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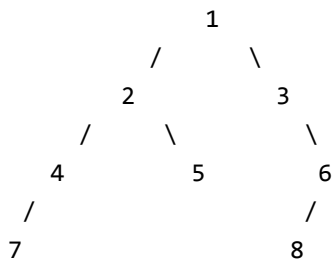
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Remove nodes on root to leaf paths of length < K

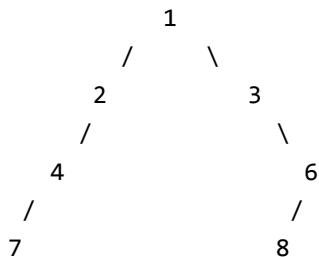
Given a Binary Tree and a number k, remove all nodes that lie only on root to leaf path(s) of length smaller than k. If a node X lies on multiple root-to-leaf paths and if any of the paths has path length $\geq k$, then X is not deleted from Binary Tree. In other words a node is deleted if all paths going through it have lengths smaller than k.

Consider the following example Binary Tree



Input: Root of above Binary Tree
k = 4

Output: The tree should be changed to following



There are 3 paths

- i) 1->2->4->7 path length = 4
- ii) 1->2->5 path length = 3
- iii) 1->3->6->8 path length = 4

There is only one path " 1->2->5 " of length smaller than 4.

The node 5 is the only node that lies only on this path, so node 5 is removed.

Nodes 2 and 1 are not removed as they are parts of other paths of length 4 as well.

If k is 5 or greater than 5, then whole tree is deleted.

If k is 3 or less than 3, then nothing is deleted.

We strongly recommend to minimize your browser and try this yourself first

The idea here is to use post order traversal of the tree. Before removing a node we need to check that all the children of that node in the shorter path are already removed.

There are 2 cases:

- i) This node becomes a leaf node in which case it needs to be deleted.
- ii) This node has other child on a path with path length $\geq k$. In that case it needs not to be deleted.

The implementation of above approach is as below :

C/C++

```
// C++ program to remove nodes on root to leaf paths of length < K
#include<iostream>
using namespace std;

struct Node
{
    int data;
    Node *left, *right;
};

//New node of a tree
Node *newNode(int data)
{
    Node *node = new Node;
    node->data = data;
    node->left = node->right = NULL;
    return node;
}

// Utility method that actually removes the nodes which are not
// on the pathLen >= k. This method can change the root as well.
Node *removeShortPathNodesUtil(Node *root, int level, int k)
{
    //Base condition
    if (root == NULL)
        return NULL;

    // Traverse the tree in postorder fashion so that if a leaf
    // node path length is shorter than k, then that node and
    // all of its descendants till the node which are not
    // on some other path are removed.
    root->left = removeShortPathNodesUtil(root->left, level + 1, k);
    root->right = removeShortPathNodesUtil(root->right, level + 1, k);

    // If root is a leaf node and it's level is less than k then
    // remove this node.
    // This goes up and check for the ancestor nodes also for the
    // same condition till it finds a node which is a part of other
    // path(s) too.
    if (root->left == NULL && root->right == NULL && level < k)
    {
        delete root;
    }
}
```

```

        return NULL;
    }

    // Return root;
    return root;
}

// Method which calls the utility method to remove the short path
// nodes.
Node *removeShortPathNodes(Node *root, int k)
{
    int pathLen = 0;
    return removeShortPathNodesUtil(root, 1, k);
}

//Method to print the tree in inorder fashion.
void printInorder(Node *root)
{
    if (root)
    {
        printInorder(root->left);
        cout << root->data << " ";
        printInorder(root->right);
    }
}

// Driver method.
int main()
{
    int k = 4;
    Node *root = newNode(1);
    root->left = newNode(2);
    root->right = newNode(3);
    root->left->left = newNode(4);
    root->left->right = newNode(5);
    root->left->left->left = newNode(7);
    root->right->right = newNode(6);
    root->right->right->left = newNode(8);
    cout << "Inorder Traversal of Original tree" << endl;
    printInorder(root);
    cout << endl;
    cout << "Inorder Traversal of Modified tree" << endl;
    Node *res = removeShortPathNodes(root, k);
    printInorder(res);
    return 0;
}

```

Run on IDE

Java

```

// Java program to remove nodes on root to leaf paths of length < k

/* Class containing left and right child of current
node and key value*/
class Node {

    int data;
    Node left, right;

    public Node(int item) {
        data = item;
    }
}

```

```

        left = right = null;
    }
}

class BinaryTree {

    static Node root;

    // Utility method that actually removes the nodes which are not
    // on the pathLen >= k. This method can change the root as well.
    Node removeShortPathNodesUtil(Node node, int level, int k) {
        //Base condition
        if (node == null) {
            return null;
        }
        // Traverse the tree in postorder fashion so that if a leaf
        // node path length is shorter than k, then that node and
        // all of its descendants till the node which are not
        // on some other path are removed.
        node.left = removeShortPathNodesUtil(node.left, level + 1, k);
        node.right = removeShortPathNodesUtil(node.right, level + 1, k);

        // If root is a leaf node and it's level is less than k then
        // remove this node.
        // This goes up and check for the ancestor nodes also for the
        // same condition till it finds a node which is a part of other
        // path(s) too.
        if (node.left == null && node.right == null && level < k) {
            return null;
        }

        // Return root;
        return node;
    }

    // Method which calls the utility method to remove the short path
    // nodes.
    Node removeShortPathNodes(Node node, int k) {
        int pathLen = 0;
        return removeShortPathNodesUtil(node, 1, k);
    }

    //Method to print the tree in inorder fashion.
    void printInorder(Node node) {
        if (node != null) {
            printInorder(node.left);
            System.out.print(node.data + " ");
            printInorder(node.right);
        }
    }

    // driver program to test for samples
    public static void main(String args[]) {
        BinaryTree tree = new BinaryTree();
        int k = 4;
        tree.root = new Node(1);
        tree.root.left = new Node(2);
        tree.root.right = new Node(3);
        tree.root.left.left = new Node(4);
        tree.root.left.right = new Node(5);
        tree.root.left.left.left = new Node(7);
        tree.root.right.right = new Node(6);
        tree.root.right.right.left = new Node(8);
        System.out.println("The inorder traversal of original tree is : ");
        tree.printInorder(root);
    }
}

```

```
Node res = tree.removeShortPathNodes(root, k);
System.out.println("");
System.out.println("The inorder traversal of modified tree is : ");
tree.printInorder(res);
}
}

// This code has been contributed by Mayank Jaiswal
```

[Run on IDE](#)

Output:

```
Inorder Traversal of Original tree
7 4 2 5 1 3 8 6
Inorder Traversal of Modified tree
7 4 2 1 3 8 6
```

Time complexity of the above solution is $O(n)$ where n is number of nodes in given Binary Tree.

This article is contributed by **Kumar Gautam**. Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above



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3.6

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AllergicToBitches • 2 months ago

Slight optimization, if level >= k, we can skip level +1 and above -

```
Node *removeShortPathNodesUtil(Node *root, int level, int k)
{
    if (root == NULL)
        return NULL;

    if(level >= k)
        return root;

    //rest of the code
}
```

^ | ▾ • Reply • Share ▸



Balu • 4 months ago

```
removeReachableLeafs(Node root, int k){
    if(root == null) return null;
    if(k>1&& isLeaf(root) return null;
    root.left = removeReachableLeafs(root.left, k-1);
    root.right = removeReachableLeafs(root.right, k-1);
    return root;
}
```

}

Please review

1 ^ | v • Reply • Share ›

**Ashish** • 5 months ago

The solution for this question will be similar to the solution for <http://www.geeksforgeeks.org/r...>, the only change being that we need to increase sum by 1 at each node as compared to increasing the sum by the node's data in that question.

^ | v • Reply • Share ›

**InnocentHacker** • 6 months ago

sorry for typo errors

```
public static BinaryTreeNode removeShortPathNodes(BinaryTreeNode root, int k){
    return removeShortPathNodesHelper(root,k,1);
}
```

```
public static BinaryTreeNode removeShortPathNodesHelper(BinaryTreeNode root, int
k,int level){
```

```
    if(root==null){
        return null;
    }
```

```
    if(level>=k){
        return root;
    }
```

```
    root.setLeft(removeShortPathNodesHelper(root.getLeft(),k,level+1));
    root.setRight(removeShortPathNodesHelper(root.getRight(),k,level+1));
    if(root.getLeft()!=null||root.getRight()!=null){
        return root;
    }
    return null;
}
```

^ | v • Reply • Share ›

**InnocentHacker** • 6 months ago

```
public class BinaryTreeUtil{
```

```
    public static BinaryTreeNode removeShortPathNodes(BinaryTreeNode root, int k){
        return removeShortPathNodesHelper(root,k,1);
    }
```

```
    public static BinaryTreeNode removeShortPathNodesHelper(BinaryTreeNode root, int
k,int level){
        if(root==null){
            return null;
```

```

}
if(level>=k){
return root;
}
root.setLeft(removeShortPathNodes(root.getLeft(),k,level+1)){
root.setRight(removeShortPathNodes(root.getRight(),k,level+1)){
if(root.getLeft()!=null||root.setRight()!=null){
return root;
}
return null;
}
}
}

```

^ | v • Reply • Share ›



pramod • 7 months ago

=====java=====

static int k=1;

public boolean remove(Node n,int count){

if(n==null) return false;

count=count+1;

boolean left=remove(n.left,count);

boolean right=remove(n.right,count);

if(left) n.left=null;

if(right) n.right=null;

if(n.left==null&& n.right==null&& count<k){ return="" true;="" }else{="" return="" false;="" }="" }="">

^ | v • Reply • Share ›



Gaurav • 8 months ago

I guess we can optimize this. In case level is greater than k, we don't need to visit the nodes

so we can use like this

if(level<k) {="" root="">left = removeShortPathNodesUtil(root->left, level + 1, k);

root->right = removeShortPathNodesUtil(root->right, level + 1, k);

}

If we check the number of recursive calls, it would be less in this case (13 instead of 17 recursive calls)

1 ^ | v • Reply • Share ›



Karan Kapoor • 8 months ago

similar to the previous post of removing nodes when node does not lie in any path $\geq k$

^ | v • Reply • Share ›



Arun Dixit • 8 months ago

```
int removeNode(Node root, int level, int k){
```

```
if(root==null)
```

```
return level;
```

```
int lh = removeNode(root.left,level+1,k);
```

```
int rh = removeNode(root.right,level+1,k);
```

```
int x = max(lh,rh);
```

```
if(root.left==null && root.right==null && x<k) root="null;" return="" x;="" }="">
```

^ | v • Reply • Share ›



saumik chaurasia • 9 months ago

```
TreeNode removeNodes(TreeNode node, int k,int len) {
```

```
if(node == null)
```

```
return null;
```

```
node.leftChild = removeNodes(node.leftChild, k,len+1) ;
```

```
node.rightChild = removeNodes(node.rightChild, k,len+1);
```

```
if(node.leftChild == null && node.rightChild == null && len < k){
```

```
System.out.println("Length " + len);
```

```
node = null;
```

```
}
```

```
return node;
```

```
}
```

^ | v • Reply • Share ›



Dman • 9 months ago



In the given code, if there exists no such path with length $\geq k$, even then root pointer is pointing to some garbage location. Hence, one condition should be added in the end that is

```
if(res==NULL)
root=NULL;
```

or the other method is pass address of pointers in the recursive function code of which is given here : <https://ideone.com/1bnDuo>

1 ^ | v • Reply • Share ›



Ayush ↗ Dman • 9 months ago

well i guess this part covers that case too..

```
if (root->left == NULL && root->right == NULL && level < k)
{
delete root;
return NULL;
}
```

^ | v • Reply • Share ›



DS+Algo • 9 months ago

No need to traverse the whole tree. i.e. when we reach to k-th node, we can return back instead of going further till leaf node.

Solution here: <http://ideone.com/f4rr3z>

1 ^ | v • Reply • Share ›



OURJAY • 10 months ago

if level of leave is <k then="" delete="" leave.="">

^ | v • Reply • Share ›



Enkesh Gupta • 10 months ago

A simple Function-

```
struct node *remove(struct node *root,int k){
if(root==NULL)
return root;
root->left=remove(root->left,k-1);
root->right=remove(root->right,k-1);
if(!root->left && !root->right){
if(k>1){
free(root);
return NULL;
}
}
return root;
```

}

^ | v • Reply • Share ›

**Sanjeev Dhankhar** • 10 months ago

Level order traversal or using stack should be the better approach here...

^ | v • Reply • Share ›

**sanju** • a year ago<http://ideone.com/vyLo9T> No need of two variables k and level. only k will do I guess

^ | v • Reply • Share ›

**Bhargava** • a year ago

May be this is of some helpful !

```
void removeK(TREE root,int count,int k)
```

{

```
TREE temp=NULL;
```

```
if(root==NULL){
```

```
return;
```

}

```
removeK(root->llink,count+1,k);
```

```
removeK(root->rlink,count+1,k); //Post-order Traversal
```

```
if((root->llink!=NULL && count<(k-1) && root->llink->llink==NULL && root->llink->rlink==NULL)){
```

[see more](#)

^ | v • Reply • Share ›

**Guest** • a year ago

don't know why it's not printing first node on my local machine . ideone however gives correct answer

```
Node *removeShortPathNodes(Node *root, int k){
```

```
if(root == NULL || k<=0) return root; // no modification required
```

```
if(!root->left && !root->right && k>0){
```

```
delete root;
```

```
return NULL;

}

root->left = removeShortPathNodes(root->left, k-1);

root->right = removeShortPathNodes(root->right, k-1);

return root;

}
```

<http://ideone.com/iobQ23>

42 ^ | v • Reply • Share ›



AllergicToBitches → Guest • 2 months ago

Ever heard of memory leak dude?? Doing preorder traversal to delete!! huh.

^ | v • Reply • Share ›



Billionaire → Guest • 7 months ago

Why so many upvotes on this one?

This is obviously wrong! You should delete its children before delete itself! i.e. post order

2 ^ | v • Reply • Share ›



AllergicToBitches → Billionaire • 2 months ago

Its 42 guest votes... seems he only upvoted his answer 42 times. lol

^ | v • Reply • Share ›



DS+Algo → Guest • 9 months ago

Ideone is not giving correct answer also.

Your first if condition should be

if(root==NULL || k<=1) return root; //Consider the case when only root is there then k=1 should not delete it

See here:<http://ideone.com/zEymcL>

^ | v • Reply • Share ›



thevagabond85 • a year ago

that obscure int pathLen = 0;

^ | v • Reply • Share ›



Abhinay Madasu • a year ago

Java Solution • <http://ideone.com/T5fnRM>



Java Solution : <http://ideone.com/10jgDm>

^ | v • Reply • Share ›



ayush katara • a year ago

can be done easily using level order traversal

^ | v • Reply • Share ›



ssk → ayush katara • a year ago

Can you explain the algorithm?

1 ^ | v • Reply • Share ›



Bhagwat Singh • a year ago

Hi

We can optimise this code .

As we are going to the depth >k which is not required .

whenever we get the path length == k return from there no need to go further because we only have to check the path with path length < k.

One change and one code addition will optimise the code:

add if(level==k) return root;

and change the if (root->left == NULL && root->right == NULL && level < k) to if (root->left == NULL && root->right == NULL) only.

1 ^ | v • Reply • Share ›



SHAILIN DESAI • a year ago

It can still be optimized further. The above program runs upto the leaf node even if that path length is greater than K.

The following C program will travel one path only till k and if the node has any one children then it will return back.

The time complexity will be O(k). (Not sure though)

```
Node *removeShortPathNodes(Node *root, int k)
```

```
{
```

```
if(root == NULL)
```

```
{
```

```
return NULL;
```

```
}
```

```
else if ( k < 2)
```

```
{
```

```
return root;
```

```
}
```

```
if(root->left != NULL)
```

[see more](#)[^](#) | [v](#) • [Reply](#) • [Share](#) ›**pavan8085** • a year ago

A simple implementation in C

```
int delete_nodes ( node_t **root, int k )
{

static int level = 1;
int max_path = 0;

if ( *root == NULL )
{
/* The non-existent children of leaf nodes are at
* level -1
*/
return -1;
}

level++;

max_path = 1 + max(delete_nodes(&((*root)->left), k),
```

[see more](#)[^](#) | [v](#) • [Reply](#) • [Share](#) ›**chicken** • a year ago

typo-> "pre-order traversal"

1 [^](#) | [v](#) • [Reply](#) • [Share](#) ›**Debanjan Chanda** • a year ago

Problem source?

[^](#) | [v](#) • [Reply](#) • [Share](#) ›**Kumar Gautam** ➔ [Debanjan Chanda](#) • a year ago

It was asked recently in a flipkart telephonic round.

[^](#) | [v](#) • [Reply](#) • [Share](#) ›**d_dll** ➔ [Kumar Gautam](#) • 9 months ago

thanks

[^](#) | [v](#) • [Reply](#) • [Share](#) ›**Debanjan Chanda** ➔ [Kumar Gautam](#) • a year ago



Thanks.

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