Good design demands compromise.

4. Make the common case fast.

### Simplicity favors regularity

Consider the following example:

| Category   | Instruction | Example   | Meaning | Comments          |
|------------|-------------|-----------|---------|-------------------|
| Arithmetic | add         | add a,b,c | $a=b+c$ | Always 3 operands |
| Arithmetic | subtract    | sub a,b,c | $a=b-c$ | Always 3 operands |

Note that each operand has *exactly* three operands.

### Smaller is faster.

MIPS has 32 32-bit registers,  $\$v0, \dots, \$v31$ , a very large number would increase the clock cycle time.

### Good design demands compromise.

The compromise represented by the MIPS design, was to make all the instructions the *same* length, thereby requiring different instruction formats.

### Make the common case fast.

The MIPS instruction set addresses this principal by making constants part of arithmetic instructions. Furthermore, by loading small constants into the upper 16-bits of a register.

## MIPS Instruction Set Summary

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Note the summary is **not** complete. Click here to see the [full list](#)

### Arithmetic Instructions

| Instruction        | Example          | Meaning           | Comments            |
|--------------------|------------------|-------------------|---------------------|
| add                | add \$1,\$2,\$3  | $\$1 = \$2 + \$3$ | Always 3 operands   |
| subtract           | sub \$1,\$2,\$3  | $\$1 = \$2 - \$3$ | Always 3 operands   |
| add immediate      | addi \$1,\$2,10  | $\$1 = \$2 + 10$  | add constant        |
| add unsigned       | addu \$1,\$2,\$3 | $\$1 = \$2 + \$3$ | Always 3 operations |
| subtract unsigned  | subu \$1,\$2,\$3 | $\$1 = \$2 - \$3$ | Always 3 operations |
| add immed.unsigned | addiu \$1,\$2,10 | $\$1 = \$2 + 10$  | Always 3 operations |

### Logical

| Instruction   | Example         | Meaning            | Comments            |
|---------------|-----------------|--------------------|---------------------|
| and           | and \$1,\$2,\$3 | $\$1 = \$2 \& \$3$ | 3 register operands |
| or            | or \$1,\$2,\$3  | $\$1 = \$2   \$3$  | 3 register operands |
| and immediate | andi \$1,\$2,10 | $\$1 = \$2 \& 10$  | AND constant        |

|                     |                |             |                         |
|---------------------|----------------|-------------|-------------------------|
| or immediate        | or \$1,\$2,10  | \$1=\$2 10  | OR constant             |
| shift left logical  | sll \$1,\$2,10 | \$1=\$2<<10 | Shift left by constant  |
| shift right logical | srl \$1,\$2,10 | \$1=\$2>>10 | Shift right by constant |

## Data Transfer

| Instruction       | Example        | Meaning                | Comments                         |
|-------------------|----------------|------------------------|----------------------------------|
| load word         | lw \$1,10(\$2) | \$1=Memory[\$2+10]     | memory to register               |
| store word        | sw \$1,10(\$2) | Memory[\$2+10]=\$1     | register to memory               |
| load upper immed. | lui \$1,10     | \$1=10x2 <sup>16</sup> | load constant into upper 16 bits |

## Conditional Branch

| Instruction         | Example         | Meaning                     | Comments          |
|---------------------|-----------------|-----------------------------|-------------------|
| branch on equal     | beq \$1,\$2,10  | if(\$1==\$2)go to PC+4+10   | Equal test        |
| branch on not equal | bne \$1,\$2,10  | if(\$1!=\$2)go to PC+4+10   | Not equal test    |
| set on less than    | slt \$1,\$2,\$3 | if(\$2<\$3)\$1=1;else \$1=0 | Less than compare |

## Unconditional Jump

| Instruction   | Example  | Meaning              | Comments                     |
|---------------|----------|----------------------|------------------------------|
| jump          | j 1000   | go to 1000           | Jump to target address       |
| jump register | jr \$31  | go to \$31           | For switch, procedure return |
| jump and link | jal 1000 | \$31=PC+4;go to 1000 | For procedure call           |

## Simple Output using MIPS I/O

Let us examine this simple program to introduce MIPS Assembler.

```
# hello.asm
# Simple routine to demo MIPS output.
# Author: R.N. Ciminero
# See Patterson & Hennessy pg. A-46 for system services.
```

```
        .text
        .globl  main
main:
    li     $v0,4          # code for print_str
    la     $a0, msg       # point to string
```

```

        syscall
        li      $v0,10          # code for exit
        syscall

        .data
msg:     .asciiz "Hello World!\n"

```

---

## Output

Hello world!

---

Comments on the program:

### # **hello.asm**

A single line *comment*, comments may also appear to the to the right of regular assembly language statements.

### **.text**

Begin program text segment.

### **.globl main**

Define a global function main.

### **main:**

A *label* **main**.

### **li \$v0,4 # code for print\_str**

Load immediate the system register,\$*v0*, with a value of four, which indicates that we intend to output an ASCII string.

### **la \$a0, msg # point to string**

Load the system register,\$*a0*, with the address of the string to be output.

### **syscall**

System routine that performs the string output based on the contents of the system registers.

### **li \$v0,10 # code for exit**

Load the system code for exit.

### **.data**

Begin the data segment of the program.

### **msg: .asciiz "Hello World!\n"**

The label *msg:* is the location of the ASCII string.

## Simple Input

---

Let us examine the following program for simple *input*.

```
# simpleio.asm
```

```

# Simple routine to demo SPIM input/output.
# Author: R.N. Ciminero
# Revision date: 10-05-93 Original def.
# See Patterson & Hennessy pg. A-46 for system services.

```

```

        .text
        .globl  main
main:
    li      $v0,4          # output msg1
    la      $a0, msg1
    syscall
    li      $v0,5          # input A and save
    syscall
    move     $t0,$v0
    li      $v0,4          # output msg2
    la      $a0, msg2
    syscall
    li      $v0,5          # input B and save
    syscall
    move     $t1,$v0
    add      $t0, $t0, $t1  # A = A + B
    li      $v0, 4         # output msg3
    la      $a0, msg3
    syscall
    li      $v0,1          # output sum
    move     $a0, $t0
    syscall
    li      $v0,4          # output lf
    la      $a0, cflf
    syscall

    li      $v0,10         # exit
    syscall
        .data
msg1:    .asciiz "\nEnter A:  "
msg2:    .asciiz "\nEnter B:  "
msg3:    .asciiz "\nA + B =  "
cflf:    .asciiz "\n"

```

---

## Output

```

Enter A: 10
Enter B: 5
A + B = 15

```

---

Comments on the program:

### **li \$v0,5 # input A and save**

Load the system register, *\$v0*, with the input code.

### **syscall**

System call to accept an *integer* from the keyboard and store it in register *\$v0*.

### **move \$t0,\$v0**

Macro to transfer the contents of system register *\$v0* to temporary register *\$t0*.

### **add \$t0, \$t0, \$t1 # A = A + B**

Add the contents of *\$t0* to the contents of *\$t1* and place the sum in *\$t0*.

## li \$v0,1 # output sum

Load the system register \$v0 with the output *integer* code.

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# Looping Structures

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Let us examine the following program for simple *looping* structures.

```
# sumit.asm
# Simple routine to sum N integers to demo a loop.
# Author: R.N. Ciminero
# Revision date: 10-06-93 Original def.
# See Patterson & Hennessy pg. A-46 for system services.
```

```
        .text
        .globl  main
main:
        li      $v0,4          # output msg1
        la      $a0, msg1
        syscall
        li      $v0,5          # input N and save
        syscall
        move    $t0,$v0

        li      $t1, 0         # initialize counter (i)
        li      $t2, 0         # initialize sum

loop:   addi     $t1, $t1, 1     # i = i + 1
        add     $t2, $t2, $t1   # sum = sum + i
        beq     $t0, $t1, exit  # if i = N, continue
        j       loop

exit:   li      $v0, 4          # output msg2
        la      $a0, msg2
        syscall

        li      $v0,1          # output sum
        move    $a0, $t2
        syscall
        li      $v0,4          # output lf
        la      $a0, lf
        syscall
        li      $v0,10         # exit
        syscall
        .data
msg1:   .asciiz  "\nNumber of integers (N)?  "
msg2:   .asciiz  "\nSum  =  "
lf:     .asciiz  "\n"
```

---

## Output

```
Number of integers (N)? 5
Sum = 15
```

---

Comments on the program:

**li \$t1, 0 # initialize counter (i)**

Temporary register *\$t1* contains the count.

**li \$t2, 0 # initialize sum**

Temporary register *\$t2* contains the sum.

**loop: addi \$t1, \$t1, 1 # i = i + 1**

Increment the counter by one.

**add \$t2, \$t2, \$t1 # sum = sum + i**

Add the counter to the sum.

**beq \$t0, \$t1, exit # if i = N, continue**

If the counter equals the number of integers, then exit the loop.

**j loop**

Else perform the summation again.

**exit: li \$v0, 4 # output msg2**

Statement to execute upon leaving the loop.