

## Convex Programming Problems

Ref:- Numerical optimization with Applications by S. Chandra, Jayadara and A. Mehra.

In the previous lecture we have seen that what is convex function? what are their properties. In this lecture, we will study some convex programming problems (CPP).

### Convex Programming Problem :-

The optimization problem of the form:

$$\begin{aligned} & \text{Min } f(x) \\ \text{s.t. } & g_i(x) \leq 0, \quad i=1, 2, \dots, n, \end{aligned}$$

} ——— (\*)

is called a convex programming problem (CPP) if  $f$  and  $g_i$  ( $i=1, 2, \dots, n$ ) are convex functions.

Remark 1:- If the optimization problem of the form:

$$\begin{aligned} & \text{Max } f(x) \\ \text{s.t. } & g_i(x) \leq 0, \quad i=1, 2, \dots, n, \end{aligned}$$

is called a concave programming problem if  $f$  is a concave function and (for  $i=1, \dots, n$ )  $g_i$ 's are convex.

## Different formats of the CPP problem:

Optimization problem	Conditions for (CPP)
(i) Min $f(x)$ s.t. $g_i(x) \leq 0, i=1, 2, \dots, n$	$f$ and $g_i (i=1, 2, \dots, n)$ both are convex.
(ii) Max $f(x)$ s.t. $g_i(x) \leq 0, i=1, 2, \dots, n$	$f$ is concave and $g_i$ 's $(i=1, 2, \dots, n)$ are convex.
(iii) Min $f(x)$ s.t. $g_i(x) \geq 0, (i=1, 2, \dots, n)$	$f$ is convex and $g_i$ 's $(i=1, 2, \dots, n)$ are concave.
(iv) Max $f(x)$ s.t. $g_i(x) \geq 0, (i=1, 2, \dots, n)$	$f$ and $g_i (i=1, 2, \dots, n)$ both are concave.

Theorem 1: Let  $\bar{x}$  be a local minimum point of the convex optimization problem  $(*)$ . Then  $\bar{x}$  is also its global min point.

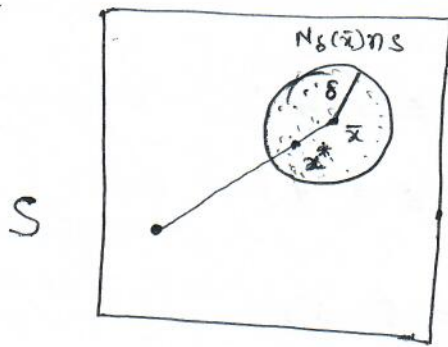
Proof: We are given that  $\bar{x}$  is a local min point of  $(*)$  (CPP).

So by definition, there exists  $\delta > 0$ , s.t.  $f(\bar{x}) \leq f(y)$ ,  
 $\forall y \in N_\delta(\bar{x}) \cap S$

$S \rightarrow$  (domain of the convex function  $f$  which is convex set).  
 $N_\delta(\bar{x})$  is a ball of radius  $\delta$  centered at  $\bar{x}$ .

Let  $x$  be an arbitrary point outside the set  $N_\delta(\bar{x}) \cap S$ .  
 Then if we show that  $f(\bar{x}) \leq f(x)$ , we are done!

Now, there certainly exists  $\lambda$ ;  $0 < \lambda < 1 \Rightarrow$   
 $x^* = \lambda x + (1-\lambda)\bar{x}$  and  $x^* \in N_\delta(\bar{x}) \cap S$ .



As  $x^* \in N_\delta(\bar{x}) \cap S$ , and  $\bar{x}$  is a local min point of  $\text{CPP}(S)$   
 we have  $f(\bar{x}) \leq f(x^*)$ .

$$\Rightarrow f(\bar{x}) \leq f(\lambda x + (1-\lambda)\bar{x})$$

But  $f$  is convex on  $S$ , and hence the above gives

$$f(\bar{x}) \leq \lambda f(x) + (1-\lambda)f(\bar{x}) = \lambda f(x) + f(\bar{x}) - \lambda f(\bar{x}).$$

$$\therefore \lambda f(\bar{x}) \leq \lambda f(x),$$

$$\Rightarrow f(\bar{x}) \leq f(x) \text{ as } 0 < \lambda < 1.$$

which implies that  $\bar{x}$  is a global min point.

Theorem 2 :- Let  $g_i$  for each  $i=1, 2, \dots, n$  be a convex function. Then  $W = \{x \in \mathbb{R}^n : g_i(x) \leq 0, i=1, 2, \dots, n\}$  is a convex set.

Proof :-  $W = \{x \in \mathbb{R}^n : g_i(x) \leq 0, i=1, 2, \dots, n\}$   
where all  $g_i$ 's are convex functions.

$$\text{Let } x_1, x_2 \in W \Rightarrow g_i(x_1) \leq 0 \quad \forall i \\ \text{ \& } g_i(x_2) \leq 0 \quad \forall i$$

Suppose  
~~let~~  $x = \lambda x_1 + (1-\lambda)x_2, \lambda \in [0, 1]$

$$g_i(x) = g_i(\lambda x_1 + (1-\lambda)x_2)$$

$$\leq \lambda g_i(x_1) + (1-\lambda)g_i(x_2) \quad (\because g_i \text{'s are convex fns}) \\ \leq \lambda \times 0 + (1-\lambda) \times 0 = 0$$

$$\Rightarrow g_i(x) \leq 0 \quad \forall i$$

$$\Rightarrow x \in W$$

$\therefore W$  is a convex set.



Theorem 3! The set of all optimal solutions of the convex optimization problem (\*) is a convex set.

Proof:-  $V = \{x^* : x^* \in S, f(x^*) \leq f(x) \forall x \in S\}$

Exercise!!

Theorem 4:- Let a non-constant convex function  $f$  be maximized over a convex set  $S$ . Then no interior point of  $S$  can be a maximizing point.  
i.e. if a maximizing point exists, then it must be a boundary point of  $S$ .

Theorem 5:- Let  $S \subseteq \mathbb{R}^n$  be convex and  $f : S \rightarrow \mathbb{R}$  be a strictly convex function. Then there is a unique minimizing point of  $f$  over  $S$ .

Proof:- Let  $\bar{x}$  and  $x^*$  ~~be~~ be two distinct minimizing points of  $f$  over  $S$ , i.e.  $\bar{x}, x^* \in S$ .

and  $f(\bar{x}) \leq f(x) \quad \forall x \in S$

also  $f(x^*) \leq f(x) \quad \forall x \in S$ .

$f(\bar{x}) = f(x^*)$

Let  $\hat{x} = \lambda \bar{x} + (1-\lambda)x^*$  for  $\lambda \in [0, 1]$ .

In particular,  
 $\hat{x} = \frac{\bar{x} + x^*}{2}$   
 $f(\hat{x}) = f\left(\frac{\bar{x} + x^*}{2}\right)$   
 $\leq \frac{1}{2}f(\bar{x}) + \frac{1}{2}f(x^*)$   
 $\leq f(\bar{x})$

Then by strict convexity of  $f$  we have

$$f(\hat{x}) = f(\lambda \bar{x} + (1-\lambda)x^*) < \lambda f(\bar{x}) + (1-\lambda)f(x^*)$$

$$\Rightarrow f(\hat{x}) < f(\bar{x}) \quad \Rightarrow \Leftarrow \quad \text{that } \bar{x} \text{ is minimizing point.}$$

[ $\because f(\bar{x}) = f(x^*)$ ]

## Examples of Convex Programming Problems:

Example 1: Check if the following problem is a convex prog. prob. (cpp).

(i) Min  $x_1 + x_2$

s.t.  $x_1^2 + x_2^2 \leq 1$   $\equiv$

$x_1^2 \leq x_2$

Min  $f = x_1 + x_2$

s.t.  $g_1(x_1, x_2) = x_1^2 + x_2^2 - 1 \leq 0$

$g_2(x_1, x_2) = x_1^2 - x_2 \leq 0$

(ii) Max  $4x_1 + 3x_2$

s.t.  $x_1 + x_2 \leq 4$

$x_1 x_2 \leq 1$

$x_1, x_2 \geq 0$ .

$\Rightarrow$   
Exercise: Max .

Exercise 2:- check if the following is c.p.p or NOT.

$$\begin{aligned} \text{(i)} \quad & \text{Max } x_2^2 \\ \text{s.t.} \quad & x_1^2 + x_2^2 \leq 4 \\ & 4x_1^2 \geq x_2, \end{aligned}$$

$$\begin{aligned} \text{(ii)} \quad & \text{Max } 4x_1 + 3x_2 \\ \text{s.t.} \quad & x_1^2 + x_2^2 \leq 1 \\ & (x_1 - 1)^2 + x_2^2 \leq 1. \end{aligned}$$

$$\begin{aligned} \text{(iii)} \quad & \text{Min } x_1^2 + x_2 \\ \text{s.t.} \quad & |x_1| + |x_2| \leq 2 \\ & x_1^2 - x_2 \geq 0. \end{aligned}$$

