Chapter 7

The Disjoint Set ADT

并查集很少考

做最小生成树时用到了并查集

记得最后1个算法

Equivalence Class

等价类. 自反、对称、传递.

1) **Definition of Equivalence Class:**

Suppose we have a set U={1,2,…,n} of n elements and a set R={(i_1,j_1), (i_2,j_2)…… (i_r,j_r)}of r relations. The relation R is an equivalence relation iff the following conditions are true(symbol'≡' represent the equivalence relation on sets, x,y,z are elements in set) :

- **Reflexive x ≡x.**
- **Symmetric x ≡y,y ≡x**
- **Transitive x ≡y and y ≡z,then x ≡z**

Equivalence Class

例如：

判别3个数a, b, c能否构成三角形的三条边？

能构成三角形的等价类：

{ (3,4,5), (4,5,6), (5,6,7), …..}

不能构成三角形的等价类：

{ (1,2,3), (2,3,5), ….. }

Equivalence Class

2) **Example:**

set s= {0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11}

pairs of equivalence:

(0 4),(3 1),(6 10),(8 9),(7 4),(6 8),(3 5),(2 11), (11 0)

Initial:{0}, {1}, {2}, {3}, {4}, {5}, {6}, {7}, {8}, {9}, {10}, {11}

0≡4 {0, 4}, {1}, {2}, {3}, {5}, {6}, {7}, {8}, {9}, {10}, {11}

3 ≡1 {0, 4}, {1, 3}, {2}, {5}, {6}, {7}, {8}, {9}, {10}, {11}

Equivalence Class

$6 \equiv 10$ {0, 4}, {1, 3}, {2}, {5}, {6, 10}, {7}, {8}, {9}, {11}

$8 \equiv 9$ {0, 4}, {1, 3}, {2}, {5}, {6, 10}, {7}, {8, 9}, {11}

$7 \equiv 4$ {0, 4, 7}, {1, 3}, {2}, {5}, {6, 10}, {8, 9}, {11}

$6 \equiv 8$ {0, 4, 7}, {1, 3}, {2}, {5}, {6, 8, 9, 10}, {11}

$3 \equiv 5$ {0, 4, 7}, {1, 3, 5}, {2}, {6, 8, 9, 10}, {11}

$2 \equiv 11$ {0, 4, 7}, {1, 3, 5}, {2, 11}, {6, 8, 9, 10}

$11 \equiv 0$ {0, 4, 7, 2, 11}, {1, 3, 5}, {6, 8, 9, 10}

Equivalence Class

3) **Online equivalence class operation**

- **Combine(a,b) : combine the equivalence classes that contains elements a and b into a single class**

- **Find(e) : determine the class that currently contains element e.**

找出e所在的等价类.

Combine(a,b) is equivalent to

i=Find(a); j=Find(b); if(i!=j) Union(i,j);

将这2个等价类做并集

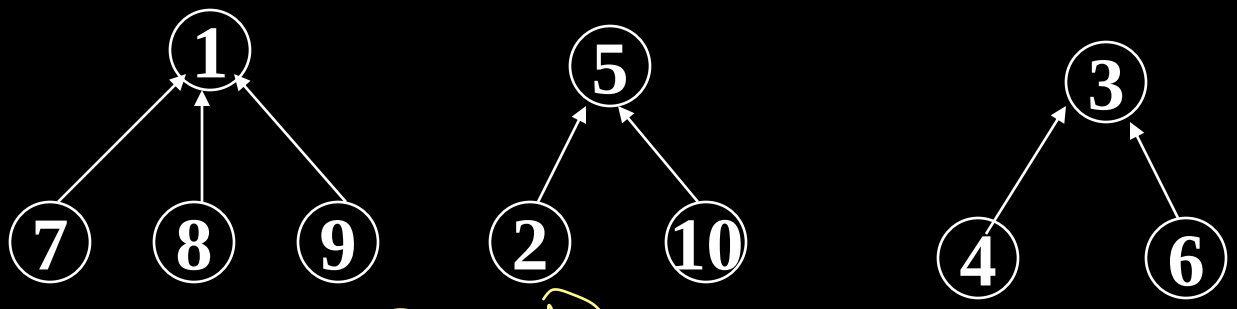
Equivalence Class

4) Tree Representation(Union-Find sets)

Example:

S_1={1,7,8,9}, S_2={5,2,10}, S_3={3,4,6},they all belong to S={1,2,3,……,10}

森林实现



	0	1	2	3	4	5	6	7	8	9	10
parent		0	5	0	3	0	3	1	1	1	5

1、5、3是树根

Equivalence Class

simple tree solution to union-find problem

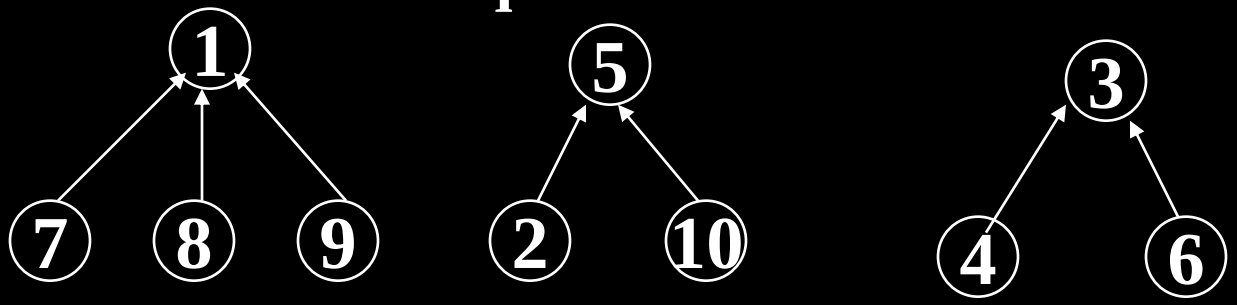
```
  void Initialize(int n)
  {  parent=new int[n+1];
     for(int e=1;e<=n;e++)
        parent[e]=0;
   }

int Find(int e)
 {  while(parent[e])
          e=parent[e];
     return e;
 }
  void Union(int i, int j)
 { parent[j]=i;
 }
```

所有的结点都是树根.
(单独的结点)

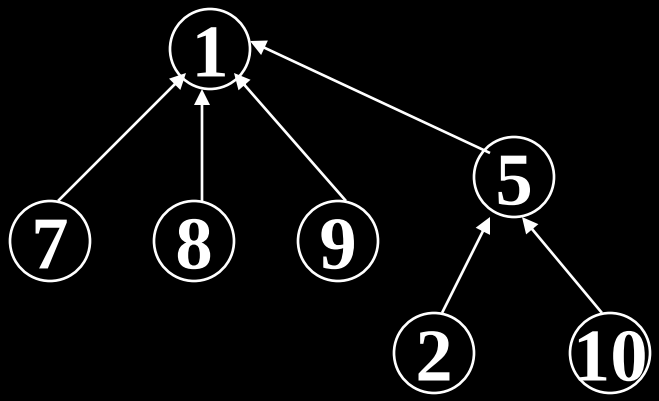
找到父结点为0时停止.

Equivalence Class



	0	1	2	3	4	5	6	7	8	9	10
parent		0	5	0	3	10	3	1	1	1	5

Union (1,5)



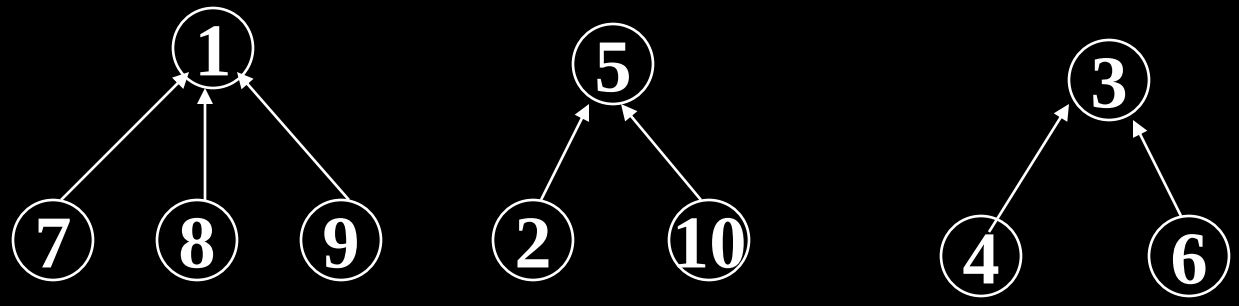
Equivalence Class

```
Java
 public class DisjSets
{  public DisjSets( int numElements )
    public void union( int root1, int root2 )
    public int find( int x )
    private int [ ] s;
}

public DisjSets( int numElements )
{   s = new int [ numElements ];
     for( int i = 0; i < s.length; i++ )
          s[ i ] = -1; //一个根结点
}
```

Equivalence Class

```
 public void union( int root1, int root2 )
{   s[ root2 ] = root1;
}
 public int find( int x )
{    if( s[x ] < 0 )
          return x;
      else
          return find( s[ x ] );
}
```



	0	1	2	3	4	5	6	7	8	9	10
S:		-1	5	-1	3	-1	3	1	1	1	5

Equivalence Class

5) **Performance Evaluation**

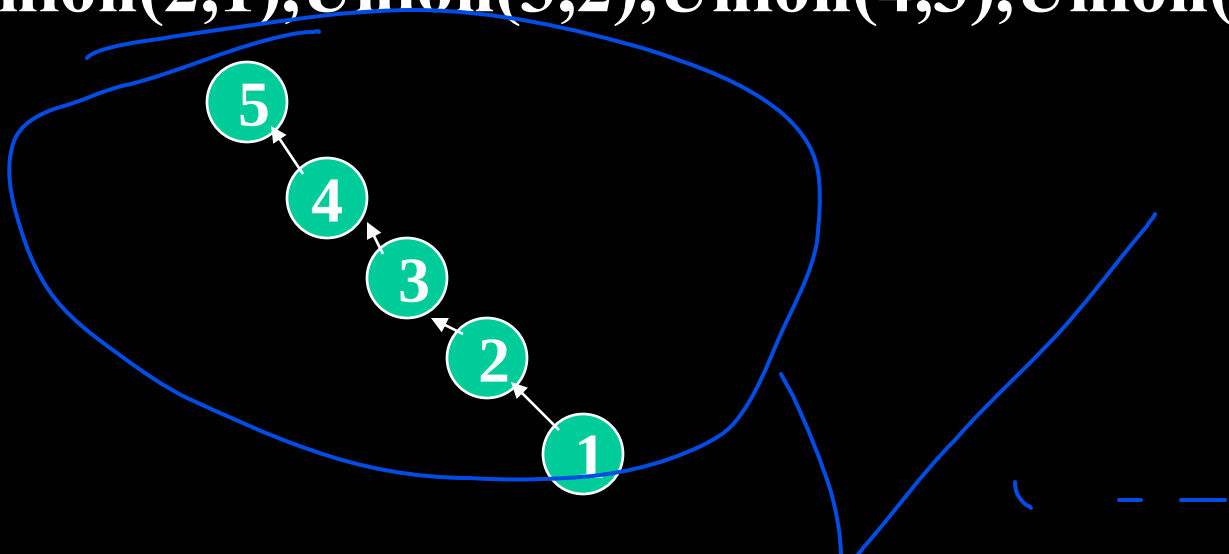
Time complexity: Find-- O(h),

Union-- θ(1)

Assume that u times unions and f times finds are to be performed, f>u,

in the worst case a tree with m elements can have a height of m:

Union(2,1),Union(3,2),Union(4,3),Union(5,4)…



Equivalence Class

★**improve Union**

合并的时候都是把矮的挂在高的上

使结果尽量矮

two rules: 节点少的做子树

- **Weight rule: if the number of nodes in tree i is less than the number in tree j, then make j the parent of i; otherwise,make i the parent of j.**

高度小的做子树.

- **Height rule: if the height of tree i is less than that of tree j, then make j the parent of i; otherwise,make i the parent of j.**

拆开,

3、下方所有结点,挂到根结点下

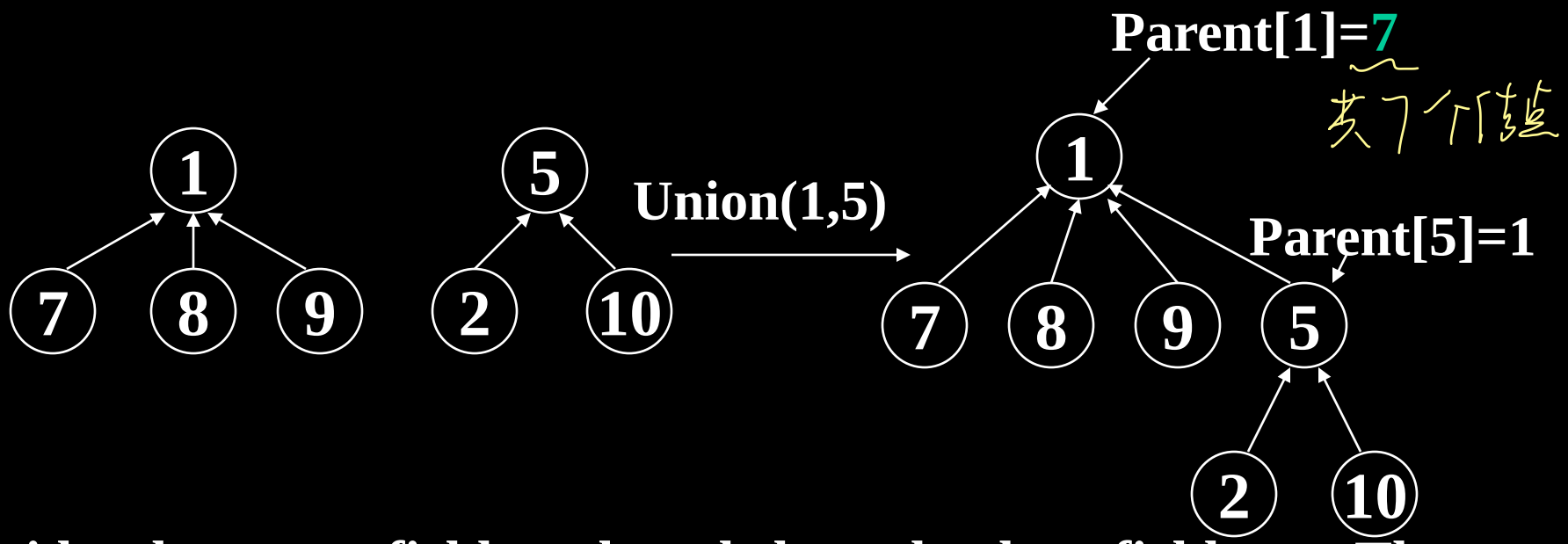
Equivalence Class

6) **Performance Enhancement**

- ✶**improve Union in order to decrease the time each find take, so that the height of tree will not increase linearly .**
- ✸**.Improvement of *Find* –path compression**

Equivalence Class

Let's discuss the weight rule(C++) :



Besides the *parent* field, each node has a boolean field *root* .The *root* field is true iff the node is presently a root node.The *parent* field of each root node is used to keep a count of the total number of nodes in the tree.

Equivalence Class

Union with the weight rule

```
  void Initialize(int n)
  { root=new bool[n+1];
    parent=new int[n+1];
    for(int e=1;e<=n;e++)
    { parent[e]=1;
      root[e]=true;
    }
  }
  int Find(int e)
  { while(!root[e])
        e=parent[e];
    return e;
  }
```

Equivalence Class

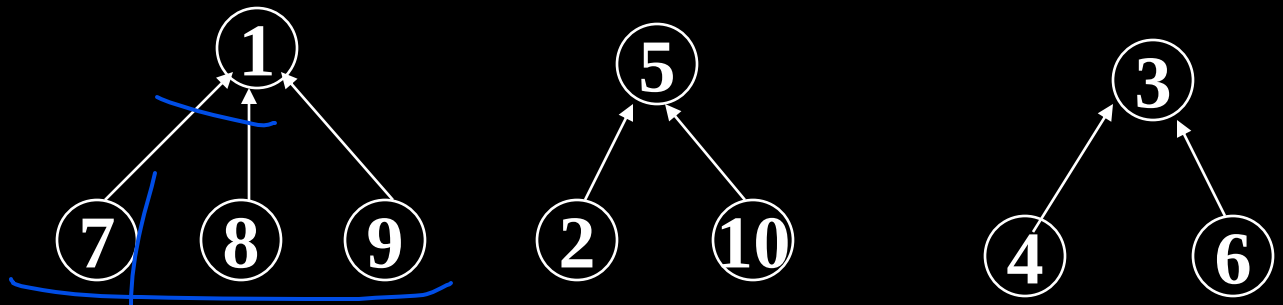
```
void Union(int i, int j)
{  if(parent[i]<parent[j])  //  i becomes subtree of j
       { parent[j]=parent[j]+parent[i];
         root[i]=false;
         parent[i]=j;
        }
    else { parent[i]=parent[i]+parent[j];
             root[j]=false;
             parent[j]=i;
           }
}
```

Equivalence Class

Java(高度规则）

用一个数组来实现，根结点中放负数，而且是代表高度。

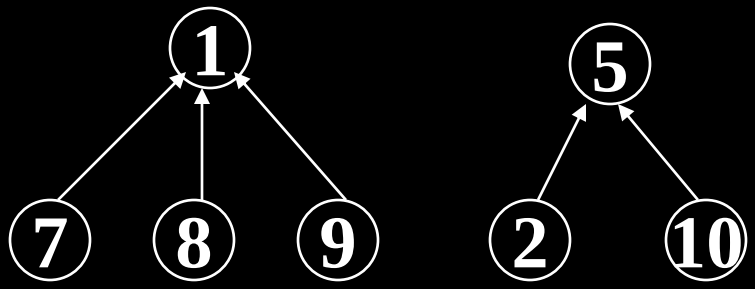
```
public void union( int root1, int root2 )
{   if(  s[ root2 ] < s[ root1 ] )
        s[ root1 ] = root2;
    else {   if( s[ root1 ] = = s[ root2 ] )
                s[ root1 ]--;
             s[ root2 ] = root1;} }
```



	0	1	2	3	4	5	6	7	8	9	10
S:		-2	5	-2	3	-2	3	1	1	1	5

h=2

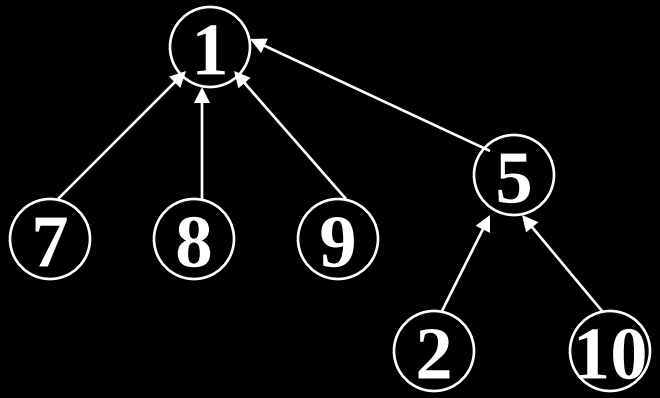
Equivalence Class



union (1, 5)

s[root1]--

s[root2] = root1



Equivalence Class

✸.Improvement of *Find* –path compression

When processing a equivalence pair, we need to operate *Find* twice, *WeightUnion* once.

Example of improvement:

Equivalence Class

```
int Find( int e) { /* C++  */ }
{  int j=e;
    while(!root[j])  j=parent[j];
    int f=e;
    while(f!=j)
      { int pf=parent[f];   parent[f]=j;  f=pf; }
}
```

Equivalence Class

Java

```
public int find( int x )
{    if( s[ x ] < 0 )
            return x;
      else
           return s[ x ] = find( s[ x ] );
}
```