

6 Explain the concept of arithmetic coding and illustrate its process with an Example of compressing the string ABABABBB.

Arithmetic coding is a lossless data compression technique used in data compression systems. Unlike fixed length or variable-length codes, arithmetic coding represents an entire message as a single fractional number between 0 and 1.

Instead of assigning separate bit codes to symbols, arithmetic coding progressively narrows a range  $[low, high]$  based on symbol probabilities until the final interval uniquely represents the message.

The cumulative probabilities of symbols are represented on a line from 0.0 to 1.0.

Each symbol occupies a sub-range proportional to its probability.

$$\text{New Range} = S + P(c) \times R$$

$S$  = cumulative probability of previous symbols

$P(c)$  = probability of current symbol

$R$  = current range (high - low)

input string ABABABBB

Step 1: R = 0

Number of A = 3

0.4 + 0.37 = 0.77

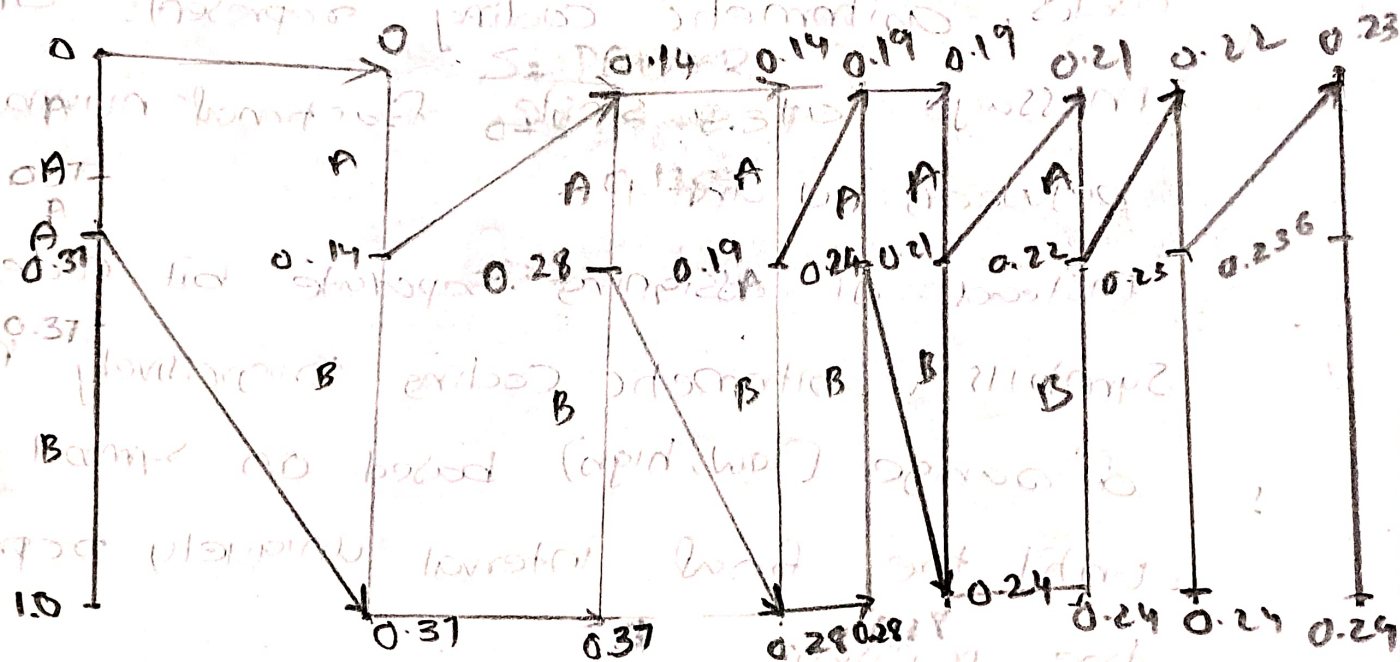
Number of B = 5

0.14

Probability of A =  $\frac{3}{8} = 0.37$

Probability of B =  $\frac{5}{8} = 0.62$

$F_1(A) = 0.37$   $F_1(B) = 1.0$



i)  $S + P(A) * R$

$= 0 + 0.37 * 0.37$

$= 0.14$

ii)  $S + P(B) * R$

$0.14 + 0.62 * 0.23$

$= 0.28$

iii)  $S + P(A) * R$

$0.14 + 0.37 * 0.14$

$= 0.19$

iv)  $S + P(B) * R$

$= 0.19 + 0.62 * 0.09 = 0.24$

v)  $S + P(A) * R$

$0.19 + 0.37 * 0.05$

$= 0.2$

vi)  $S + P(B) * R$

$0.21 + 0.62 * 0.03$

$= 0.22$

vii)  $S + P(B) * R$

$0.22 + 0.62 * 0.02 = 0.23$

viii)  $S + P(B) * R$

$0.23 + 0.62 * 0.01$

$= 0.2362$

tap =  $\frac{0.23 + 0.24}{2} = 0.235$

~~Arithmetic~~ Arithmetic coding generates a unique tag for a sequence without building codes for all sequences of length  $n$  unlike Huffman coding.

The interval is repeatedly subdivided into same proportions as the original range.

The advantages are high compression efficiency than Huffman coding and close to entropy limit.