

# Quantum Dynamics, Logistic Map, and Digit Distribution of Special Math Constants

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**Abstract**—Using the logistic map instead of the base quadratic system as in my previous paper [11], I obtain very similar quantum dynamics, this time for the function  $\sin^2(\sqrt{x})$  instead of  $\exp(x)$ . When  $x$  is a small integer or a product of consecutive primes, my framework reveals new insights on the digit distribution of major math constants. I also discuss deep findings about the chaotic nature of dynamical systems and several applications including in AI.

## I. INTRODUCTION

Let  $n$  be a fixed, large integer, say  $n = 10^6$ . The **dynamical system**  $S_{k+1} = S_k^2$ , starting with  $S_0 = 2^n + 1$ , that is, a string consisting of  $n - 1$  zeros with a one at both ends, has this particular property: the first  $n$  binary digits of  $S_n$  match those of  $\exp(1)$ , give or take. To make it a true mapping and without changing the conclusions, each  $S_k$  is rescaled: it is multiplied by an integer power of 2 (negative or positive), so that it stays in  $[1, 2]$  for all  $k = 0, 1, 2$  and so on.

If you replace  $S_0 = 2^n + 1$  by  $S_0 = 2^n + x$  where  $x$  is a small integer, say  $x = \pm 1$  or  $x = 3$ , then the first  $n$  binary digits of  $S_n$  match those of  $\exp(x)$ . This remains true if each  $S_k$  is truncated to a precision of  $2n$  bits. I discuss the details in [9], [10] and [11].

I call the dynamical system in question the **base quadratic map**. Other quadratic dynamical systems include the **logistic map**  $S_{k+1} = 4S_k(1 - S_k)$  and the standard **quadratic map**  $S_{k+1} = S_k^2 + c$  where  $c$  is a constant. The latter is defined in the complex plane and leads to the **Mandelbrot set** featured in Figure 1.

The goal here is to obtain similar results for the logistic map, thus expanding my investigations about the **digit sum function** in [9]. For the logistic map, starting with the **seed string**  $S_0 = 2^{-2n}$  consisting of a single ‘1’, we get  $S_n = \sin^2(1)$  correct to  $2n$  binary digits if we keep a  $2n$ -bit precision at all times. Even better: if  $S_0 = x \cdot 2^{-2n}$ , then  $S_n = \sin^2(\sqrt{x})$  correct to  $2n$  bits. The **quantum dynamics** of the digit sum are very similar to those observed with the base quadratic map. This framework offers new directions to reach our ultimate goal: proving deep results about the digit distribution of special math constants.

## II. LOGISTIC MAP AND THE DIGIT SUM FUNCTION

For the logistic map  $S_{k+1} = 4S_k(1 - S_k)$  with  $0 \leq S_k \leq 1$ , there is a closed-form expression for the  $k$ -th iterate:

$$S_k = \sin^2(2^k \arcsin \sqrt{S_0}). \quad (1)$$

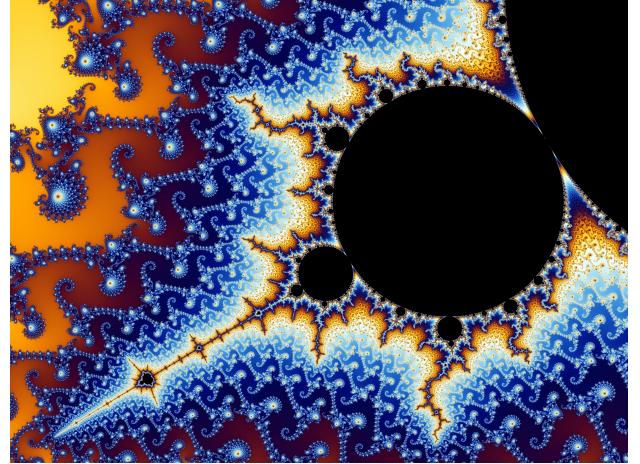


Fig. 1. Mandelbrot set linked to the standard quadratic map

In particular, if  $S_0 = x \cdot 2^{-2n}$  with  $x > 0$ , then using a Taylor expansion for  $\arcsin(\sqrt{S_0})$ , we get

$$S_k = \sin^2\left(\frac{x^{1/2}}{2^{n-k}} + \frac{1}{6} \cdot \frac{x^{3/2}}{2^{3n-k}} + \dots\right). \quad (2)$$

Thus, when  $k = n$  and  $n \rightarrow \infty$ , we get  $S_n \rightarrow \sin^2(\sqrt{x})$ . By contrast, for the base quadratic map with the seed  $S_0 = 2^n + x$  and proper rescaling (multiplication by an integer power of 2), we obtained the asymptotic formula

$$S_k = \left(1 + \frac{x}{2^n}\right)^{2^k} \sim \exp\left(\frac{x}{2^{n-k}}\right) \quad (3)$$

converging to  $\exp(x)$  when  $k = n$  and  $n \rightarrow \infty$ . The right part in (3) has an accuracy of about  $2n - k$  bits if the precision on the left part is kept to  $2n$  bits at all times.

### A. Model comparison, with illustrations

Here  $n = 10^5$  is fixed. The digit sum function  $\zeta_k$  counts the number of ‘1’ in the first  $n$  binary digits of  $S_k$ , for  $k = 0, 1, 2$  and so on. Since we start with a seed  $S_0$  close to 0 with the proportion of ‘1’ typically increasing over time until reaching about 50% at  $k = n$ , I use the **adjusted digit sum**  $\zeta'_k$  instead. It counts the number of ‘1’ in the first  $n$  digits of  $S_k$ , starting at position  $n - k$  in the digit expansion of  $S_k$ .

Let  $\rho = k/n$ . The behavior of  $\zeta'_k$  is trivial when  $\rho < 0.50$ . Up until  $\rho = 0.75$ , patterns are usually strong and obvious. It starts to get somewhat chaotic as  $\rho > 0.90$ , and when  $\rho \geq 1$ , we are in full **chaotic phase**. Thus I focus on  $0.75 < \rho < 1$ .

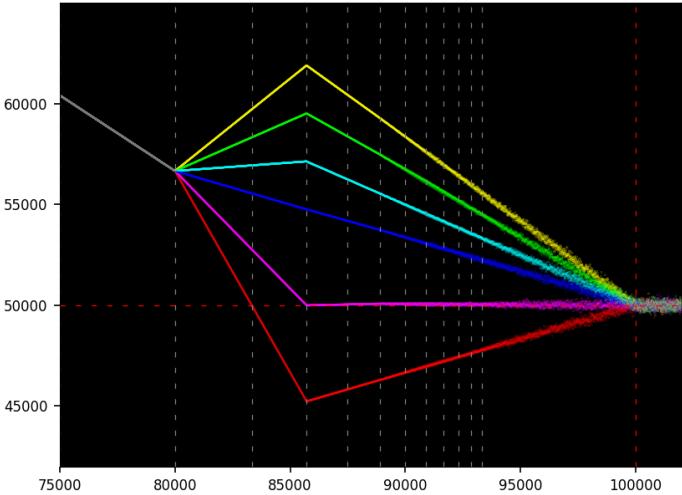


Fig. 2. Logistic map:  $\zeta'_k$  with  $x = 1, n = 10^5, k$  on X-axis

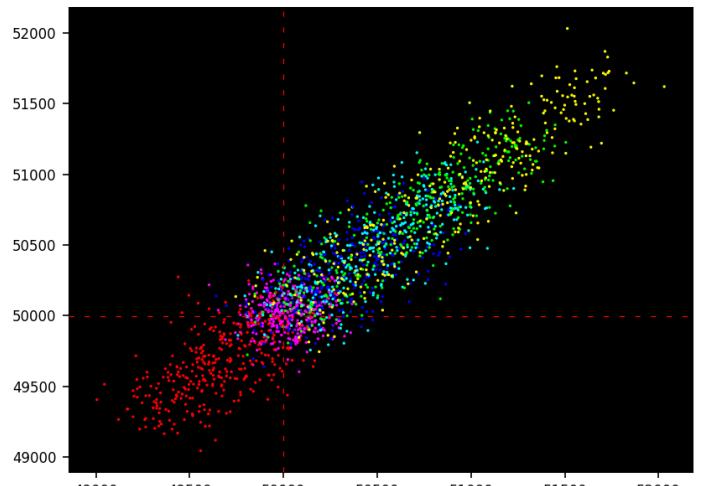


Fig. 5. Log. map:  $(\zeta'_{k-12}, \zeta'_k)$  with  $0.98 < \frac{k}{n} < 1, x = 1, n = 10^5$

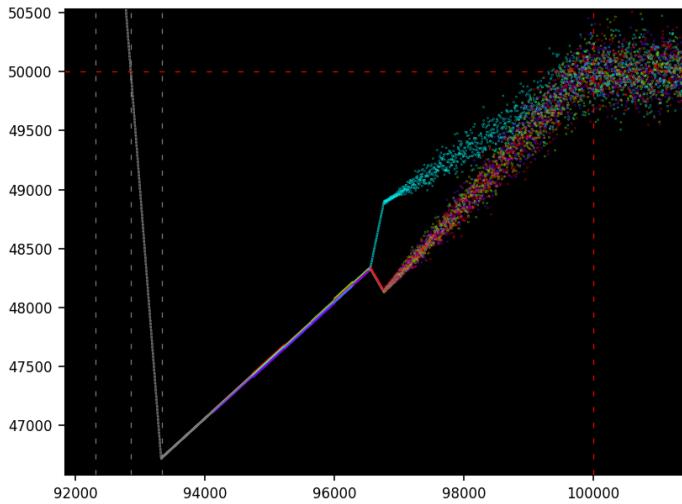


Fig. 3. Logistic map:  $\zeta'_k$  with  $x = \pi_\kappa, n = 10^5 - 1, k$  on X-axis

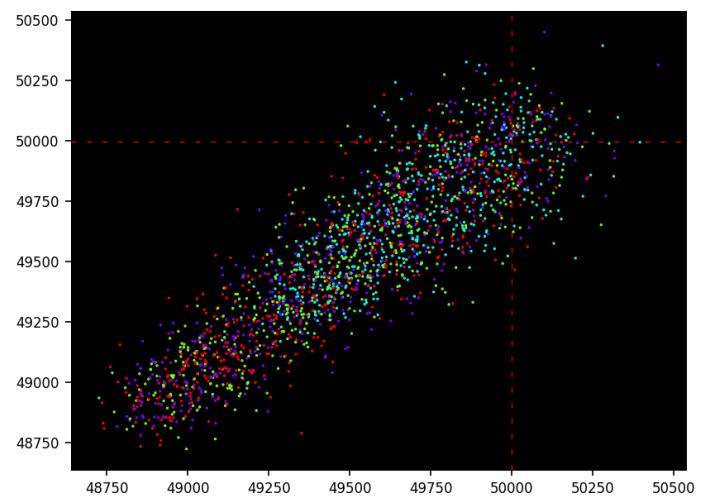


Fig. 6. Log. map:  $(\zeta'_{k-12}, \zeta'_k)$  with  $0.98 < \frac{k}{n} < 1, x = \pi_\kappa, n = 10^5 - 1$

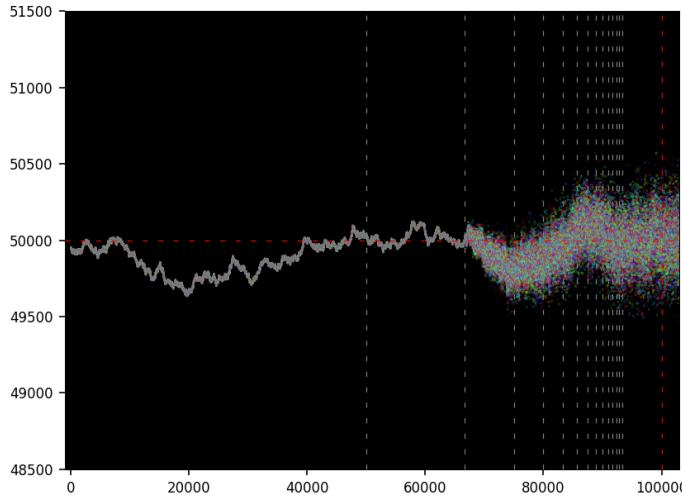


Fig. 4. Logistic map:  $\zeta'_k$  with  $x = \sqrt{2}, n = 10^5, k$  on X-axis

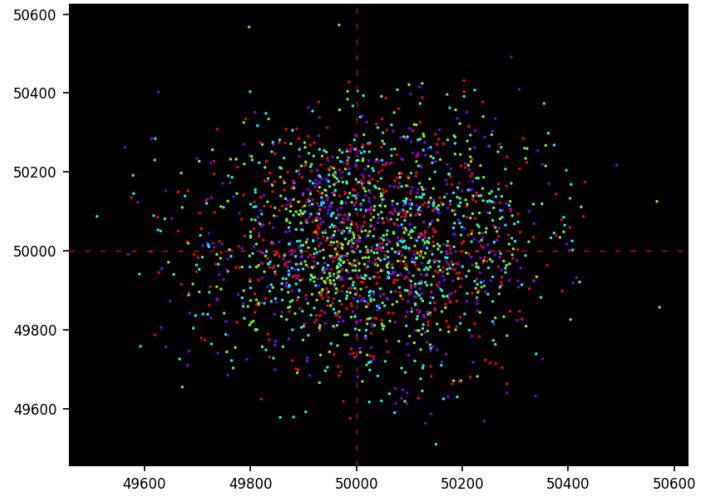


Fig. 7. Log. map:  $(\zeta'_{k-12}, \zeta'_k)$  with  $0.98 < \frac{k}{n} < 1, x = \sqrt{2}, n = 10^5$

From now on, **BQ** denotes the base quadratic map with illustrations in [11]. With the seeds  $S_0 = x \cdot 2^{-2n}$  for the logistic map and  $S_0 = 2^n + x$  rescaled to  $1 + x \cdot 2^{-n}$  for the BQ map, we observe the following 3 types of behavior for  $\zeta'_k$ .

- **Chaotic.** This represents the vast majority of cases, for instance if  $x$  is a random number in  $[0, 1]$ . See Figures 4 and 7.
- **Quantic.** See Figures 2 and 5 for the logistic map, and Figure 8 for the BQ map. It happens when  $x$  is a small integer, say  $x = 1$ . The color indicates the **congruential class** that  $k$  belongs to, modulo 6.
- **Hybrid.** You don't have multiple branches depending on the **congruential class** that  $k$  belongs to as in the quantic case, at least until  $k \approx 0.97 \cdot n$ . See Figure 3 and 6. Here,  $x = \pi_\kappa$  is the  $\kappa$ -th **primorial**, that is, the product of the first  $\kappa$  primes ( $\kappa = 9$ ).

Zoom in on any figure to see the details. In particular, the vertical dashed lines indicate the abscissa of change points (forking or slope change) in the function  $\zeta'_k$ . The values seen in the pictures (specific  $k$ 's on the X-axis) are those of the BQ map, but remain valid for the logistic map. They correspond to  $k = \rho n$ , with  $\rho = \frac{1}{2}, \frac{2}{3}, \frac{3}{4}, \frac{4}{5}, \frac{5}{6}$  and so on.

The peculiar shape in the quantic and hybrid cases, both for the BQ and logistic maps, is explained by the unusual Taylor series when fully expanding formulas (2) and (3). See section II.B in [11] for details.

Figures 2, 3, and 4 represent time series with **quantum states** for the digit sum  $\zeta'_k$  with  $k$  on the X-axis and  $\zeta'_k$  on the Y-axis. By contrast, Figures 5, 6 and 7 are scatterplots representing the vector  $(\zeta'_{k-w}, \zeta'_k)$  for  $0.98 \cdot n < k < n$ . It shows a **spectral view** when we are approaching chaos; full chaos starts at  $k \geq n$  in the quantic and hybrid cases, and at about  $k = 0$  in almost all other cases. The parameter  $w$  is called the **time lag**. An alternative view consists of showing the autocorrelation function computed on the  $\zeta'_k$  sequence when  $k$  is close to  $n$ , for various time lag values  $w = 1, 2$  and so on. However, in the quantic and hybrid cases, the process is **non-stationary** until  $k \geq n$ .

With a random seed, the scatterplot in Figure 7 shows a Gaussian distribution centered at  $(\frac{n}{2}, \frac{n}{2})$ . Curiously, the seed with  $x = \pi^2/4$  leads to  $\zeta'_k$  hovering around  $n/2$  as in Figure 4 when  $k < n$ , but suddenly dropping to 0 at  $k = n$ , since  $S_n \approx \sin^2(\sqrt{x}) = 1$ . Keep in mind that  $\zeta'_k$  counts the '1' in the first  $n$  digits of  $S_k$ , starting at position  $n - k$  in the digit expansion of  $S_k$ .

See also Figures 8, 9, 10 and 11. For the BQ map,  $x$  is the parameter in the seed  $S_0 = 2^n + x$ ; for the logistic map, it comes from  $S_0 = x \cdot 2^{-2n}$ . In Figures 3 and 10 featuring the logistic map,  $x$  is a multiple of 3; to see the congruential classes in 6 colors, I had to replace  $n = 10^5$  by  $n = 10^5 - 1$ , which is a multiple of 3.

Note that Figures 10 and 11 feature the non-adjusted digit sum  $\zeta_k$  instead of  $\zeta'_k$ . Discontinuities in  $\zeta'_k$  are sometimes observed. Each figure has its own scale: the X- and Y-axis are custom-truncated depending on the case, to provide the best view, in particular near  $k = n$  where the real action is taking

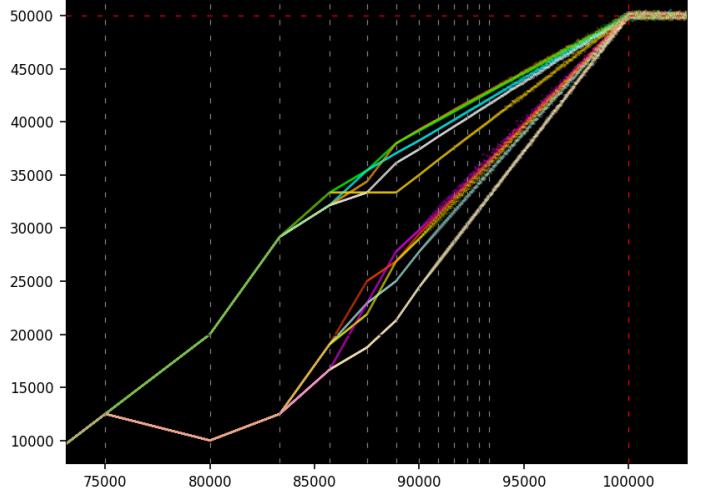


Fig. 8. BQ map:  $\zeta'_k$  with  $x = 1, n = 10^5, k$  on X-axis

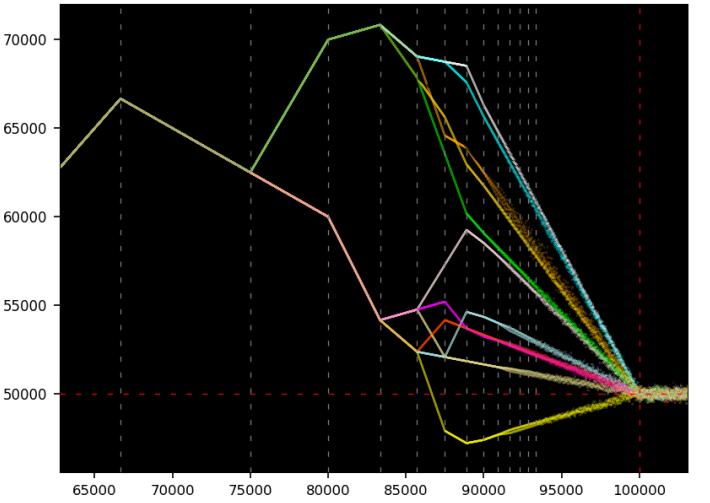


Fig. 9. BQ map:  $\zeta'_k$  with  $x = -1, n = 10^5, k$  on X-axis

place. Pictures featuring moving averages that collapse multi-branches into a single one with a constant slope, are posted in [11].

Finally, in [11], I discuss the inverse iterations for the BQ map, moving backward from  $S_{k+1}$  to  $S_k$  rather than forward from  $S_k$  to  $S_{k+1}$ . It is possible to do the same for the logistic map. However, in both maps, the inverse mapping  $S_{k+1}$  to  $S_k$  is not one-to-one, but one-to-two. In the Python code in section IV,  $S_k$  is denoted as `prod`. The link to **fractals** is visible in Figures 1 and 2 in [9].

#### B. Normality of special math constants

My framework was first designed for the BQ map to answer this question: are the binary digits of  $e$  evenly distributed? That is, is  $e$  **simply normal** in base 2? This famous multi-century old conjecture is still an open question. For the logistic map, replace  $e$  by  $\sin^2(1)$ . In this section, I share the progress that I made recently thanks to the new methodology discussed in this article.

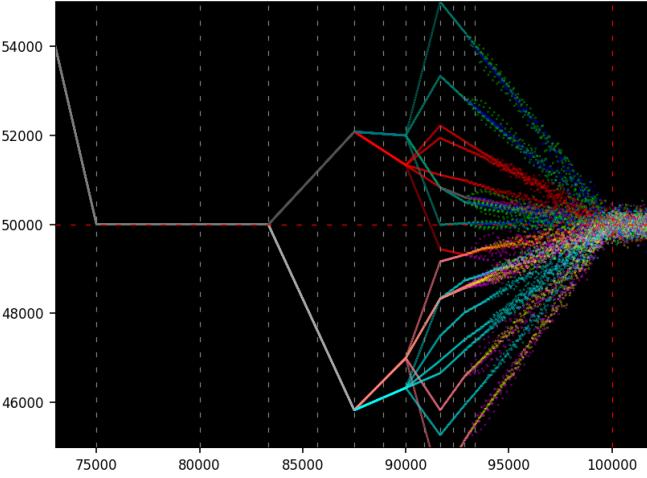


Fig. 10. Logistic map:  $\zeta_k$  with  $x = 9, n = 10^5 - 1, k$  on X-axis

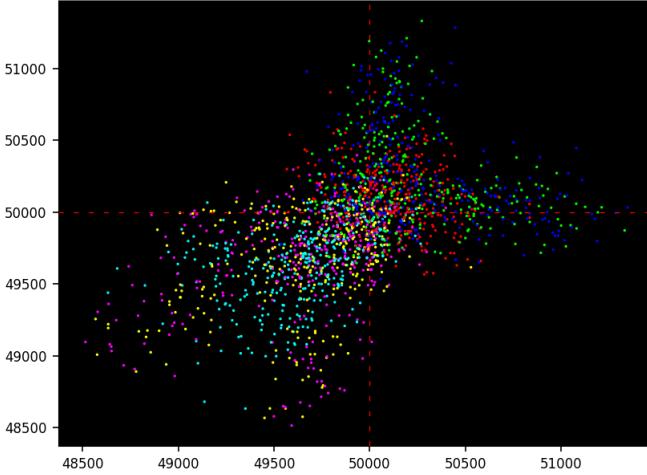


Fig. 11. Log. map:  $(\zeta_{k-12}, \zeta_k)$  with  $0.98 < \frac{k}{n} < 1, x = 9, n = 10^5 - 1$

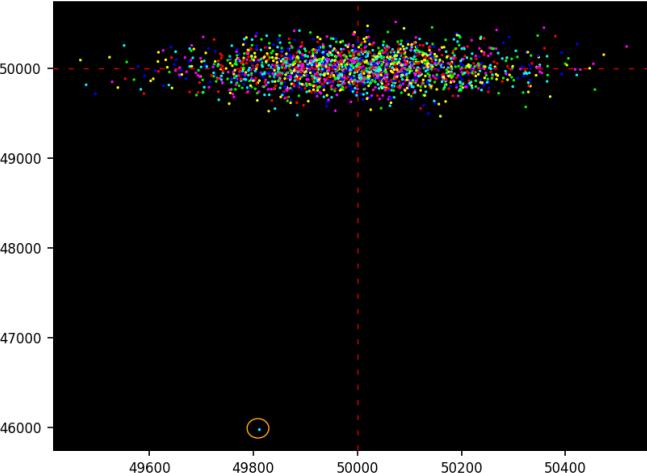


Fig. 12. Same as Fig. 7 but with  $x$  leading to  $\zeta'_n/n \approx 0.46$  instead of 0.50

Let  $n = 10^5$ . Using a random  $S_n$  in  $[0, 1]$  with a proportion  $p = 46\%$  of ‘1’ in its binary digit expansion rather than  $p = 50\%$  as in Figure 4, how would the sequence  $(\zeta'_k)$  look like, for  $k = 0, 1$  and so on? The answer: exactly as in Figure 4. But the spectral view would be very different from the corresponding Figure 7. Instead, it would look like Figure 12.

The explanation is as follows: near to  $k = n$  with  $n = 10^5$ ,  $\zeta'_k/n$  is always close to 50%, with a sudden drop exactly and only at  $k = n$ , where  $\zeta'_n/n \approx 46\%$ . Thus, the single outlier in Figure 12, circled in orange. If you increase  $n$  from  $10^5$  to (say)  $10^{50}$ , you would still get an outlier even if increasing  $p = 0.46$  to  $p = 0.48$ .

It sounds as if you cannot get the **quantic** behavior observed in Figure 2 if the digits of  $S_n$  follow a **Bernoulli process** of parameter  $p \neq \frac{1}{2}$ , for  $n$  large enough. Instead, it would have the **chaotic** behavior pictured in Figure 4

That is, the binary digits of  $\sin^2(1)$  either do not follow a Bernoulli process, or if they do, then the probability of ‘1’ is  $p = \frac{1}{2}$ . Of course, almost everyone believe the latter to be true, not the former. There is no formal proof yet, but the progress made here is very encouraging.

A side effect (conjecture) is the following: if the digits of  $S_n$  follow a Bernoulli process with  $p \neq \frac{1}{2}$ , then the digits of

$$S_{n-1} = \frac{1 \pm \sqrt{1 - S_n}}{2} \quad (4)$$

follow a Bernoulli process with  $p = \frac{1}{2}$ . By “following”, I mean that the empirical joint digit distribution converges to that of a Bernoulli process as  $n \rightarrow \infty$ .

Empirical evidence is easy to obtain. Generate  $y = S_n$ , a number in  $[0, 1]$  with digits simulated to follow a Bernoulli process with  $p = 0.46$ . Since  $y = \sin^2(\sqrt{x})$ , we have  $x = \arcsin^2(\sqrt{y})$ . Let’s use  $S_0 = x \cdot 2^{-2n}$ . Note that the inverse map is not one-to-one, see (4). Set the precision to  $2n$  bits and proceed as in all other cases ( $x = 1, 9$ , and so on) using the code in section IV to generate  $S_k$  and compute  $\zeta'_k$  for  $k = 0, 1$  and so on. I did the test, with the following Python code to generate  $y$  and get  $x$ .

---

```

import gmpy2
import numpy as np

np.random.seed(410)
n = 100000
u = np.random.binomial(n=1, p=0.46, size=2*n)
stri = [str(bit) for bit in u]
stri = ''.join(stri)
ctx = gmpy2.get_context()
ctx.precision = 2*n
y = gmpy2.mpz(stri, 2)
y /= 2**((2*n))
x = gmpy2.asin(gmpy2.sqrt(y))
x = x*x

```

---

### C. Applications and references

Here I compiled a list of useful references related to the topic, broken down by application, with a focus on literature recently published.

- The framework presented here relies on discrete **quadratic dynamical systems**. This family also includes the **logistic**

map and the example discussed in [19]. For additional references, see my book on chaos and dynamical systems [7].

- Showing that the binary digits are evenly distributed is the first step towards proving that  $e$  is a **normal number**. Andrew Granville and David Bailey [3] are good references on this topic. For recent publications on normal numbers, see Verónica Becher [4] and [2], [12]. One of best results known for any major math constant is the fact that the proportion of ones in the first  $n$  binary digits of  $\sqrt{2}$  is larger than  $\sqrt{2n}$ , see [18].
- The digit sum or digit count functions (both are identical for binary digits) is also known as the **Hamming weight**, with a fast algorithm described [here](#) and a full chapter in [20]. The Wolfram entry for the **digit sum** (see [here](#)) features an exact closed-form formula for the number of digits equal to 1 in the binary expansion of any integer, with more references. For a discussion on the **carry digit** function (a **2-cocycle**) that propagates 1's from right to left in the successive iterations  $S_k$ , see [1], [5].
- An interesting application of the digit sum is featured in [13] in the context of genotype maps, with processes not unlike the dynamical systems discussed in this article, and **blancmange curves** almost identical to Figure 3.3 in my book on numeration systems [7].
- There is a connection to **quantum maps** and **quantum cryptography** [6], [17]. For **PRNGs** (pseudo-random generators) based on irrational numbers, see chapter 13 in [8] or chapter 4 in [7]. Finally, if you use an arbitrary seed  $S_0$  with about 50% of '0' and '1', you obtain strings  $S_k$  that look random after very few iterations.
- **Deep neural networks** have been used to identify the underlying model of dynamical systems, based on available data produced by simulations or from real life observations, see [15], [16], [21]. In our case, the model would be a simple formula that generates the values of the digit sum function, to study its asymptotic properties. See also [14].

### III. CONCLUSION

This article is the fourth in this series, preceded by [9], [10] and [11]. The previous articles focus on the simplest quadratic map, while this one shows how we can attain similar results with the most well-known dynamical system: the logistic map.

Orbits in chaotic dynamical systems are very sensitive to the seed. Starting with close seeds  $S_0 = x \cdot 2^{-2n}$  and  $S'_0 = 0$  I computed  $\Delta_k = |S_k - S'_k|$  and showed that  $\Delta_n \approx \sin^2(\sqrt{x})$  with a precision of about  $2n$  bits, also getting good approximations to  $\Delta_k$  when  $k \leq n$ . In the logistic map,  $S'_0$  is a **fixed point**.

I also discussed the unique quantic behavior of the digit sum function when using special values of  $x$ . This opens up new directions to study the digit distribution of special math constants,  $e$  in particular. Most importantly, it leads to many applications including cryptography, synthetic data, pattern

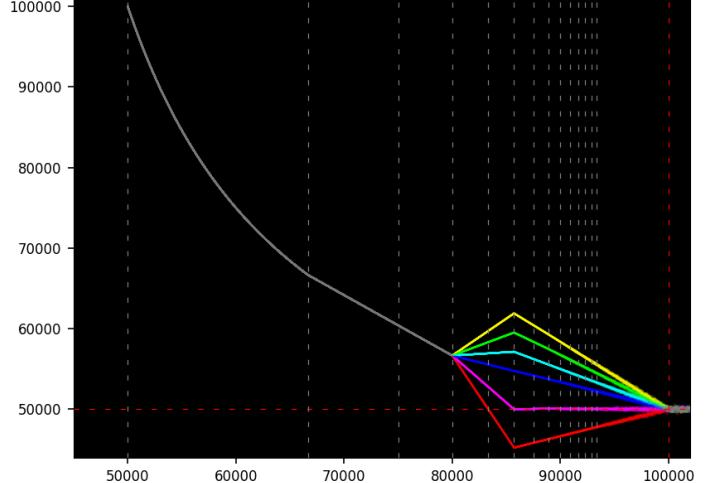


Fig. 13. Unusual precision parameters producing curve and discontinuity

detection or proof automation with LLMs [10], agent based modeling, and more.

The material presented here can also be used to complement a course on dynamical systems, for scientific research, or to start a PhD thesis on the subject.

### IV. PYTHON CODE

In the code, the variable `prod` represents  $S_k$ . I also use external functions to check how many of the binary digits of  $S_n$  match those of  $\sin^2(\sqrt{x})$ . The code is on GitHub, [here](#).

The fact that  $0.111\dots = 1.000\dots$  in base 2 can have side effects on the values of  $\zeta_k$  and  $\zeta'_k$ . Whether the rightmost digits of  $S_k$  are '1000' or '0111' has no real impact on the accuracy. However it has an impact – usually small – on the digit sum, but potentially large especially when  $k$  is small. It remains to be seen if it contributes to the increased chaos near  $k = n$ .

A spectacular manifestation of this side effect is as follows. To compute  $\zeta'_k$ , I look at the first  $n$  digits, starting at position  $2n - k$ , of the integer number  $\lfloor 2^{2n} S_k \rfloor$ . The result is shown in Figure 13, with  $k$  on the X-axis and  $\zeta'_k$  on the Y-axis, with  $n = 10^5$  and starting with the seed  $S_0 = x \cdot 2^{-2n}$  with  $x = 1$ .

When  $k < \frac{n}{2}$ ,  $\zeta'_k = 0$ . However at  $k = \frac{n}{2}$ ,  $\zeta'_k \approx n$ . So there is a discontinuity at  $k = \frac{n}{2}$ . And then, the path from  $k = \frac{n}{2}$  to  $k = \frac{2}{3}n$  is not a straight line, but a curve. I did not expect that kind of behavior. To the contrary, if you use  $\lfloor 2^{4n} S_k \rfloor$  rather than  $\lfloor 2^{2n} S_k \rfloor$  by setting `pow` to 4 rather than 2 in the code, the behavior is different until  $k = \frac{2}{3}n$ , though not thereafter: when  $k < \frac{n}{2}$ ,  $\zeta'_k \approx n$  and there is no discontinuity at  $k = \frac{n}{2}$ . Also the curvy section is replaced by a straight line. Both behaviors are correct.

```
import gmpy2
import numpy as np

n = 99999 # choose n divisible by 3 if x = 9 and
           ncolors = 6
H = int(1.1*n)

import colorsys

def hsv_to_rgb(h, s, v):
```

```

return tuple(round(i * 255) for i in
            colorsys.hsv_to_rgb(h, s, v))

def generate_contrasting_colors(ncolors):
    colors = []
    for i in range(ncolors):
        hue = i / ncolors
        col = hsv_to_rgb(hue, 1.0, 1.0)
        color = (col[0]/255, col[1]/255, col[2]/255)
        colors.append(color)
    return colors

ncolors = 6 # 4 colors for hybrid case, 6 for
            # quantic
colorTable = generate_contrasting_colors(ncolors)

#--- 1. Main

import gmpy2
import numpy as np

kmin = 0.00 * n # do not compute digit count if k
                <= kmin
kmax = 1.15 * n # do not compute digit count if k
                >= kmax
kmax = min(H, kmax)

# precision set to L bits to keep at least about n
# correct bits till k=n
ctx = gmpy2.get_context()
ctx.precision = 2*n

# x = gmpy2.mpz(2*3*5*7*11*13*19*23*29)
x = gmpy2.mpz(1)

prod = gmpy2.mpfr(x/2**((2*n)))

# local variables
arr_count1 = []
arr_colors = []
xvalues = []
ecnt1 = -1
e_approx = "N/A"

OUT = open("digit_sum.txt", "w")

for k in range(1, H+1):

    prod = 4*prod*(1 - prod)
    pow = 2*n # 2n --> 0.1111.. | 4n --> 1.0000..
    pstri = bin(gmpy2.mpz(2**pow) * prod)
    stri = pstri[0:2*n]

    if k > kmin and k < kmax:
        stri = stri[2:]
        lstri = len(stri)
        if k == n:
            e_approx = stri
        # estri = stri[0:n]      # for digit sum
        estri = stri[max(0,n-k):2*n-k] # for adjusted
                                        # digit sum
        ecnt1 = estri.count('1') * n / (1+len(estri))
        arr_count1.append(ecnt1)
        color = colorTable[k % ncolors]

        arr_colors.append(color)
        xvalues.append(k)
        OUT.write(str(k)+"\t"+str(ecnt1)+"\t"
                  +str(lstri)+"\n")

    if k%1000 == 0:
        print("%d %d %d" %(k, ecnt1, lstri))

OUT.close()

#--- 2. Compute bits of sin^2(sqrt(x_)) and count
# correct bits in my computation

# Set precision to L binary digits
gmpy2.get_context().precision = 4*n
e_value = gmpy2.sin(gmpy2.sqrt(x))
e_value = e_value * e_value

# Convert e_value to binary string
e_binary = gmpy2.digits(e_value, 2)[0]

k = 0
while k < len(e_approx) and e_approx[k] ==
        e_binary[k]:
    k += 1
# e_binary should be equal to e_approx up to about
# n bits
print("\n%d correct digits (n = %d)" %(k, n))

e_approx_decimal = 0
for k in range(0,80):
    e_approx_decimal += int(e_approx[k])/(2***(k+1))
print("e_exact : %16.14f" % (e_value))
print("e_approx: %16.14f" % (e_approx_decimal))
print("Up to factor 2 at integer power of 2.")

#--- 3. Create the main plot

import matplotlib.pyplot as plt
import matplotlib as mpl
import numpy as np

mpl.rcParams['axes.linewidth'] = 0.5
plt.rcParams['xtick.labelsize'] = 8
plt.rcParams['ytick.labelsize'] = 8
plt.rcParams['axes.facecolor'] = 'black'

plt.scatter(xvalues, arr_count1,s=0.005,
            c=arr_colors)
plt.axhline(y=n/2,color='red',linestyle='--',
            linewidth=0.6,dashes=(5,10))
plt.axhline(y=n/5, color='black', linestyle='--',
            linewidth = 0.6, dashes=(5, 10))
plt.axvline(x=n, color='red', linestyle='-' ,
            linewidth = 0.6, dashes=(5, 10))

for k in range(1,15):
    plt.axvline(x=k*n/(k+1),c='gray',linestyle='--',
                linewidth=0.6,dashes=(5, 10))

# we start with about 0% of 1 going up to about 50%
plt.ylim([0.44*n, 1.01*n])
plt.xlim([0.45*n, 1.02*n])
plt.show()

#--- 4. Create AR scatterplot

nv = n
lag = 12
tail = 2000
plt.scatter(arr_count1[n-tail:lag:n-lag],
            arr_count1[n-tail:n], s=0.4,
            c=arr_colors[n-tail:lag:n-lag])
# plt.plot(arr_count1[n-tail:lag:n-lag],
#           arr_count1[n-tail:n], linewidth=0.6) ##,
#           c=arr_colors[n-tail:lag:n-lag])
plt.axhline(y=n/2,color='red',linestyle='--',
            linewidth=0.6,dashes=(5,10))
plt.axvline(x=n/2, color='red', linestyle='-' ,
            linewidth = 0.6, dashes=(5, 10))

plt.show()

```

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