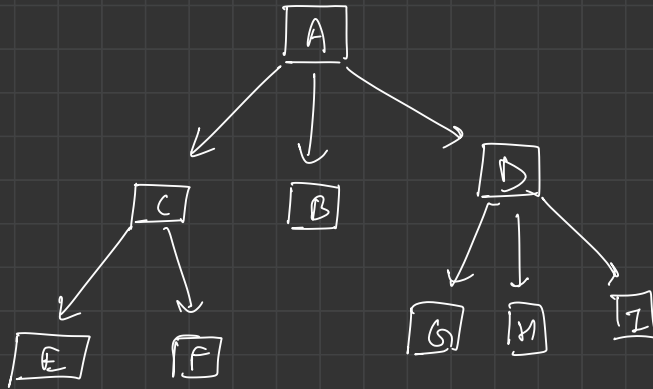


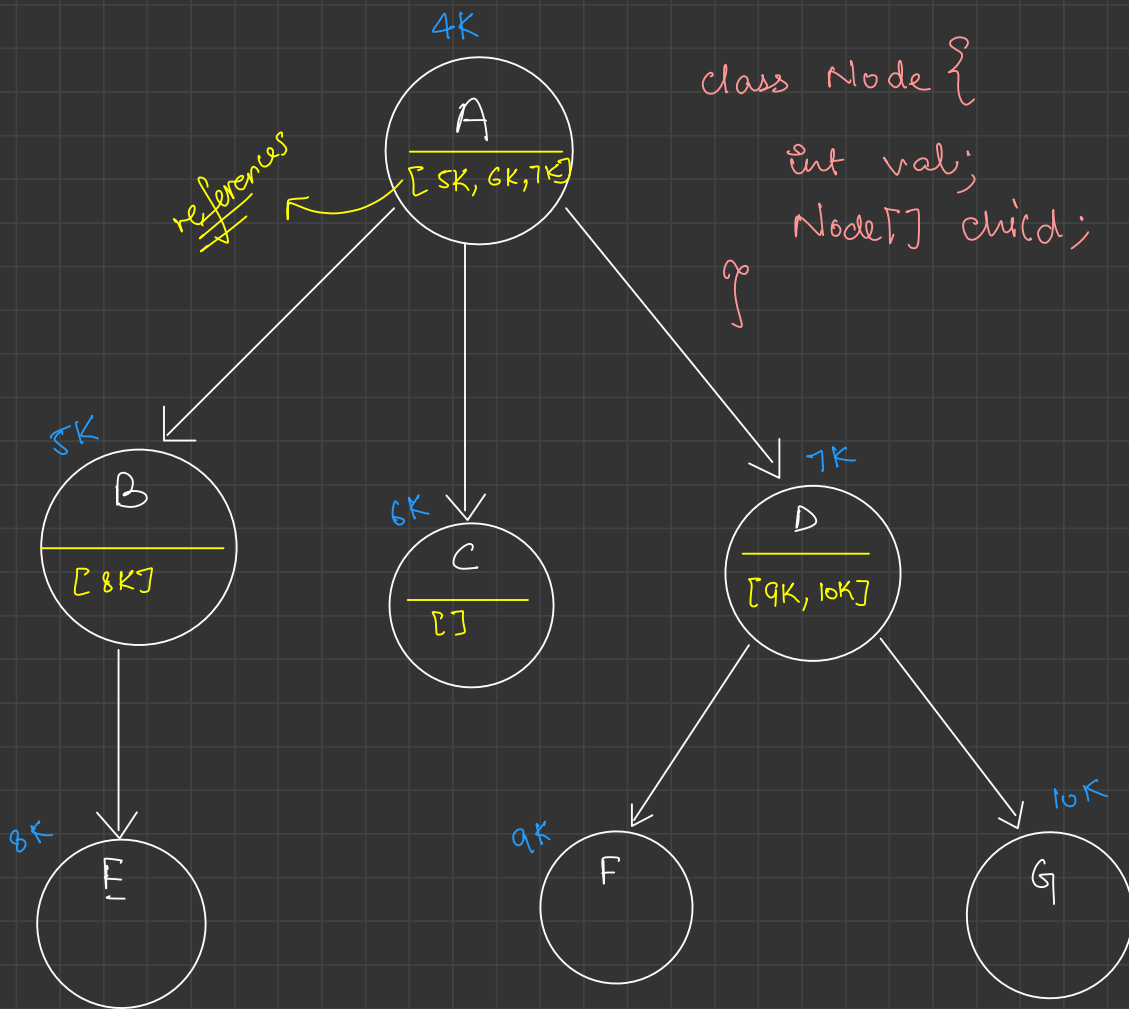


# Binary Trees

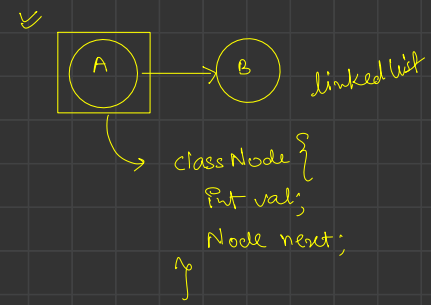
store your office employees data!

Trees!





```
class Node {  
    int val;  
    Node[] child;  
}
```



# Binary Trees

(Atmost 2 childs)

2 options

Siblings

B, C  
D, E  
F, G

left subtree

Node

degree = 2

left child

degree 2

A

Node / root node

right child

Node

Node

right subtree

degree

No. of childs of a Node

degree of leaf Node = 0

leaf

{ Node with 0 child }

D

degree = 1

E

F

G

leaf

H

degree = 0

I

leaf

leaf

for leaf Node

node.left == null  
node.right == null

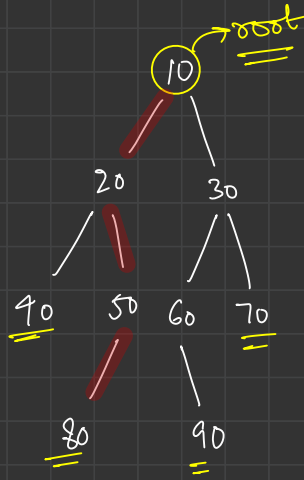
class Node {

int val;

Node left;

Node right;

}

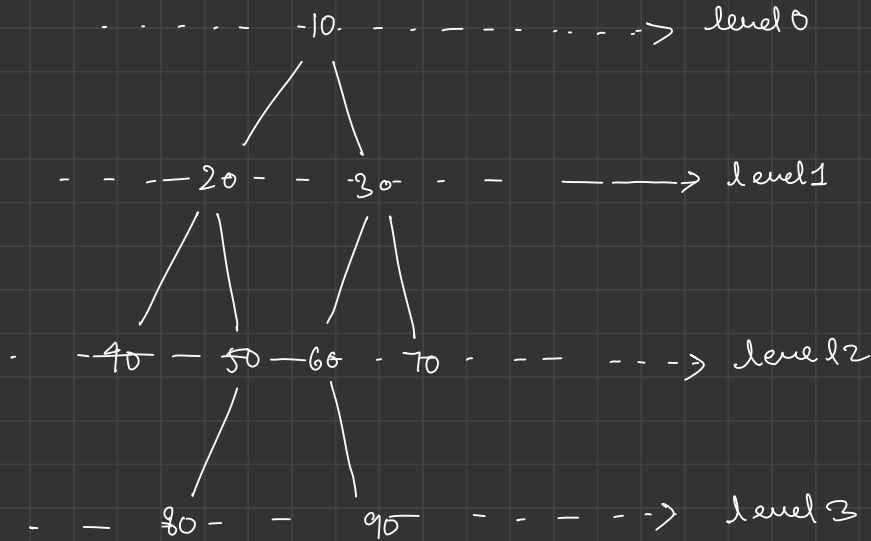


## height of a binary tree

{ dist b/w root node and deepest leaf  
Node in term of edges

height = 3

{Think of level as generation}



Perfect Binary Tree } A tree where no. of nodes at  $h^{\text{th}}$  level is  $2^h$

height = 3

--- 10 --- level 0

{ 1 Node }  $\rightarrow 2^0$

--- 20 --- 30 --- level 1

{ 2 Nodes }  $\rightarrow 2^1$

--- 40 --- 50 --- 60 --- 70 --- level 2

{ 4 Nodes }  $\rightarrow 2^2$

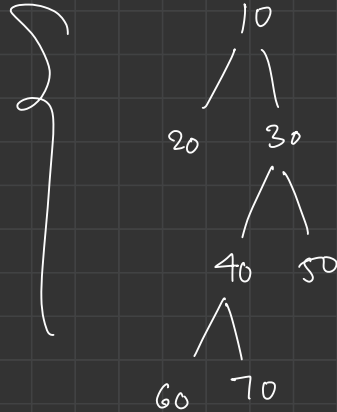
--- 80 --- 90 --- 100 --- 110 --- 120 --- 130 --- 140 --- 150 --- level 3

{ 8 Nodes }  $\rightarrow 2^3$

--- h level h  $\rightarrow$  {  $2^h$  nodes }

# Full Binary Tree

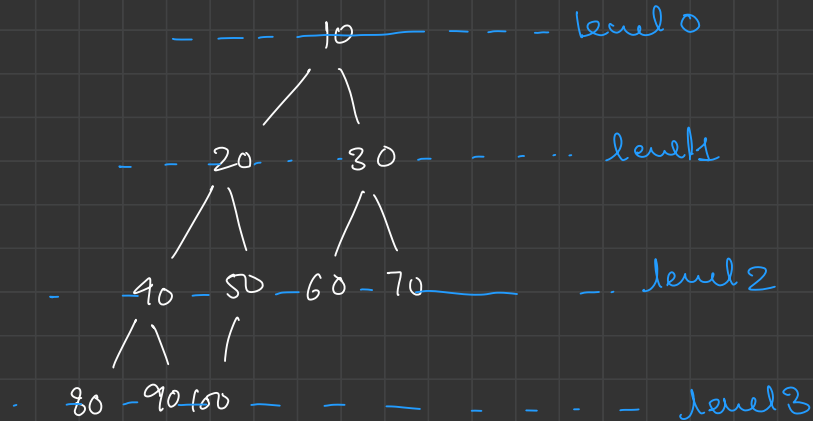
→ A BT where each Node has either zero or two children.





## Complete Binary Tree

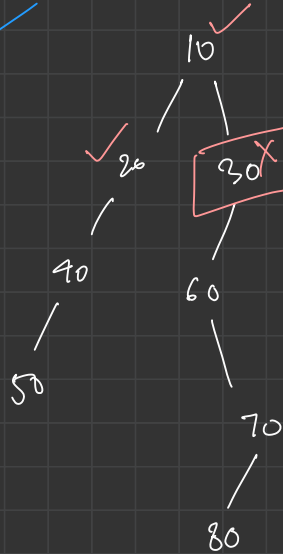
where each level is completely filled, except the last level,  
and the nodes in last level are as left as possible



# Balanced Binary Trees

$$| \text{height of left subtree} - \text{height of right subtree} | \leq 1$$

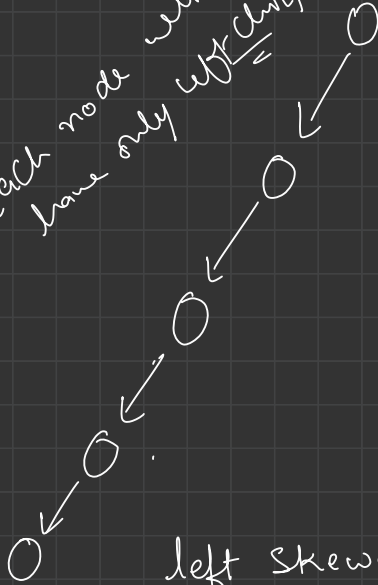
for Every  
Node



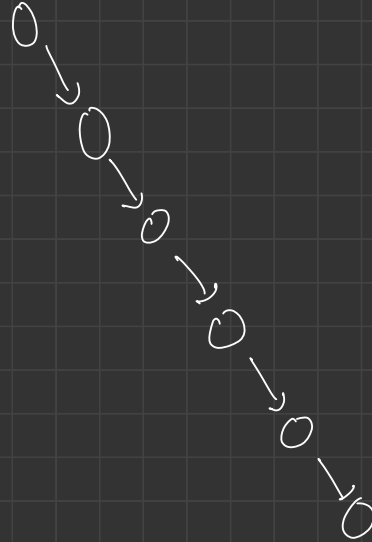
Not Balanced

## Skew tree

each node will  
have only left child



left skewed tree



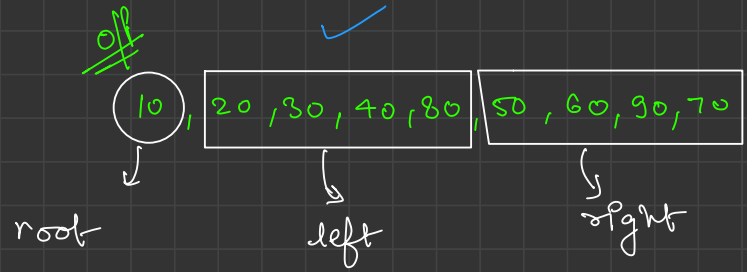
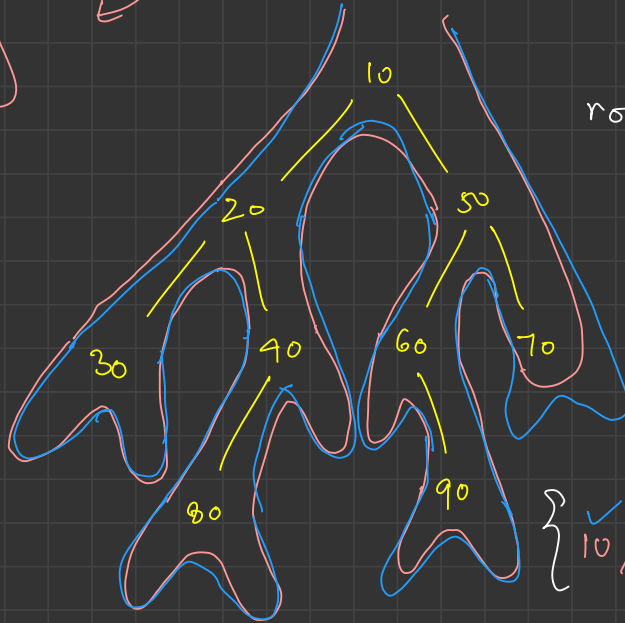
right skewed tree

# Traversal over tree

Q.

① pre order traversal

✓ print (node)  
✓ left subtree  
✓ right subtree



Euler path

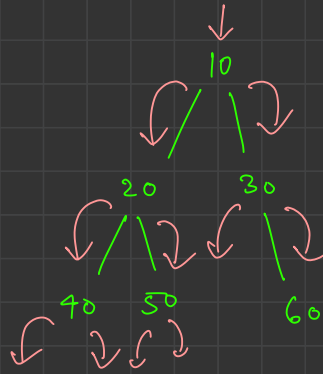
{ 10, 20, 30, 40, 80, 50, 60, 90, 70 }

```
void preOrder (TreeNode root)
{
    ① if (root == null)
        return;
```

```
    ② print (root.val);
```

```
    ③ preOrder (root.left);
```

```
    ④ preOrder (root.right);
}
```

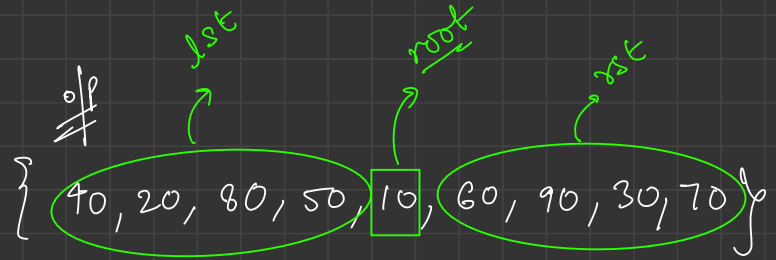
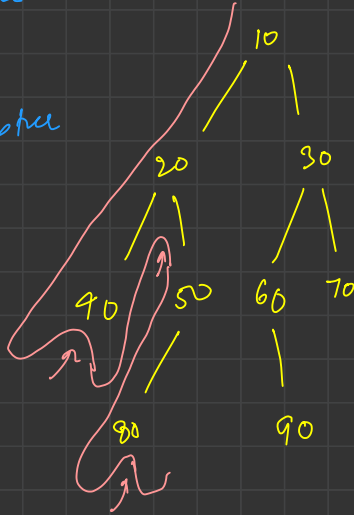


10, 20, 40, 50, 30, 60

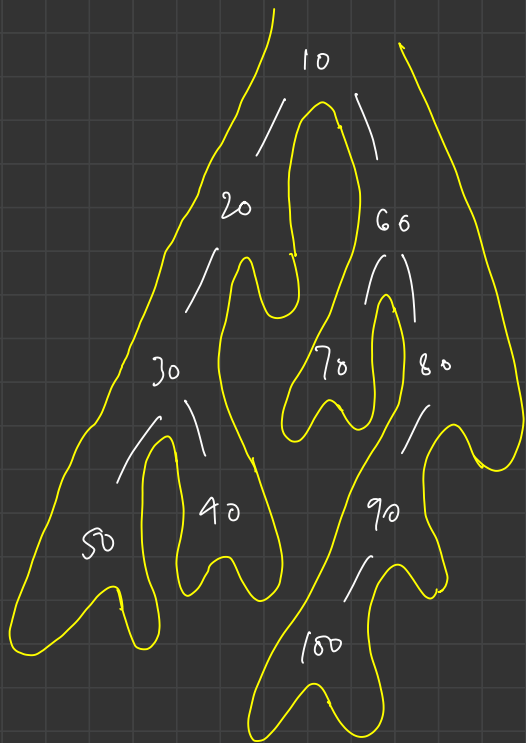
callstack

# In Order Traversal

- ↳ left subtree
- ↳ root
- ↳ right subtree



{ 40, 20, 80, 50, 10, 60, 90, 30, 70 }



Inorder

50, 30, 40, 20, 10, 70, 60, 100, 90, 80

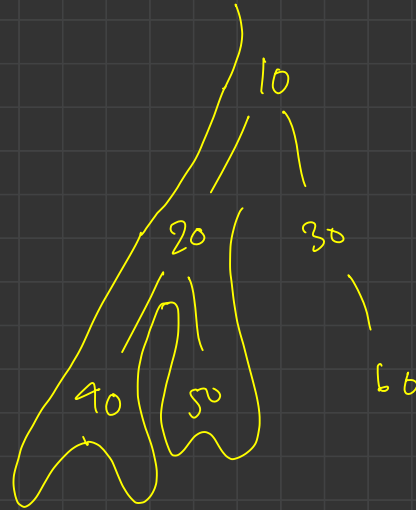
```
void Inorder (TreeNode root)
{
    if (root == null)
        return;
```

```
    Inorder (root.left);
```

```
    print (root.val);
```

```
    Inorder (root.right);
```

```
}
```

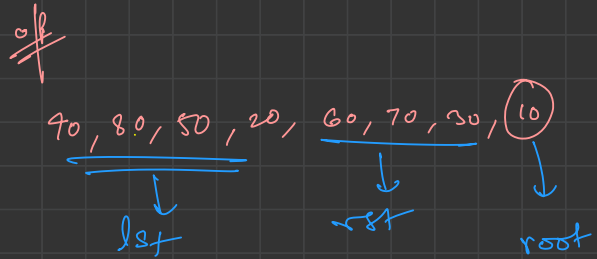
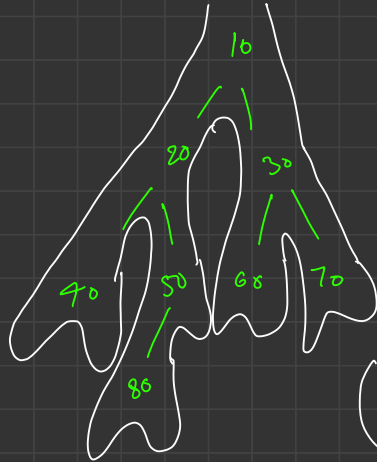


✓ ✓  
40, 20, 50, 10, 30, 60



# post Order Traversal

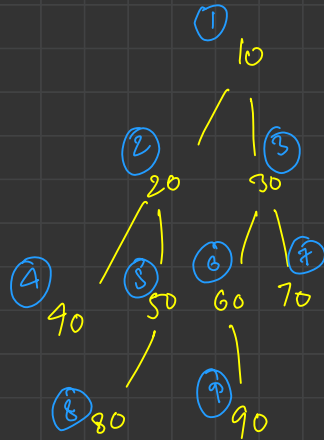
- ↳ left subtree
- ↳ right subtree
- ↳ print()



(40, 80, 50, 20, 60, 70, 30, 10)

# Size of a Binary Tree

(No. of Nodes in Binary Tree)



Size = 9

```
int size (TreeNode root)
{
    if (root == null) return 0;

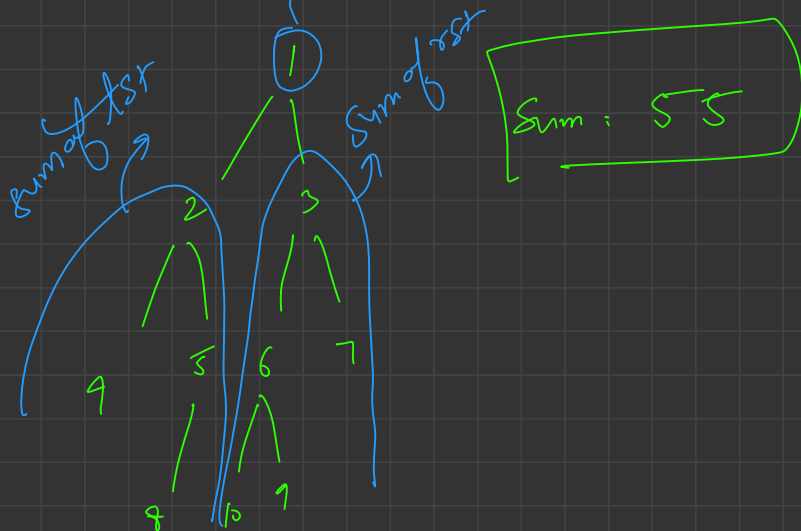
    int lft = size (root.left);
    int rht = size (root.right);

    return lft + 1 + rht;
}
```

# Sum of tree

↳ Sum of all nodes value in a tree

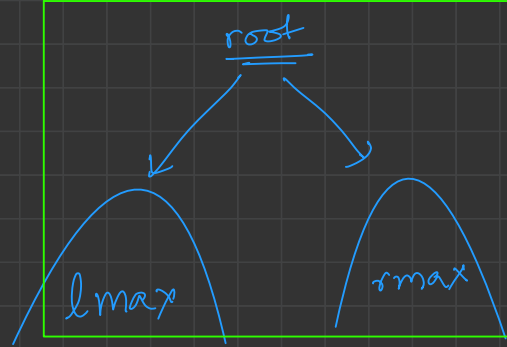
$\text{Sum of lsr} + \text{root.val} + \text{Sum of rsr}$



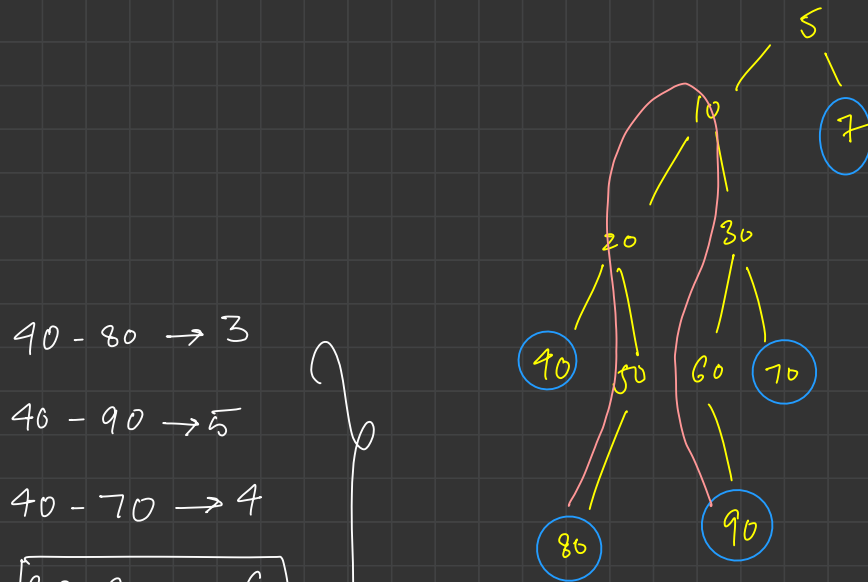
# Max of tree

→ Max<sup>m</sup> node value present in a tree

$\max(\text{lmax}, \text{root.val}, \text{rmax})$



diameter of tree { max<sup>m</sup> dist. bet<sup>w</sup> any two leaf Nodes }



40 - 80 → 3

40 - 90 → 5

40 - 70 → 4

80 - 90 → 6

80 - 70 → 5

90 - 70 → 3

diameter

