Gajendra Processor

Lab Report

submitted by

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Under the supervision

 \mathbf{of}

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1 General CPU Register

1.1 Diagram

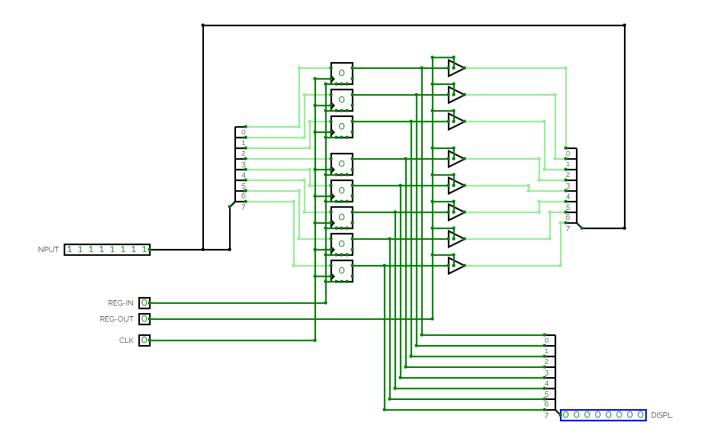


Figure 1: 8-bit Register

1.2 I/O Control

The register consists of a bidirectional common ' I/O' line which stores the 8 bit input from common bus. The 8 bit value can later be sent to output. The flow of data is controlled by the signals REG_in and REG_out.

1.3 Working

The register consists of D flip flops, responsible for the storage of memory. REG_acts as the control signal for input, when it is 1 the input is read from the common bus while REG_out is the control signal for output in the same way.

2 Program Counter

2.1 Diagram

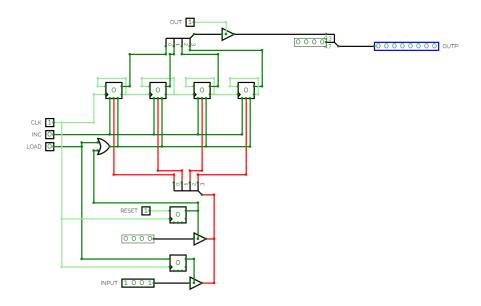


Figure 2: Program Counter

2.2 Working

Program counter is used to execute the program. It is a 4 bit data signal used to access the memory locations of the ROM. ROM consists of instructions and data of the program that has to be run . The counter starts the program at 0000 thereby accessing the instruction set in that location of the ROM via MAR. This is stored in Instruction Register which executes the program. The counter keeps incrementing by 1 and executes all the instructions in order. This cycle keep continuing until the program ends.

The following are the functionalities that are present in the Program Counter.

2.3 PC_INC

If PC INC = 1 then the value of the address stored in the Program Counter is incremented by 1 in the next positive trigger of the clock.

2.4 PC_OUT

If PC_OUT = 1 then the Program Counter outputs the address stored in it into the **common** bus of the **Gajendra Processor** as Least Significant Bits (0-3).

2.5 PC_LOAD

If $PC_LOAD = 1$ then the Program Counter takes the Least Significant Bits (0-3) of the **common bus** into it through the **input** and updates the counter to that value (which is the next instruction for the ROM).

3 MEMORY

3.1 Diagram

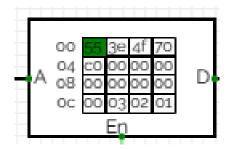


Figure 3: Memory

Memory is the ROM which contains both instruction and data required for a program. The following are the functionalities that are present in the Memory.

3.2 MEM_OUT

The ROM's output is controlled by a tri-state buffer. MEM _OUT acts as the enable line for this tri-state buffer. If MEM_OUT = 1 then the memory outputs a 8-bit output into the common bus.

4 MEMORY ADDRESS REGISTER

During a computer run , the 4-bit address in the Programming Counter is latched into the MAR .Then the MAR applies this 4-bit address to the Memory where a read operation is performed.

The following are the functionalities that are present in the MAR .

4.1 Diagram

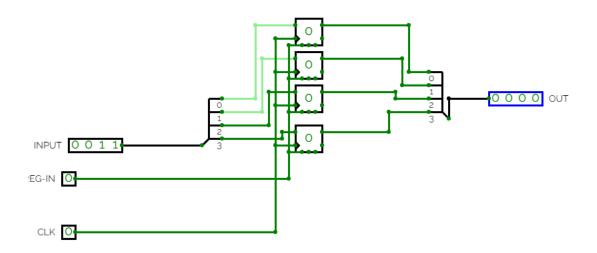


Figure 4: Memory Address Register

4.2 MAR IN

If MAR $_{
m IN}=1$, then the MAR takes the 4-bit address (Least Significant in common bus) from the Program Counter through common bus.

5 INSTRUCTION REGISTER

5.1 Diagram

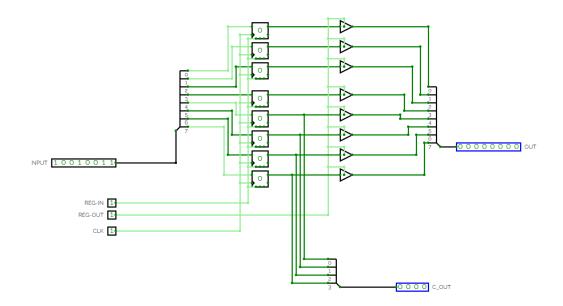


Figure 5: Instruction Register

To fetch an instruction from the memory the computer does a memory read operation . This places the contents of the addressed memory location on the common bus. The instruction register loads this content into it on the next positive clock edge .

The contents of the instruction register are split into two nibbles of 4-bits each (least 4 significant bits , most 4-significant bits) , The most 4-significant bits are sent into **Controller** and it will also send the 8-bit content into the **Instruction decoder** .

The following are the functionalities that are present in the Instruction register.

5.2 INS_REG_IN

If INS_REG $_{-}$ IN = 1 , then the INS register takes the 8-bit content as input from the common bus .

5.3 INS_REG_OUT

If INS_REG $_{-}$ OUT = 1 , then INS register loads its 8-bit content into Instruction decoder and thus into common bus.

5.4 CTRL

The Most Significant 4-bits(4-7) load in to the Controller and INS Decoder.

6 INSTRUCTION DECODER

6.1 Diagram

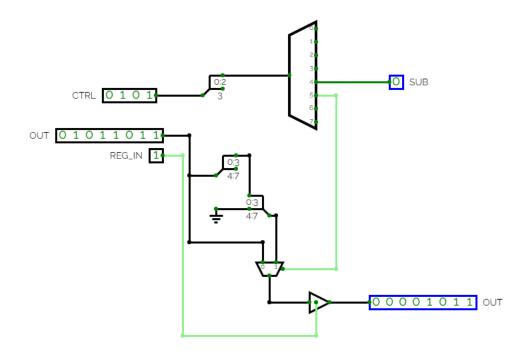


Figure 6: Instruction Decoder

It helps the ALU by giving the instruction whether to add or subtract, by using the most significant 4-bits (i.e CTRL). It also helps us with the instruction **LDI** by making the 4-most significant bits into 0000, so that the remaining 4-bits gets stored in our **accumulator** (i.e register A).

The following are the functionalities that are present in the Instruction Decoder.

6.2 INS _DEC _OUT

If INS $_\mathrm{DEC}\ _\mathrm{OUT}=1$, then it takes the 8-bit content from OUT of INS register .

6.3 ALU_OUT

The INS_DEC generates a 1-bit input for the **ALU**.

It generates the following outputs for the corresponding instructions .

 $0 \implies Addition$

 $1 \implies Subtraction$

7 Arithmetic and Logic Unit(ALU)

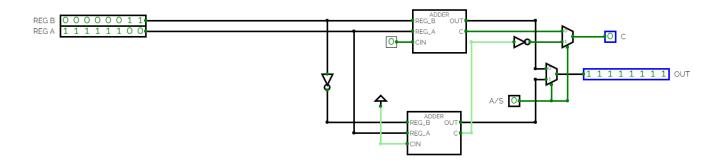


Figure 7: Gajendra ALU

The ALU performs the arithmetic operations, executed by the computer. Supported operations are Addition(ADD) and Subtraction(SUB). The final 8-bit data is then loaded into Accumulator or REG $_A$.

7.1 ADD

The ADD operation is used to add two 8-bit numbers using a 8-bit full adder which adds without carry .

7.2 SUB

The SUB operation is used to subtract two 8-bit numbers using a 8-bit full adder with initial carry in as 1. This one is supplied by INS_DEC.

8 Status Register

8.1 Diagram

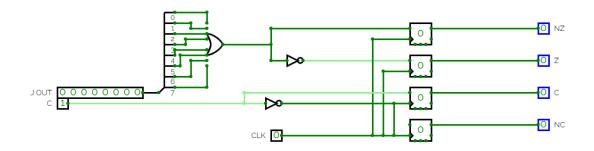


Figure 8: Status Register

8.2 I/O

It takes 2 inputs ALU $_$ OUT and C (carry out from ALU) and gives four outputs namely: C (1 if carry = 1), NC(negation of C), Z(1 if(ALU $_$ OUT = 1) and NZ (negation of Z). Tise NZ line is used as input to the controller to aid in the JNZ function.

8.3 Uses

It is used to implement JUMP and JUMP if not zero conditions.

9 Control Processor

9.1 Diagram

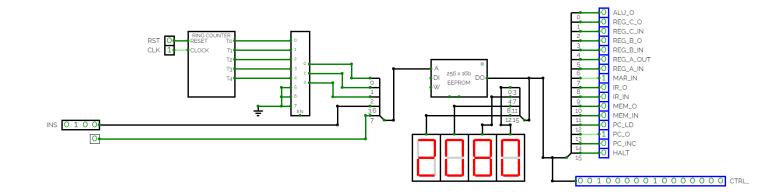


Figure 9: Control Processor

9.2 I/O

The control processor takes 4 bit instruction from INS_REG .We use a ring counter of length 5 indicating the 5 T-states. The EEPROM in the control processor takes 8-bit address is split into 3 parts. The first bit zero and next 4 bits are the instruction and the last 3 bits are the T-state that is to be implemented. The EEPROM gives 16-bit output which is the control word which is also the output of the Control Processor.

9.3 Working

The control word consists of all inputs of the program which are: REG_in, REG_out, PC_load, PC_out, PC_inc, flag, MEM_in, MEM_out, MAR_in, INS_out, INS_out and ALU_out. It controls all the other components of the program and is hence the most important component of the program .

9.4 Ring Counter

Every instruction has micro instructions called 'T states' that are needed to be implemented in a particular sequence one after another. This is done using a ring counter. We use a ring counter of length 5 bits in the program which governs the micro instructions.

9.5 EEPROM Programming

Each memory location in EEPROM contains 8 bits necessary to execute its micro instructions. The each memory location contains the control word of that particular micro instruction.

10 Micro instructions and Controller Logic Design

10.1 NOP:

Figure 10: Micro Instruction for NOP

10.2 LDA:

```
LDA :1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<IR_OUT|1<<MAR_IN T2

1<<MEM_OUT|1<<REGA_IN T3

0 T4
```

Figure 11: Micro instruction for LDA

10.3 ADD:

```
      STA :1<<PC_OUT|1<<MAR_IN</td>
      T0

      1<<PC_INC|1<<MEM_OUT|1<<IR_IN</td>
      T1

      1<<MEM_IN|1<<REGA_OUT</td>
      T2

      0
      T3

      0
      T4
```

Figure 12: Micro instruction for ADD

10.4 SUB:

```
      SUB :1<<PC_OUT|1<<MAR_IN</td>
      T0

      1<<PC_INC|1<<MEM_OUT|1<<IR_IN</td>
      T1

      1<<IR_OUT|1<<MAR_IN</td>
      T2

      1<<MEM_IN|1<<REGA_OUT</td>
      T3

      0
      T4
```

Figure 13: Micro instruction for SUB

10.5 LDI:

```
LDI :1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<IR_OUT|1<<REGA_IN T2

0 T3

0 T4
```

Figure 14: Micro instruction for LDI

10.6 JMP:

```
JMP :1<<PC_OUT|1<<MAR_IN T0

1<<PC_INC|1<<MEM_OUT|1<<IR_IN T1

1<<IR_OUT|1<<PC_LOAD T2

0 T3

0 T4
```

Figure 15: Micro instruction for JMP

10.7 JNZ:

```
JNZ :1<<PC_OUT|1<<MAR_IN
                                                       T0
    1<<PC_INC|1<<MEM_OUT|1<<IR_IN
                                                       T1
    1<<IR_OUT | 1<<PC_LOAD
                                                       T2
    0
                                                       Т3
    0
                                                       T4
```

Figure 16: Micro instruction for JNZ

11 Instruction Set

NOP: 11.1

Description

Loads the Instruction Register and increments the Program counter. No other operation is carried out.

Operation (i) None Operands: **Program Counter** Syntax: $PC \leftarrow PC + 1$ (i) NOP X None 8-bit op-code: 0000 XXXX

Figure 17: Instruction Set for NOP

11.2 LDA:

Description

Loads Register A with the value of the common bus (this value is obtained from the memory).

Operation

(i) $R_A \leftarrow (X)$

Syntax: Operands: Program Counter

(i) LDA X $0 \le X \le 15$ PC \leftarrow PC + 1

8-bit op-code:

0001 0xX(xxxx)

Figure 18: Instruction Set for LDA

11.3 ADD:

Description

Add the value at the given address to accumulator.

Operation

(i) $R_B \leftarrow (X)$

(ii) $R_A \leftarrow R_A + R_B$

Syntax: Operands: Program Counter

(i) ADD X $0 \le X \le 15$ PC \leftarrow PC + 1

8-bit op-code:

0011 0xX(xxxx)

Figure 19: Instruction Set for ADD $\,$

11.4 SUB:

Description

(i)

Subtract the value at the given address from the accumulator.

Operation

(i) $R_B \leftarrow (X)$ (ii) $R_A \leftarrow R_A - R_B$

Syntax: Operands: Program Counter

ADD X $0 \le X \le 15$ $PC \leftarrow PC + 1$

8-bit op-code:

0100 0xX(xxxx)

Figure 20: Instruction Set for SUB

11.5 LDI:

Description

(i)

Loads the value given along with the instruction to the Register A.

Operation

(i) $R_A \leftarrow X$

Syntax: Operands: Program Counter

LDI X $0 \le X \le 15$ PC \leftarrow PC + 1

8-bit op-code:

0101 0xX(xxxx)

Figure 21: Instruction Set for LDI

11.6 JMP:

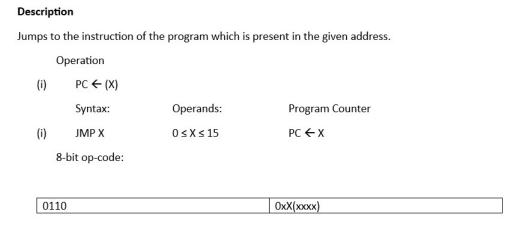


Figure 22: Instruction Set for JMP

12 Sample Tasks

12.1 For Addition

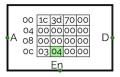


Figure 23: Rom data for addition of numbers 3,4

It first loads the data which is stored in the address 0x0c (i.e 3) and then as the program counter is increased it will take the instruction addition stored in the address 0x01 and load the number stored in the memory address 0x0d (i.e 4), then after it takes the instruction out and gives the output 7 at the display of the register C.

12.2 For Subtraction

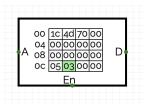


Figure 24: Rom data for subtraction of 2 from 5

It first loads the data which is stored in the address 0x0c (i.e 5) and then as the program counter is increased it will take the instruction subtraction stored in the address 0x01 and load

the number stored in the memory address $0x0d(i.e\ 3)$, then after it takes the instruction out and gives the output 2 at the display of the register C.

12.3 For both addition and subtraction

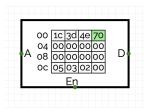


Figure 25: Rom data for 5+3-2