

Mathematical Software Programming (02635)

Lecture 5 — October 2, 2025

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This week

Topics

- ▶ Abstract data structures
- ▶ Dynamic memory

Learning objectives

- ▶ Describe and use **data structures** such as arrays, linked lists, stacks, and queues.
- ▶ Choose appropriate data types and **data structures** for a given problem.
- ▶ Design, implement, and document a program that solves a mathematical problem.

Abstract data types

An abstract data type is a high-level abstraction: the *behavior* is defined, not the implementation.

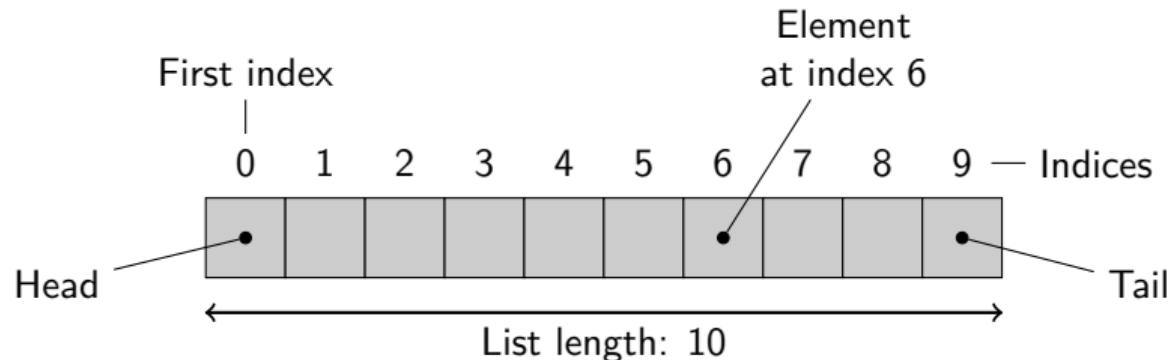
Examples

- ▶ sets and lists
- ▶ vectors, matrices, tensors
- ▶ polynomials
- ▶ complex numbers, rational numbers
- ▶ points, circles, triangles, rectangles, polygons, etc.

Lists

A *list* is an ordered set of elements with the following properties

- ▶ an element can be accessed, inserted, or deleted at any position
- ▶ a list can be split into sublists
- ▶ two lists can be concatenated



Indices can be 0-based or 1-based.

Example

Suppose $L_1 = (5, 3, 9)$, $L_2 = (9, 4, 7, 0)$, and $L_3 = (10, -2, 5, 8)$.

- ▶ concatenation of L_1 and L_2 (denoted $L_1 \cap L_2$)

$$(5, 3, 9) \cap (9, 4, 7, 0) = (5, 3, 9, 9, 4, 7, 0)$$

- ▶ splitting L_3 at index 1 (0-based)

$$\text{split}(L_3, 1) \rightsquigarrow (10), (-2, 5, 8)$$

- ▶ accessing the 3rd element of L_2 (i.e., index 2)

$$\text{access}(L_2, 2) \rightsquigarrow 7$$

- ▶ inserting the value 8 at index 2 of L_1

$$\text{insert}(L_1, 2, 8) \rightsquigarrow L_1 = (5, 3, 8, 9),$$

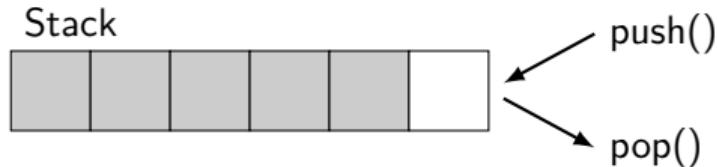
- ▶ deleting the head of L_2

$$\text{delete}(L_2, 0) \rightsquigarrow L_2 = (4, 7, 0).$$

Stacks and queues

Stack

- ▶ a list where elements are inserted and deleted at one end only
- ▶ *first-in-last-out* (FILO)



Queue

- ▶ a list where elements are inserted at one end and deleted/extracted at the other
- ▶ *first-in-first-out* (FIFO)



Structures in C

A *structure* is a type declaration that *groups* a set of variables.

```
struct complex {  
    double real;  
    double imag;  
};
```

```
struct complex a; // a is uninitialized  
struct complex b = {1.0,0.0};  
struct complex c = {.real=0.0, .imag=1.0};
```

- ▶ a, b and c are variables of type struct complex
- ▶ real and imag are so-called *members* of the structure
- ▶ period (.) is the *member access operator* (e.g., c.real = 5;)

Example: sum of elements in array

Data structure with array length and pointer to first element:

```
struct array {  
    size_t length;  
    double *x;  
};
```

Sum function:

```
double sum(struct array a) {  
    double sum = 0.0;  
    for (size_t k=0; k<a.length; k++) sum += a.x[k];  
    return sum;  
}
```

```
double data[] = {1.0,2.0,3.0};  
struct array a = {3, data};  
printf("sum = %g\n", sum(a));
```

Nested structures

A struct member can be another structure, an array, etc.

Example

Structure representing a point in \mathbb{R}^2 :

```
struct point {  
    double x,y;  
};
```

Structure representing a triangle:

```
struct triangle {  
    struct point v1,v2,v3;  
};
```

Arrays of structures

```
struct point point_array[10];           // uninitialized array of length 10
point_array[0].x = 1.0;                  // init. member x of first element
point_array[0].y = 2.0;                  // init. member y of first element
point_array[9].x = 4.0;                  // init. member x of last element
point_array[9].y = 7.0;                  // init. member y of last element
```

```
struct triangle {
    struct point vertices[3];
};
```

Pointer to a structure

- ▶ A pointer to a structure is useful as function input/output
- ▶ Use `p->member` to access a member of a structure via a pointer `p`

Example

```
double data[10] = {0.0};  
struct array a = {10, data};  
struct array * p = &a;  
p->x[0] = 1.0;
```

Remark: `p->x` is equivalent to `(*p).x` whereas `*p.x` is equivalent to `*(p.x)`.

Structure as return value

A function can return a structure or a pointer to a structure.

Examples

```
struct point midpoint(struct point a, struct point b) {
    struct point c = {(a.x+b.x)/2.0,(a.y+b.y)/2.0};
    return c;
}
```

```
struct triangle * largest(struct triangle * a, struct triangle * b) {
    return (area(*a) >= area(*b)) ? a : b;
}
```

Never return a pointer to an *automatic* local variable.

Type definitions

Type definitions are used to create an alias for an existing type:

```
typedef <type> <alias>;
```

Example

```
typedef struct point { double x,y; } point_t;
typedef struct triangle { point_t vertices[3]; } triangle_t;
```

point_t and triangle_t are now aliases for struct point and struct triangle:

```
point_t a={.x=1.0,.y=-2.0}, b={.x=2.0,.y=8.0}, c={.x=5.0,.y=2.0};
triangle_t T={.vertices={a,b,c}};
```

Dynamic memory allocation

Prototypes (stdlib.h)

```
void *malloc(size_t size);
void *calloc(size_t nelements, size_t elementSize);
void *realloc(void *pointer, size_t size);
void free(void *pointer);
```

Allocating storage for an array of length N

```
double *p = malloc(N*sizeof(double));
if (p == NULL) {
    // Code to deal with memory allocation failure ...
}
```

The pointer p must be used to *free/deallocate* storage.

Extending dynamically allocated memory

```
size_t N = 100;
double *p = malloc(N*sizeof(double));
if (p == NULL) { /* Code to handle memory allocation failure */ }

...
// Request more memory (2*N)
N *= 2;
double *ptmp = realloc(p, N*sizeof(double));
if (ptmp == NULL) {
    // Code to handle reallocation failure ...
    // p is still a valid pointer
}
else
    p = ptmp;
```

Freeing/deallocating memory

```
double *p = malloc(100*sizeof(double));  
if (p==NULL) { /* Handle memory allocation failure */ }  
...  
free(p);
```

Common errors

- ▶ Freeing memory twice
- ▶ Attempting to free automatically allocated memory (p was not returned by malloc)
- ▶ Using pointer to access memory *after* freeing memory
- ▶ Forgetting to free memory (may result in a *memory leak*)

Memory: stack vs heap

Stack (automatic allocation)

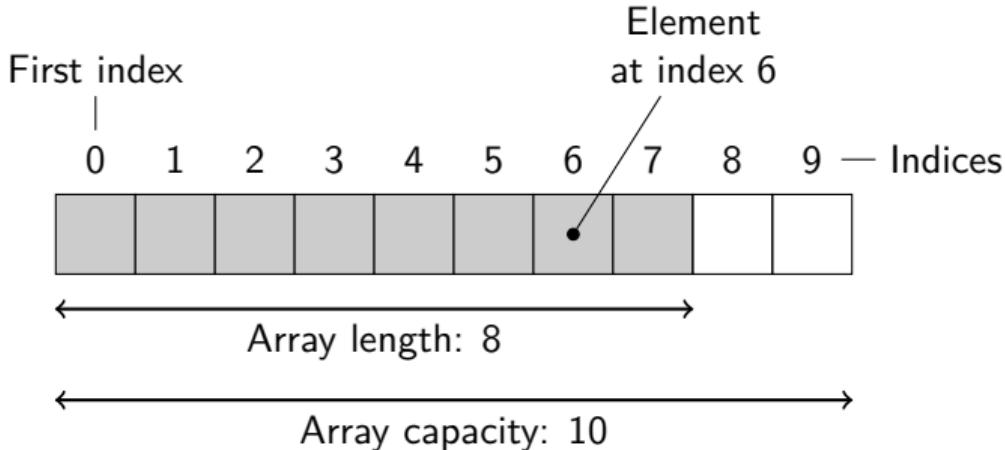
- ▶ Memory layout fixed at compile-time (no resizing of variables or arrays)
- ▶ Used for local variables (function parameters, temporary variables)
- ▶ Fast access (due to contiguous memory and “CPU cache friendliness”)
- ▶ No manual memory management (automatically allocated/deallocated)
- ▶ Limited size (typically 2-10 MB per thread)

Heap (dynamic allocation)

- ▶ Memory allocated at runtime (using `malloc`, etc.)
- ▶ Used for dynamic/global data (accessible across functions)
- ▶ Manual memory management (must explicitly allocate and free memory)
- ▶ Flexible size (only limited by system)
- ▶ Slower access (due to pointer indirection and possible fragmentation)
- ▶ Risk of memory leaks (if not properly deallocated)

Remark: Static and global variables are stored in the “data segment”, not the stack or heap.

Dynamic array



```
typedef struct array {
    size_t len;          // Length of array
    size_t capacity;    // Capacity
    double * val;        // Pointer to first element
} array_t;
```

Example: allocating/deallocating a dynamic array

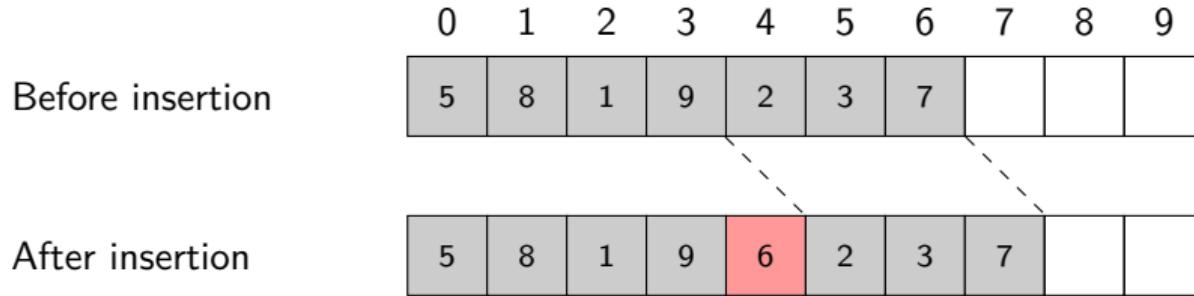
```
array_t *array_alloc(const size_t capacity) {  
    array_t *a = malloc(sizeof(*a));  
    if (a == NULL) return NULL;  
    a->capacity = (capacity > 0) ? capacity : 1; // Minimum capacity: 1  
    a->len = 0;  
    a->val = malloc((a->capacity) * sizeof(*(a->val)));  
    if (a->val == NULL) { free(a); return NULL; }  
    return a;  
}
```

```
void array_dealloc(array_t *a) {  
    if (a) { free(a->val); free(a); }  
}
```

Important: `free(a->val)` must precede `free(a)`

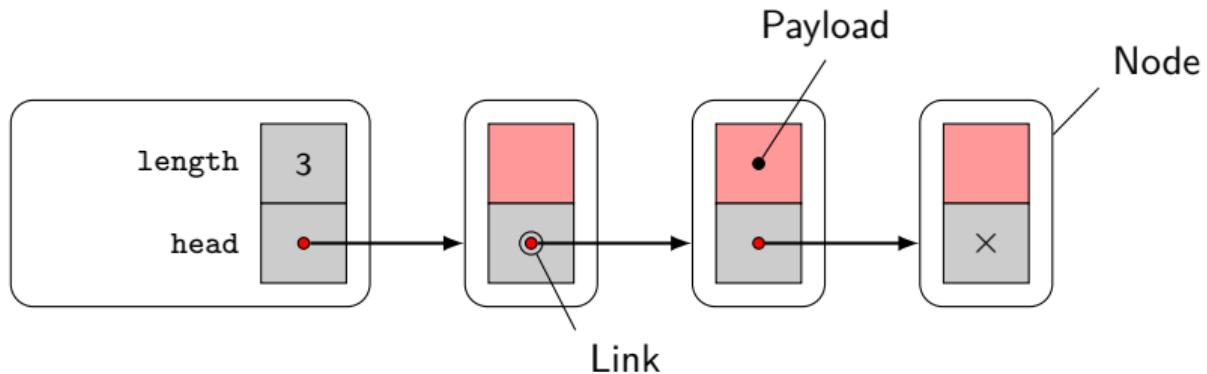
Inserting an element in a dynamic array

Insert “6” at index 4:



- ▶ Inserting at index 0 requires that all entries are relocated
- ▶ Appending (inserting at index len) may involve relocation if length is equal to capacity
- ▶ Increasing capacity may involve relocation
 - ▶ design choice: increase capacity by constant amount/factor/...

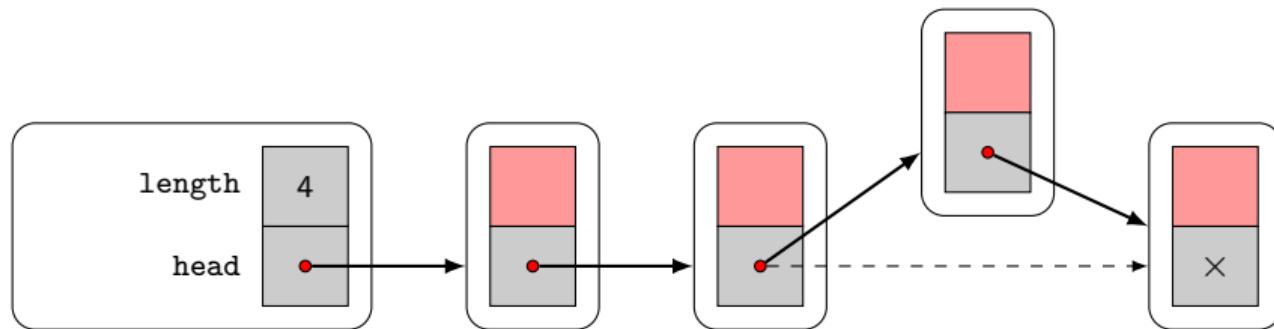
Singly linked list



```
struct snode {
    double x;
    struct snode * next;  };
struct sllist {
    size_t length;
    struct snode * head;  };
```

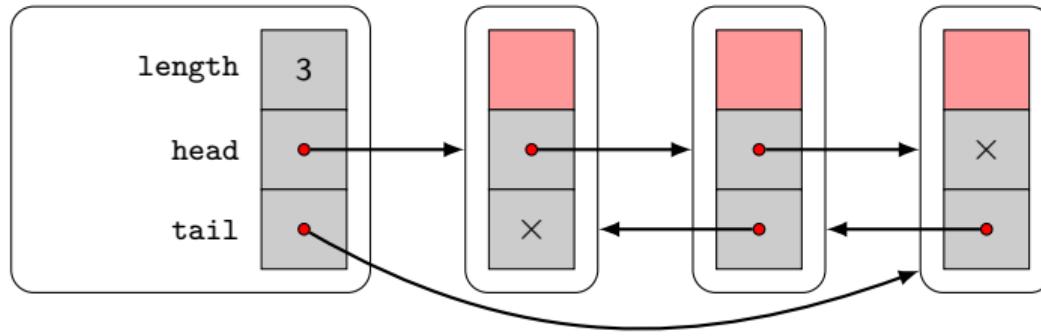
Inserting an element in a singly linked list

Insert node at index 2:



- ▶ Inserting a node at index 0 is fast
- ▶ Inserting a node at any other index:
 - ▶ find preceding node (up to n operations if list has length n)
 - ▶ insert new node (does not depend on length)

Doubly linked list



```
struct dnode {
    double x;
    struct dnode * next;
    struct dnode * prev;  };
struct dllist {
    size_t length;
    struct dnode * head;
    struct dnode * tail;  };
```

Implementing a list

Array-based implementation

- ▶ cost of finding/accessing an element does not depend on list length n
- ▶ worst-case cost of inserting/deleting an element is proportional to n

Linked list implementation

- ▶ worst-case cost of finding/accessing an element is proportional to n
- ▶ cost of inserting/deleting an element does not depend on n

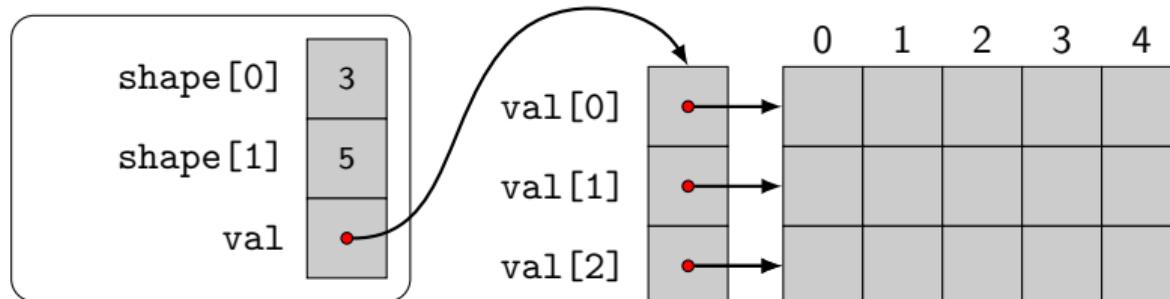
Two-dimensional arrays: linear indexing

```
enum storage_order {RowMajor, ColMajor};  
typedef struct array2d {  
    size_t shape[2];  
    enum storage_order order;    // RowMajor or ColMajor  
    double * val;  
} array2d_t;
```

```
array2d_t *a = array2d_alloc((size_t []){3,5}, RowMajor);  
a->val[i*5+j] = 2.0;    // set element at index (i,j) to 2.0
```

Two-dimensional arrays as an array of arrays

```
typedef struct carray2d {  
    size_t shape[2];  
    double ** val;  
} carray2d_t;
```



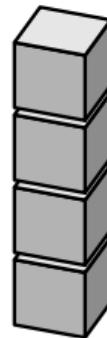
```
carray2d_t * a = carray2d_alloc((size_t []){3,5});  
a->val[i][j] = 2.0; // set element at index (i,j) to 2.0
```

Multidimensional arrays

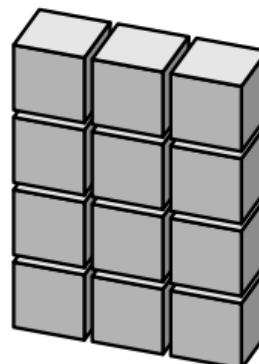
k -dimensional array

- ▶ array *shape* is a k -tuple: (m_1, \dots, m_k)
- ▶ *strides* tuple (s_1, \dots, s_k) determines storage order (generalization of row or column major)
- ▶ indexing requires a *multi-index* which is a k -tuple of the form (i_1, \dots, i_k)

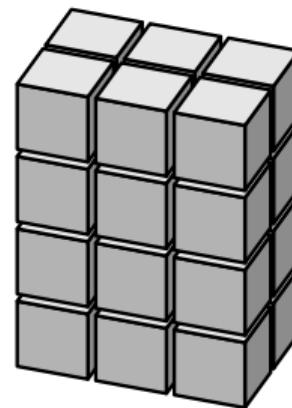
one-dimensional array



two-dimensional array



three-dimensional array



four-dimensional array



MSP Tools

Small, educational C library that is based on the data structures introduced in today's note.

```
CC=gcc
CPPFLAGS=-Imsptools/include
CFLAGS=-Wall -std=c17
LDLIBS=-lm -lmsptools
LDFLAGS=-Lmsptools/lib

.PHONY: all clean
all: msptools myprogram
msptools:
    git clone https://gitlab.gbar.dtu.dk/mskan/msptools.git
    make --directory=msptools
clean:
    -$(RM) myprogram
```

Exercises

- ▶ Abstract data types
- ▶ Structures and implementation
- ▶ Memory allocation

Remember that every `malloc()` call should have a matching `free()` call.