# Session 03: Recursion Applications

## What is Recursion?

**Recursion** is a process where a function **calls itself** to solve a smaller version of the same problem.

### Real-Life Analogy:

Looking in two mirrors facing each other – reflection inside reflection – similar to function calling itself again and again.

#### How Does Recursion Work?

- 1. A recursive function calls itself with a **smaller problem**.
- 2. There must be a **base case** to **stop** the recursion.
- 3. Each call is **stored in memory** using a **function call stack**.

## Application 1: Factorial (n!)

### Definition:

Factorial of n is the product of all numbers from n to 1.

```
Formula: n! = n \times (n-1) \times (n-2) \times ... \times 1
```

## Iterative Version

```
int factorialIter(int n) {
    int result = 1;
    for(int i = 1; i <= n; i++) {
        result *= i;
    }
    return result;
}</pre>
```



# 1. Factorial Using Recursion

```
#include <iostream>
using namespace std;
// Recursive function to calculate factorial
int factorial(int n) {
    if (n == 1) return 1; // Base case
    return n * factorial(n - 1); // Recursive case
}
int main() {
    int n;
    cout << "Enter a number to find factorial: ";</pre>
    cin >> n;
    if (n < 1) {
        cout << "Factorial not defined for numbers less than 1.";</pre>
    } else {
        cout << "Factorial of " << n << " is: " << factorial(n) << endl;</pre>
    }
    return 0;
}
```

## Recursive Version

```
int factorialRec(int n) {
    if (n == 0 || n == 1)
        return 1; // base case
    else
        return n * factorialRec(n - 1); // recursive call
}
```

#### Dry Run: factorialRec(4)

Call

Value Returned

```
factorialRec(4) \rightarrow 4 * factorialRec(3)
factorialRec(3) \rightarrow 3 * factorialRec(2)
factorialRec(2) \rightarrow 2 * factorialRec(1)
factorialRec(1) \rightarrow 1 (base case)
Returning: 2×1=2 \rightarrow 3×2=6 \rightarrow 4×6=24
```

#### Memory Stack Diagram

## Application 2: Fibonacci Sequence

## **Definition:**

```
F(0) = 0, F(1) = 1,

F(n) = F(n-1) + F(n-2)
```

## Recursive Version

```
int fib(int n) {
    if (n <= 1)
        return n;
    return fib(n-1) + fib(n-2);
}</pre>
```

#### Memoized Version (Optimized)

```
int fibMemo(int n, int memo[]) {
    if (n <= 1)
        return n;
    if (memo[n] != -1)
        return memo[n];

    memo[n] = fibMemo(n-1, memo) + fibMemo(n-2, memo);
    return memo[n];
}

Usage:
int main() {
    int n = 10;
    int memo[n+1];
    fill_n(memo, n+1, -1);
    cout << "Fibonacci(" << n << ") = " << fibMemo(n, memo) << endl;
}</pre>
```

### Dry Run: fib(4)

#### Problem with Recursive Fibonacci

Calls same values repeatedly

• Time complexity = O(2<sup>n</sup>) Memoization reduces to O(n)

## 2. Fibonacci Series Using Recursion

```
#include <iostream>
using namespace std;
// Recursive function to get nth Fibonacci number
int fibonacci(int n) {
    if (n == 0) return 0; // Base case
    if (n == 1) return 1; // Base case
    return fibonacci(n - 1) + fibonacci(n - 2); // Recursive call
}
int main() {
    int n;
    cout << "Enter the number of terms in Fibonacci series: ";</pre>
    cin >> n;
    cout << "Fibonacci series: ";</pre>
    for (int i = 0; i < n; i++) {
        cout << fibonacci(i) << " ";</pre>
    }
    cout << endl;</pre>
    return 0;
}
```

## Application 3: GCD Using Recursion (Euclidean Algorithm)

### **★** GCD: Greatest Common Divisor

The largest number that divides both numbers.

#### Recursive Version

```
int gcd(int a, int b) {
    if (b == 0)
        return a;
    return gcd(b, a % b);
}
```

#### Ory Run: gcd(48, 18)

```
      Call
      a
      b
      a % b

      gcd(48, 18)
      48
      18
      12

      gcd(18, 12)
      18
      12
      6

      gcd(12, 6)
      12
      6
      0

      gcd(6, 0)
      6
      0
      —
```

#### **☑** Final GCD = 6

## Memory Stack for gcd (48, 18)



# 3. GCD (Greatest Common Divisor) **Using Recursion**

```
#include <iostream>
using namespace std;
// Recursive function to calculate GCD
int gcd(int a, int b) {
                               // Base case
    if (b == 0) return a;
    return gcd(b, a % b);
                                // Recursive call
}
int main() {
    int a, b;
    cout << "Enter two numbers to find GCD: ";</pre>
    cin >> a >> b;
    cout << "GCD of " << a << " and " << b << " is: " << gcd(a, b) << endl;</pre>
    return 0;
}
```

## Practice Worksheet

- A. Factorial
  - 1. Find factorialRec(5)  $\rightarrow$
  - 2. Fill stack diagram for factorialRec(3)
- B. Fibonacci
  - 1. Find fib(5) (no memo)  $\rightarrow$  \_\_\_\_\_

2. What's the time complexity of basic recursive Fibonacci?

#### C. GCD

- 1.  $gcd(36, 24) \rightarrow$ \_\_\_\_\_
- 2. Trace the recursive calls for gcd(30, 12)

## 4. Tower of Hanoi Using Recursion

```
#include <iostream>
using namespace std;
// Recursive function to solve Tower of Hanoi
void towerOfHanoi(int n, char source, char helper, char destination) {
    if (n == 1) {
         cout << "Move disk 1 from " << source << " to " << destination <<
endl:
        return;
    }
     towerOfHanoi(n - 1, source, destination, helper); // Move n-1 disks to
helper
    cout << "Move disk " << n << " from " << source << " to " << destination
     towerOfHanoi(n - 1, helper, source, destination); // Move n-1 disks to
destination
}
int main() {
    int n:
    cout << "Enter number of disks: ";</pre>
    cin >> n;
    cout << "Steps to solve Tower of Hanoi:" << endl;</pre>
       towerOfHanoi(n, 'A', 'B', 'C'); // A = source, B = helper, C =
destination
```

```
return 0;
}
```



## Quiz: Recursion Applications



## Section A: True or False

- Q1. A recursive function must always have a base case.
- **Q2.** Recursive Fibonacci is more efficient than iterative.
- **Q3.** The GCD of two numbers can be found using recursion.
- Q4. Recursion uses a call stack to track function calls.

## Section B: Conceptual Questions

-	
	What will happen if a recursive function does not reach its base case?
Q7.	Why is memoization used in Fibonacci?

**Q8.** Write the recursive formula for the Fibonacci sequence.

## Section C: Output Prediction

**Q9.** What will be the output of this code?

```
int factorial(int n) {
    if (n <= 1)
       return 1;
    return n * factorial(n - 1);
```

```
}
cout << factorial(4);

→ _____
Q10. What is the output of gcd(48, 18) using recursion?
 → _____
Q11. How many recursive calls will be made in factorial(3)?</pre>
```

## Section D: Dry Run Stack Trace

Fill in the blanks to trace the recursive stack.

**Q12.** Dry run for factorial(3)

#### **Function Call Returned Value**

factorial(3)  $\rightarrow$  factorial(2)  $\rightarrow$  factorial(1)  $\rightarrow$  Return values:

**Q13.** Dry run for gcd(30, 12)

#### **Function Call Returned Value**

 $\gcd(30, 12) \rightarrow$  $\gcd(12, 6) \rightarrow$  $\gcd(6, 0) \rightarrow$ 

Return values:

## **☑** Section E: Code Completion

**Q14.** Complete the recursive GCD function:

```
int gcd(int a, int b) {
    if (_____)
        return a;
    return ____;
}
```

# Section F: Multiple Choice

**Q15.** Which of the following has time complexity  $O(2^n)$ ?

- a) Iterative factorial
- b) Recursive factorial
- c) Basic recursive Fibonacci
- d) Memoized Fibonacci

Q16. What is the base case for Fibonacci?

```
a) F(0) = 1
```

b) 
$$F(0) = 0, F(1) = 1$$

c) 
$$F(n) = 0$$

d) None of the above

## Answer Key (Teacher Only)

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```
A1. True
```

A2. False

A3. True

A4. True

**Q5.** 
$$n == 0$$
 or  $n == 1$ 

**Q6.** Infinite recursion → stack overflow

Q7. To avoid recomputing same values → improves speed

**Q8.** 
$$F(n) = F(n-1) + F(n-2)$$

**Q9**. 24

**Q10**. 6

**Q11.** 3 calls

Q12.

**Function Call Returned Value** 

```
factorial(3) \rightarrow 6
factorial(2) \rightarrow 2
factorial(1) \rightarrow 1

Q13.

Function Call Returned Value \gcd(30, 12) \rightarrow 6
\gcd(12, 6) \rightarrow 6
\gcd(6, 0) \rightarrow 6

Q14.
```

cpp CopyEdit

if (b == 0)

return gcd(b, a % b);

**Q15.** c) Basic recursive Fibonacci **Q16.** b) 
$$F(\emptyset) = \emptyset$$
,  $F(1) = 1$