# ACM 常用算法模板



## Vijay

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## 1 搜索 Search

## 1.1 A\*算法

#### 1.1.1 八数码

```
struct matrix {
   vector<vector<int>> g;
   bool operator<(matrix a) const {</pre>
       return g < a.g;
   }
} st, v;
int h(matrix a) {
                     //h 函数定义:不在标准位置的数字个数
   int res{};
   for (int i = 0; i < 3; ++i) for (int j = 0; j < 3; ++j) if (a.g[i][j] !=
st.g[i][j]) ++res;
   return res;
struct node {
   matrix m;
   int cnt;
                  //移动次数
   bool operator<(node a) const { return cnt + h(m) > a.cnt + h(a.m); }
};
void solution() {
   st.g = vector<vector<int>>{{1, 2, 3}, {8, 0, 4}, {7, 6, 5}};
                                                                    //
定义标准矩阵
   v.g = vector<vector<int>>(3, vector<int>(3));
   for (auto& x : v.g) for (int i = 0; i < 3; ++i) {
                                                          //输入
       char c;
       cin >> c;
       x[i] = c - '0';
   priority_queue<node> q;
                      //防止重复
   set<matrix> s;
   int fx = 0, fy = 0;
                          //空格的位置
   q.push({v, 0});
   while (!q.empty()) {
       auto t = q.top();
       q.pop();
       if (!h(t.m)) {
                          //判断与标准矩阵是否一阵
           cout << t.cnt << '\n';
           return;
        for (int i = 0; i < 3; ++i) for (int j = 0; j < 3; ++j) if
|(!t.m.g[i][j])| fx = i, fy = j;
                                   //查找空格位置
       for (auto& dir : DIRS) {
                                      //向四种方向移动
           int x = fx + dir.first, y = fy + dir.second;
           if (x >= 0 && y >= 0 && x < 3 && y < 3) {
              swap(t.m.g[fx][fy], t.m.g[x][y]);
              if (!s.count(t.m)) {
```

```
s.emplace(t.m);
q.push({t.m, t.cnt + 1});
}
swap(t.m.g[fx][fy], t.m.g[x][y]); //撤销操作
}
}
}
```

#### 1.1.2 K 短路

```
vector<double> dist;
                         //从终点到各点的最短距离
void solution() {
   int n, m, u, v;
                      //点数,边数
   double e, s;
                      //边权,总权值
   cin >> n >> m >> s;
   dist = vector<double>(n + 1, INF);
   dist[n] = 0;
   vector<bool> vis(n + 1, false);
                                     //判断节点是否已被访问
   vector<int> cnt(n + 1, 0);
                                     //每个点被访问的次数
   vector<vector<pair<int,double>>> adj(n + 1), adj2(n + 1);
                                                          //邻接表;
   for (int i = 0; i < m; ++i) {
       cin >> u >> v >> e;
       adj[u].emplace_back(v, e);
       adj2[v].emplace_back(u, e);
   }
   struct node {
                      //使用 A*所需的结构体
       int idx{};
                      //节点编号
       double d{};
                      //表示从起点到该节点的实际距离
       bool operator<(node a) const {</pre>
          return d + dist[idx] > a.d + dist[a.idx];
   };
   struct node2 {
                      //计算终点到所有节点的最短路
       int idx{};
       double d{};
       bool operator<(node2 a) const { return d > a.d; }
   };
   priority_queue<node> q;
   priority_queue<node2> q2;
   int ans{};
   q2.push({n, 0});
   while (!q2.empty()) {
       auto t = q2.top();
       q2.pop();
       if (vis[t.idx]) continue;
       vis[t.idx] = true;
       dist[t.idx] = t.d;
       for (auto& e : adj2[t.idx]) q2.push({e.first, e.second + t.d});
   int k = (int)s / dist[1];
   q.push({1, 0});
```

```
//使用 A*算法
while (!q.empty()) {
    auto t = q.top();
    q.pop();
    ++cnt[t.idx];
    if (t.idx == n) {
        s -= t.d;
        if (s < 0) {
            cout << ans << '\n';</pre>
            return;
        }
        ++ans;
    for (auto& e : adj[t.idx])
        if (cnt[e.first] <= k && t.d + e.second <= s)</pre>
           q.push({e.first, e.second + t.d});
cout << ans << '\n';
```

## 1.2 Dancing Links X 算法

#### 1.2.1 精确覆盖问题

```
const int MX = 1e5 + 10;
int n, m, ans, stk[MX];
struct DLX {
   int n{}, m{}, idx{};
                             //行数,列数,当前列的索引
   int L[MX], R[MX], U[MX], D[MX], col[MX], row[MX], head[MX], sz[MX];
   //指针,行列,问题解,表头,每列数量
   void init (const int& r, const int& c) { //初始化循环双向链表
       n = r, m = c;
       for (int i = 0; i <= c; ++i) {
          L[i] = i - 1;
          R[i] = i + 1;
          U[i] = D[i] = i;
       }
       L[0] = c, R[c] = 0;
       idx = c;
       memset(head, 0 ,sizeof(head));
       memset(sz, 0, sizeof(sz));
   void insert(const int& r, const int& c) {//插入节点
       col[++idx] = c, row[idx] = r, ++sz[c];
       D[idx] = D[c], U[D[c]] = idx, U[idx] = c, D[c] = idx;
       if (!head[r]) head[r] = L[idx] = R[idx] = idx;
       else {
          R[idx] = R[head[r]], L[R[head[r]]] = idx;
          L[idx] = head[r], R[head[r]] = idx;
       }
   }
   void remove(const int& c) { // 移除 c 列以及所有与其相交的行
       L[R[c]] = L[c], R[L[c]] = R[c];
```

```
for (int i = D[c]; i != c; i = D[i])
           for (int j = R[i]; j != i; j = R[j])
               U[D[j]] = U[j], D[U[j]] = D[j], --sz[col[j]];
   void resume(const int& c) { //恢复 c 列以及所有与其相交的行
       L[R[c]] = R[L[c]] = c;
       for (int i = U[c]; i != c; i = U[i])
           for (int j = L[i]; j != i; j = L[j])
               U[D[j]] = D[U[j]] = j, ++sz[col[j]];
   }
   bool dance(int dep) { //递归求解
       int c = R[0];
       if (!c) {
           ans = dep;
           return true;
       for (int i = c; i != 0; i = R[i]) {
           if (sz[i] < sz[c]) c = i;
       remove(c);
       for (int i = D[c]; i != c; i = D[i]) {
           stk[dep] = row[i];
           for (int j = R[i]; j != i; j = R[j]) remove(col[j]);
           if (dance(dep + 1)) return true;
           for (int j = L[i]; j != i; j = L[j]) resume(col[j]);
       resume(c);
       return false;
} solver;
void solution() {
   cin >> n >> m;
   solver.init(n, m);
   for (int i = 1; i <= n; ++i) {
       for (int j = 1; j <= m; ++j) {
           int x;
           cin >> x;
           if (x) solver.insert(i, j);
       }
                       //从第一行开始递归求解
   solver.dance(0);
   if (ans) for (int i = 0; i < ans; ++i) cout << stk[i] << ' ';
   else cout << "No Solution!" << '\n';</pre>
1.2.2 重复覆盖问题
const int MX = 1e5 + 10;
int n, m, ans;
struct DLX {
   bool vis[MX];
                              //行数,列数,当前列的索引
   int n{}, m{}, idx{};
   int L[MX], R[MX], U[MX], D[MX], col[MX], row[MX], head[MX], sz[MX];
```

```
//指针,行列,问题解,表头,每列数量
void init (const int& r, const int& c) {
                                          //初始化循环双向链表
   n = r, m = c;
   for (int i = 0; i <= c; ++i) {
       L[i] = i - 1;
       R[i] = i + 1;
       U[i] = D[i] = i;
   L[0] = c, R[c] = 0;
   idx = c;
   memset(head, 0 ,sizeof(head));
   memset(sz, 0, sizeof(sz));
}
void insert(const int& r, const int& c) {//插入节点
   col[++idx] = c, row[idx] = r, ++sz[c];
   D[idx] = D[c], U[D[c]] = idx, U[idx] = c, D[c] = idx;
   if (!head[r]) head[r] = L[idx] = R[idx] = idx;
   else {
       R[idx] = R[head[r]], L[R[head[r]]] = idx;
       L[idx] = head[r], R[head[r]] = idx;
   }
}
void remove(const int& c) {
                             // 删除和恢复和精确覆盖有所不同
   for(int i = D[c]; i != c; i = D[i])
       L[R[i]] = L[i], R[L[i]] = R[i];
void resume(const int& c) {
   for(int i = U[c]; i != c; i = U[i])
       L[R[i]] = R[L[i]] = i;
}
int h() { // 预估函数
   int ret = 0;
   memset(vis, 0, sizeof vis);
   for(int i = R[0]; i; i = R[i]) {
       if(vis[i]) continue;
       ++ret; vis[i] = true;
       for(int j = D[i]; j != i; j = D[j])
           for(int k = R[j]; k != j; k = R[k])
              vis[col[k]] = true;
   return ret;
void dance(int dep) { //递归求解
   if (dep + h() >= ans) return;
   int c = R[0];
   if (!c) {
       if (dep < ans) ans = dep;</pre>
       return;
   for (int i = c; i != 0; i = R[i]) {
       if (sz[i] < sz[c]) c = i;
   }
   remove(c);
```

```
for (int i = D[c]; i != c; i = D[i]) {
            remove(i);
            for (int j = R[i]; j != i; j = R[j]) remove(j);
            dance(dep + 1);
            for (int j = L[i]; j != i; j = L[j]) resume(j);
            resume(i);
        }
       return;
   }
} solver;
void solution() {
   while (cin >> n >> m) {
        solver.init(n, n);
        for (int i = 1; i <= n; ++i) solver.insert(i, i);</pre>
       for (int i = 1; i <= m; ++i) {
            int x, y;
            cin >> x >> y;
            solver.insert(x, y);
            solver.insert(y, x);
        solver.dance(1);
        cout << ans << '\n';</pre>
   }
```

## 2 字符串 String

## 2.1 KMP 算法

```
vector<int> prefix_function(string s) { //前缀函数
   int n = (int)s.length();
   vector<int> pi(n);
   for (int i = 1; i < n; i++) {
       int j = pi[i - 1];
       while (j > 0 \&\& s[i] != s[j]) j = pi[j - 1];
       if (s[i] == s[j]) j++;
       pi[i] = j;
   return pi;
vector<int> find_occurrences(string text, string pattern) {
                                                                //KMP 算
法
   string cur = pattern + '#' + text;
   int sz1 = text.size(), sz2 = pattern.size();
   vector<int> v;
   vector<int> lps = prefix_function(cur);
   for (int i = sz2 + 1; i \le sz1 + sz2; i++) {
       if (lps[i] == sz2)
       v.push_back(i - 2 * sz2);
   }
   return v;
```

## 2.2 Z 算法

## 2.3 字典树

#### 2.3.1 查询是否出现

```
struct Trie {
    int next[MX][26], cnt;
    bool exist[MX]; // 该结点结尾的字符串是否存在

void insert(string s) { // 插入字符串
    int p = 0;
    for (int i = 0; i < s.size(); ++i) {
        int c = s[i] - 'a';
        if (!next[p][c]) next[p][c] = ++cnt; // 如果没有,就添加结点
        p = next[p][c];
    }
    exist[p] = true;
}

bool find(string s) { // 查找字符串
```

```
int p = 0;
    for (int i = 0; i < s.size(); ++i) {
        int c = s[i] - 'a';
        if (!next[p][c]) return false;
        p = next[p][c];
    }
    return exist[p];
}</pre>
```

### 2.3.2 查询前缀出现次数

```
struct Trie {
   int next[MX][26],cnt,pre[MX];
   bool p;
   int n, m;
   void insert(string s) {
       int p = 0;
       for(int i=0;i<s.size();i++)</pre>
           int c = s[i]-'a';
           if(!next[p][c]) next[p][c] = ++cnt;
           pre[next[p][c]]++;//前缀保存
           p = next[p][c];
       }
   int search(string s) {
       int p = 0;
       for(int i = 0; i < s.size();i++)</pre>
           int c = s[i] - 'a';
           if(!next[p][c]) return 0;
           p=next[p][c];
       }//root 经过此循环后变成前缀最后一个字母所在位置
       return pre[p];
   }
};
```

## 2.4 AC 自动机

```
#include <deque>
#include <iostream>

void promote()
{
    std::ios::sync_with_stdio(0);
    std::cin.tie(0);
    std::cout.tie(0);
    return;
}

typedef char chr;
typedef std::deque<int> dic;
```

```
const int maxN = 2e5;
const int maxS = 2e5;
const int maxT = 2e6;
int n;
chr s[maxS + 10];
chr t[maxT + 10];
int cnt[maxN + 10];
struct AhoCorasickAutomaton
   struct Node
       int son[30];
       int val;
       int fail;
       int head;
       dic index;
   } node[maxS + 10];
   struct Edge
   {
       int head;
       int next;
    } edge[maxS + 10];
   int root;
   int ncnt;
   int ecnt;
   void Insert(chr *str, int i)
       int u = root;
       for (int i = 1; str[i]; i++)
           if (node[u].son[str[i] - 'a' + 1] == 0)
               node[u].son[str[i] - 'a' + 1] = ++ncnt;
           u = node[u].son[str[i] - 'a' + 1];
       }
       node[u].index.push_back(i);
       return;
   }
   void Build()
       dic q;
       for (int i = 1; i \le 26; i++)
           if (node[root].son[i])
               q.push_back(node[root].son[i]);
       while (!q.empty())
       {
           int u = q.front();
           q.pop_front();
           for (int i = 1; i <= 26; i++)
               if (node[u].son[i])
```

```
{
                   node[node[u].son[i]].fail = node[node[u].fail].son[i];
                   q.push_back(node[u].son[i]);
               }
               else
               {
                   node[u].son[i] = node[node[u].fail].son[i];
               }
           }
       }
       return;
   }
   void Query(chr *str)
       int u = root;
       for (int i = 1; str[i]; i++)
           u = node[u].son[str[i] - 'a' + 1];
           node[u].val++;
       return;
   }
   void addEdge(int tail, int head)
   {
       ecnt++;
       edge[ecnt].head = head;
       edge[ecnt].next = node[tail].head;
       node[tail].head = ecnt;
       return;
   }
   void DFS(int u)
       for (int e = node[u].head; e; e = edge[e].next)
           int v = edge[e].head;
           DFS(v);
           node[u].val += node[v].val;
       for (auto i : node[u].index)
           cnt[i] += node[u].val;
       return;
   }
   void FailTree()
       for (int u = 1; u <= ncnt; u++)
           addEdge(node[u].fail, u);
       DFS(root);
       return;
} ACM;
int main()
```

```
{
    std::cin >> n;
    for (int i = 1; i <= n; i++)
    {
        std::cin >> (s + 1);
        ACM.Insert(s, i);
    }
    ACM.Build();
    std::cin >> (t + 1);
    ACM.Query(t);
    ACM.FailTree();
    for (int i = 1; i <= n; i++)
        std::cout << cnt[i] << '\n';
    return 0;
}</pre>
```

## 2.5 Manacher 算法

```
void Manacher(char s[], int len)
{ // 原字符串和串长
   int 1 = 0;
   String[l++] = '$'; // 0 下标存储为其他字符, 防止越界
   String[l++] = '#';
   for (int i = 0; i < len; i++)
       String[l++] = s[i];
       String[1++] = '#';
   String[1] = 0; // 空字符
   int MaxR = 0;
   int flag = 0;
   for (int i = 0; i < 1; i++)
        cnt[i] = MaxR > i ? min(cnt[2 * flag - i], MaxR - i) : 1; //
2*flag-i 是i 点关于flag 的对称点
       while (String[i + cnt[i]] == String[i - cnt[i]])
           cnt[i]++;
       if (i + cnt[i] > MaxR)
          MaxR = i + cnt[i];
          flag = i;
       }
   }
```

## 2.6 字符串哈希

```
int count_unique_substrings(string const &s)
{
   int n = s.size();
```

```
const int b = 31;
const int m = 1e9 + 9;
vector<long long> b_pow(n);
b_pow[0] = 1;
for (int i = 1; i < n; i++)
   b_pow[i] = (b_pow[i - 1] * b) % m;
vector<long long> h(n + 1, 0);
for (int i = 0; i < n; i++)
   h[i + 1] = (h[i] + (s[i] - 'a' + 1) * b_pow[i]) % m;
int cnt = 0;
for (int l = 1; l <= n; l++)
    set<long long> hs;
   for (int i = 0; i <= n - 1; i++)
        long cur_h = (h[i + 1] + m - h[i]) % m;
       cur_h = (cur_h * b_pow[n - i - 1]) % m;
       hs.insert(cur_h);
   cnt += hs.size();
return cnt;
```

## 2.7 后缀自动机

```
void sam_extend(char c)
   int cur = sz++;
   st[cur].len = st[last].len + 1;
   int p = last;
   while (p != -1 && !st[p].next.count(c))
       st[p].next[c] = cur;
       p = st[p].link;
   if (p == -1)
       st[cur].link = 0;
   }
   else
   {
       int q = st[p].next[c];
       if (st[p].len + 1 == st[q].len)
       {
           st[cur].link = q;
       }
       else
       {
           int clone = sz++;
           st[clone].len = st[p].len + 1;
           st[clone].next = st[q].next;
```

```
st[clone].link = st[q].link;
    while (p != -1 && st[p].next[c] == q)
    {
        st[p].next[c] = clone;
        p = st[p].link;
    }
        st[q].link = st[cur].link = clone;
    }
}
last = cur;
}
```

## 3 数据结构 Data Structure

## 3.1 并查集

```
#include <bits/stdc++.h>
using namespace std;
struct dsu
   vector<size_t> pa, size, sum;
   explicit dsu(size_t size_)
       : pa(size_ * 2), size(size_ * 2, 1), sum(size_ * 2)
       // size 与 sum 的前半段其实没有使用,只是为了让下标计算更简单
       iota(pa.begin(), pa.begin() + size_, size_);
       iota(pa.begin() + size_, pa.end(), size_);
       iota(sum.begin() + size_, sum.end(), 0);
   }
   void unite(size_t x, size_t y)
   {
       x = find(x), y = find(y);
       if (x == y)
           return;
       if (size[x] < size[y])</pre>
           swap(x, y);
       pa[y] = x;
       size[x] += size[y];
       sum[x] += sum[y];
   }
   void move(size_t x, size_t y)
   {
       auto fx = find(x), fy = find(y);
       if (fx == fy)
           return;
       pa[x] = fy;
       --size[fx], ++size[fy];
```

```
sum[fx] -= x, sum[fy] += x;
   }
   size_t find(size_t x) { return pa[x] == x ? x : pa[x] = find(pa[x]); }
};
int main()
   size_t n, m, op, x, y;
   while (cin >> n >> m)
       dsu dsu(n + 1); // 元素范围是 1...n
       while (m--)
        {
            cin >> op;
            switch (op)
            case 1:
               cin >> x >> y;
               dsu.unite(x, y);
               break;
            case 2:
               cin >> x >> y;
               dsu.move(x, y);
               break;
            case 3:
               cin >> x;
               x = dsu.find(x);
               cout << dsu.size[x] << ' ' << dsu.sum[x] << '\n';</pre>
               break;
            default:
               assert(false); // not reachable
            }
       }
   }
```

## 3.2 线段树

#### 3.2.1 区间加/求和线段树

```
#include <bits/stdc++.h>
using namespace std;

template <typename T>
class SegTreeLazyRangeAdd
{
   vector<T> tree, lazy;
   vector<T> *arr;
   int n, root, n4, end;

   void maintain(int cl, int cr, int p)
```

```
{
    int cm = cl + (cr - cl) / 2;
    if (cl != cr && lazy[p])
        lazy[p * 2] += lazy[p];
        lazy[p * 2 + 1] += lazy[p];
        tree[p * 2] += lazy[p] * (cm - cl + 1);
        tree[p * 2 + 1] += lazy[p] * (cr - cm);
        lazy[p] = 0;
    }
}
T range_sum(int l, int r, int cl, int cr, int p)
    if (1 <= c1 && cr <= r)
       return tree[p];
    int m = cl + (cr - cl) / 2;
    T sum = 0;
    maintain(cl, cr, p);
    if (1 <= m)
        sum += range_sum(1, r, cl, m, p * 2);
    if (r > m)
       sum += range_sum(1, r, m + 1, cr, p * 2 + 1);
    return sum;
}
void range_add(int 1, int r, T val, int cl, int cr, int p)
    if (1 <= c1 && cr <= r)
    {
        lazy[p] += val;
        tree[p] += (cr - cl + 1) * val;
       return;
    int m = cl + (cr - cl) / 2;
    maintain(cl, cr, p);
    if (1 <= m)
        range_add(l, r, val, cl, m, p * 2);
    if (r > m)
       range_add(1, r, val, m + 1, cr, p * 2 + 1);
    tree[p] = tree[p * 2] + tree[p * 2 + 1];
}
void build(int s, int t, int p)
{
    if (s == t)
        tree[p] = (*arr)[s];
       return;
    }
    int m = s + (t - s) / 2;
    build(s, m, p * 2);
    build(m + 1, t, p * 2 + 1);
    tree[p] = tree[p * 2] + tree[p * 2 + 1];
}
```

```
public:
   explicit SegTreeLazyRangeAdd<T>(vector<T> v)
       n = v.size();
       n4 = n * 4;
       tree = vector<T>(n4, 0);
       lazy = vector<T>(n4, 0);
       arr = &v;
       end = n - 1;
       root = 1;
       build(0, end, 1);
       arr = nullptr;
   }
   void show(int p, int depth = 0)
       if (p > n4 \mid | tree[p] == 0)
           return;
       show(p * 2, depth + 1);
       for (int i = 0; i < depth; ++i)
           putchar('\t');
       printf("%d:%d\n", tree[p], lazy[p]);
       show(p * 2 + 1, depth + 1);
   }
   T range_sum(int 1, int r) { return range_sum(1, r, 0, end, root); }
   void range_add(int 1, int r, int val) { range_add(l, r, val, 0, end,
root); }
};
```

#### 3.2.2 区间修改/求和线段树

```
#include <bits/stdc++.h>
using namespace std;
template <typename T>
class SegTreeLazyRangeSet
   vector<T> tree, lazy;
   vector<T> *arr;
   int n, root, n4, end;
   void maintain(int cl, int cr, int p)
       int cm = cl + (cr - cl) / 2;
       if (cl != cr && lazy[p])
       {
           lazy[p * 2] = lazy[p];
           lazy[p * 2 + 1] = lazy[p];
           tree[p * 2] = lazy[p] * (cm - cl + 1);
           tree[p * 2 + 1] = lazy[p] * (cr - cm);
           lazy[p] = 0;
       }
   }
```

```
T range_sum(int l, int r, int cl, int cr, int p)
       if (1 <= c1 && cr <= r)
           return tree[p];
       int m = cl + (cr - cl) / 2;
       T sum = 0;
       maintain(cl, cr, p);
       if (1 <= m)
           sum += range_sum(1, r, cl, m, p * 2);
       if (r > m)
           sum += range_sum(1, r, m + 1, cr, p * 2 + 1);
       return sum;
   }
   void range_set(int 1, int r, T val, int cl, int cr, int p)
       if (1 <= c1 && cr <= r)
       {
           lazy[p] = val;
           tree[p] = (cr - cl + 1) * val;
           return;
       }
       int m = cl + (cr - cl) / 2;
       maintain(cl, cr, p);
       if (1 <= m)
           range_set(1, r, val, cl, m, p * 2);
       if (r > m)
           range_set(1, r, val, m + 1, cr, p * 2 + 1);
       tree[p] = tree[p * 2] + tree[p * 2 + 1];
   }
   void build(int s, int t, int p)
       if (s == t)
           tree[p] = (*arr)[s];
           return;
       }
       int m = s + (t - s) / 2;
       build(s, m, p * 2);
       build(m + 1, t, p * 2 + 1);
       tree[p] = tree[p * 2] + tree[p * 2 + 1];
   }
public:
   explicit SegTreeLazyRangeSet<T>(vector<T> v)
       n = v.size();
       n4 = n * 4;
       tree = vector<T>(n4, 0);
       lazy = vector<T>(n4, 0);
       arr = &v;
       end = n - 1;
       root = 1;
       build(0, end, 1);
```

```
arr = nullptr;
}

void show(int p, int depth = 0)
{
    if (p > n4 || tree[p] == 0)
        return;
    show(p * 2, depth + 1);
    for (int i = 0; i < depth; ++i)
        putchar('\t');
    printf("%d:%d\n", tree[p], lazy[p]);
    show(p * 2 + 1, depth + 1);
}

T range_sum(int l, int r) { return range_sum(l, r, 0, end, root); }

void range_set(int l, int r, int val) { range_set(l, r, val, 0, end, root); }
};</pre>
```

#### 3.3 主席树

```
#include <algorithm>
#include <cstdio>
#include <cstring>
using namespace std;
const int maxn = 1e5; // 数据范围
int tot, n, m;
int sum[(maxn << 5) + 10], rt[maxn + 10], ls[(maxn << 5) + 10],
   rs[(maxn << 5) + 10];
int a[maxn + 10], ind[maxn + 10], len;
inline int getid(const int &val)
{ // 离散化
   return lower_bound(ind + 1, ind + len + 1, val) - ind;
int build(int 1, int r)
{ // 建树
   int root = ++tot;
   if (1 == r)
       return root;
   int mid = l + r >> 1;
   ls[root] = build(1, mid);
   rs[root] = build(mid + 1, r);
   return root; // 返回该子树的根节点
int update(int k, int l, int r, int root)
{ // 插入操作
   int dir = ++tot;
   ls[dir] = ls[root], rs[dir] = rs[root], sum[dir] = sum[root] + 1;
   if (1 == r)
```

```
return dir;
   int mid = l + r \gg 1;
   if (k <= mid)</pre>
       ls[dir] = update(k, 1, mid, ls[dir]);
   else
       rs[dir] = update(k, mid + 1, r, rs[dir]);
   return dir;
int query(int u, int v, int l, int r, int k)
{ // 查询操作
   int mid = l + r \gg 1,
        x = sum[ls[v]] - sum[ls[u]]; // 通过区间减法得到左儿子中所存储的数
值个数
   if (1 == r)
       return 1;
   if (k <= x) // 若 k 小于等于 x ,则说明第 k 小的数字存储在在左儿子中
       return query(ls[u], ls[v], l, mid, k);
   else // 否则说明在右儿子中
       return query(rs[u], rs[v], mid + 1, r, k - x);
inline void init()
   scanf("%d%d", &n, &m);
   for (int i = 1; i <= n; ++i)
       scanf("%d", a + i);
   memcpy(ind, a, sizeof ind);
   sort(ind + 1, ind + n + 1);
   len = unique(ind + 1, ind + n + 1) - ind - 1;
   rt[0] = build(1, len);
   for (int i = 1; i <= n; ++i)
       rt[i] = update(getid(a[i]), 1, len, rt[i - 1]);
int 1, r, k;
inline void work()
   while (m--)
       scanf("%d%d%d", &1, &r, &k);
       printf("%d\n", ind[query(rt[l - 1], rt[r], 1, len, k)]); // 回答
询问
   }
int main()
   init();
   work();
   return 0;
```

## 3.4 划分树

```
\#define \_mid(a, b) ((a + b) / 2)
using namespace std;
typedef long long 11;
const int maxn = 1e5 + 10;
const int INF = 0x3f3f3f3f;
int sorted[maxn];
int cnt[20][maxn];
int tree[20][maxn];
void build(int 1, int r, int k)
   if (r == 1) // 如果区间内只有一个数,返回
       return;
   int mid = _{mid}(l, r), flag = mid - l + 1; // 求出flag
   for (int i = 1; i <= r; i++)
       if (tree[k][i] < sorted[mid]) // sorted 代表排序好了的数组
          flag--;
   int bufl = l, bufr = mid + 1;
   for (int i = 1; i <= r; i++)
       cnt[k][i] = (i == 1) ? 0 : cnt[k][i - 1]; // 初始化
        if (tree[k][i] < sorted[mid] || tree[k][i] == sorted[mid] &&</pre>
flag > 0
       { // 如果有多个中值
          tree[k + 1][bufl++] = tree[k][i];
          cnt[k][i]++; // 进入左子树
          if (tree[k][i] == sorted[mid])
              flag--;
       }
       else // 进入右子树
          tree[k + 1][bufr++] = tree[k][i];
   build(l, mid, k + 1);
   build(mid + 1, r, k + 1);
int ask(int k, int sl, int sr, int l, int r, int x)
   if (sl == sr)
       return tree[k][s1];
   cntl = (l == sl) ? 0 : cnt[k][l - 1]; // 是否和查询区间重合
   int cntl2r = cnt[k][r] - cntl;
                                       // 计算l到r有cntl2r个数进入左
子树
   if (cnt12r >= x)
                                       // 如果大于当前查询的 k 则进入左子
树(因为左子树中最大的数大于第 k 大的数)
       return ask(k + 1, sl, \_mid(sl, sr), sl + cnt[k][r] -
1, x);
   else
```

```
{ // 否则进入右子树
int lr = _mid(sl, sr) + 1 + (l - sl - cntl);
return ask(k + 1, _mid(sl, sr) + 1, sr, lr, lr + r - l - cntl2r,
x - cntl2r);
}
}
```

## 3.5 单调栈

```
const int N = 100010;
stack<int> s;
int a[N], n;
int main()
{
    cin >> n;
    for (int i = 1; i <= n; i++) cin >> a[i];

    for (int i = 1; i <= n; i++) {
        while(!s.empty() && a[s.top()] >= a[i]) s.pop();
        if(!s.empty()) cout << a[s.top()] << ' ';
        else cout << -1 << ' ';
        s.push(i);
    }
    return 0;
}</pre>
```

## 3.6 伸展树

```
#include <cstdio>
const int N = 100005;
int rt, tot, fa[N], ch[N][2], val[N], cnt[N], sz[N];

struct Splay
{
    void maintain(int x) { sz[x] = sz[ch[x][0]] + sz[ch[x][1]] + cnt[x]; }

    bool get(int x) { return x == ch[fa[x]][1]; }

    void clear(int x)
    {
        ch[x][0] = ch[x][1] = fa[x] = val[x] = sz[x] = cnt[x] = 0;
    }

    void rotate(int x)
    {
        int y = fa[x], z = fa[y], chk = get(x);
        ch[y][chk] = ch[x][chk ^ 1];
        if (ch[x][chk ^ 1])
            fa[ch[x][chk ^ 1]] = y;
        ch[x][chk ^ 1] = y;
```

```
fa[y] = x;
    fa[x] = z;
    if (z)
        ch[z][y == ch[z][1]] = x;
    maintain(y);
    maintain(x);
}
void splay(int x)
    for (int f = fa[x]; f = fa[x], f; rotate(x))
        if (fa[f])
            rotate(get(x) == get(f) ? f : x);
    rt = x;
}
void ins(int k)
{
    if (!rt)
    {
        val[++tot] = k;
        cnt[tot]++;
        rt = tot;
        maintain(rt);
        return;
    int cur = rt, f = 0;
    while (1)
    {
        if (val[cur] == k)
        {
            cnt[cur]++;
            maintain(cur);
            maintain(f);
            splay(cur);
            break;
        f = cur;
        cur = ch[cur][val[cur] < k];</pre>
        if (!cur)
        {
            val[++tot] = k;
            cnt[tot]++;
            fa[tot] = f;
            ch[f][val[f] < k] = tot;
            maintain(tot);
            maintain(f);
            splay(tot);
            break;
        }
    }
}
int rk(int k)
{
    int res = 0, cur = rt;
```

```
while (1)
    {
        if (k < val[cur])</pre>
        {
            cur = ch[cur][0];
        }
        else
        {
            res += sz[ch[cur][0]];
            if (k == val[cur])
                splay(cur);
                return res + 1;
            }
            res += cnt[cur];
            cur = ch[cur][1];
        }
    }
}
int kth(int k)
{
    int cur = rt;
    while (1)
        if (ch[cur][0] && k <= sz[ch[cur][0]])</pre>
            cur = ch[cur][0];
        }
        else
            k -= cnt[cur] + sz[ch[cur][0]];
            if (k \ll 0)
                splay(cur);
                return val[cur];
            cur = ch[cur][1];
        }
    }
}
int pre()
    int cur = ch[rt][0];
    if (!cur)
        return cur;
    while (ch[cur][1])
        cur = ch[cur][1];
    splay(cur);
    return cur;
}
int nxt()
{
    int cur = ch[rt][1];
```

```
if (!cur)
           return cur;
       while (ch[cur][0])
           cur = ch[cur][0];
       splay(cur);
       return cur;
   }
   void del(int k)
       rk(k);
       if (cnt[rt] > 1)
           cnt[rt]--;
           maintain(rt);
           return;
       if (!ch[rt][0] && !ch[rt][1])
           clear(rt);
           rt = 0;
           return;
       if (!ch[rt][0])
           int cur = rt;
           rt = ch[rt][1];
           fa[rt] = 0;
           clear(cur);
           return;
       if (!ch[rt][1])
       {
           int cur = rt;
           rt = ch[rt][0];
           fa[rt] = 0;
           clear(cur);
           return;
       }
       int cur = rt;
       int x = pre();
       fa[ch[cur][1]] = x;
       ch[x][1] = ch[cur][1];
       clear(cur);
       maintain(rt);
} tree;
int main()
   int n, opt, x;
   for (scanf("%d", &n); n; --n)
       scanf("%d%d", &opt, &x);
       if (opt == 1)
           tree.ins(x);
```

## 3.7 替罪羊树

```
const int MAXN = 1500005, INF = 0x7f7f7f7f;
const double ALPHA = 0.7;
int L[MAXN], R[MAXN], N[MAXN], val[MAXN], size[MAXN], cnt = 1, FV[MAXN],
FN[MAXN];
int flatten(int pos, int *fv, int *fn) // 拉平
   int 1 = 0, r = 0, unempty = N[pos] != 0;
   if (L[pos])
       1 = flatten(L[pos], fv, fn);
   if (unempty)
       fv[1] = val[pos];
       fn[1] = N[pos];
   if (R[pos])
       r = flatten(R[pos], fv + 1 + unempty, fn + 1 + unempty);
   return 1 + r + unempty;
void rebuild(int pos, int l, int r) // 重建
   int mid = (1 + r) / 2, sz1 = 0, sz2 = 0;
   if (1 < mid)
       L[pos] = ++cnt;
       rebuild(L[pos], 1, mid - 1);
       sz1 = size[L[pos]];
   }
   else
       L[pos] = 0;
   if (mid < r)
       R[pos] = ++cnt;
       rebuild(R[pos], mid + 1, r);
       sz2 = size[R[pos]];
   }
   else
       R[pos] = 0;
   val[pos] = FV[mid];
```

```
N[pos] = FN[mid];
   size[pos] = sz1 + sz2 + N[pos];
void try_restructure(int pos) // 尝试重构当前子树
   double k = max(size[L[pos]], size[R[pos]]) / double(size[pos]);
   if (k > ALPHA)
       int sz = flatten(pos, FV, FN);
       rebuild(pos, 0, sz - 1);
   }
void insert(int v, int pos = 1) // 插\lambda
   size[pos]++;
   if (N[pos] == 0 \&\& L[pos] == 0 \&\& R[pos] == 0)
       val[pos] = v;
       N[pos] = 1;
   else if (v < val[pos])
       if (L[pos] == 0)
           L[pos] = ++cnt;
       insert(v, L[pos]);
   else if (v > val[pos])
       if (R[pos] == 0)
           R[pos] = ++cnt;
       insert(v, R[pos]);
   }
   else
       N[pos]++;
   try_restructure(pos);
void remove(int v, int pos = 1) // 删除
   size[pos]--;
   if (v < val[pos])</pre>
       remove(v, L[pos]);
   else if (v > val[pos])
       remove(v, R[pos]);
   else
       N[pos]--;
   try restructure(pos);
int countl(int v, int pos = 1) // 统计比 v 小的数的个数
   if (v < val[pos])</pre>
       return L[pos] ? countl(v, L[pos]) : 0;
   else if (v > val[pos])
       return size[L[pos]] + N[pos] + (R[pos] ? countl(v, R[pos]) : 0);
   else
       return size[L[pos]];
```

```
int countg(int v, int pos = 1) // 统计比 v 大的数的个数
   if (v > val[pos])
       return R[pos] ? countg(v, R[pos]) : 0;
   else if (v < val[pos])</pre>
       return size[R[pos]] + N[pos] + (L[pos] ? countg(v, L[pos]) : 0);
   else
       return size[R[pos]];
int rank(int v) // 求排名
   return countl(v) + 1;
int kth(int k, int pos = 1) // 求指定排名的数
   if (size[L[pos]] + 1 > k)
       return kth(k, L[pos]);
   else if (size[L[pos]] + N[pos] < k)</pre>
       return kth(k - size[L[pos]] - N[pos], R[pos]);
       return val[pos];
int pre(int v) // 求前驱
   int r = countl(v);
   return kth(r);
int suc(int v) // 求后继
   int r = size[1] - countg(v) + 1;
   return kth(r);
```

## 3.8 一维树状数组

#### 3.8.1 单点修改 + 区间查询

```
inline int lowbit(int x) // 求二进制下最低位的1
{
    return x & -x;
}

inline void add(int x, int d) //单点修改 给第x号元素 + d
{
    while (x <= n) {
        aa[x] += d;
        x += lowbit(x);
    }
}</pre>
```

```
inline int ask(int x) // 求从第一个元素 到 第 x 个元素的总和
{
    int res = 0;
    while (x) {
        res += aa[x];
        x -= lowbit(x);
    }
    return res;
}
inline int getsum(int x, int y) // 求区间[l, r]内的元素总和
{
    return ask(r) - ask(l - 1);
}
```

#### 3.8.2 区间修改 + 单点查询

```
inline int lowbit(int x) // 求二进制下最低位的1
   return x \& -x;
void add(int x, int d) // 修改点
   while (x <= n)
      p[x] += d;
      x += lowbit(x);
   }
void range_add(int l, int r, int d) // 经过差分处理后, p 数组修改只要在第 l
个元素的 +d 和第r + 1 个元素-d 即可完成区间修改
   add(1, d);
   add(r + 1, -d);
int ask(int x) // 求第x 个位置的元素值
   int res = 0;
   while (x)
      res += p[x];
      x -= lowbit(x);
   return res;
```

#### 3.8.3 区间修改 + 区间查询

typedef long long 11;

```
11 sum1[N], sum2[N];
inline int lowbit(int x) // 求二进制下最低位的1
   return x \& -x;
inline void add(ll x, ll d) // 对 sum1 数组和 sum2 数组进行维护
   for (int i = x; i \le n; i += lowbit(i))
   {
       sum1[i] += d;
       sum2[i] += x * d;
   }
inline void range_add(ll l, ll r, ll d) // 差分思想进行区间维护
   add(1, d);
   add(r + 1, -d);
inline ll ask(ll x) // 求[1, x]内的所有元素之和
   11 \text{ res} = 0;
   for (int i = x; i > 0; i -= lowbit(i))
       res += (x + 1) * sum1[i] - sum2[i];
   return res;
inline ll range_ask(ll l, ll r) // 求区间[l, r]内所有元素之和
   return ask(r) - ask(l - 1);
```

#### 3.8.4 求逆序对数量

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;
const int N = 5e5 + 7;

int n;
int tr[N], a[N];
vector<int> mp;
ll ans;

inline int lowbit(int x)
{
    return x & -x;
}
```

```
void add(int x, int d)
   while (x <= n)
       tr[x] += d;
       x += lowbit(x);
   }
int ask(int x)
   int res = 0;
   while (x)
       res += tr[x];
       x -= lowbit(x);
   return res;
inline int get_id(int x)
   return lower_bound(mp.begin(), mp.end(), x) - mp.begin() + 1;
int main()
   scanf("%d", &n);
   for (int i = 1; i <= n; ++i)
       scanf("%d", &a[i]), mp.push_back(a[i]);
   // 离散化
   sort(mp.begin(), mp.end());
   mp.erase(unique(mp.begin(), mp.end()), mp.end());
   // 先序遍历找这个点前有多少个点比它大
   for (int i = 1; i <= n; ++i)
       int u = get_id(a[i]);
       add(u, 1);
       ans = ans + i - ask(u);
   }
   /*
   //后序遍历,看在这个点后有多少个点比它小
   for (int i = n; i > 0; --i) {
       int u = get_id(a[i]);
       add(u, 1);
       ans += ask(u - 1);
   printf("%lld\n", ans);
   return 0;
```

## 3.9 二维树状数组

### 3.9.1 单点修改 + 二维区间查询

```
inline int lowbit(int x)
{
    return x & -x;
}

void add(int x, int y, int d)
{
    for (int i = x; i <= n; i += lowbit(i))
        for (int j = y; j <= n; j += lowbit(j))
            a[i][j] += d;
}

int ask(int x, int y)
{
    int res = 0;
    for (int i = x; i > 0; i -= lowbit(i))
        for (int j = y; j > 0; j -= lowbit(j))
            res += a[i][j];
    return res;
}

int range_ask(int x1, int y1, int x2, int y2)
{
    return ask(x2, y2) - ask(x1 - 1, y2) - ask(x2, y1 - 1) + ask(x1 - 1, y1 - 1);
}
```

#### 3.9.2 二维区间修改 + 单点查询

```
inline int lowbit(int x)
{
    return x & -x;
}

void update(int x, int y, int d) //单点修改
{
    for (int i = x; i <= n; i += lowbit(i))
        for (int j = y; j <= m; j += lowbit(j));
            tree[i][j] += d;
}

void ranger_update(int x1, int y1, int x2, int y2, int d) //二维区间修改
{
    update(x1, y1, d);
    update(x1, y2 + 1, -d);
    update(x2 + 1, y1, -d);
    update(x2 + 1, y2 + 1, d);
}

void ask(int x, int y)
{</pre>
```

```
int res = 0;
for (int i = x; i > 0; i -= lowbit(i))
    for (int j = y; j > 0; j -= lowbit(j))
    res += tree[i][j];
}
```

#### 3.9.3 二维区间修改 + 二维区间查询

```
inline int lowbit(int x)
   return x \& -x;
|void update(int x, int y, int d) // 单点修改
   for (int i = x; i <= n; i += lowbit(i))
       for (int j = y; j <= m; j += lowbit(j));</pre>
   tree[i][j] += d;
void ranger_update(int x1, int y1, int x2, int y2, int d) // 二维区间修改
   update(x1, y1, d);
   update(x1, y2 + 1, -d);
   update(x2 + 1, y1, -d);
   update(x2 + 1, y2 + 1, d);
int ask(int x, int y) // 单点查询
   int res = 0;
   for (int i = x; i > 0; i -= lowbit(i))
       for (int j = y; j > 0; j -= lowbit(j))
           res += tree[i][j];
   return res;
int ranger_ask(int x1, int y1, int x2, int y2) // 利用二维差分进行二维区间
查询
   return ask(x2, y2) - ask(x2, y1 - 1) - ask(x1 - 1, y2) + ask(x1 - 1, y2)
y1 - 1);
```

## 4 数学 Number Theory

#### 4.1 Euler 筛法

```
void Euler(int n)
{
```

```
for (int i = 2; i <= n; ++i)
{
    if (!vis[i])
    {
       pri[cnt++] = i;
    }
    for (int j = 0; j < cnt; ++j)
    {
       if (1ll * i * pri[j] > n) break;
       vis[i * pri[j]] = 1;
       if (i % pri[j] == 0) break;
    }
}
```

### 4.2 分解质因数

# 4.3 快速幂

```
long long qpow(double x, long long n) {
    long long ans = 1.0;
    long long y = x;
    while (n > 0) {
        if (n & 1) ans = ans * y % MOD;
        y = y * y % MOD;
        n >>= 1;
    }
    return ans;
}
```

# 4.4 牛顿迭代法

double sqrt\_newton(double n)

```
{
    const double eps = 1E-15;
    double x = 1;
    while (true)
    {
        double nx = (x + n / x) / 2;
        if (abs(x - nx) < eps)
            break;
        x = nx;
    }
    return x;
}</pre>
```

## 4.5 快速数论变换

```
//大数相乘
inline int read()
   int x = 0, f = 1;
   char ch = getchar();
   while (ch < '0' || ch > '9')
       if (ch == '-')
          f = -1;
       ch = getchar();
   while (ch <= '9' && ch >= '0')
       x = 10 * x + ch - '0';
       ch = getchar();
   return x * f;
void print(int x)
   if (x < 0)
       putchar('-'), x = -x;
   if (x >= 10)
       print(x / 10);
   putchar(x % 10 + '0');
const int N = 300100, P = 998244353;
inline int qpow(int x, int y)
   int res(1);
   while (y)
       if (y & 1)
          res = 1ll * res * x % P;
       x = 111 * x * x % P;
       y >>= 1;
```

```
}
   return res;
int r[N];
void ntt(int *x, int lim, int opt)
   int i, j, k, m, gn, g, tmp;
   for (i = 0; i < lim; ++i)
        if (r[i] < i)
           swap(x[i], x[r[i]]);
   for (m = 2; m \leftarrow 1im; m \leftarrow 1)
       k = m \gg 1;
       gn = qpow(3, (P - 1) / m);
       for (i = 0; i < lim; i += m)
           for (j = 0; j < k; ++j, g = 111 * g * gn % P)
               tmp = 111 * x[i + j + k] * g % P;
               x[i + j + k] = (x[i + j] - tmp + P) \% P;
               x[i + j] = (x[i + j] + tmp) % P;
            }
       }
   if (opt == -1)
       reverse(x + 1, x + lim);
       int inv = qpow(lim, P - 2);
       for (i = 0; i < lim; ++i)
           x[i] = 111 * x[i] * inv % P;
   }
int A[N], B[N], C[N];
char a[N], b[N];
int main()
   int i, lim(1), n;
   scanf("%s", &a);
   n = strlen(a);
   for (i = 0; i < n; ++i)
       A[i] = a[n - i - 1] - '0';
   while (\lim < (n << 1))
       lim <<= 1;
   scanf("%s", &b);
   n = strlen(b);
   for (i = 0; i < n; ++i)
       B[i] = b[n - i - 1] - '0';
   while (lim < (n << 1))
       lim <<= 1;
   for (i = 0; i < lim; ++i)
```

```
r[i] = (i \& 1) * (lim >> 1) + (r[i >> 1] >> 1);
ntt(A, lim, 1);
ntt(B, lim, 1);
for (i = 0; i < lim; ++i)
    C[i] = 111 * A[i] * B[i] % P;
ntt(C, lim, -1);
int len(0);
for (i = 0; i < lim; ++i)
    if (C[i] >= 10)
       len = i + 1, C[i + 1] += C[i] / 10, C[i] %= 10;
    if (C[i])
       len = max(len, i);
while (C[len] >= 10)
   C[len + 1] += C[len] / 10, C[len] %= 10, len++;
for (i = len; \sim i; --i)
    putchar(C[i] + '0');
puts("");
return 0;
```

# 4.6 线性同余方程

```
int ex_gcd(int a, int b, int &x, int &y)
   if (b == 0)
   {
       x = 1;
       y = 0;
       return a;
   int d = ex_gcd(b, a \% b, x, y);
   int temp = x;
   x = y;
   y = temp - a / b * y;
   return d;
bool liEu(int a, int b, int c, int &x, int &y)
   int d = ex_gcd(a, b, x, y);
   if (c % d != 0)
       return 0;
   int k = c / d;
   x *= k;
   y *= k;
   return 1;
```

### 4.7 中国剩余定理

```
LL CRT(int k, LL *a, LL *r)
{
    LL n = 1, ans = 0;
    for (int i = 1; i <= k; i++)
        n = n * r[i];
    for (int i = 1; i <= k; i++)
    {
        LL m = n / r[i], b, y;
        exgcd(m, r[i], b, y); // b * m mod r[i] = 1
        ans = (ans + a[i] * m * b % n) % n;
    }
    return (ans % n + n) % n;
}</pre>
```

### 4.8 扩展欧几里得算法

```
void exgcd(int a, int b, int& x, int& y) {
    if (b == 0) {
        x = 1, y = 0;
        return;
    }
    exgcd(b, a % b, y, x);
    y -= a / b * x;
}
```

## 4.9 高斯消元

## 4.10 高精度计算

#include <cstdio>

```
#include <cstring>
static const int LEN = 1004;
int a[LEN], b[LEN], c[LEN], d[LEN];
void clear(int a[])
   for (int i = 0; i < LEN; ++i)
       a[i] = 0;
void read(int a[])
   static char s[LEN + 1];
   scanf("%s", s);
   clear(a);
   int len = strlen(s);
   for (int i = 0; i < len; ++i)
       a[len - i - 1] = s[i] - '0';
void print(int a[])
   int i;
   for (i = LEN - 1; i >= 1; --i)
       if (a[i] != 0)
           break;
   for (; i >= 0; --i)
       putchar(a[i] + '0');
   putchar('\n');
void add(int a[], int b[], int c[])
   clear(c);
   for (int i = 0; i < LEN - 1; ++i)
       c[i] += a[i] + b[i];
       if (c[i] >= 10)
           c[i + 1] += 1;
           c[i] -= 10;
       }
   }
void sub(int a[], int b[], int c[])
   clear(c);
   for (int i = 0; i < LEN - 1; ++i)
```

```
c[i] += a[i] - b[i];
       if (c[i] < 0)
           c[i + 1] -= 1;
           c[i] += 10;
       }
   }
void mul(int a[], int b[], int c[])
   clear(c);
   for (int i = 0; i < LEN - 1; ++i)
       for (int j = 0; j <= i; ++j)
           c[i] += a[j] * b[i - j];
       if (c[i] >= 10)
           c[i + 1] += c[i] / 10;
           c[i] %= 10;
       }
   }
inline bool greater_eq(int a[], int b[], int last_dg, int len)
   if (a[last_dg + len] != 0)
       return true;
   for (int i = len - 1; i >= 0; --i)
       if (a[last_dg + i] > b[i])
           return true;
       if (a[last_dg + i] < b[i])
           return false;
   return true;
void div(int a[], int b[], int c[], int d[])
   clear(c);
   clear(d);
   int la, lb;
   for (la = LEN - 1; la > 0; --la)
       if (a[la - 1] != 0)
           break;
   for (lb = LEN - 1; lb > 0; --lb)
       if (b[lb - 1] != 0)
           break;
   if (1b == 0)
       puts("> <");</pre>
       return;
```

```
}
   for (int i = 0; i < la; ++i)
        d[i] = a[i];
   for (int i = la - lb; i >= 0; --i)
        while (greater_eq(d, b, i, lb))
            for (int j = 0; j < lb; ++j)
                d[i + j] -= b[j];
                if (d[i + j] < 0)
                    d[i + j + 1] -= 1;
                    d[i + j] += 10;
            }
            c[i] += 1;
        }
   }
int main()
   read(a);
   char op[4];
   scanf("%s", op);
   read(b);
   switch (op[0])
    case '+':
        add(a, b, c);
        print(c);
        break;
    case '-':
        sub(a, b, c);
        print(c);
        break;
    case '*':
        mul(a, b, c);
        print(c);
        break;
    case '/':
        div(a, b, c, d);
        print(c);
        print(d);
        break;
   default:
        puts("> <");</pre>
    }
   return 0;
```

# 5 动态规划 Dynamic Programming

### 5.1 背包 DP

```
// 0-1 背包问题母代码(二维)
void bags()
   vector<int> weight = {1, 3, 4}; //各个物品的重量
   vector<int> value = {15, 20, 30}; //对应的价值
   int bagWeight = 4;
                              //背包最大能放下多少重的物品
   // 二维数组: 状态定义:dp[i][j]表示从 0-i 个物品中选择不超过 j 重量的物品的最
大价值
   vector<vector<int>> dp(weight.size() + 1, vector<int>(bagWeight + 1,
0));
   // 初始化:第一列都是 0,第一行表示只选取 0 号物品最大价值
   for (int j = bagWeight; j >= weight[0]; j--)
      dp[0][j] = dp[0][j - weight[0]] + value[0];
   // weight 数组的大小 就是物品个数
   for (int i = 1; i < weight.size(); i++) // 遍历物品(第 0 个物品已经初始
化)
   {
      for (int j = 0; j <= bagWeight; j++) // 遍历背包容量
         if (j < weight[i])</pre>
                                 //背包容量已经不足以拿第 i 个物品了
            dp[i][j] = dp[i - 1][j]; //最大价值就是拿第 i-1 个物品的最大
价值
          //背包容量足够拿第 i 个物品,可拿可不拿: 拿了最大价值是前 i-1 个物品
扣除第 1 个物品的 重量的最大价值加上 1 个物品的价值
         //不拿就是前 i-1 个物品的最大价值,两者进行比较取较大的
         else
             dp[i][j] = max(dp[i - 1][j], dp[i - 1][j - weight[i]] +
value[i]);
      }
   cout << dp[weight.size() - 1][bagWeight] << endl;</pre>
// 首先是背包分类的模板:
// 1、0/1 背包: 外循环 nums,内循环 target,target 倒序且 target>=nums[i];
// 2、完全背包:外循环 nums,内循环 target,target 正序且 target>=nums[i];
// 3、组合背包:外循环 target,内循环 nums,target 正序且 target>=nums[i];
// 4、分组背包: 这个比较特殊,需要三重循环: 外循环背包 bags,内部两层循环根据题目
的要求转化为 1,2,3 三种背包类型的模板
// 然后是问题分类的模板:
```

```
// 1、最值问题: dp[i] = max/min(dp[i], dp[i-nums]+1)或 dp[i] = max/min(dp[i], dp[i-num]+nums);
// 2、存在问题(bool): dp[i]=dp[i]||dp[i-num];
// 3、组合问题: dp[i]+=dp[i-num];
```

### 5.2 数位 DP

```
//给定两个正整数 a,b, 求在 [a,b] 中的所有整数中,每个数码(digit) 各出现了多少
次。
#include <bits/stdc++.h>
using namespace std;
const int N = 15;
typedef long long 11;
11 1, r, dp[N], sum[N], mi[N];
ll ans1[N], ans2[N];
int a[N];
inline void solve(ll n, ll *ans) {
 11 \text{ tmp} = n;
 int len = 0;
 while (n) a[++len] = n \% 10, n \neq 10;
 for (int i = len; i >= 1; --i) {
   for (int j = 0; j < 10; j++) ans[j] += dp[i - 1] * a[i];
   for (int j = 0; j < a[i]; j++) ans[j] += mi[i - 1];
   tmp -= mi[i - 1] * a[i], ans[a[i]] += tmp + 1;
   ans[0] -= mi[i - 1];
 }
int main() {
 scanf("%11d%11d", &1, &r);
 mi[0] = 111;
 for (int i = 1; i <= 13; ++i) {
   dp[i] = dp[i - 1] * 10 + mi[i - 1];
   mi[i] = 1011 * mi[i - 1];
 solve(r, ans1), solve(l - 1, ans2);
 for (int i = 0; i < 10; ++i) printf("%lld ", ans1[i] - ans2[i]);</pre>
 return 0:
```

#### 5.3 插头 DP

```
//在 N * M 的棋盘内铺满 1 * 2 或 2 * 1 的多米诺骨牌,求方案数。
#include <bits/stdc++.h>
using namespace std;
const int N = 11;
long long f[2][1 << N], *f0, *f1;
int n, m;
int main() {</pre>
```

```
while (cin >> n >> m && n) {
   f0 = f[0];
   f1 = f[1];
   fill(f1, f1 + (1 << m), 0);
   f1[0] = 1;
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < m; ++j) {
       swap(f0, f1);
       fill(f1, f1 + (1 << m), 0);
#define u f0[s]
       for (int s = 0; s < 1 << m; ++s)
         if (u) {
            if (j != m - 1 \&\& (!(s >> j \& 3))) f1[s ^ 1 << j + 1] +=
u; // 横放
           f1[s ^ 1 << j] += u; // 竖放或不放
         }
     }
   cout << f1[0] << endl;</pre>
```

## 5.4 线性 DP

```
//最长上升子序列 LIS
const int N = 1e4 + 1;
int a[N];
int d[N];
int main() {
 int n;
 cin >> n;
 for (int i = 1;i <= n;i++)
   cin >> a[i];
 d[1] = a[1];//将第一个元素存入
 int len = 1;
 for (int i = 2; i <= n; i++) {
     if (a[i] > d[len]) d[++len] = a[i];
     else {
       int j = lower_bound(d+1,d + len + 1,a[i]) - d;
       d[j] = a[i];//替换序列中比他第一个大的元素
 }
 cout << len;</pre>
 return 0;
```

# 6 计算几何 Computational Geometry

## 6.1 常数定义相关

```
//long double 的输入输出
scanf("%Lf" , &a);
printf("%.10Lf" , a);
//常用函数:fabsl(a),cosl(a).....
//即在末尾加上了字母 l
//常数定义
const double eps = 1e-8;
const double PI = acos(-1.0);

int sgn(double x)//符号函数, eps 使用最多的地方
{
    if (fabs(x) < eps)
        return 0;
    if (x < 0)
        return -1;
    else
        return 1;
}
```

### 6.2 点类

```
struct Point
                  //点类
   double x, y;
   Point() {}
   Point(double \underline{x}, double \underline{y}) : x(\underline{x}), y(\underline{y}) {}
    Point operator-(const Point &b) const { return Point(x - b.x, y -
b.y); }
    Point operator+(const Point &b) const { return Point(x + b.x, y +
b.y); }
   double operator^(const Point &b) const { return x * b.y - y * b.x; }
   double operator*(const Point &b) const { return x * b.x + y * b.y; }
//点积
   bool operator (const Point &b) const { return x < b.x \mid | (x == b.x \&\&
y < b.y); }
    bool operator==(const Point &b) const { return sgn(x - b.x) == 0 \&\&
sgn(y - b.y) == 0; }
   Point Rotate(double B, Point P) //绕着点P, 逆时针旋转角度B(弧度)
   {
       Point tmp;
       tmp.x = (x - P.x) * cos(B) - (y - P.y) * sin(B) + P.x;
       tmp.y = (x - P.x) * sin(B) + (y - P.y) * cos(B) + P.y;
        return tmp;
   }
};
double dist(Point a, Point b) { return sqrt((a - b) * (a - b)); } //两点
```

```
间距离
double len(Point a){return sqrt(a.x * a.x + a.y * a.y);}//向量的长度
```

### 6.3 直线类

```
struct Line
              //直线类
   Point s, e;
   Line() {}
   Line(Point _s, Point _e) : s(_s), e(_e) {}
   //两直线相交求交点
   //第一个值为0表示直线重合,为1表示平行,为2是相交
   //只有第一个值为2时,交点才有意义
   pair<int, Point> operator&(const Line &b) const
   {
       Point res = s;
       if (sgn((s - e) ^ (b.s - b.e)) == 0)
          if (sgn((s - b.e) ^ (b.s - b.e)) == 0)
              return make_pair(0, res); //重合
          else
              return make_pair(1, res); //平行
       double t = ((s - b.s) ^ (b.s - b.e)) / ((s - e) ^ (b.s - b.e));
       res.x += (e.x - s.x) * t;
       res.y += (e.y - s.y) * t;
       return make_pair(2, res);
};
```

## 6.4 相交关系判断

### 6.5 距离相关计算

```
Point PointToLine(Point P, Line L) //求点到直线的距离
   Point result;
   double t = ((P - L.s) * (L.e - L.s)) / ((L.e - L.s) * (L.e - L.s));
   result.x = L.s.x + (L.e.x - L.s.x) * t;
   result.y = L.s.y + (L.e.y - L.s.y) * t;
   return result;
Point NearestPointToLineSeg(Point P, Line L) // 求点到线段的距离
   Point result;
   double t = ((P - L.s) * (L.e - L.s)) / ((L.e - L.s) * (L.e - L.s));
   if (t >= 0 \&\& t <= 1)
       result.x = L.s.x + (L.e.x - L.s.x) * t;
       result.y = L.s.y + (L.e.y - L.s.y) * t;
   }
   else
   {
       if (dist(P, L.s) < dist(P, L.e))</pre>
           result = L.s;
       else
           result = L.e;
   return result;
```

## 6.6 点和直线相关

```
//计算多边形面积,点的编号从 0~n-1
double CalcArea(Point p[], int n)
{
    double res = 0;
    for (int i = 0; i < n; i++)
        res += (p[i] ^ p[(i + 1) % n]) / 2;
    return fabs(res);
}
//*判断点在线段上
bool OnSeg(Point P, Line L)
{
    return sgn((L.s - P) ^ (L.e - P)) == 0 &&
        sgn((P.x - L.s.x) * (P.x - L.e.x)) <= 0 &&
        sgn((P.y - L.s.y)) * (P.y - L.e.y)) <= 0;
}</pre>
```

# 6.7 凸包相关

```
int ConvexHull(Point *p, int n, Point *ch) //求凸包Andrew 算法
   sort(p, p + n);
   n = unique(p, p + n) - p; // 去重
   int m = 0;
   for (int i = 0; i < n; ++i)
       while (m > 1 \&\& sgn((ch[m - 1] - ch[m - 2]) \land (p[i] - ch[m - 1]))
<= 0)
            --m;
       ch[m++] = p[i];
   int k = m;
   for (int i = n - 2; i >= 0; i--)
       while (m > k \&\& sgn((ch[m - 1] - ch[m - 2]) \land (p[i] - ch[m - 1]))
<= 0)
            --m;
       ch[m++] = p[i];
   if (n > 1)
       m--;
   return m;
```

### 6.8 极角排序

```
// 叉积: 对于 tmp = a x b
// 如果 b 在 a 的逆时针(左边):tmp > 0
// 顺时针(右边): tmp < 0
// 同向: tmp = 0
// 相对于原点的极角排序
// 如果是相对于某一点 x,只需要把 x 当作原点即可
bool mycmp(Point a, Point b)
{
   if (atan2(a.y, a.x) != atan2(b.y, b.x))
      return atan2(a.y, a.x) < atan2(b.y, b.x);
   else
      return a.x < b.x;
}</pre>
```

# 6.9 点和多边形的位置关系

```
// 点形成一个凸包,而且按逆时针排序
// 如果是顺时针把里面的<0 改为>0
// 点的编号:0~n-1
// 返回值:
// -1:点在凸多边形外
// 0:点在凸多边形边界上
```

```
// 1:点在凸多边形内
int inConvexPoly(Point a, Point p[], int n)
   for (int i = 0; i < n; i++)
   {
       if (sgn((p[i] - a) ^ (p[(i + 1) % n] - a)) < 0)
          return -1;
       else if (OnSeg(a, Line(p[i], p[(i + 1) % n])))
          return 0;
   return 1;
1//判断点是否在凸包内
bool inConvex(Point A, Point *p, int tot)
   int l = 1, r = tot - 2, mid;
   while (1 <= r)
       mid = (l + r) >> 1;
       double a1 = (p[mid] - p[0]) ^ (A - p[0]);
       double a2 = (p[mid + 1] - p[0]) ^ (A - p[0]);
       if (a1 >= 0 \&\& a2 <= 0)
       {
           if (((p[mid + 1] - p[mid]) ^ (A - p[mid])) >= 0)
              return true;
          return false;
       }
       else if (a1 < 0)
          r = mid - 1;
       else
           1 = mid + 1;
   return false;
// 判断点在任意多边形内
// 射线法, poly[]的顶点数要大于等于 3,点的编号 0~n-1
// 返回值
// -1:点在凸多边形外
// 0:点在凸多边形边界上
// 1:点在凸多边形内
int inPoly(Point p, Point poly[], int n)
   int cnt;
   Line ray, side;
   cnt = 0;
   ray.s = p;
   ray.e.y = p.y;
   ray.e.x = -1000000000000.0; //-INF,注意取值防止越界
   for (int i = 0; i < n; i++)
   {
       side.s = poly[i];
```

```
side.e = poly[(i + 1) % n];
       if (OnSeg(p, side))
           return 0;
       //如果平行轴则不考虑
       if (sgn(side.s.y - side.e.y) == 0)
           continue;
       if (OnSeg(side.s, ray))
           if (sgn(side.s.y - side.e.y) > 0)
               cnt++;
       else if (OnSeg(side.e, ray))
           if (sgn(side.e.y - side.s.y) > 0)
               cnt++;
       else if (inter(ray, side))
           cnt++;
   if (cnt % 2 == 1)
       return 1;
   else
       return -1;
// 判断凸多边形
bool isconvex(Point poly[], int n)
   bool s[3];
   memset(s, false, sizeof(s));
   for (int i = 0; i < n; i++)
         s[sgn((poly[(i + 1) % n] - poly[i]) ^ (poly[(i + 2) % n] -
poly[i])) + 1] = true;
       if (s[0] && s[2])
           return false;
   return true;
//判断凸包是否相离 bool isConvexHullSeparate(int n, int m, Point a[], Point
b[])
   for (int i = 0; i < n; i++)
       if (inPoly(a[i], b, m) != -1)
           return false;
   for (int i = 0; i < m; i++)
       if (inPoly(b[i], a, n) != -1)
           return false;
   for (int i = 0; i < n; i++)
```

### 6.10 闵可夫斯基和

```
const int MAX = 2e5 + 5;
Point s1[MAX], s2[MAX];
int Minkowski(Point A[], int n, Point B[], int m, Point M[])
   int tot = 0;
   for (int i = 0; i < n; i++)
       s1[i] = A[(i + 1) \% n] - A[i];
   for (int i = 0; i < m; i++)
       s2[i] = B[(i + 1) \% m] - B[i];
   M[tot] = A[0] + B[0];
   int p1 = 0, p2 = 0;
   while (p1 < n \&\& p2 < m)
       ++tot, M[tot] = M[tot - 1] + ((s1[p1] ^ s2[p2]) >= 0 ? s1[p1++] :
s2[p2++]);
   while (p1 < n)
       ++tot, M[tot] = M[tot - 1] + s1[p1++];
   while (p2 < m)
       ++tot, M[tot] = M[tot - 1] + s2[p2++];
   return tot + 1;
```

# 7 图论 Graph Theory

### 7.1 拓扑排序

```
int n, m;
vector<int> G[MAXN];
int in[MAXN]; // 存储每个结点的入度
bool toposort()
{
vector<int> L;
```

```
queue<int> S;
for (int i = 1; i <= n; i++)
    if (in[i] == 0)
        S.push(i);
while (!S.empty())
{
    int u = S.front();
    S.pop();
    L.push_back(u);
    for (auto v : G[u])
        if (--in[v] == 0)
           S.push(v);
    }
if(L.size() == n)
    for (auto i : L)
       cout << i << ' ';
    return true;
}
else
{
    return false;
}
```

# 7.2 最短路

#### 7.2.1 Bellman-Ford 算法

```
if (dis[u] == inf)
         continue;
      // 无穷大与常数加减仍然为无穷大
      // 因此最短路长度为 inf 的点引出的边不可能发生松弛操作
      for (auto ed : e[u])
         int v = ed.v, w = ed.w;
         if (dis[v] > dis[u] + w)
            dis[v] = dis[u] + w;
            flag = true;
         }
      }
   }
   // 没有可以松弛的边时就停止算法
   if (!flag)
      break;
}
// 第 n 轮循环仍然可以松弛时说明 s 点可以抵达一个负环
return flag;
```

#### 7.2.2 Dijkstra 算法

```
struct edge
   int v, w;
};
struct node
   int dis, u;
   bool operator>(const node &a) const { return dis > a.dis; }
};
vector<edge> e[maxn];
int dis[maxn], vis[maxn];
priority_queue<node, vector<node>, greater<node>> q;
void dijkstra(int n, int s)
   memset(dis, 63, sizeof(dis));
   dis[s] = 0;
   q.push({0, s});
   while (!q.empty())
       int u = q.top().u;
       q.pop();
       if (vis[u])
           continue;
       vis[u] = 1;
       for (auto ed : e[u])
       {
```

```
int v = ed.v, w = ed.w;
if (dis[v] > dis[u] + w)
{
          dis[v] = dis[u] + w;
          q.push({dis[v], v});
          }
}
```

## 7.3 强连通分量

```
int dfn[N], low[N], dfncnt, s[N], in_stack[N], tp;
int scc[N], sc; // 结点 i 所在 SCC 的编号
int sz[N];
               // 强连通 i 的大小
void tarjan(int u)
   low[u] = dfn[u] = ++dfncnt, s[++tp] = u, in_stack[u] = 1;
   for (int i = h[u]; i; i = e[i].nex)
       const int &v = e[i].t;
       if (!dfn[v])
           tarjan(v);
           low[u] = min(low[u], low[v]);
       else if (in_stack[v])
           low[u] = min(low[u], dfn[v]);
       }
   if (dfn[u] == low[u])
       ++sc;
       while (s[tp] != u)
           scc[s[tp]] = sc;
           sz[sc]++;
           in_stack[s[tp]] = 0;
           --tp;
       }
       scc[s[tp]] = sc;
       sz[sc]++;
       in_stack[s[tp]] = 0;
       --tp;
   }
```

## 7.4 割点和桥

#### 7.4.1 割点

```
#include <bits/stdc++.h>
using namespace std;
int n, m; // n: 点数 m: 边数
int dfn[100001], low[100001], inde, res;
// dfn:记录每个点的时间戳
// Low: 能不经过父亲到达最小的编号,inde: 时间戳,res: 答案数量
bool vis[100001], flag[100001]; // flag: 答案 vis: 标记是否重复
vector<int> edge[100001];
                          // 存图用的
void Tarjan(int u, int father)
                        // u 当前点的编号, father 自己爸爸的编号
   vis[u] = true;
                         // 标记
   low[u] = dfn[u] = ++inde; // 打上时间戳
   int child = 0;
                         // 每一个点儿子数量
   for (auto v : edge[u])
   { // 访问这个点的所有邻居 (C++11)
      if (!vis[v])
      {
         child++;
                                  // 多了一个儿子
         Tarjan(v, u);
                                   // 继续
         low[u] = min(low[u], low[v]); // 更新能到的最小节点编号
         if (father != u \&\& low[v] >= dfn[u] \&\&
             !flag
                [u]) // 主要代码
                     // 如果不是自己,且不通过父亲返回的最小点符合割点的
要求,并且没有被标记过
                    // 要求即为: 删了父亲连不上去了, 即为最多连到父亲
         {
             flag[u] = true;
             res++; // 记录答案
      else if (v != father)
         low[u] =
             min(low[u], dfn[v]); // 如果这个点不是自己,更新能到的最小节
点编号
   if (father == u \&\& child >= 2 \&\&
      !flag[u])
   { // 主要代码,自己的话需要 2 个儿子才可以
      flag[u] = true;
      res++; // 记录答案
   }
int main()
```

```
cin >> n >> m; // 读入数据
for (int i = 1; i <= m; i++)
{ // 注意点是从 1 开始的
   int x, y;
   cin >> x >> y;
   edge[x].push_back(y);
   edge[y].push back(x);
                          // 使用 vector 存图
}
for (int i = 1; i <= n; i++) // 因为 Tarjan 图不一定连通
   if (!vis[i])
   {
                    // 时间戳初始为 0
       inde = 0;
       Tarjan(i, i); // 从第 i 个点开始,父亲为自己
   }
cout << res << endl;</pre>
for (int i = 1; i <= n; i++)
   if (flag[i])
       cout << i << " "; // 输出结果
return 0;
```

### 7.4.2 割边

```
int low[MAXN], dfn[MAXN], dfs_clock;
bool isbridge[MAXN];
vector<int> G[MAXN];
int cnt_bridge;
int father[MAXN];
void tarjan(int u, int fa)
   father[u] = fa;
   low[u] = dfn[u] = ++dfs_clock;
   for (int i = 0; i < G[u].size(); i++)
       int v = G[u][i];
       if (!dfn[v])
           tarjan(v, u);
           low[u] = min(low[u], low[v]);
           if (low[v] > dfn[u])
           {
               isbridge[v] = true;
               ++cnt_bridge;
           }
       }
       else if (dfn[v] < dfn[u] && v != fa)
           low[u] = min(low[u], dfn[v]);
       }
   }
```

### 7.5 欧拉图

```
// 欧拉回路:通过图中每条边恰好一次的回路
// 欧拉通路:通过图中每条边恰好一次的通路
// 欧拉图: 具有欧拉回路的图
// 半欧拉图: 具有欧拉通路但不具有欧拉回路的图
// 给定一张有 500 个顶点的无向图,求这张图的一条欧拉路或欧拉回路。如果有多组解,
输出最小的那一组。
// 在本题中,欧拉路或欧拉回路不需要经过所有顶点。
#include <algorithm>
#include <cstdio>
#include <stack>
#include <vector>
using namespace std;
struct edge {
 int to;
 bool exists;
 int revref;
 bool operator<(const edge& b) const { return to < b.to; }</pre>
};
vector<edge> beg[505];
int cnt[505];
const int dn = 500;
stack<int> ans;
|void Hierholzer(int x) { // 关键函数
 for (int& i = cnt[x]; i < (int)beg[x].size();) {
   if (beg[x][i].exists) {
     edge e = beg[x][i];
     beg[x][i].exists = 0;
     beg[e.to][e.revref].exists = 0;
     ++i;
     Hierholzer(e.to);
   } else {
     ++i;
   }
 ans.push(x);
int deg[505];
int reftop[505];
int main() {
 for (int i = 1; i <= dn; ++i) {
   beg[i].reserve(1050); // vector 用 reserve 避免动态分配空间,加快速度
 }
```

```
int m;
scanf("%d", &m);
for (int i = 1; i <= m; ++i) {
  int a, b;
  scanf("%d%d", &a, &b);
  beg[a].push_back((edge){b, 1, 0});
  beg[b].push_back((edge){a, 1, 0});
  ++deg[a];
 ++deg[b];
}
for (int i = 1; i <= dn; ++i) {
 if (!beg[i].empty()) {
    sort(beg[i].begin(), beg[i].end()); // 为了要按字典序贪心,必须排序
  }
}
for (int i = 1; i <= dn; ++i) {
 for (int j = 0; j < (int)beg[i].size(); ++j) {</pre>
    beg[i][j].revref = reftop[beg[i][j].to]++;
  }
}
int bv = 0;
for (int i = 1; i <= dn; ++i) {
  if (!deg[bv] && deg[i]) {
   bv = i;
  } else if (!(deg[bv] & 1) && (deg[i] & 1)) {
    bv = i;
  }
}
Hierholzer(bv);
while (!ans.empty()) {
 printf("%d\n", ans.top());
 ans.pop();
}
```

## 7.6 最小环

```
// 给出一个图,问其中的有 n 个节点构成的边权和最小的环 (n >= 3) 是多大。
int val[maxn + 1][maxn + 1]; // 原图的邻接矩阵

inline int floyd(const int &n)
{
    static int dis[maxn + 1][maxn + 1]; // 最短路矩阵
    for (int i = 1; i <= n; ++i)
        for (int j = 1; j <= n; ++j)
        dis[i][j] = val[i][j]; // 初始化最短路矩阵
    int ans = inf;</pre>
```

### 7.7 树哈希

```
// 给定一棵以点1为根的树,你需要输出这棵树中最多能选出多少个互不同构的子树。
// 两棵有根树 T1,T2 同构当且仅当他们的大小相等,且存在一个顶点排列 σ 使得在 T1 中
i 是 j 的祖先当且仅当在 T2 中 σi 是 σj 的祖先。
#include <cctype>
#include <chrono>
#include <cstdio>
#include <random>
#include <set>
#include <vector>
typedef unsigned long long ull;
const ull mask =
std::chrono::steady_clock::now().time_since_epoch().count();
ull shift(ull x)
   x ^= mask;
   x ^= x << 13;
   x ^= x >> 7;
   x ^= x << 17;
   x ^= mask;
   return x;
const int N = 1e6 + 10;
int n;
ull hash[N];
std::vector<int> edge[N];
std::set<ull> trees;
void getHash(int x, int p)
   hash[x] = 1;
   for (int i : edge[x])
```

```
{
       if (i == p)
           continue;
       getHash(i, x);
       hash[x] += shift(hash[i]);
   trees.insert(hash[x]);
int main()
   scanf("%d", &n);
   for (int i = 1; i < n; i++)
       int u, v;
       scanf("%d%d", &u, &v);
       edge[u].push_back(v);
       edge[v].push_back(u);
   }
   getHash(1, 0);
   printf("%lu", trees.size());
```

## 7.8 最小树形图

```
bool solve()
   ans = 0;
   int u, v, root = 0;
   for (;;)
       f(i, 0, n) in[i] = 1e100;
       f(i, 0, m)
           u = e[i].s;
           v = e[i].t;
           if (u != v \&\& e[i].w < in[v])
           {
               in[v] = e[i].w;
               pre[v] = u;
            }
       f(i, 0, m) if (i != root && in[i] > 1e50) return 0;
       int tn = 0;
       memset(id, -1, sizeof id);
       memset(vis, -1, sizeof vis);
       in[root] = 0;
       f(i, 0, n)
       {
           ans += in[i];
           v = i;
           while (vis[v] != i && id[v] == -1 && v != root)
```

```
{
           vis[v] = i;
           v = pre[v];
       if (v != root && id[v] == -1)
           for (int u = pre[v]; u != v; u = pre[u])
               id[u] = tn;
           id[v] = tn++;
        }
    if (tn == 0)
       break;
    f(i, 0, n) if (id[i] == -1) id[i] = tn++;
   f(i, 0, m)
       u = e[i].s;
       v = e[i].t;
       e[i].s = id[u];
       e[i].t = id[v];
       if (e[i].s != e[i].t)
           e[i].w -= in[v];
    }
   n = tn;
   root = id[root];
return ans;
```

## 7.9 最小直径生成树

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 502;
typedef long long 11;
typedef pair<int, int> pii;
11 d[MAXN][MAXN], dd[MAXN][MAXN], rk[MAXN][MAXN], val[MAXN];
const ll INF = 1e17;
int n, m;
bool cmp(int a, int b) { return val[a] < val[b]; }</pre>
void floyd()
   for (int k = 1; k <= n; k++)
       for (int i = 1; i <= n; i++)
           for (int j = 1; j <= n; j++)
               d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
struct node
   11 u, v, w;
} a[MAXN * (MAXN - 1) / 2];
```

```
void solve()
   // 求图的绝对中心
   floyd();
   for (int i = 1; i <= n; i++)
       for (int j = 1; j <= n; j++)
       {
           rk[i][j] = j;
           val[j] = d[i][j];
       sort(rk[i] + 1, rk[i] + 1 + n, cmp);
   11 P = 0, ans P = INF;
   // 在点上
   for (int i = 1; i <= n; i++)
   {
       if (d[i][rk[i][n]] * 2 < ansP)</pre>
           ansP = d[i][rk[i][n]] * 2;
           P = i;
       }
   }
   // 在边上
   int f1 = 0, f2 = 0;
   11 disu = INT_MIN, disv = INT_MIN, ansL = INF;
   for (int i = 1; i <= m; i++)
   {
       ll u = a[i].u, v = a[i].v, w = a[i].w;
       for (int p = n, i = n - 1; i >= 1; i--)
       {
           if (d[v][rk[u][i]] > d[v][rk[u][p]])
           {
               if (d[u][rk[u][i]] + d[v][rk[u][p]] + w < ansL)
               {
                   ansL = d[u][rk[u][i]] + d[v][rk[u][p]] + w;
                   f1 = u, f2 = v;
                     disu = (d[u][rk[u][i]] + d[v][rk[u][p]] + w) / 2 -
d[u][rk[u][i]];
                   disv = w - disu;
               }
               p = i;
           }
       }
   cout << min(ansP, ansL) / 2 << '\n';</pre>
   // 最小路径生成树
   vector<pii> pp;
   for (int i = 1; i <= 501; ++i)
       for (int j = 1; j <= 501; ++j)
           dd[i][j] = INF;
   for (int i = 1; i <= 501; ++i)
       dd[i][i] = 0;
   if (ansP <= ansL)</pre>
```

```
{
   for (int j = 1; j <= n; j++)
       for (int i = 1; i <= m; ++i)
           ll u = a[i].u, v = a[i].v, w = a[i].w;
           if (dd[P][u] + w == d[P][v] && dd[P][u] + w < dd[P][v])
               dd[P][v] = dd[P][u] + w;
               pp.push_back({u, v});
           u = a[i].v, v = a[i].u, w = a[i].w;
           if (dd[P][u] + w == d[P][v] && dd[P][u] + w < dd[P][v])
               dd[P][v] = dd[P][u] + w;
               pp.push_back({u, v});
           }
        }
   }
   for (auto [x, y] : pp)
       cout << x << ' ' << y << '\n';
}
else
{
   d[n + 1][f1] = disu;
   d[f1][n + 1] = disu;
   d[n + 1][f2] = disv;
   d[f2][n + 1] = disv;
    a[m + 1].u = n + 1, a[m + 1].v = f1, a[m + 1].w = disu;
   a[m + 2].u = n + 1, a[m + 2].v = f2, a[m + 2].w = disv;
   n += 1;
   m += 2;
   floyd();
   P = n;
   for (int j = 1; j <= n; j++)
        for (int i = 1; i <= m; ++i)
           ll u = a[i].u, v = a[i].v, w = a[i].w;
           if (dd[P][u] + w == d[P][v] && dd[P][u] + w < dd[P][v])
               dd[P][v] = dd[P][u] + w;
               pp.push_back({u, v});
           u = a[i].v, v = a[i].u, w = a[i].w;
           if (dd[P][u] + w == d[P][v] && dd[P][u] + w < dd[P][v])
               dd[P][v] = dd[P][u] + w;
               pp.push_back({u, v});
           }
        }
    }
    cout << f1 << ' ' << f2 << '\n';
    for (auto [x, y] : pp)
       if (x != n \&\& y != n)
           cout << x << ' ' << y << '\n';
```

```
}
void init()
   for (int i = 1; i <= 501; ++i)
       for (int j = 1; j <= 501; ++j)
           d[i][j] = INF;
   for (int i = 1; i <= 501; ++i)
       d[i][i] = 0;
int main()
   init();
   cin >> n >> m;
   for (int i = 1; i <= m; ++i)
       11 u, v, w;
       cin >> u >> v >> w;
       w *= 2;
       d[u][v] = w, d[v][u] = w;
       a[i].u = u, a[i].v = v, a[i].w = w;
   }
   solve();
   return 0;
```

# 8 网络流 Flow Network

# 8.1 最大流

# 8.1.1 Edmonds-Karp 算法

```
int a[maxn], p[maxn]; // a: 点 x -> BFS 过程中最近接近点 x 的边给它的最
大流
                      // p: 点 x -> BFS 过程中最近接近点 x 的边
   void init(int n)
      for (int i = 0; i < n; i++)
          G[i].clear();
      edges.clear();
   }
   void AddEdge(int from, int to, int cap)
      edges.push_back(Edge(from, to, cap, 0));
      edges.push back(Edge(to, from, 0, 0));
      m = edges.size();
      G[from].push_back(m - 2);
      G[to].push_back(m - 1);
   }
   int Maxflow(int s, int t)
      int flow = 0;
      for (;;)
      {
          memset(a, 0, sizeof(a));
          queue<int> Q;
          Q.push(s);
          a[s] = INF;
          while (!Q.empty())
              int x = Q.front();
              Q.pop();
              for (int i = 0; i < G[x].size(); i++)
              { // 遍历以 x 作为起点的边
                 Edge &e = edges[G[x][i]];
                 if (!a[e.to] && e.cap > e.flow)
                     p[e.to] = G[x][i]; // G[x][i] 是最近接近点 e.to 的
训
                     a[e.to] =
                         min(a[x], e.cap - e.flow); // 最近接近点 e.to
的边赋给它的流
                     Q.push(e.to);
                 }
              }
              if (a[t])
                 break; // 如果汇点接受到了流,就退出 BFS
          if (!a[t])
               break; // 如果汇点没有接受到流,说明源点和汇点不在同一个连通
分量上
          for (int u = t; u != s;
```

## 8.1.2 Dinic 算法

```
struct MF
   struct edge
       int v, nxt, cap, flow;
   } e[N];
   int fir[N], cnt = 0;
   int n, S, T;
   11 maxflow = 0;
   int dep[N], cur[N];
   void init()
   {
       memset(fir, -1, sizeof fir);
       cnt = 0;
   }
   void addedge(int u, int v, int w)
       e[cnt] = \{v, fir[u], w, 0\};
       fir[u] = cnt++;
       e[cnt] = {u, fir[v], 0, 0};
       fir[v] = cnt++;
   }
   bool bfs()
       queue<int> q;
       memset(dep, 0, sizeof(int) * (n + 1));
       dep[S] = 1;
       q.push(S);
       while (q.size())
           int u = q.front();
           q.pop();
           for (int i = fir[u]; \sim i; i = e[i].nxt)
```

```
{
               int v = e[i].v;
               if ((!dep[v]) && (e[i].cap > e[i].flow))
                   dep[v] = dep[u] + 1;
                   q.push(v);
                }
            }
       }
       return dep[T];
   }
   int dfs(int u, int flow)
   {
       if ((u == T) || (!flow))
           return flow;
       int ret = 0;
       for (int \&i = cur[u]; \sim i; i = e[i].nxt)
           int v = e[i].v, d;
           if ((dep[v] == dep[u] + 1) \&\&
               (d = dfs(v, min(flow - ret, e[i].cap - e[i].flow))))
               ret += d;
               e[i].flow += d;
               e[i ^ 1].flow -= d;
               if (ret == flow)
                   return ret;
           }
        }
       return ret;
   }
   void dinic()
       while (bfs())
           memcpy(cur, fir, sizeof(int) * (n + 1));
           maxflow += dfs(S, INF);
       }
   }
} mf;
8.1.3 MPM 算法
struct MPM
   struct FlowEdge
   {
       int v, u;
       long long cap, flow;
       FlowEdge() {}
```

```
FlowEdge(int _v, int _u, long long _cap, long long _flow)
        : v(_v), u(_u), cap(_cap), flow(_flow) {}
    FlowEdge(int _v, int _u, long long _cap)
        : v(_v), u(_u), cap(_cap), flow(011) {}
};
const long long flow_inf = 1e18;
vector<FlowEdge> edges;
vector<char> alive;
vector<long long> pin, pout;
vector<list<int>> in, out;
vector<vector<int>> adj;
vector<long long> ex;
int n, m = 0;
int s, t;
vector<int> level;
vector<int> q;
int qh, qt;
void resize(int _n)
{
    n = _n;
    ex.resize(n);
    q.resize(n);
    pin.resize(n);
    pout.resize(n);
    adj.resize(n);
    level.resize(n);
    in.resize(n);
    out.resize(n);
}
MPM() {}
MPM(int _n, int _s, int _t)
{
    resize(_n);
    s = _s;
    t = _t;
}
void add_edge(int v, int u, long long cap)
{
    edges.push back(FlowEdge(v, u, cap));
    edges.push_back(FlowEdge(u, v, 0));
    adj[v].push_back(m);
    adj[u].push_back(m + 1);
    m += 2;
}
bool bfs()
{
    while (qh < qt)
        int v = q[qh++];
```

```
for (int id : adj[v])
               if (edges[id].cap - edges[id].flow < 1)</pre>
                   continue;
               if (level[edges[id].u] != -1)
                   continue;
               level[edges[id].u] = level[v] + 1;
               q[qt++] = edges[id].u;
           }
       }
       return level[t] != -1;
   }
   long long pot(int v) { return min(pin[v], pout[v]); }
   void remove node(int v)
   {
       for (int i : in[v])
           int u = edges[i].v;
           auto it = find(out[u].begin(), out[u].end(), i);
           out[u].erase(it);
           pout[u] -= edges[i].cap - edges[i].flow;
       for (int i : out[v])
           int u = edges[i].u;
           auto it = find(in[u].begin(), in[u].end(), i);
           in[u].erase(it);
           pin[u] -= edges[i].cap - edges[i].flow;
       }
   }
   void push(int from, int to, long long f, bool forw)
       qh = qt = 0;
       ex.assign(n, 0);
       ex[from] = f;
       q[qt++] = from;
       while (qh < qt)
           int v = q[qh++];
           if (v == to)
               break;
           long long must = ex[v];
           auto it = forw ? out[v].begin() : in[v].begin();
           while (true)
           {
               int u = forw ? edges[*it].u : edges[*it].v;
                      long long pushed = min(must, edges[*it].cap -
edges[*it].flow);
               if (pushed == 0)
                   break;
               if (forw)
               {
                   pout[v] -= pushed;
```

```
pin[u] -= pushed;
               }
               else
                {
                   pin[v] -= pushed;
                   pout[u] -= pushed;
               if (ex[u] == 0)
                   q[qt++] = u;
               ex[u] += pushed;
               edges[*it].flow += pushed;
               edges[(*it) ^ 1].flow -= pushed;
               must -= pushed;
               if (edges[*it].cap - edges[*it].flow == 0)
               {
                    auto jt = it;
                   ++jt;
                    if (forw)
                    {
                            in[u].erase(find(in[u].begin(), in[u].end(),
*it));
                       out[v].erase(it);
                    }
                    else
                    {
                         out[u].erase(find(out[u].begin(), out[u].end(),
*it));
                        in[v].erase(it);
                    it = jt;
               }
               else
                    break;
                if (!must)
                   break;
           }
       }
   }
   long long flow()
        long long ans = 0;
       while (true)
        {
           pin.assign(n, 0);
           pout.assign(n, 0);
           level.assign(n, -1);
            alive.assign(n, true);
            level[s] = 0;
           qh = 0;
           qt = 1;
           q[0] = s;
            if (!bfs())
               break;
           for (int i = 0; i < n; i++)
            {
```

```
out[i].clear();
               in[i].clear();
           for (int i = 0; i < m; i++)
               if (edges[i].cap - edges[i].flow == 0)
                   continue;
               int v = edges[i].v, u = edges[i].u;
               if (level[v] + 1 == level[u] && (level[u] < level[t] || u
== t))
               {
                   in[u].push_back(i);
                   out[v].push_back(i);
                   pin[u] += edges[i].cap - edges[i].flow;
                   pout[v] += edges[i].cap - edges[i].flow;
           pin[s] = pout[t] = flow_inf;
           while (true)
           {
               int v = -1;
               for (int i = 0; i < n; i++)
                   if (!alive[i])
                       continue;
                   if (v == -1 \mid | pot(i) < pot(v))
                       v = i;
               }
               if (v == -1)
                   break;
               if (pot(v) == 0)
                   alive[v] = false;
                   remove_node(v);
                   continue;
               long long f = pot(v);
               ans += f;
               push(v, s, f, false);
               push(v, t, f, true);
               alive[v] = false;
               remove_node(v);
           }
       return ans;
    }
};
8.1.4 ISAP 算法
```

```
struct Edge
{
   int from, to, cap, flow;
   Edge(int u, int v, int c, int f) : from(u), to(v), cap(c), flow(f) {}
```

```
};
bool operator<(const Edge &a, const Edge &b)
   return a.from < b.from || (a.from == b.from && a.to < b.to);</pre>
struct ISAP
   int n, m, s, t;
   vector<Edge> edges;
   vector<int> G[maxn];
   bool vis[maxn];
   int d[maxn];
   int cur[maxn];
   int p[maxn];
   int num[maxn];
   void AddEdge(int from, int to, int cap)
       edges.push_back(Edge(from, to, cap, 0));
       edges.push back(Edge(to, from, 0, 0));
       m = edges.size();
       G[from].push back(m - 2);
       G[to].push_back(m - 1);
   }
   bool BFS()
   {
       memset(vis, 0, sizeof(vis));
       queue<int> Q;
       Q.push(t);
       vis[t] = 1;
       d[t] = 0;
       while (!Q.empty())
            int x = Q.front();
           Q.pop();
           for (int i = 0; i < G[x].size(); i++)
               Edge &e = edges[G[x][i] ^ 1];
               if (!vis[e.from] && e.cap > e.flow)
                   vis[e.from] = 1;
                   d[e.from] = d[x] + 1;
                   Q.push(e.from);
               }
            }
       return vis[s];
   }
   void init(int n)
       this->n = n;
       for (int i = 0; i < n; i++)
```

```
G[i].clear();
   edges.clear();
}
int Augment()
{
    int x = t, a = INF;
   while (x != s)
       Edge &e = edges[p[x]];
       a = min(a, e.cap - e.flow);
       x = edges[p[x]].from;
    }
   x = t;
   while (x != s)
       edges[p[x]].flow += a;
       edges[p[x] ^ 1].flow -= a;
       x = edges[p[x]].from;
    }
   return a;
}
int Maxflow(int s, int t)
   this->s = s;
   this->t = t;
   int flow = 0;
   BFS();
   memset(num, 0, sizeof(num));
    for (int i = 0; i < n; i++)
       num[d[i]]++;
    int x = s;
   memset(cur, 0, sizeof(cur));
   while (d[s] < n)
       if (x == t)
           flow += Augment();
           x = s;
       int ok = 0;
       for (int i = cur[x]; i < G[x].size(); i++)
           Edge &e = edges[G[x][i]];
           if (e.cap > e.flow && d[x] == d[e.to] + 1)
               ok = 1;
               p[e.to] = G[x][i];
               cur[x] = i;
               x = e.to;
               break;
            }
       if (!ok)
```

```
int m = n - 1;
               for (int i = 0; i < G[x].size(); i++)
                   Edge &e = edges[G[x][i]];
                   if (e.cap > e.flow)
                       m = min(m, d[e.to]);
                if (--num[d[x]] == 0)
                   break;
               num[d[x] = m + 1]++;
               cur[x] = 0;
               if (x != s)
                   x = edges[p[x]].from;
           }
        }
       return flow;
   }
};
```

#### 8.2 最小割

```
//最大流最小割定理
#include <algorithm>
#include <cstdio>
#include <cstring>
#include <queue>
const int N = 1e4 + 5, M = 2e5 + 5;
int n, m, s, t, tot = 1, lnk[N], ter[M], nxt[M], val[M], dep[N], cur[N];
void add(int u, int v, int w) {
 ter[++tot] = v, nxt[tot] = lnk[u], lnk[u] = tot, val[tot] = w;
void addedge(int u, int v, int w) { add(u, v, w), add(v, u, 0); }
int bfs(int s, int t) {
 memset(dep, 0, sizeof(dep));
 memcpy(cur, lnk, sizeof(lnk));
 std::queue<int> q;
 q.push(s), dep[s] = 1;
 while (!q.empty()) {
   int u = q.front();
   q.pop();
   for (int i = lnk[u]; i; i = nxt[i]) {
     int v = ter[i];
     if (val[i] \&\& !dep[v]) q.push(v), dep[v] = dep[u] + 1;
   }
 return dep[t];
int dfs(int u, int t, int flow) {
 if (u == t) return flow;
```

```
int ans = 0;
 for (int &i = cur[u]; i && ans < flow; i = nxt[i]) {</pre>
   int v = ter[i];
   if (val[i] \&\& dep[v] == dep[u] + 1) {
      int x = dfs(v, t, std::min(val[i], flow - ans));
     if (x) \ val[i] -= x, \ val[i ^ 1] += x, \ ans += x;
 if (ans < flow) dep[u] = -1;
 return ans;
int dinic(int s, int t) {
 int ans = 0;
 while (bfs(s, t)) {
   int x;
   while ((x = dfs(s, t, 1 << 30))) ans += x;
 return ans;
int main() {
 scanf("%d%d%d%d", &n, &m, &s, &t);
 while (m--) {
   int u, v, w;
   scanf("%d%d%d", &u, &v, &w);
   addedge(u, v, w);
 printf("%d\n", dinic(s, t));
 return 0;
```

### 8.3 费用流

#### 8.3.1 基于 EK 算法

```
struct qxx
{
    int nex, t, v, c;
};

qxx e[M];
int h[N], cnt = 1;

void add_path(int f, int t, int v, int c)
{
    e[++cnt] = (qxx){h[f], t, v, c}, h[f] = cnt;
}

void add_flow(int f, int t, int v, int c)
{
    add_path(f, t, v, c);
    add_path(t, f, 0, -c);
```

```
}
int dis[N], pre[N], incf[N];
bool vis[N];
bool spfa()
   memset(dis, 0x3f, sizeof(dis));
   queue<int> q;
   q.push(s), dis[s] = 0, incf[s] = INF, incf[t] = 0;
   while (q.size())
       int u = q.front();
       q.pop();
       vis[u] = 0;
       for (int i = h[u]; i; i = e[i].nex)
            const int &v = e[i].t, &w = e[i].v, &c = e[i].c;
           if (!w \mid | dis[v] \leftarrow dis[u] + c)
               continue;
           dis[v] = dis[u] + c, incf[v] = min(w, incf[u]), pre[v] = i;
           if (!vis[v])
               q.push(v), vis[v] = 1;
       }
   }
   return incf[t];
int maxflow, mincost;
void update()
   maxflow += incf[t];
   for (int u = t; u != s; u = e[pre[u] ^ 1].t)
       e[pre[u]].v -= incf[t], e[pre[u] ^ 1].v += incf[t];
       mincost += incf[t] * e[pre[u]].c;
// 调用: while(spfa())update();
```

#### 8.3.2 基于 Dinic 算法

```
#include <algorithm>
#include <cstdio>
#include <cstring>
#include <queue>

const int N = 5e3 + 5, M = 1e5 + 5;
const int INF = 0x3f3f3f3f;
int n, m, tot = 1, lnk[N], cur[N], ter[M], nxt[M], cap[M], cost[M], dis[N], ret;
bool vis[N];
```

```
void add(int u, int v, int w, int c)
     ter[++tot] = v, nxt[tot] = lnk[u], lnk[u] = tot, cap[tot] = w,
cost[tot] = c;
void addedge(int u, int v, int w, int c) { add(u, v, w, c), add(v, u, 0,
-c); }
bool spfa(int s, int t)
   memset(dis, 0x3f, sizeof(dis));
   memcpy(cur, lnk, sizeof(lnk));
   std::queue<int> q;
   q.push(s), dis[s] = 0, vis[s] = 1;
   while (!q.empty())
   {
       int u = q.front();
       q.pop(), vis[u] = 0;
       for (int i = lnk[u]; i; i = nxt[i])
           int v = ter[i];
           if (cap[i] && dis[v] > dis[u] + cost[i])
               dis[v] = dis[u] + cost[i];
               if (!vis[v])
                   q.push(v), vis[v] = 1;
           }
       }
   }
   return dis[t] != INF;
int dfs(int u, int t, int flow)
   if (u == t)
       return flow;
   vis[u] = 1;
   int ans = 0;
   for (int &i = cur[u]; i && ans < flow; i = nxt[i])
       int v = ter[i];
       if (!vis[v] && cap[i] && dis[v] == dis[u] + cost[i])
           int x = dfs(v, t, std::min(cap[i], flow - ans));
           if (x)
               ret += x * cost[i], cap[i] -= x, cap[i ^ 1] += x, ans +=
x;
       }
   vis[u] = 0;
   return ans;
int mcmf(int s, int t)
```

```
int ans = 0;
   while (spfa(s, t))
       int x;
       while ((x = dfs(s, t, INF)))
           ans += x;
   return ans;
int main()
   int s, t;
   scanf("%d%d%d%d", &n, &m, &s, &t);
   while (m--)
       int u, v, w, c;
       scanf("%d%d%d%d", &u, &v, &w, &c);
       addedge(u, v, w, c);
   int ans = mcmf(s, t);
   printf("%d %d\n", ans, ret);
   return 0;
```

# 9 STL 标准模板库 Standard Template Library

## 9.1 Priority queue

```
top() 访问队首元素
push() 入队
pop() 堆顶(队首)元素出队
size() 队列元素个数
empty()是否为空
注意没有 clear()! 不提供该方法
优先队列只能通过 top()访问队首元素(优先级最高的元素)
```

### 9.2 Deque

```
push_back(x)/push_front(x) 把x压入后/前端 back()/front() 访问(不删除)后/前端元素 pop_back() pop_front() 删除后/前端元素 erase(iterator it)删除双端队列中的某一个元素 erase(iterator first,iterator last) 删除双端队列中[first,last) 中的元素 empty()判断 deque 是否空 size() 返回 deque 的元素数量 clear()清空 deque
```

# 10 其他 Other

## 10.1 竞赛模板

```
#include <bits/stdc++.h>
using namespace std;
using ll = long;
const int INF = 0x3f3f3f3f;
const int MOD = 1e9 + 7;
const int MX = 1e5 + 10;
const vector<pair<int, int>> DIRS = {{1, 0}, {-1, 0}, {0, 1}, {0, -1}};
template <typename T>
istream& operator>>(istream& is, vector<T>& v) {
   for (auto& x : v) cin >> x;
   return is;
template <typename T>
ostream &operator<<(ostream& os, vector<T>& v) {
   for (auto& x : v) cout << x << ' ';
   return os;
void solution() {
signed main() {
   cin.tie(nullptr);
   ios::sync_with_stdio(false);
   int t{1};
   cin >> t;
   while (t--) solution();
   return 0;
```