

# COMP9020

## Foundations of Computer Science

# COMP9020 19T3 Staff

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Lectures: Mondays and Wednesdays 4-6pm  
Consults: Wednesdays, 2-3pm, Rm204 K17  
Research: Theoretical CS: Algorithms, Formal verification

# What you can expect from me

# What is this course about?

## What is Computer Science?

*“Computer science no more about computers than astronomy is about telescopes”*

– E. Dijkstra



# Course Aims

The actual content is taken from a list of subjects that constitute the basis of the tool box of every serious practitioner of computing:

- |                                 |           |
|---------------------------------|-----------|
| • number theory                 | week 1    |
| • sets, relations and functions | weeks 2–3 |
| • big-O notation                | week 3    |
| <hr/>                           |           |
| • recursion                     | week 4    |
| • graph theory                  | week 5    |
| • logic                         | week 6    |
| • induction                     | week 7    |
| <hr/>                           |           |
| • combinatorics                 | week 8    |
| • probability and expectation   | week 9    |

# Course Material

All course information is placed on the course website

[www.cse.unsw.edu.au/~cs9020/](http://www.cse.unsw.edu.au/~cs9020/)

Content includes:

- Lecture slides and recordings
- Quizzes and Assignments
- Course Forums
- Practice questions
- Challenge questions

# Course Material

Textbook:

- KA Ross and CR Wright: [Discrete Mathematics](#)

Supplementary textbook:

- E Lehman, FT Leighton, A Meyer:  
[Mathematics for Computer Science](#)



# Assessment Summary

60% exam, 30% assignments, 10% quizzes:

- 16 quizzes, worth up to 1 mark each
- 3 assignments, worth up to 10 marks each
- final exam (2 hours) worth up to 60 marks

Quizzes are available for 2 days before each lecture. Assignments due at the end of weeks 4, 7 and 10.

**You must achieve 40% on the final exam to pass**

Your final score will be taken from your 10 best quiz results, 3 assignments and final exam.

## More information

View the course outline at:

[www.cse.unsw.edu.au/~cs9020/outline.html](http://www.cse.unsw.edu.au/~cs9020/outline.html)

Particularly the sections on **Student conduct** and **Plagiarism**.

**What I will expect from you**

# Assessments

To achieve good marks in this course you need to demonstrate:

- Your understanding of the material
- Your ability to work with the material

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# Mathematical communication

## Guidelines for good mathematical writing

Mathematical writing should be:

- Clear
- Logical
- Sensical

# Examples

## Example

Ex 1 a) ~~300~~ 51 b) 72 c) 12

$$\begin{aligned}\text{Ex 2: } (A \setminus B) \cup (B \setminus A) &= (A \cap B^c) \cup (B \cap A^c) = (A \cup B) \cap (A \cup A^c) \cap (B \cup B^c) \cap (\cancel{A \cup A^c}) \\ &= (A \cup B) \cap (A^c \cup B^c) = (A \cup B) \cap (A \cap B)^c = (A \cup B) \setminus (A \cap B) \text{ by DeM, Dist}\end{aligned}$$

Ex 3 a) Yes b) No c) Yes d) No e) Yes Ex 4 a) True b) False

~~Ex 4 a) True b) False~~

# Examples

## Example

Ex. 2

$$\begin{aligned}(A \setminus B) \cup (B \setminus A) &= (A \cap B^c) \cup (B \cap A^c) && \text{(Def.)} \\ &= ((A \cap B^c) \cup B) \cap ((A \cap B^c) \cup A^c) && \text{(Dist.)} \\ &= (A \cup B) \cap (B^c \cup B) \\ &\quad \cap (A \cup A^c) \cap (B^c \cup A^c) && \text{(Dist.)} \\ &= (A \cup B) \cap (A^c \cup B^c) && \text{(Ident.)} \\ &= (A \cup B) \cap (A \cap B)^c && \text{(DeM.)} \\ &= (A \cup B) \setminus (A \cap B) && \text{(Def.)}\end{aligned}$$



# Examples

## Example

Ex. 4a

We will show that if  $R_1$  and  $R_2$  are symmetric, then  $R_1 \cap R_2$  is symmetric.

Suppose  $(a, b) \in R_1 \cap R_2$ . Then  $(a, b) \in R_1$  and  $(a, b) \in R_2$ . Because  $R_1$  is symmetric,  $(b, a) \in R_1$ . Because  $R_2$  is symmetric,  $(b, a) \in R_2$ . Therefore  $(b, a) \in R_1 \cap R_2$ . Therefore  $R_1 \cap R_2$  is symmetric.

# Proofs

A large component of your work in this course is giving **proofs** of **propositions**.

A **proposition** is a statement that is either true or false.

## Example

Propositions:

- $3 + 5 = 8$
- All integers are either even or odd
- There exist  $a, b, c$  such that  $1/a + 1/b + 1/c = 4$

Not propositions:

- $3 + 5$
- $x$  is even or  $x$  is odd
- $1/a + 1/b + 1/c = 4$

# Proposition structure

Common proposition structures include:

If A then B  $(A \Rightarrow B)$

A if and only if B  $(A \Leftrightarrow B)$

For all x, A  $(\forall x.A)$

There exists x such that A  $(\exists x.A)$

$\forall$  and  $\exists$  are known as **quantifiers**.

# Proofs

A large component of your work in this course is giving **proofs** of **propositions**.

A proof of a proposition is an argument to convince the reader/marker that the proposition is true.

A **proof** of a proposition is a finite sequence of logical steps, starting from base assumptions (**axioms** and **hypotheses**), leading to the proposition in question.

# Proofs

## Example

Prove:  $3 \times 2 = 2 \times 3$

$$\begin{aligned} 3 \times 2 &= (2 + 1) \times 2 \\ &= (2 \times 2) + (1 \times 2) \\ &= (1 \times 2) + (2 \times 2) \\ &= 2 + (2 \times 2) \\ &= (2 \times 1) + (2 \times 2) \\ &= 2 \times (1 + 2) \\ &= 2 \times 3. \end{aligned}$$

# Proofs: pitfalls

Starting from the proposition and deriving true **is not valid**.

## Example

Prove:  $1 = -1$

$$\begin{array}{rcl} & 1 & = -1 \\ \text{So } (1)^2 & = & (-1)^2 \\ \text{So } 1 & = & 1 \quad \text{which is true.} \end{array}$$

Does this mean that  $1 = -1$ ?

## Proofs: pitfalls

Make sure each step is logically valid: for example,  $x = y$  implies  $x^2 = y^2$  but  $x^2 = y^2$  does not imply  $x = y$ .

### Example

Suppose  $a = b$ . Then,

$$\begin{array}{rcl} & a^2 & = ab \\ \text{So} & a^2 - b^2 & = ab - b^2 \\ \text{So} & (a - b)(a + b) & = (a - b)b \\ \text{So} & a + b & = b \\ \text{So} & a & = 0 \end{array}$$

This is true no matter what value  $a$  is given at the start, so does that mean everything is equal to 0?

# Proofs: pitfalls

For propositions of the form  $\forall x.A$  where  $x$  can have infinitely many values:

- You cannot enumerate infinitely many cases in a proof.
- Only considering a finite number of cases is not sufficient.

## Example

For all  $n$ ,  $n^2 + n + 41$  is prime

True for  $n = 0, 1, 2, \dots, 39$ . Not true for  $n = 40$ .



# Proofs: pitfalls

The order of quantifiers matters when it comes to propositions:

## Example

- For every number  $x$ , there is a number  $y$  such that  $y$  is larger than  $x$
- There is a number  $y$  such that for every number  $x$ ,  $y$  is larger than  $x$

## Proof strategies: direct proof

Proposition form	You need to do this
$A \Rightarrow B$	Assume A and prove B
$A \Leftrightarrow B$	Prove “If A then B” and “If B then A”
$\forall x.A$	Show A holds for every possible value of x
$\exists x.A$	Find a value of x that makes A true

# Proof strategies: contradiction

To prove  $A$  is true, assume  $A$  is false and derive a contradiction.  
That is, start from the negation of the proposition and derive false.

## Example

Prove:  $\sqrt{2}$  is irrational

Proof: Assume  $\sqrt{2}$  is rational ...

# Negating propositions

Proposition form	Its negation
$A$ and $B$	Not $A$ or not $B$
$A$ or $B$	Not $A$ and Not $B$
$A \Rightarrow B$	$A$ and Not $B$
$A \Leftrightarrow B$	$(A$ and Not $B)$ or $($ not $A$ and $B)$
$\forall x.A$	$\exists x.$ not $A$
$\exists x.A$	

# Negating propositions

Proposition form	Its negation
$A$ and $B$	?
$A$ or $B$	
$A \Rightarrow B$	
$A \Leftrightarrow B$	
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## Negating propositions

Proposition form	Its negation
$A$ and $B$	?
$A$ or $B$	?
$A \Rightarrow B$	?
$A \Leftrightarrow B$	?
$\forall x.A$	?
$\exists x.A$	

## Negating propositions

Proposition form	Its negation
$A$ and $B$	? not a or not b
$A$ or $B$	? not a and not b
$A \Rightarrow B$	? a and not b
$A \Leftrightarrow B$	? a and not b, or b and not a
$\forall x.A$	? E x not a
$\exists x.A$	? floor x not

## Proof strategies: contrapositive

To prove a proposition of the form “If A then B” you can prove “If not B then not A”

### Example

Prove: If  $m + n \geq 73$  then  $m \geq 37$  or  $n \geq 37$ .

## Proof strategies: dealing with $\forall$

How can we check infinitely many cases?

- Choose an **arbitrary** element: an object with no assumptions about it (may have to check several cases)
- Induction (see week 7)

### Example

Prove: For every integer  $n$ ,  $n^2$  will have remainder 0 or 1 when divided by 4.

**Note:** “Arbitrary” is not the same as “random”.