

# Summary of COMP523 Advanced Algorithm

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# Chapter 1

## Symmetry Notation

### 1.1 Asymptotic Notation

Asymptotic notation is a way of describing the limiting behavior of a function when the argument tends towards a particular value or infinity. In computer science, asymptotic notation is frequently used to describe the running time or space usage of an algorithm.

- $O$ -notation:  $f(n) = O(g(n))$  if there exist constants  $c$  and  $n_0$  such that  $0 \leq f(n) \leq cg(n)$  for all  $n \geq n_0$ .
- $\Omega$ -notation:  $f(n) = \Omega(g(n))$  if there exist constants  $c$  and  $n_0$  such that  $0 \leq cg(n) \leq f(n)$  for all  $n \geq n_0$ .
- $\Theta$ -notation:  $f(n) = \Theta(g(n))$  if there exist constants  $c_1, c_2$  and  $n_0$  such that  $0 \leq c_1g(n) \leq f(n) \leq c_2g(n)$  for all  $n \geq n_0$ .
- $o$ -notation:  $f(n) = o(g(n))$  if for any constant  $c > 0$ , there exists a constant  $n_0$  such that  $0 \leq f(n) < cg(n)$  for all  $n \geq n_0$ .
- $\omega$ -notation:  $f(n) = \omega(g(n))$  if for any constant  $c > 0$ , there exists a constant  $n_0$  such that  $0 \leq cg(n) < f(n)$  for all  $n \geq n_0$ .

### 1.2 Comparing Functions

#### 1.2.1 Transitivity

- $f(n) = O(g(n))$  and  $g(n) = O(h(n))$  implies  $f(n) = O(h(n))$ .
- $f(n) = \Omega(g(n))$  and  $g(n) = \Omega(h(n))$  implies  $f(n) = \Omega(h(n))$ .
- $f(n) = \Theta(g(n))$  and  $g(n) = \Theta(h(n))$  implies  $f(n) = \Theta(h(n))$ .

For example,  $n^2 = O(n^3)$  and  $n^3 = O(n^4)$  implies  $n^2 = O(n^4)$ .

#### 1.2.2 Reflexivity

- $f(n) = O(f(n))$ .
- $f(n) = \Omega(f(n))$ .
- $f(n) = \Theta(f(n))$ .

For example,  $n^2 = O(n^2)$ .

### 1.2.3 Symmetry

- $f(n) = O(g(n))$  implies  $g(n) = O(f(n))$ .
- $f(n) = \Omega(g(n))$  implies  $g(n) = \Omega(f(n))$ .
- $f(n) = \Theta(g(n))$  implies  $g(n) = \Theta(f(n))$ .
- $f(n) = o(g(n))$  implies  $g(n) = \omega(f(n))$ .
- $f(n) = \omega(g(n))$  implies  $g(n) = o(f(n))$ .

For example,  $n^2 = O(n^3)$  implies  $n^3 = \Omega(n^2)$ .

### 1.2.4 Transpose Symmetry

- $f(n) = O(g(n))$  if and only if  $g(n) = \Omega(f(n))$ .
- $f(n) = \Theta(g(n))$  if and only if  $g(n) = \Theta(f(n))$ .
- $f(n) = o(g(n))$  if and only if  $g(n) = \omega(f(n))$ .
- $f(n) = \omega(g(n))$  if and only if  $g(n) = o(f(n))$ .

For example,  $n^2 = O(n^3)$  if and only if  $n^3 = \Omega(n^2)$ .

### 1.2.5 sum and maximum

$$f_1(n) + f_2(n) + \cdots + f_k(n) = \Theta(\max(f_1(n), f_2(n), \dots, f_k(n)))$$

where  $k$  is a constant positive integer.

Let  $f_j(n) = j$ ,  $k = n$ , then

$$f_1(n) + f_2(n) + \cdots + f_k(n) = n(n+1)/2 = \Theta(n^2)$$

### 1.2.6 Running time hierarchy

- logarithmic:  $O(\log n)$
- linear:  $O(n)$
- $n \log n$ :  $O(n \log n)$
- quadratic:  $O(n^2)$
- polynomial:  $O(n^k)$
- exponential:  $O(c^n)$
- constant:  $O(1)$
- superconstant:  $\omega(1)$
- sublinear:  $o(n)$
- superlinear:  $\omega(n)$
- superpolynomial:  $\omega(n^k)$
- subexponential:  $o(c^n)$

## 1.3 Expect of algorithms

**Correctness:** An algorithm is correct if it halts with the correct output for every input instance.

**Termination:** An algorithm is terminating if it halts for every input instance.

**Efficiency:** An algorithm is efficient if it halts with the correct output for every input instance and runs in polynomial time.

## Chapter 2

# Recursion and Divide and Conquer techniques

### 2.1 Finding Majority in array

The pseudocode of the algorithm is shown in Algorithm 2.1.

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**Algorithm 1** Finding Majority in array

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```
1: procedure MAJORITY( $A$ )
2:    $n \leftarrow$  length of  $A$ 
3:   if  $n = 0$  then
4:     return  $-1$ 
5:   end if
6:   if  $n = 1$  then
7:     return  $A[1]$ 
8:   end if
9:   if  $n \neq 1$  and  $n$  is odd then
10:
11:   end if
12:   Array  $B$  of size  $n/2$ 
13:   set  $j=0$ 
14:   for  $i = 1$  to  $n/2$  do
15:     if  $A[2i - 1] = A[2i]$  then
16:        $B[j] \leftarrow A[2i - 1]$ 
17:        $j \leftarrow j + 1$ 
18:     end if
19:   end for
20:    $m \leftarrow$  MAJORITY( $B$ )
21:    $count \leftarrow 0$ 
22:   for  $i = 1$  to  $n$  do
23:     if  $A[i] = m$  then
24:        $count \leftarrow count + 1$ 
25:     end if
26:   end for
27:   if  $count > n/2$  then
28:     return  $m$ 
29:   else
30:     return  $-1$ 
31:   end if
32: end procedure
```

---

**Correctness:**

Lemma: If  $A$  has a majority element, then the majority element of  $A$  is also the majority element of  $B$ .

Base case:  $n = 1$ , the majority element is  $A[1]$ .

Induction hypothesis: Assume that the lemma is true for  $n = k$ , we will prove that the lemma is true for  $n = k + 1$ .

Induction step: If  $A$  has a majority element, then the majority element of  $A$  is also the majority element of  $B$ .

Case 1 ( $A$  has a majority element  $m$ ): Then by the lemma, it is also the majority element of  $B$ . Then  $m$  appears more than  $k/2$  times in  $B$ . Then  $m$  appears more than  $(k + 1)/2$  times in  $A$ .

Case 2 ( $A$  has no majority element): Then  $B$  has no majority element. Then  $A$  has no majority element.

**Proof the lemma:**

proof by contradiction. Assume that  $A$  has a majority element  $m$  and  $B$  has a majority element  $m'$ , but  $m \neq m'$ .

Let  $x$  be the numbers of occurrence of  $m$  in  $A$ .

Let  $y$  be the numbers of occurrence of  $m'$  in  $B$ .

Then  $2y$  times from pairs that are represented in  $B$  by a value different from  $m'$ , and  $x - 2y$  times, since each occurrence of  $m$  in  $A$  that is not paired with another occurrence of  $m$  in  $A$  is paired with an occurrence of  $m'$  in  $B$ .

In total, this gives  $2y + x - 2y = x$  occurrences of  $m$  in  $A$ , which is a contradiction.

**Running time:**

Recursive formula for the running time:

$$T(n) \leq T(n/2) + cn$$

where  $c$  is a constant.

The solution to the recurrence is  $T(n) = O(n)$ .

## 2.2 Searching in logarithmic time

Searching faster with BinarySearch.

It is a particular case of the divide-and-conquer paradigm.

**Input:** A sorted array  $A$  of  $n$  elements and a value  $x$ .

**Output:** An index  $i$  such that  $A[i] = x$  or the special value  $-1$  if  $x$  does not appear in  $A$ .

**Pseudocode** is shown in Algorithm 2.2.

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**Algorithm 2** BinarySearch
 

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```

1: procedure BINARYSEARCH( $x, i, j$ )
2:   if  $i = j$  then
3:     if  $A[i] = x$  then
4:       return  $i$ 
5:     else
6:       return  $-1$ 
7:     end if
8:   else
9:     if  $x = A[\lfloor (i + j)/2 \rfloor]$  then
10:      return  $\lfloor (i + j)/2 \rfloor$ 
11:    else if  $x < A[\lfloor (i + j)/2 \rfloor]$  then
12:      return BINARYSEARCH( $x, i, \lfloor (i + j)/2 \rfloor$ )
13:    else
14:      return BINARYSEARCH( $x, \lfloor (i + j)/2 \rfloor + 1, j$ )
15:    end if
16:  end if
17: end procedure

```

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**Running time:**

The number of comparisons performed by BinarySearch is:

$$T(n) \leq T(n/2) + 4$$

Keep calculate:

$$\begin{aligned}
T(n) &\leq T(n/2) + 4 \\
&\leq T(n/4) + 4 + 4 \\
&\leq T(n/8) + 4 + 4 + 4 \\
&\leq T(n/2^k) + 4k \\
&\leq T(n/2^{\log(n-1)}) + 4\log(n-1) \\
&= T(2) + 4(\log n - 1) \\
&\leq 4\log n - 4 \\
&= 4\log n
\end{aligned}$$

proof  $T(n) \leq 4\log n$ :

Base case:  $n = 1, T(1) = 0 \leq 4\log 1 = 0$ .

Induction hypothesis: Assume that the lemma is true for  $n = k$ , we will prove that the lemma is true for  $n = k + 1$ .

Induction step:  $T(k + 1) \leq 4\log(k + 1)$ .

$$\begin{aligned}
T(k + 1) &\leq T(k/2) + 4 \\
&\leq 4\log(k/2) + 4 \\
&= 4\log k - 4 + 4 \\
&= 4\log k \\
&\leq 4\log(k + 1)
\end{aligned}$$

**Memory usage:**

The memory usage of BinarySearch is:

$$M(n) = O(\log n)$$

**Comparing BinarySearch and LinearSearch:**

$$\begin{aligned}
T_{\text{BinarySearch}}(n) &= O(\log n) \\
T_{\text{LinearSearch}}(n) &= O(n) \\
T_{\text{BinarySearch}}(n) &= O(\log n) < O(n) = T_{\text{LinearSearch}}(n) \\
M_{\text{BinarySearch}}(n) &= O(\log n) < O(1) = M_{\text{LinearSearch}}(n)
\end{aligned}$$

## 2.3 Running time of Divide and Conquer algorithms

The Master Theorem:

Suppose that  $T(n)$  satisfies the recurrence:

$$T(n) \leq aT(n/b) + cn^d$$

where  $a \geq 1, b > 1, c > 0$  and  $d \geq 0$  are constants.

Then  $T(n)$  has the following asymptotic bounds:

$$T(n) = \begin{cases} O(n^d) & \text{if } d > \log_b a \\ O(n^d \log n) & \text{if } d = \log_b a \\ O(n^{\log_b a}) & \text{if } d < \log_b a \end{cases}$$

This theorem is useful for solving recurrences of the form:

$$T(n) = aT(n/b) + f(n)$$

where  $a \geq 1$ ,  $b > 1$  and  $f(n)$  is an asymptotically positive function.

**Example:**

$$T(n) = 8T(n/2) + 100n^2$$

$a = 8$ ,  $b = 2$ ,  $f(n) = 100n^2$ ,  $d = 2$ ,  $\log_b a = \log_2 8 = 3$ .

$d = 2 < \log_b a = 3$ , so  $T(n) = O(n^{\log_b a}) = O(n^3)$ .

## 2.4 Finding pair of points closest to each other

**Input:** A set  $P$  of  $n$  points in the plane.

**Output:** The pair of points in  $P$  that are closest to each other.

**Pseudocode** is shown in Algorithm 2.4. **Running time:**

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### Algorithm 3 ClosestPair

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```

1: procedure CLOSESTPAIR( $P_1, \dots, P_n$ )
2:   Construct  $P_x$  and  $P_y$ .  $P_x$  is sorted by  $x$ -coordinate,  $P_y$  is sorted by  $y$ -coordinate.
3:   return CLOSESTPAIRREC( $P_x, P_y$ )
4: end procedure

```

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### Algorithm 4 ClosestPairRec

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```

1: procedure CLOSESTPAIRREC( $P_x, P_y$ )
2:   if  $|P_x| = |P_y| \leq 3$  then
3:     For each pair of points  $(P_i, P_j)$ , compute  $d(P_i, P_j)$ 
4:     return the pair of points with the smallest distance
5:   end if
6:   Construct  $Q_x, Q_y, R_x$  and  $R_y$ .
7:    $(l_1, l_2) = \text{CLOSESTPAIRREC}(Q_x, Q_y)$ 
8:    $(r_1, r_2) = \text{CLOSESTPAIRREC}(R_x, R_y)$ 
9:    $\delta = \min\{d(l_1, l_2), d(r_1, r_2)\}$ 
10:   $x^* =$  the largest  $x$ -coordinate in  $Q_x$ 
11:   $L = \{(x, y) : x = x^*\}$ 
12:   $S = \{p \in P : p \in L \text{ and } p \text{ is within } \delta \text{ of } L\}$ 
13:  Construct  $S_v$ 
14:  for  $p \in S$  do
15:    Let  $q$  be the point in  $S_v$  closest to  $p$ 
16:    if  $d(p, q) < \delta$  then
17:       $\delta = d(p, q)$ 
18:       $(s_1, s_2) = (p, q)$ 
19:    end if
20:  end for
21:  if  $d(s_1, s_2) < \min\{d(l_1, l_2), d(r_1, r_2)\}$  then
22:    return  $(s_1, s_2)$ 
23:  end if
24:  if  $d(l_1, l_2) < d(r_1, r_2)$  then
25:    return  $(l_1, l_2)$ 
26:  else
27:    return  $(r_1, r_2)$ 
28:  end if
29: end procedure

```

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$$T(n) \leq 2T(n/2) + O(n \log n) = O(n \log n)$$

**Example:**



## Chapter 3

# Graph Algorithms

### 3.1 Graph Definitions

**Graph:** A graph  $G$  consists of a set  $V$  of vertices and a set  $E$  of edges, where each edge is associated with a pair of vertices.

**Directed Graph:** A directed graph  $G$  consists of a set  $V$  of vertices and a set  $E$  of directed edges, where each directed edge is associated with an ordered pair of vertices.

**Undirected Graph:** An undirected graph  $G$  consists of a set  $V$  of vertices and a set  $E$  of undirected edges, where each undirected edge is associated with an unordered pair of vertices.

**Neighbours of a vertex  $v$ :** Set of vertices that are connected to  $v$  by an edge.

**Degree of a vertex  $v$ :** number of neighbours of  $v$ , denoted by  $deg(v)$ .

**Path:** A sequence of (non-repeating) nodes with consecutive nodes being connected by an edge.  
length = node count - 1 = edge count.

**Distance between two nodes:** The number of edges in the shortest path between the two nodes.

**Graph diameter:** The maximum distance between any two nodes in the graph.

**Lines, cycles, trees and cliques:**

**Line:** A graph with  $n$  vertices and  $n - 1$  edges.

**Cycle:** A graph with  $n$  vertices and  $n$  edges.

**cliques:** A graph with  $n$  vertices and  $n(n - 1)/2$  edges.

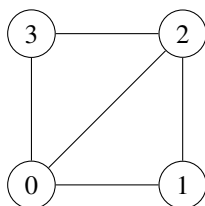
**Tree:** A graph with  $n$  vertices and  $n - 1$  edges.

**Graph representations:**

**Adjacency matrix:** A  $n \times n$  matrix  $A$  where  $A_{ij} = 1$  if there is an edge between  $i$  and  $j$ , and  $A_{ij} = 0$  otherwise.

examples of adjacency matrices:

Given the following graph:



The adjacency matrix is:

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix}$$

**Adjacency matrix for directed graphs:** A  $n \times n$  matrix  $A$  where  $A_{ij} = 1$  if there is an edge from  $i$  to  $j$ , and  $A_{ij} = 0$  otherwise.

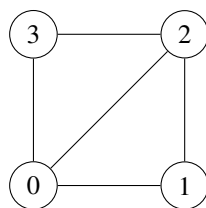
examples of adjacency matrices for directed graphs:  
 Given the following graph:



The adjacency matrix is:

$$\begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

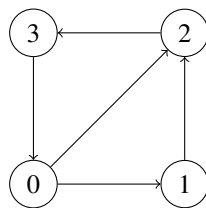
**Adjacency list:** A list of lists, where the  $i$ th list contains the neighbours of vertex  $i$ .  
 Given the following graph:



The adjacency list is:

$$\begin{bmatrix} 1 & 2 & 3 \\ 0 & 2 & \\ 0 & 1 & 3 \\ 0 & 2 & \end{bmatrix}$$

**Adjacency list for directed graphs:** A list of lists, where the  $i$ th list contains the neighbours of vertex  $i$ .  
 Given the following graph:



The adjacency list is:

$$\begin{bmatrix} 1 & 2 \\ 2 & \\ 3 & \\ 0 & \end{bmatrix}$$

**Adjacency matrix vs adjacency list:**

Adjacency matrix	Adjacency list
$O(1)$ to check if there is an edge between $i$ and $j$	$O(\min(\deg(i), \deg(j)))$ to check if there is an edge between $i$ and $j$
$O(n)$ to find the neighbours of $i$	$O(\deg(i))$ to find the neighbours of $i$
$O(n^2)$ space	$O(n + m)$ space

## 3.2 Depth-first search

**Depth-first search:** A graph search algorithm that explores the neighbours of a vertex before exploring the neighbours of its neighbours.

example of depth-first search:



The depth-first search sequence is:

0, 1, 2, 3, 5, 4

**Depth-first search algorithm:**

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**Algorithm 5** Depth-first search algorithm

---

```
1: procedure DFS( $G, v$ )
2:   for  $e \in V$  do
3:     if  $e$  is unexplored then
4:        $u = \text{head of } e$ 
5:       if  $u$  is unexplored then
6:          $e$  is a tree edge
7:         DFS( $G, u$ )
8:       else
9:          $e$  is a back edge
10:      end if
11:    end if
12:  end for
13: end procedure
```

---

**Running time of depth-first search:**  $O(n + m)$

## 3.3 Breadth-first search

**Breadth-first search:** A graph search algorithm that explores the neighbours of a vertex before exploring the neighbours of its neighbours.

example of breadth-first search:



The breadth-first search sequence starting from vertex 0 is 0, 1, 2, 3, 4, 5.

**Breadth-first search algorithm:**

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**Algorithm 6** Breadth-first search algorithm

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```
1: procedure BFS( $G, s$ )
2:   initial empty list  $L$ 
3:    $L \leftarrow s$ 
4:    $i \leftarrow 0$ 
5:   while  $L[i] \neq \emptyset$  do
6:      $L_{i+1} \leftarrow \text{emptylist}$ 
7:     for  $v \in L[i]$  do
8:       for edges  $(e)$  incident to  $v$  do
9:         if  $e$  is unexplored then
10:            $w \leftarrow$  the other end of  $e$ 
11:           if  $w$  is unexplored then
12:             label  $e$  as a tree edge
13:             add  $w$  to  $L_{i+1}$ 
14:           else
15:             label  $e$  as a cross edge
16:           end if
17:         end if
18:       end for
19:     end for
20:      $i \leftarrow i + 1$ 
21:   end while
22: end procedure
```

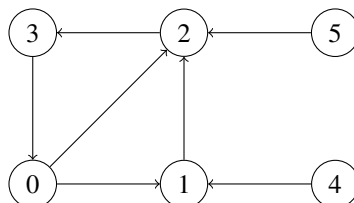
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**Running time of breadth-first search:**  $O(n + m)$

### 3.4 Strong Connectivity

**Directed graph:** A graph where the edges have a direction.

Examples:



#### DFS and BFS on directed graphs:

Very similar to undirected graphs, except that we only consider edges that go out of a vertex.

Running time is  $O(n + m)$

For example graph above the DFS sequence is 0, 1, 2, 3.

The BFS sequence is 0, 1, 2, 3.

#### 3.4.1 Connectivity

**Weak connectivity:** If we ignore the direction for all edges, there would be a path from any vertex to any other vertex.

**Strong Connectivity:** For every two nodes  $u$  and  $v$ , there is a path from  $u$  to  $v$  and a path from  $v$  to  $u$ .

### 3.4.2 Mutual Reachability

Two nodes  $u$  and  $v$  are mutually reachable if there is a path from  $u$  to  $v$  and a path from  $v$  to  $u$ .

**Strong connectivity:** For every pair of nodes  $u$  and  $v$ , these two nodes are mutually reachable.

**Transitivity:** If  $u$  is mutually reachable with  $v$  and  $v$  is mutually reachable with  $w$ , then  $u$  is mutually reachable with  $w$ .

### 3.4.3 Testing strong connectivity

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**Algorithm 7** Testing strong connectivity

---

```
1: procedure TESTSTRONGCONNECTIVITY( $G$ )
2:   define  $G^R$  to be the graph with the same vertices as  $G$  but with all edges reversed
3:   Select a node  $s$  in  $G$ 
4:   BFS( $G, s$ ), BFS( $G^R, s$ )
5:   for each node  $v$  do
6:     if  $v$  is unexplored in either BFS then
7:       return False
8:     end if
9:   end for
10:  return True
11: end procedure
```

---

## 3.5 Testing bipartiteness

**Bipartite graph:** A graph  $G = (V, E)$  is bipartite if and only if the vertices can be partitioned into two sets  $V_1$  and  $V_2$  such that every edge has one end in  $V_1$  and the other end in  $V_2$ .

A Graph  $G = (V, E)$  is bipartite if and only if it has no odd cycles. (odd cycle: a cycle with odd number of edges)

**Testing bipartiteness:**

Given a graph  $G = (V, E)$ , we want to test if  $G$  is bipartite.

Given a graph  $G = (V, E)$ , decide if it is 2-colourable.

Given a graph  $G = (V, E)$ , decide if it has an odd cycle.

**Colouring the nodes** It is quite familiar with BFS:

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**Algorithm 8** Colouring the nodes

---

```
1: procedure COLOURING( $G, s$ )
2:   initial empty list  $L$ 
3:   initial empty list  $C$ 
4:    $L \leftarrow s$ 
5:    $C[s] \leftarrow red$ 
6:    $i \leftarrow 0$ 
7:   while  $L[i] \neq \emptyset$  do
8:      $L_{i+1} \leftarrow emptylist$ 
9:     for  $v \in L[i]$  do
10:      for edges ( $e$ ) incident to  $v$  do
11:        if  $e$  is unexplored then
12:           $w \leftarrow$  the other end of  $e$ 
13:          if  $w$  is unexplored then
14:            label  $e$  as a tree edge
15:            add  $w$  to  $L_{i+1}$ 
16:            if  $i + 1$  is odd then
17:               $C[w] \leftarrow green$ 
18:            else
19:               $C[w] \leftarrow red$ 
20:            end if
21:          else
22:            label  $e$  as a cross edge
23:            if  $C[v] = C[w]$  then
24:              return False
25:            end if
26:          end if
27:        end if
28:      end for
29:    end for
30:     $i \leftarrow i + 1$ 
31:  end while
32:  for  $e(v, w) \in G$  do
33:    if  $C[v] = C[w]$  then
34:      return False
35:    end if
36:  end for
37:  return True
38: end procedure
```

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**Running time of colouring the nodes:**  $O(n + m)$

**Correctness of colouring the nodes:**

Proof by contradiction.

Suppose that  $G$  is not bipartite.

Then  $G$  has an odd cycle.

Suppose to the contrary that the algorithm return True.

That means that the algorithm did not detect the odd cycle.

## 3.6 DAGs and Topological Ordering

**DAG:** A directed acyclic graph (DAG) is a directed graph with no directed cycles.  
examples of DAGs:



**Topological ordering:** Given a graph  $G = (V, E)$ , a topological ordering of  $G$  is an ordering of the nodes  $u_1, u_2, \dots, u_n$  such that for every edge  $(u_i, u_j)$ , we have  $i < j$ .

Intuitively, a topological ordering is an ordering of the nodes such that every edge goes from left to right.

example of topological ordering based on given graph above:

3, 0, 1, 2, 4, 5

**Topological ordering implies DAG:**

- If  $G$  has a topological ordering, then  $G$  is a DAG.
- Suppose by contradiction that  $G$  has a topological ordering  $u_1, u_2, \dots, u_n$  but  $G$  also has a cycle  $C$ .
- Let  $u_j$  be the smallest element of  $C$  in the topological ordering.
- Let  $u_i$  be its predecessor in  $C$ .
- $u_i$  must appear before  $u_j$  in the topological ordering.
- This contradicts the fact that  $u_j$  is the smallest element of  $C$  in the topological ordering.

**DAG implies topological ordering:**

Proof by induction: Base case: If  $G$  has one or two nodes, then  $G$  has a topological ordering.

Induction steps: Assume that a DAG up to  $k$  nodes has a topological ordering (induction hypothesis). we will prove that a DAG with  $k + 1$  nodes has a topological ordering.

- By our lemma, there is at least one source node in  $G$ , and let  $u$  be the node.
- Put  $u$  at the beginning of the topological ordering.
- Consider the graph  $G'$ , obtained by  $G$  by removing  $u$  and its incident edges.
- $G'$  is a DAG with  $k$  nodes.
- It has a topological ordering  $u_1, u_2, \dots, u_k$  by the induction hypothesis.
- Append this ordering to  $u$  to get a topological ordering of  $G$ .

Here is the algorithm:

---

**Algorithm 9** Topological Sorting

---

```

1: procedure TOPOLOGICALSORTING( $G$ )
2:   find a source vertex  $u$ 
3:   set  $u$  as the first element of the topological ordering
4:    $G' \leftarrow G$  with  $u$  and its incident edges removed
5:    $L \leftarrow$  TOPOLOGICALSORTING( $G'$ )
6:   append  $L$  to  $u$ 
7: end procedure

```

---

Running time of the algorithm is  $O(n^2)$

**Modified Topological Sorting:**

Running time of the algorithm is  $O(n + m)$

---

**Algorithm 10** Modified Topological Sorting

---

```
1: procedure MODIFIEDTOPOLOGICALSORTING( $G$ )
2:    $L \leftarrow \text{emptylist}$ 
3:    $S \leftarrow$  set of all source vertices
4:   while  $S \neq \emptyset$  do
5:     remove a vertex  $u$  from  $S$ 
6:     append  $u$  to  $L$ 
7:     for each edge  $(u, v)$  do
8:       remove edge  $(u, v)$  from  $G$ 
9:       if  $v$  is a source vertex then
10:        add  $v$  to  $S$ 
11:       end if
12:     end for
13:   end while
14:   if  $G$  has edges then
15:     return  $G$  has a cycle
16:   else
17:     return  $L$ 
18:   end if
19: end procedure
```

---

### 3.7 Finding strongly connected components

**connected components:** A connected component of an undirected graph is subgraph of the graph where any two nodes are connected by a path.

**strongly connected components:** A strongly connected component of a directed graph is a subgraph of the graph where any two nodes are mutually reachable.(mutually reachable: there is a path from  $u$  to  $v$  and a path from  $v$  to  $u$ )

**Finding strongly connected components:**

**Kosaraju's algorithm:**

---

**Algorithm 11** Kosaraju's algorithm

---

```
1: procedure KOSARAJU( $G$ )
2:   Initialise stack  $S$ 
3:   Select a arbitrary node  $s$ 
4:   DFS_tree=DFS( $G, s$ )
5:    $S \leftarrow$  nodes in DFS_tree
6:    $G^R \leftarrow$  nodes in order of  $S$ 
7:   DFS( $G^R, s$ )
8:   return the nodes in the DFS tree
9: end procedure
```

---

**Running time of Kosaraju's algorithm:**  $O(n + m)$