Ref: Numerical optimization with applications by S. Chandra, Jayadera, &

In the previous Lectures we have seen some conven programming problems. We also discussed un the Constraints optimization with using with equality Constraints using Lagrange's optimization technique. Now, in this lecture we will see KKT conditions. Remember when we have equality type constraints then we use Layrange's Method to find out the optimal solution of that type of problems, But, suppose We have a constraint optimization problem with inequality type constraints, then karaush-kuhn-Tucker proposed some condition to find optimal solutions of the problems. We ealled that condition as KKT-condition in short. Let's see the following 600 form of the KKT- Problem:

A Karaush-Kuhn-Tueker (KKT) Conditions:

Consider the optimization problem as

Min :
$$f(x)$$

S.1. $g_1(x) \neq 0$, $i = 1, 2, ..., m$,

where f: IR" > IR, gi: 1R" > IR.

Also, assume that I and of are continuously differentiable functions.

Necessary Condition:

Let x* be en local min point of the problem (*) at which basic constraint qualification holds.

Then there exist multipliers (called KKT-multi-pliers) . It, i21, 2, ..., m such that the following conditions hold:

- (i) $\nabla f(x^*) + \frac{m}{2} \stackrel{\star}{\lambda}_i \nabla g_i(x^*) = 0$ (optimality Cond?)
- (ii) g: (x*) = 0, i=1,2,..., m, (Feasibility Cond+)
- (ii) $\lambda^{i}_{i}(x^{*}) \pm 0$, i_{2}, \ldots, m , (Complementary slackness)
- (iv) i > 0 + i (Non-negativity)

These (above four) conditions of cire called KKT-Conditions.

Note: Here, 1ti's are called KKT- multipliers, x* called KKT - point.

Let (x*, 1, 12, ..., 1m) satisfy the KKT- conditions (i) to (iv). Let I amd gifor i= 1,2,-, m be differentiable conven functions. Then * * is a global min point of the problem .

Note that f and di (+i) are diff. convex for that means the optimization problem @ becomes Conver optimization problem.

Proof: Flut us suppose S= fx ein: g; (n) = 0, i=1,2,..., mg Now, find and gend is diff. convent for.

then $f(n) - f(n^*) > (n-x^*)^T \nabla f(n) + x \in S - 0$ \$ 91°(n) - 9;(n*) > (x-x*) T \q;(n) + x \ \ S.

Now, from necessary condition (in it >0 (non-negative) Therefor we can multiply 1; in the above inequally and adding them we get

Z 1, g; (n) - Z 1, g; (n) > (n-xx) T Z 1, vg; (n) - 2)

Adding equations 10 and 10, we get

dding equations () ma $f(n) - f(n^*) + \frac{m}{2} \lambda^* g(n) - \frac{m}{2} \lambda^* g(n^*) \geq (n - n^*)^T (\nabla fen) + \frac{m}{2} \lambda^* g(n)$ $f(n) - f(n^*) + \frac{m}{2} \lambda^* g(n) - \frac{m}{2} \lambda^* g(n^*) \geq (n - n^*)^T (\nabla fen)$ $f(n) - f(n^*) + \frac{m}{2} \lambda^* g(n) - \frac{m}{2} \lambda^* g(n^*) \geq (n - n^*)^T (\nabla fen)$ $f(n) - f(n^*) + \frac{m}{2} \lambda^* g(n) - \frac{m}{2} \lambda^* g(n^*) \geq (n - n^*)^T (\nabla fen)$

Therefore, $f(n) - f(n^*) > - \overline{Z} \lambda_i^* g(n) > 0$

=) x* is the global minimum point.

Example 1: Min $f(x_1, x_2) = 2x_1 + x_2$ S.t. $x_1^2 + x_2^2 \leq 4$ $x_1 - x_2 \leq 0$

$$=$$
) Min $f(x)$ ≤ 0

Where $2 = (n_1, n_2)$ $f(n_1, x_2) = 2n_1 + n_2$ $g_1(x_1, n_2) = n_1^2 + n_2^2 - 4 \le 0$ $g_2(n_1, n_2) = n_1^2 + n_2 \le 0$

Step-I

Wits first cheek this problem is a CPP or NOT?

If conven, then KKT-conditions will be sufficient.

Step. 2: Calculate Hessian Matrin for the constraints (which for y_1 , $\nabla_{y_1}^2 = \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix}$, $\lambda_1 = 2 > 6$ dinear)

i. $\nabla^2 g_1$ is positive-definite, hence convex.

step-3: Try to write all the KKT- conditions

(i)
$$\nabla f(\vec{x}) + \sum_{i=1}^{2} \lambda_{i}^{*} g_{i}(x^{*}) = 0$$

=>
$$(2 1) + \frac{1}{2} (2x4 2x2) + \frac{1}{2} (1-1) = (0,0)$$

(ii)
$$g_1(x^*) \leq 6 \quad \forall i \quad , g_1(x^*) = x_1^2 + x_2^2 - 4 \leq 6$$

 $g_2(x^*) = x_1 - x_2 \leq 6$

(iii)
$$\lambda_{1}^{*} g_{1}(x^{*}) = 0 = \lambda_{1}^{*} (x_{1}^{2} + x_{2}^{2} - 4) = 0$$

$$\lambda_{2}^{*} g_{2}(x^{*}) = 0 = \lambda_{1}^{*} (x_{1} - x_{2}) = 0$$

$$\lambda_1^{\times}, \lambda_2^{\times} \geq 0$$

Step-4: Try to find x*

Case-I: $\lambda_1^* = \lambda_2^* = 0$

ease 11)- 1,20, x=x, 12 >0.

Cone-11: $\chi_1^2 + \chi_2^2 - 4 = 0$, $\lambda_1^* > 0$, $\lambda_1^* = 0$.

Cone IF; x1-12=0, x12+12-4=0, 1120 1210.

Case-I:

Case- II ;

Cone-III +
$$\chi_1 \lambda_1^* = -1$$

 $\chi_2 \lambda_2^* = -1/2$
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theek it satisfy all the KKT worditions or NOT?

If yes, then (x1, x2) is the global min for the CPP.

4

E. Similars to the presion optimization (min) problem, we can extend the logic the sufficient condition for the maximization problem on

if x* is the local max for the corresponding

max f(x) $s.t. g_i(x) \leq 0.$

Then that local maxima would be the global maxima if we the objective function is fern) is concave and the associted constraints, i.e. gick) (feasible space) is convert. In the other way we can say that KKT conditions are the sufficient word?, if for is concave and feasible space is convert for the maximization problem.

Exercise: Find the solution of the following optimization problem. Example: min $f(\underline{x}) = xu^2 + x_2^2 - 2xu$ S.t. $x_1^2 + x_2 - 1 = 0$