



## CSP2348-5243 Data Structures

### Solutions to Tutorial 03

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#### Exercise 1:

code	##	cost
int ascending(int[] array, boolean unique) {	01	$O(1)$
if (array==null) return -1;	02	$O(1)$
int n=array.length;	03	$O(1)$
for (i=1; i<n; i++){	04	$O(1) \times \#e$
if (array[i-1]>array[i]) return i;	05	$O(1) \times \#e$
if (unique && array[i-1]==array[i]) return i;	06	$O(1) \times \#e$
}	07	-
return n;	08	$O(1)$
}	09	-
loop-control=1,2,3,...,n-1 #e = number of executions $\approx O(n)$ maximum cost = $O(n)$		

#### Exercise 2:

code	#	cost
int searchLinear(int[] array, int item) {	01	$O(1)$
if (array==null) return -1;	02	$O(1) \times \#e$
for (int i=0; i<array.length; i++)	03	$O(1) \times \#e$
if (array[i]>=item)	04	$O(1) \times \#e$
return i;	05	$O(1) \times \#e$
return array.length;	06	$O(1)$
}		
loop-control = 0,1, 2, 3,..., n-1 #e = number of executions $\approx O(n)$ maximum cost = $O(n)$		

### Exercise 3:

code	#	cost
public void merge(int[] a1, int l1, int r1, int[] a2, int l2, int r2, int[] a3, int l3) {	01	O(1)
//merge existing a1[l1,...,r1] and existing a2[l2,...,r2]	02	-
//into existing a3[l3], where a1 and a2 are sorted	03	-
int i = l1, j = l2, k = l3;	04	O(1)
while (i <= r1 && j <= r2) {	05	O(1)x#e1
if (a1[i] <= a2[j]) {	06	O(1)x#e1
a3[k++] = a1[i++];	07	O(1)x#e1
} else {	08	-
a3[k++] = a2[j++];	09	O(1)x#e1
}	10	-
}	11	-
while (i <= r1) {	12	O(1)x#e2
a3[k++] = a1[i++];	13	O(1)x#e2
}	14	-
while (j <= r2) {	15	O(1)x#e3
a3[k++] = a2[j++];	16	O(1)x#e3
}	17	-
}	18	-
<p>1st while loop: Loop-control1 = 1, 2, ... <math>n_1+n_2</math>                    #e1 = number of execution <math>\approx O(n_1+n_2)</math></p> <p>(Note: 2<sup>nd</sup> &amp; 3<sup>rd</sup> while loops would execute one of them only)</p> <p>2<sup>nd</sup> while loop: Loop-control2 = ..., <math>n_1</math>                    #e2 = number of execution <math>\approx O(n_1)</math></p> <p>3<sup>rd</sup> while loop: Loop-control2 = ..., <math>n_2</math>                    #e3 = number of execution <math>\approx O(n_2)</math></p> <p>Maximum cost = <math>O(n_1+ n_2) + O(n_1) + O(n_2) = O(n_1+ n_2)</math></p>		

### Exercise 4 ~ 7:

See separate codes in separate documents