

Quadratic Programming Problems.

Ref: Numerical optimization with Applications by S. Chandra, Jayadeva, and A. Mehra.

In the earlier ~~class~~ lectures we studied convex optimization problems, the Lagrang optimization technique when we solve the optimization problems with equality type constraints. In this lecture we will study "Quadratic Optimization Problems".

Quadratic Programming problems:

A quadratic programming problem (QPP) is the special class of non-linear optimization problems in which the objective function is quadratic and all the constraints are linear.

The general mathematical formulation of a QPP is as follows:

$$\text{Min } f(x) = x^T Q x + c^T x$$

$$\text{s.t. } Ax \leq b,$$

$$x \geq 0.$$

where $Q = [q_{ij}]_{n \times n}$ symmetric positive definite matrix,

- $c \in \mathbb{R}^n$
- $x \in \mathbb{R}^n$
- $b \in \mathbb{R}^m$
- $A = [a_{ij}]_{m \times n}$

Example 1:- Min $f(x) = 3x_1^2 + 4x_2^2 + 2x_1x_2 - 2x_1 - 3x_2$,

$$\text{s.t. } 3x_1 + 2x_2 \leq 6$$

$$x_1 + x_2 \leq 2$$

$$x_1, x_2 \geq 0.$$

$$\text{OR } \min f(x) = \overbrace{(x_1 \ x_2)}^{x^T} \overbrace{\begin{pmatrix} 3 & 1 \\ 1 & 4 \end{pmatrix}}^Q \overbrace{\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}}^x + \overbrace{(-2 \ -3)}^{c^T} \overbrace{\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}}^x$$
$$\text{s.t. } \overbrace{\begin{pmatrix} 3 & 2 \\ 1 & 1 \end{pmatrix}}^A \overbrace{\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}}^x \leq \overbrace{\begin{pmatrix} 6 \\ 2 \end{pmatrix}}^b.$$

$$\overbrace{\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}}^x \geq 0.$$

$$Q = \begin{pmatrix} 3 & 1 \\ 1 & 4 \end{pmatrix}$$

$$D_1 = 3 > 0, \quad D_2 = 12 - 1 = 11 > 0$$

\therefore The matrix Q is positive definite.

Hence, it is positive semi-definite also.

Therefore, the above example is a Quadratic programming problem.

Solution of QPP by graphical Method:-

Example 2:- Min $(x_1 - 2)^2 + (x_2 - 1)^2$

$$\text{s.t. } x_1 + x_2 \leq 2$$

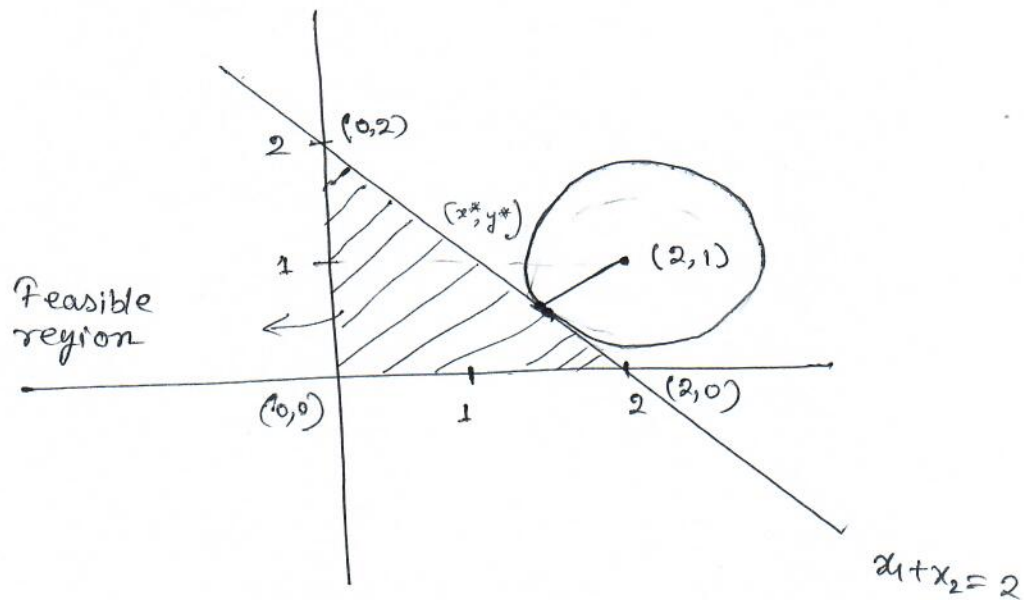
$$x_1, x_2 \geq 0$$

Some observations are :-

(2)

- (i) Feasible region of (QPP) is always a convex set,
(since all the constraints are linear)
- (ii) Unlike (LPP), the optimal solution of QPP/NLPP may ~~be~~ not be attained at the vertices of the feasible region.

Solution of Example 2:



$$x^* + y^* = 2 \quad \text{--- (i)}$$

$$\text{slope } \frac{y^* - 1}{x^* - 2} = 1$$

$$\Rightarrow x^* - y^* = 1 \quad \text{--- (ii)}$$

Solve equations (i) & (ii), $x^* = 3/2$, $y^* = 1/2$

Exercise - check if the following problem is a QPP or NOT?

$$\min f = 3x_1^2 + 2x_2^2 + x_1x_2 - 4x_1 - 2x_2$$

$$\text{s.t.} \quad x_1 + 2x_2 \leq 6$$

$$x_1, x_2 \geq 0.$$

□ KKT conditions for QPP :-

(3)

The problem (QPP) can be re-written as

$$\begin{aligned} \text{Min } f(x) &= x^T Q x + c^T x \\ \text{s.t. } Ax &\leq b, \quad \text{--- (i)} \\ -x &\leq 0, \quad \text{--- (ii)} \end{aligned} \quad \left. \vphantom{\begin{aligned} \text{Min } f(x) &= x^T Q x + c^T x \\ \text{s.t. } Ax &\leq b, \quad \text{--- (i)} \\ -x &\leq 0, \quad \text{--- (ii)} \end{aligned}} \right\} \text{--- (1)}$$

Let the KKT multiplier associated with constraints (i) and (ii) be $u \in \mathbb{R}^m$ and $v \in \mathbb{R}^n$, respectively. Then the KKT conditions for the problem (1) are as follows :

$$c^T + 2x^T Q + u^T A - v^T = 0$$

$$u^T (Ax - b) - v^T x = 0,$$

$$Ax - b \leq 0,$$

$$x \geq 0, u \geq 0, v \geq 0.$$

From (i)

□ $Ax - b \leq 0$. we need $u \in \mathbb{R}^m$, i.e. m number of variables.

From (ii) $-x \leq 0$ we need $v \in \mathbb{R}^n$, i.e., n number of variables.

Therefore, the total number of KKT multipliers are $(m+n)$.

(i) Optimality Condition :-

$$\nabla f(x) + \sum_{i=1}^m \lambda_i \nabla g_i(x) = 0$$

Prob.
For (1)

$$2x^T Q + c^T + u^T A + v^T (-I) = 0 \quad \text{--- (a)}$$

(ii) Feasibility Condⁿ :-

$$g_i(x) \leq 0$$

For problem ① $Ax - b \leq 0$, ——— ⑤

(iii) Complementary slackness Condⁿ :-

$$\lambda_i^* g_i(x^*) = 0$$

For prob ① $u^T(Ax - b) - v^T x = 0$ ——— ⑥

(iv) Non-Negative Condⁿ :-

$$\lambda_i^* \geq 0$$

for problem ① $x \geq 0, u \geq 0, v \geq 0$, ——— ⑦

▣ Rewrite the above conditions :-

For ① $2x^T Q = 2Q^T x = 2Qx$ ($\because Q$ is symmetric)

\therefore taking transpose of eq^s ①, ②, ③, ④

$$2Qx + c + A^T u - v = 0$$

$$u^T(Ax - b) - v^T x = 0$$

$$Ax - b + s = 0 \quad (\text{slack variable added})$$

$$x \geq 0, u \geq 0, v \geq 0 \quad (\text{all variables are non-negative})$$

Thm: Let Q be a symmetric, ^{positive} semi-definite matrix of order n . Then for any $x, y \in \mathbb{R}^n$

$$2x^T Q y \leq (x^T Q x + y^T Q y).$$

Proof: $z^T Q z \geq 0 \quad \forall z \in \mathbb{R}^n$ ($\because Q$ is +ve semi-definite)
By definition.

$$\Rightarrow (x-y)^T Q (x-y) \geq 0, \quad \forall x, y \in \mathbb{R}^n \quad (\because \text{it's true for all } z \in \mathbb{R}^n)$$

$$\Rightarrow x^T Q x - x^T Q y - y^T Q x + y^T Q y \geq 0.$$

$$\Rightarrow x^T Q y + y^T Q x \leq x^T Q x + y^T Q y$$

$$\Rightarrow \boxed{2x^T Q y \leq (x^T Q x + y^T Q y)}$$

$$x_{1 \times n}^T Q_{n \times n} y_{n \times 1} = [d]_{1 \times 1} \Rightarrow \text{scalar.}$$

$$\begin{aligned} x^T Q y &= (x^T Q y)^T \\ &[\because (\text{scalar})^T = \text{scalar}] \\ \downarrow \\ x^T Q y &= y^T Q^T x \\ &= y^T Q x \quad [\because Q^T = Q] \end{aligned}$$