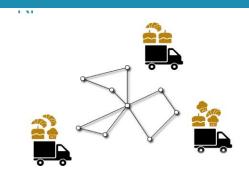
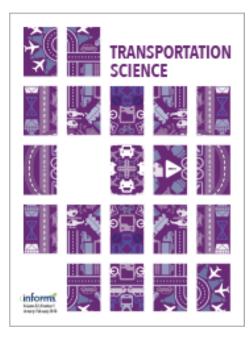


Vehicle routing A focus on heuristic design

schedulingseminar.com

Jan Christiaens
Greet Vanden Berghe





TRANSPORTATION SCIENCE

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Fifty Years of Vehicle Routing

Gilbert Laporte

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The Vehicle Routing Problem (VRP) was introduced 50 years ago by Dantzig and Ramser under the title "The Truck Dispatching Problem." The study of the VRP has given rise to major developments in the fields of exact algorithms and heuristics. In particular, highly sophisticated exact mathematical programming decomposition algorithms and powerful metaheuristics for the VRP have been put forward in recent years. The purpose of this article is to provide a brief account of this development.

Key words: vehicle routing problem; traveling salesman problem; exact algorithms; heuristics; metaheuristics; survey

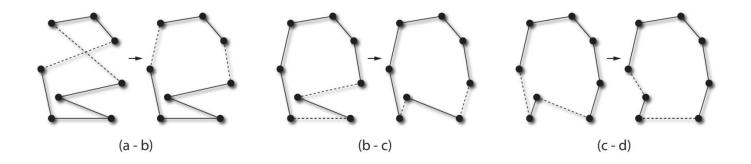
History: Received: August 2009; revision received: September 2009; accepted: September 2009. Published online in Articles in Advance October 21, 2009.



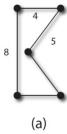
Heuristics

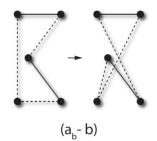


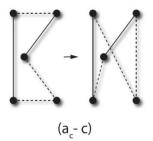
TSP – 2-opt

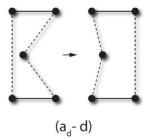


TSP - 3-opt







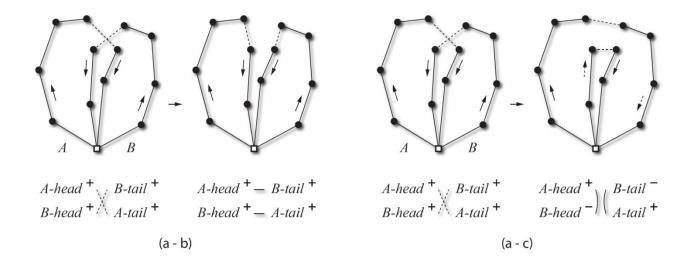


TSP - More heuristics

- k-opt (Lin and Kernighan, 1973)
 - Remove k edges, reconnect the possibly reversed strings
- *or-opt* (Or, 1976)
 - Relocate a string with maximum cardinality 3



VRP – Inter-route 2-opt

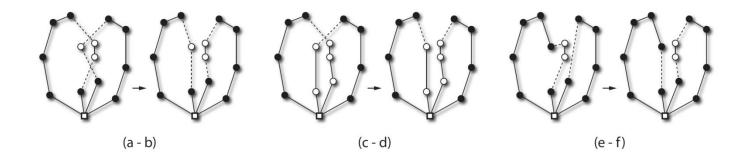


VRP - More neighborhoods

- *Or-opt* (Or, 1976)
 - Relocate 1 string with maximum cardinality 3, possibly to a different route
- Single, double, pair+single, double pair (Waters, 1987)
 - Relocate 2 strings with maximum cardinality 2, possibly to different routes
- Chain-exchange (Fahrion and Wrede, 1990)
 - Relocate 2 strings with maximum cardinality |t| / 2



VRPTW - Exchange, Cross, Relocate



VRPTW - More heuristics

- 2-opt* (Potvin and Rousseau, 1995)
 - o inter-route 2-opt without string reversal, or Cross
- CROSS-exchange (Taillard et al., 1997)
 - Savelsbergh's Exchange, Cross, Relocate
 - (+) Relocate may operate as intra-route operator



Ruin & recreate (R&R)

- Term introduced by Schrimpf et al.
 - Record Breaking Optimization Results Using the Ruin and Recreate Principle (2000)
 - Optimization with Ruin Recreate (2002)

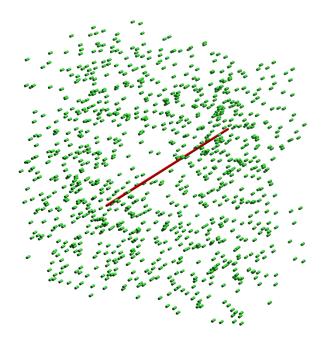
(United States Patent, IBM)

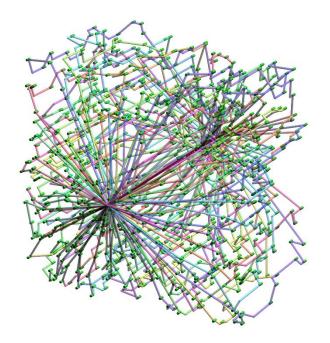
- # ruins, 1 recreate
- Concept established in older publications
 - Performance of Interconnection Rip-Up and Reroute Strategies (1981)



VRPTW heuristics

| VRPTW heuristics | Reference | | | |
|------------------------------|--------------------------|--|--|--|
| Exchange, cross and relocate | Savelsbergh 1988 | | | |
| 2-opt* | Potvin and Rousseau 1995 | | | |
| CROSS-exchange | Taillard et al. 1997 | | | |





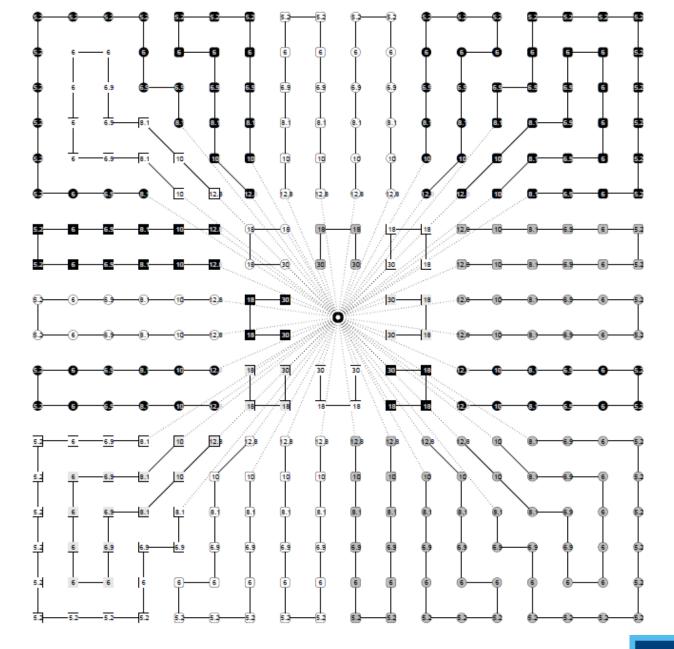
Ruin & recreate heuristics

| | Reference | Ruin | Recreate |
|-------|-------------------------|--|---|
| LNS | Shaw 1998 | related removal | Branch & Bound |
| R&R | Schrimpf et al. 2000 | random, radial or string removal | greedy insertion |
| ALNS | Pisinger and Røpke 2007 | random, worst, related, cluster, time-oriented or historical removal | greedy or regret insertion with(out) noise function |
| SISRs | THIS TALK | adjacent string removal | greedy insertion with blinks |



Data





CVRPLIB

Capacitated Vehicle Routing Problem Library

| All Instances | Plots | New Instances | CVRP Challenge | Updates | Posts | Links | About |
|---------------|-------|---------------|----------------|---------|-------|-------|-------|
|---------------|-------|---------------|----------------|---------|-------|-------|-------|

You are here: Home

| Benchmark | Instance | n | K | Q | UB | Opt | Features |
|---|----------|---|---|---|----|-----|----------|
| Set A (Augerat, 1995) | | , | | | , | | |
| Set B (Augerat, 1995) | | | | | | | € 🗐 |
| ▶ Set E (Christofides and Eilon, 1969) | | | | | | | € 🗐 |
| Set F (Fisher, 1994) | | | | | | | ₽ 🗐 |
| Set M (Christofides, Mingozzi and Toth, 1979) | | | | | | | ₽ 🗐 |
| Set P (Augerat, 1995) | | | | | | | € 🗐 |
| Christofides, Mingozzi and Toth (1979) | | | | | | | € 🗐 |
| Rochat and Taillard (1995) | | | | | | | ₽ |
| ▶ Golden et al. (1998) | | | | | | | ₽ 🗐 |
| Li et al. (2005) | | | | | | | € 🗐 |
| ▶ Uchoa et al. (2014) | | | | | | | € 🖨 |
| | | | | | | | |





February, 23rd 2018- Improved BKSs reported by Keld Helsga in using <u>LKH-3</u>: X-n524 k153 X-n837-k142 (193810), X-n856-k95 (89002), X-n936-k151 (13292 X X-n957-k87 (85-78) and x-n857-k142 (193810), X-n856-k95 (89002), X-n936-k151 (13292 X X-n957-k87 (85-78))

August, 14th 2017 - Improved BKSs reported (Toffolo, Vidal, Wauters (2017), Heuristics for paper): X-n294-k50 (47161), X-n322-k28 (29834), X-n327-k20 (27532), X-n344-k43 (42056),



B19-k171 (158265),

appear as a working

■ November, 23rd 2016 - Improved BKSs reported by Jan Christiaens obtained using ASB-RF (55269), X-n459-k26 (24173), X-n491-k59 (66510), X-n502-k39 (69230), X-n536-k96 (94988), X-n501-k42 (42722), X-n599-k92 (100490), X-n627-k43 (62210), X-n670-k130 (146451), X-n685-k75 (68261), X-n701-k44 (81934), X-n716-k35 (43414), X-n749-k98 (77365), X-n766-k71 (114525), X-n783-k48 (72445), X-n801-k40 (73331), X-n837-k142 (193813), X-n876-k59 (99331), X-n895-k37 (53946), X-n916-k207 (329247), X-n936-k151 (132926) and X-n957-k87 (85482).

- November, 6th 2016 Improved BKS to X-n256-k16 (18839) reported by Túlio Toffolo and Thibaut Vidal. This contradicts a previous claim that a solution with value 18880 was optimal. After investigation, it was found that a typo in a script made the BCP method (Pecin et al., 2014) to be run with a minimum of 17 routes. The improving solution has 16 routes. The status of that instance is also corrected to "open".
- April, 25th 2016 Proven optimal solutions by the BCP method (Pecin et al., 2014): X-n214-k11 (10856) and X-n233-k16 (19230).
- April, 4th 2016 -Improved and proven optimal solutions by the BCP method (Pecin et al., 2014): X-n331-k15 (31102) and X-n439-k37 (36391).

April, 1st 2016 - Improved BKSs reported by Jan Christiaens: X-n322-k28 (29848), X-n336-k84 (139135), X-n344-k43 (42068), X-n351-k40 (25928), X-n384-k52 (65981), X-n401-k29 (66202), X-n449-k29 (55302), X-n459-k26 (24179), X-n480-k70 (89458), X-n491-k59 (66520), X-n502-k39 (69232), X-n536-k96 (94991), X-n548-k50 (86701), X-n573-k30 (50719), X-n586-k159 (190423), X-n599-k92 (108541), X-n613-k62 (59556), X-n627-k43 (62217), X-n641-k35 (63737), X-n670-k130 (146477), X-n685-k75 (68276), X-n701-k44 (81962), X-n716-k35 (43441), X-n733-k159 (136250), X-n749-k98 (77402), X-n766-k71 (114534), X-n783-k48 (72453), X-n801-k40 (73344), X-n819-k171 (158267), X-n837-k142 (193836), X-n856-k95 (89007), X-n876-k59 (99360), X-n895-k37 (53948), X-n916-k207 (329299), X-n957-k87 (85517), X-n979-k58 (119008) and X-n1001-k43 (72404).

■ June, 22nd 2015 - Improved BKSs reported by Jan Christiaens: X-n322-k28 (29854), X-n336-k84 (139165), X-n344-k43 (42073), X-n351-k40 (25936), X-n384-k52 (66021), X-n401-k29 (66219), X-n480-k70 (89488), X-n491-k59 (66523), X-n536-k96 (95062), X-n561-k42 (42754), X-n573-k30 (50726), X-n586-k159 (190454), X-n599-k92 (108600), X-n627-k43 (62264), X-n641-k35 (63760), X-n670-k130 (146570), X-n685-k75 (68291), X-n701-k44 (81997), X-n716-k35 (43491), X-n733-k159 (136313), X-n749-k98 (77423), X-n766-k71 (114566), X-n783-k48 (72547), X-n801-k40 (73367), X-n819-k171 (158298), X-n837-k142 (193933), X-n856-k95 (89040), X-n876-k59 (99424), X-n895-k37 (54030), X-n916-k207 (329394), X-n936-k151 (132946), X-n957-k87 (85566), X-n979-k58 (119072) and X-n1001-k43 (72477).



SISRs

- 1. Ruin
- 2. Recreate
- 3. Fleet minimization



1. SISRs ruin

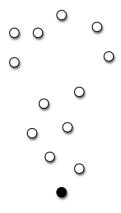
Remove customers – Adjacent string removal

- capacity slack
- spatial slack

- vehicles gain free capacity
- vehicles are detached from regions

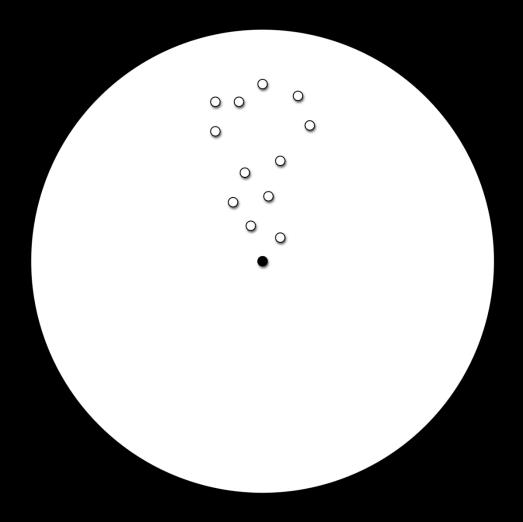


Spatial slack



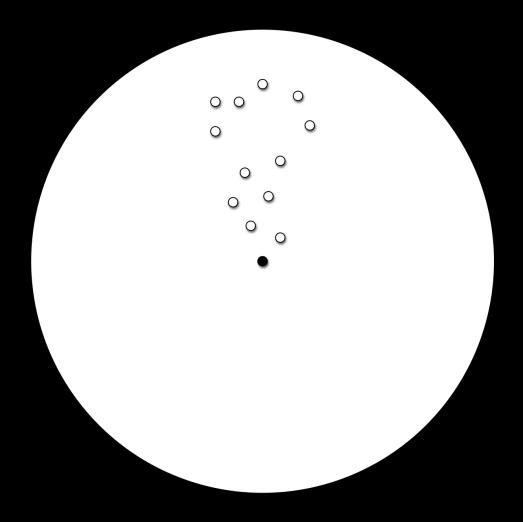


Spatial slack – Maximum distance





Spatial slack – Maximum distance





Spatial slack – Total distance

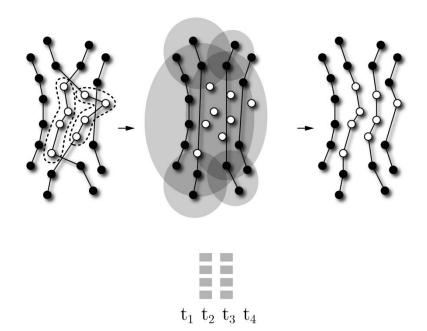




Spatial slack – Total distance

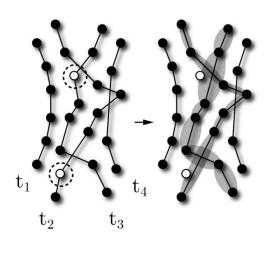




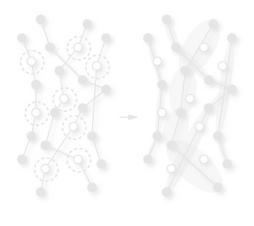


Ruin methods - Alternatives

Small number

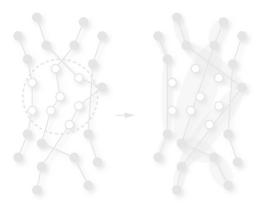


Random





Radial

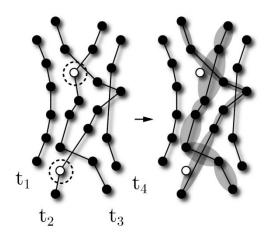






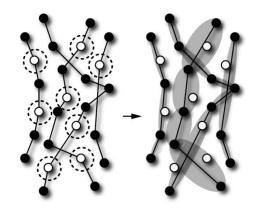
Ruin methods - Alternatives

Small number



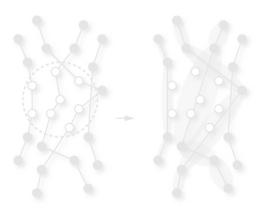
 t_1 t_2 t_3 t_4

Random





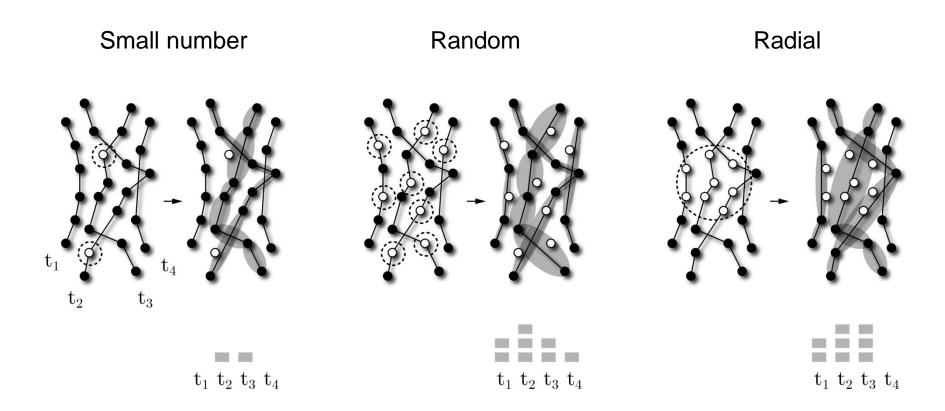
Radial

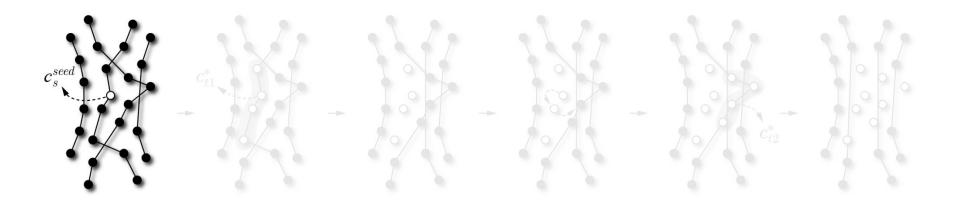


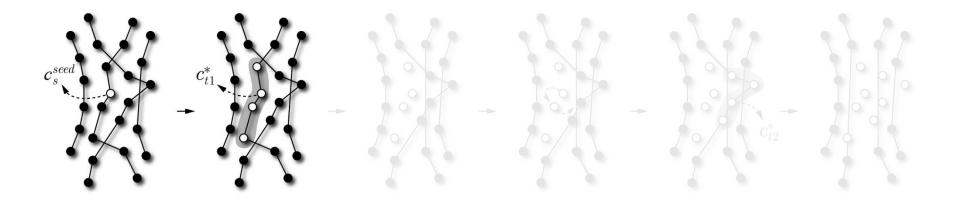


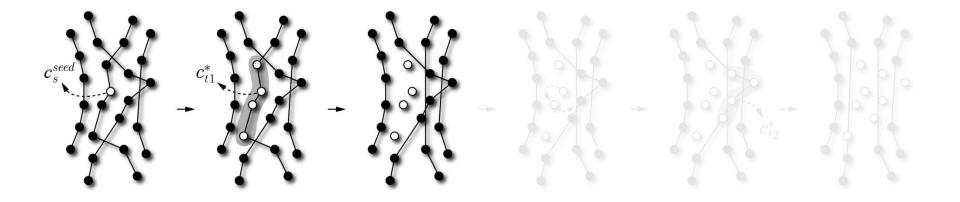


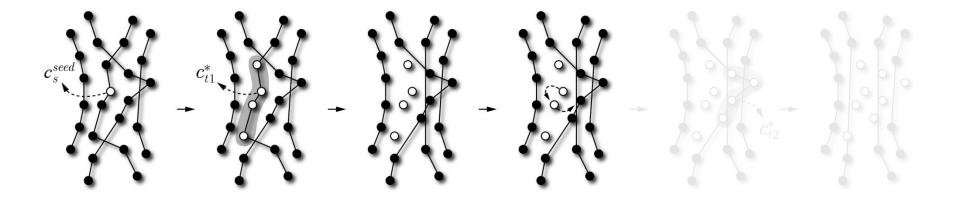
Ruin methods - Alternatives

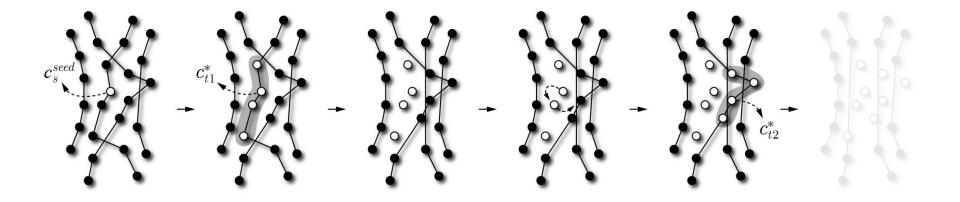


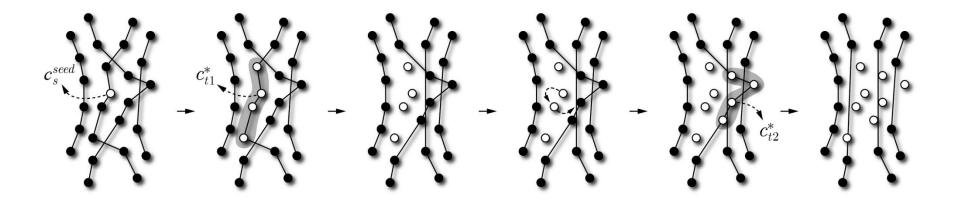












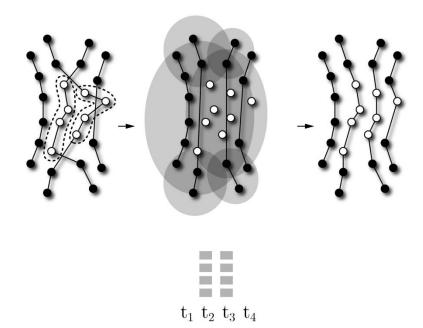
2. SISRs recreate

- insert customers
- options enabled by slack

- greedy insertion with blinks
- strongly dependent on ruin



SISRs recreate



SISRs recreate – Greedy insertion with blinks

Greedy insertion:

- always at best position
- very greedy ⊗

Greedy insertion with blinks

- •
- •

- •
- •
- •

```
procedure RECREATE(s)
     sort(A)
     for c \in A do
          P \leftarrow \text{NULL}
          for t \in T (which can serve c) do
                for P_t in t do
                           if P = \text{NULL or } costAt(P_t) < costAt(P) then
                                 P \leftarrow P_t
          if P = \text{NULL then}
                T \leftarrow T \cup \text{new tour } t
                P \leftarrow \text{position in } t
          insert c at P
           A \leftarrow A \setminus \{c\}
end procedure
```

SISRs recreate – Greedy insertion with blinks

Greedy insertion:

- always at best position
- very greedy ⊗

Greedy insertion with blinks:

- mostly at best position
- less greedy ☺
- •

- •
- •

```
procedure RECREATE(s)
     sort(A)
     for c \in A do
          P \leftarrow \text{NULL}
          for t \in T (which can serve c) do
                for P_t in t do
                     if U(0,1) < 1-\beta then
                          if P = \text{NULL or } costAt(P_t) < costAt(P) then
                                P \leftarrow P_t
          if P = \text{NULL then}
                T \leftarrow T \cup \text{new tour } t
                P \leftarrow \text{position in } t
          insert c at P
          A \leftarrow A \setminus \{c\}
end procedure
```

SISRs Recreate – Greedy insertion with blinks

Greedy insertion:

- always at best position
- very greedy ⊗

Greedy insertion with blinks:

- mostly at best position
- less greedy ☺

- HBSS (Bresina 1996)
- rank all options (best to worst)
- exponentially decreasing probability

```
procedure RECREATE(s)
     sort(A)
     for c \in A do
          P \leftarrow \text{NULL}
          for t \in T (which can serve c) do
                for P_t in t do
                     if U(0,1) < 1-\beta then
                           if P = \text{NULL or } costAt(P_t) < costAt(P) then
                                P \leftarrow P_t
          if P = \text{NULL then}
                T \leftarrow T \cup \text{new tour } t
                P \leftarrow \text{position in } t
          insert c at P
          A \leftarrow A \setminus \{c\}
end procedure
```

SISRs recreate – Greedy insertion with blinks

Greedy insertion:

- always at best position
- very greedy ⊗

Greedy insertion with blinks:

- mostly at best position
- less greedy ☺
- rank-based selection probabilities
- without ranking all options!! ©

- HBSS (Bresina 1996)
- rank all options (best to worst)
- exponentially decreasing probability

```
procedure RECREATE(s)
     sort(A)
     for c \in A do
          P \leftarrow \text{NULL}
          for t \in T (which can serve c) do
                for P_t in t do
                     if U(0,1) < 1-\beta then
                           if P = \text{NULL or } costAt(P_t) < costAt(P) then
                                P \leftarrow P_t
          if P = \text{NULL then}
                T \leftarrow T \cup \text{new tour } t
                P \leftarrow \text{position in } t
          insert c at P
          A \leftarrow A \setminus \{c\}
end procedure
```

3. SISRs fleet minimization

• Fleet minimization - absences-based acceptance criterion





SISRs fleet minimization

- VRPTW, PDPTW ...
- Not strictly related to distance → two-stage approach
 - 1. Fleet minimization
 - 2. Distance minimization
- Nagata, Bräysy, and Dullaert (2010)
 Nagata and Kobayashi (2010)
 PDPTW
 - 1. Remove vehicle
 - "Squeeze" until all customers are served
 - 3. Repeat



Fleet minimization – 1st attempt

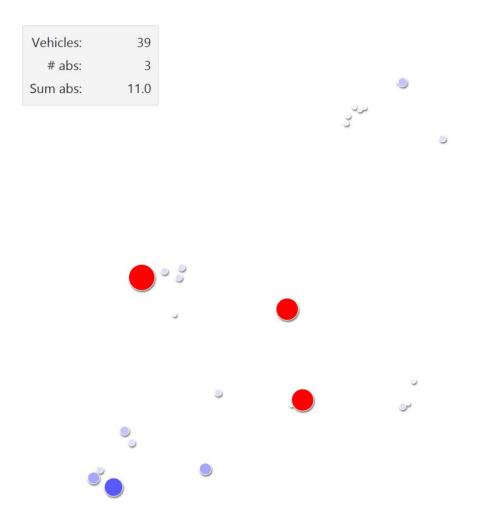
- Consider partial solutions
- Let A = set of absent customers
 - 1. Apply SISRs
 - 2. Accept s* if

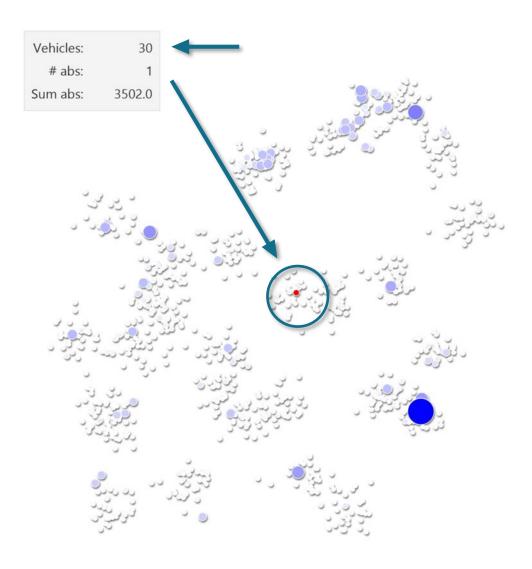


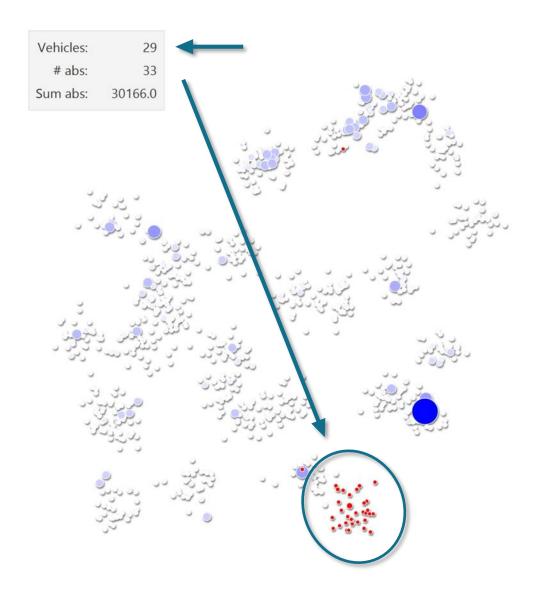
3. Repeat

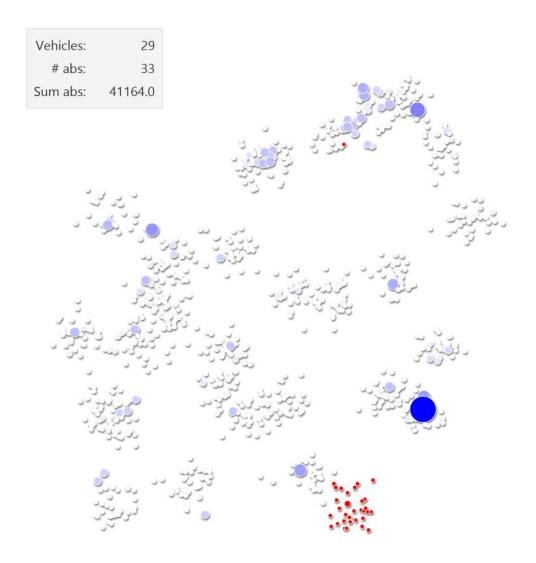


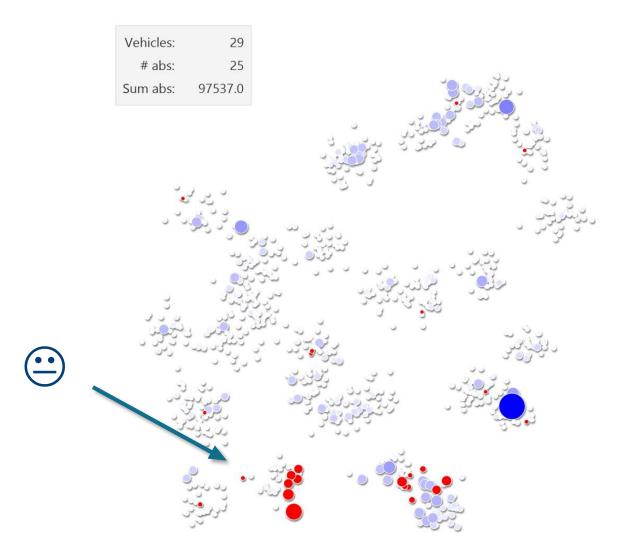
Vehicles: # abs: Sum abs:

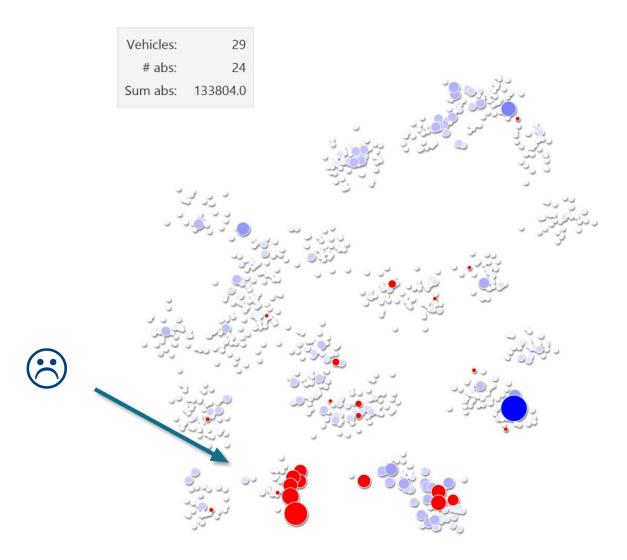


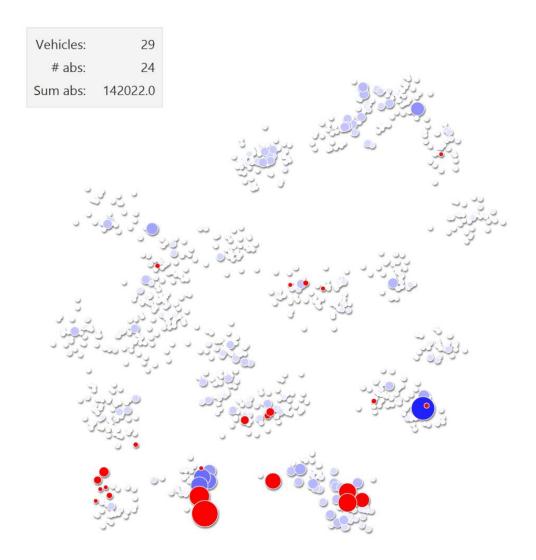


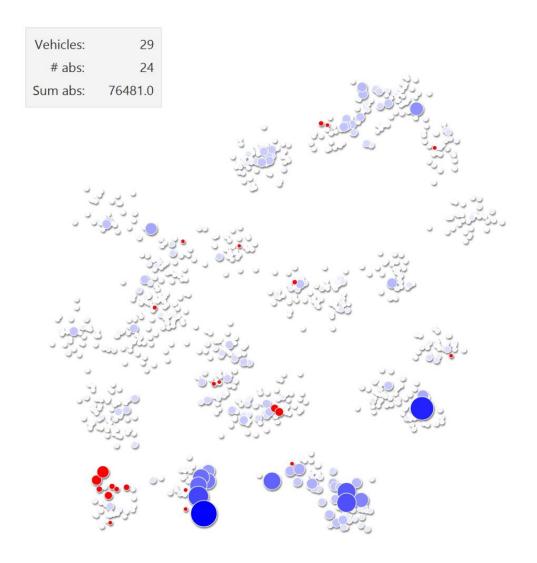


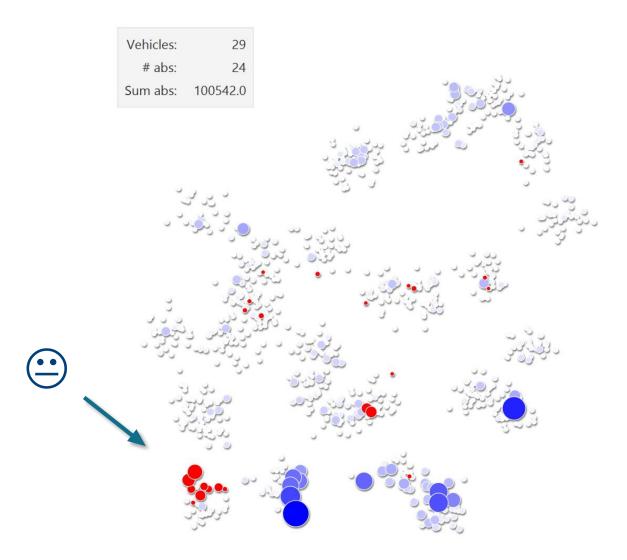


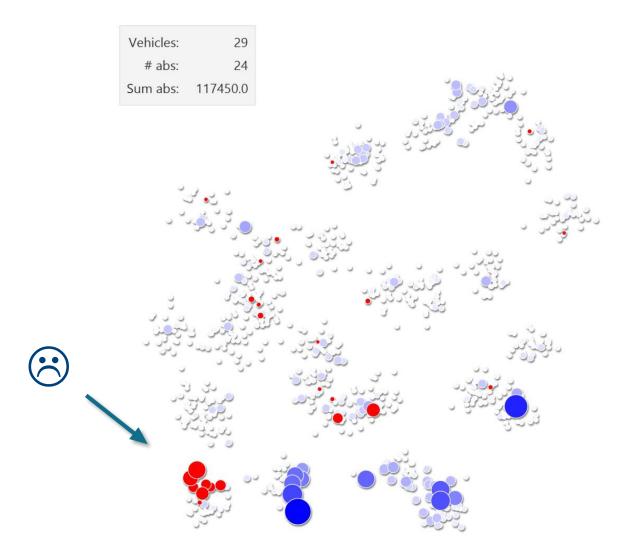


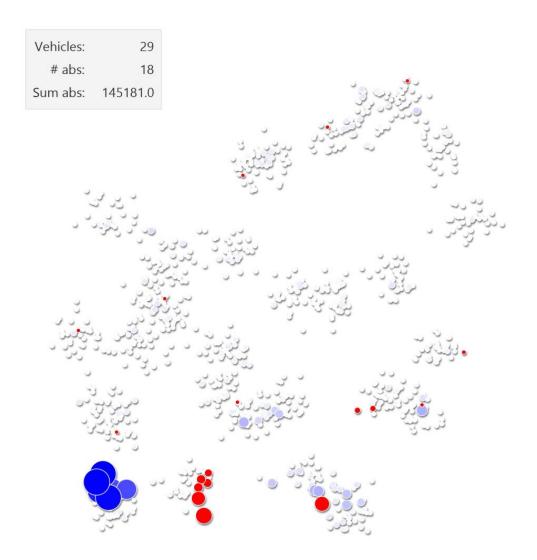


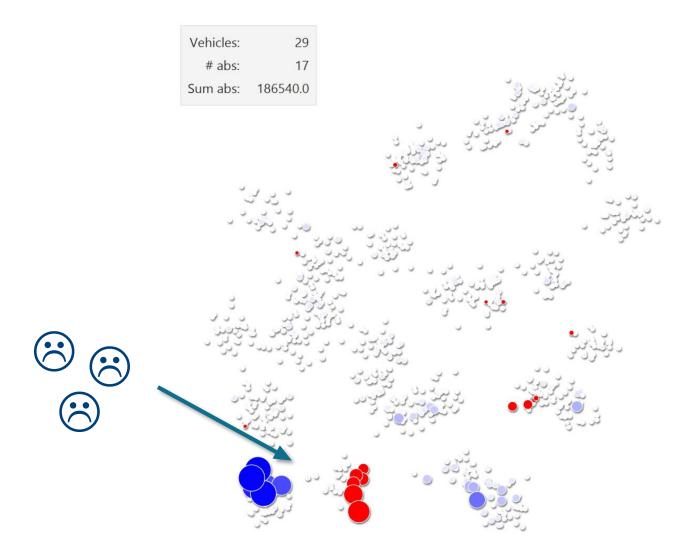












Absences-based acceptance criterion

Accept s* if



Absences-based acceptance criterion

Accept s* if





Absences-based acceptance criterion

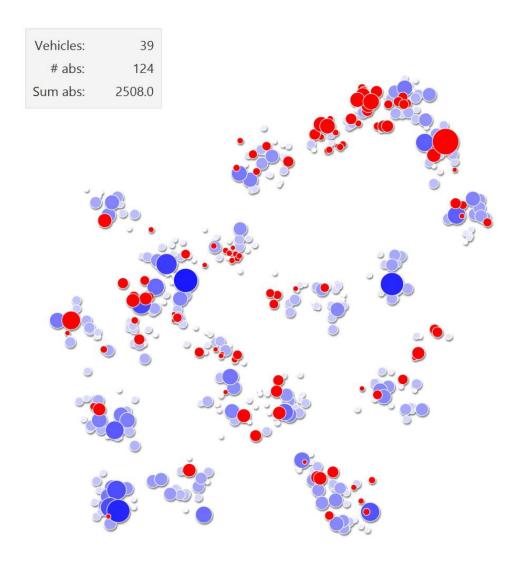
Accept s* if



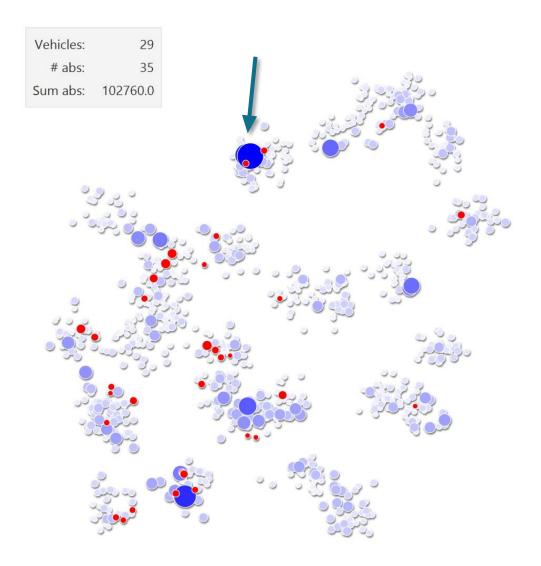
or if

$$\sum_{c \in A} abs_c > \sum_{c \in A^*} abs_c$$

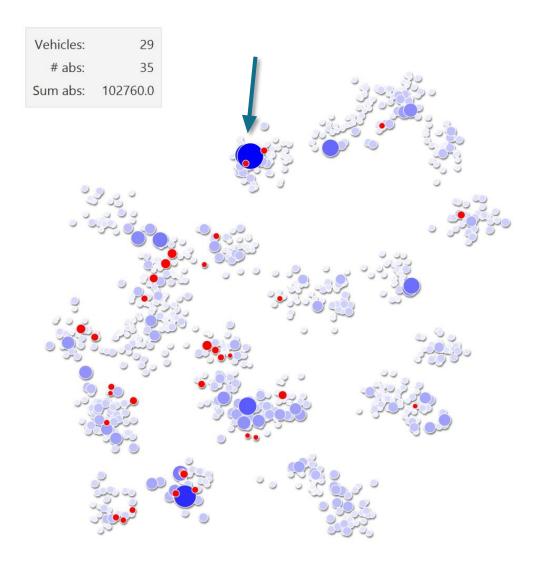
Vehicles: # abs: Sum abs:

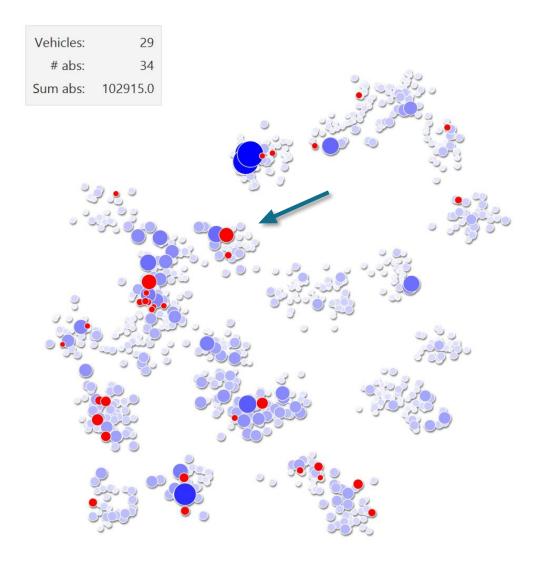


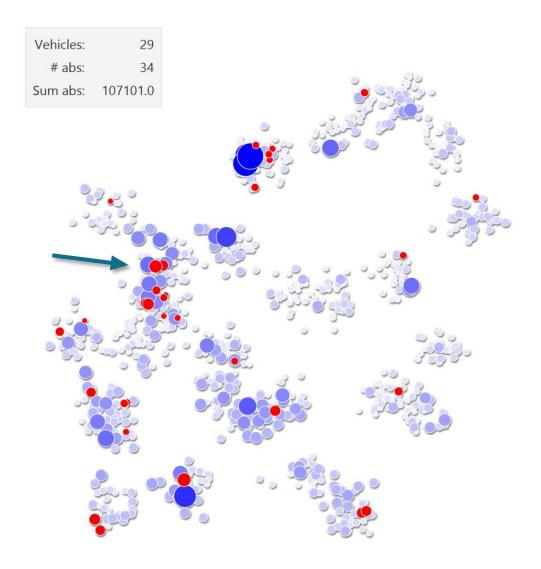
Skip to 29 vehicles

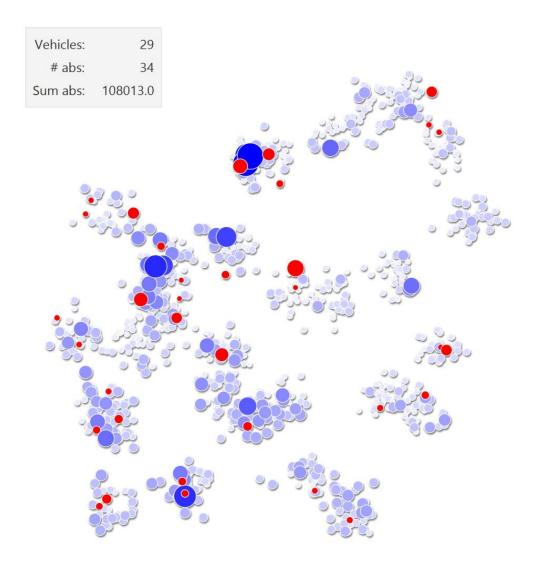


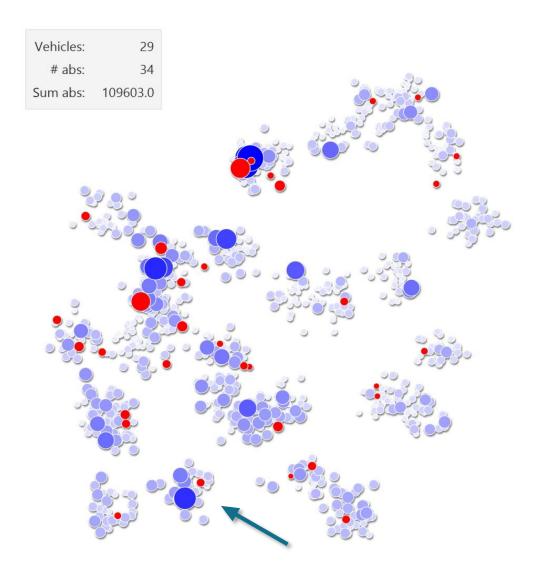


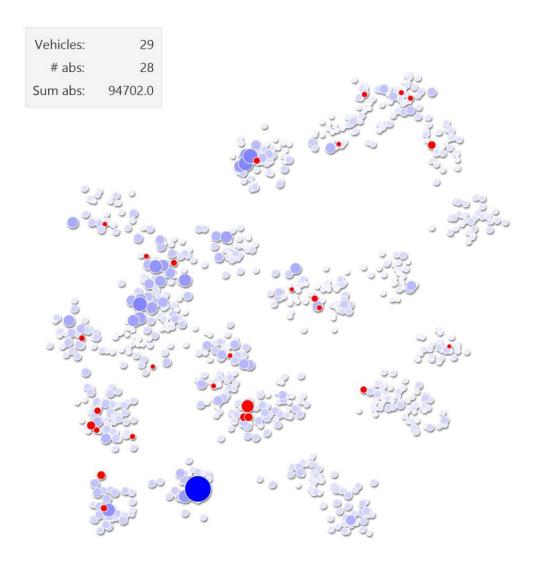








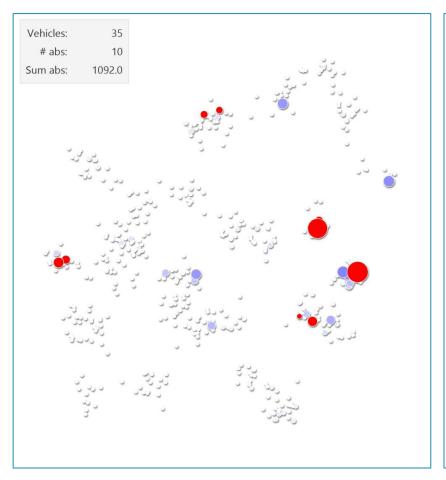


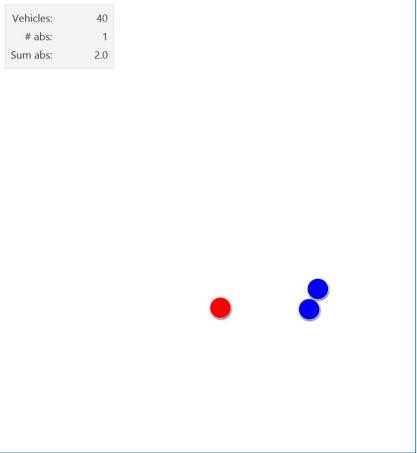


Acceptance criteria compared – Start



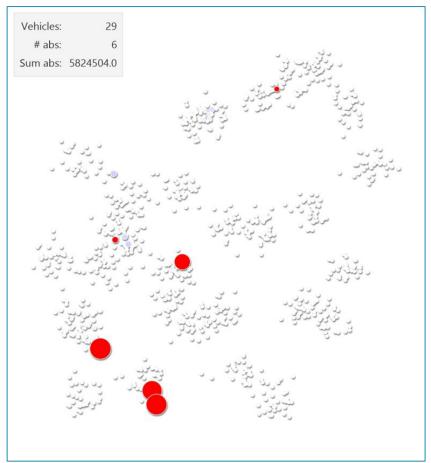
Acceptance criteria compared – Start

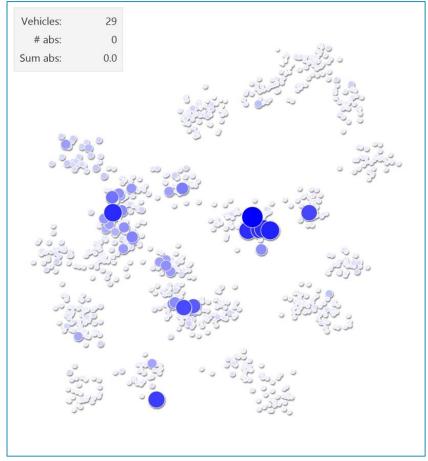




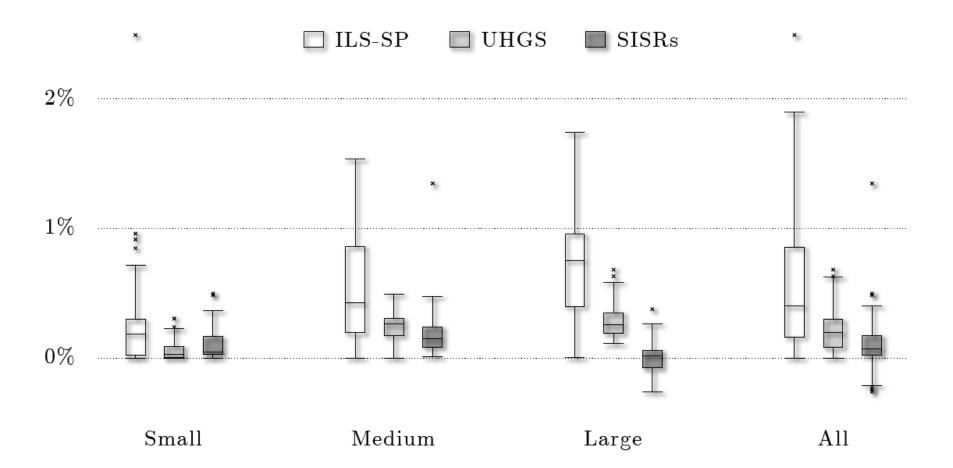


Acceptance criteria compared – Final solutions



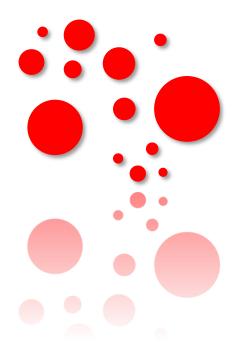












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