Last time:

GD:
$$x^{(t+1)} = x^{(t)} - \mu \nabla f(x^{(t)})$$

Dhe gave an interpretation of the above GD update step whereby

oct+1) minimizes the f =

$$g(z) = f(x^{(t)}) + \nabla f(x^{(t)})^{T}(z - x^{(t)})$$

$$+ \frac{1}{2\mu} ||z - x^{(t)}||^{2}$$

2 We talked about the method of steepest descent:

GD kelongs to a class of algorithms where $\alpha^{(t+1)} = \alpha^{(t)} - \mu P^{(t)}$

1 chosen iteratively

· Specifically GD selects p(ts)
to minimize

min \frac{\mathfaire}{\psi \left(\pi^{(t)})^T p}{\psi \left(\psi^{(t)})^T p}

· Other algorithms can be derived by replacing 1/191/2 by other norms

e.g. Ne obtain coordinate descent when we pick P to minimize

min $DF(x^{(t)})^T P$ $P = \frac{||P||_2}{||P||_2}$

Here we end up with

it = index of largest 1 | entry of 7F(x(n))

P*= -sign (J; f(x(t))) // Tf(x(t)) // E;*

vector of all zeros

except at j entry where

it is I

 $= -\frac{\partial f(x^{(t)})}{\partial x_{j}^{*}} \vec{e_{j}}^{*} \qquad \left(= -\nabla_{j}^{*} f(x^{(t)}) \vec{e_{j}}^{*} \right)$

Gradient Descent under Constraints	:
We have been solving	
$\begin{cases} mi & f(x) \\ xepn \end{cases}$	
by running the iterations	
$\int x^{(t+1)} = x^{(t)} - n \nabla F(x^{(t)})$	
But what change do we need to	

But what changes do we need to make to GD if we now want to solve:

min f(x) subject to xESZ

2CRⁿ

The issue is that even if $x^{(0)}$ or any $x^{(0)} \in SZ$, there is no quarantee that $x^{(t+1)}$ as given by $x^{(0)}$ is also in SZ.

Me'll need the following defin

The projection of a point of anto a set so is defined as the closest pt in so, to one of the closest pt is defined as

 $\prod_{\Omega} (x) = \underset{Y \in \Omega}{\operatorname{aramin}} \|x - y\|$

Examples: $Z = T_{\mathcal{R}}(z)$ $T_{\mathcal{R}}(x)$

If $\Omega = B_2^n := \{x : ||x|| \le 1\}$ then $T_2(x) = \{x, ||x|| \le 1\}$ $\frac{x}{||x||}$

Hen
$$T_{x}(x) = \int x \quad \text{if} \quad x_{1} \geq 0$$

$$\begin{cases} (0, x_{2}, x_{3}, \dots, x_{n}) \\ \text{if} \quad x_{1} < 0 \end{cases}$$

Projected GD:

In order to solve where 2 convex min f(x) subject to $x \in S^2$ we run the iteration

$$\alpha^{(t+1)} = \pi \left(x^{(t)} - u^{(t)} \nabla F(x^{(t)}) \right)$$
Standard
GD step

Projected onto 2

$$\chi^{(t+1)} = \chi^{(t)} - \mu^{(t)} \nabla f(\chi^{(t)})$$

$$\chi^{(t+1)} = \underset{x \in S_{2}}{\operatorname{argmi}} \| \chi^{(t+1)} - \chi \|$$

In general, solving $x^{(t+1)} = \operatorname{argmi} \|y^{(t+1)} - x\|$

is itself a constrained opt. Problem that may not have a nice closed form solution.

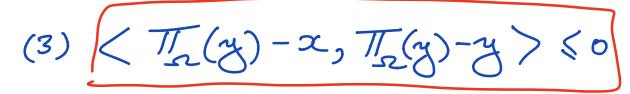
That said, there are interesting cases, where 52 is nice (like the exemples above) and where $T_2(x)$ is easy to compute.

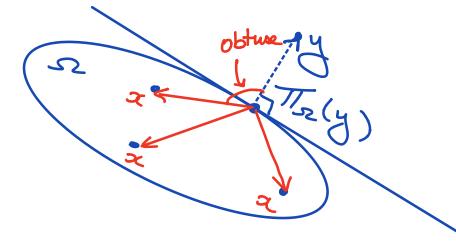
Some facts about projections:

Lemma: If $SZCR^n$ is convex and closed (and not empty) and if $x \in SZ$, $y \in R^n$

then (1) The (y) is unique.

(2) Ta (Ta(y)) = Te(y)





So (对- 死(y), 不(y)) > 数-死(y), x>

(5) $\| \sqrt{\|x\|_{2}} (y) - x \| \le \|y - x\|$

Ex:
$$\max_{\|x\|_{\infty} \le 1} \langle y, x \rangle$$

$$\|x\|_{\infty} \le 1$$

$$\leq \max_{\|x\|_{\infty} \le 1} \|y\|_{1} \|x\|_{\infty}$$

$$|x|_{\infty} \le 1$$
Find x s.t. $\langle y, x \rangle = \|y\|_{1}$

$$\|y\|_{1} = \sum_{i} |y_{i}|_{1}$$

$$\langle y, x \rangle = \sum_{i} |y_{i}|_{2} \langle y_{i}, x \rangle + hm$$

$$\langle y, x \rangle = \sum_{i} |y_{i}|_{2} \langle y_{i}, x \rangle + hm$$

$$\langle y, x \rangle = \sum_{i} |y_{i}|_{2} \langle y_{i}, x \rangle + hm$$

Application: Fund opt.

C: vector of expected returns for assets

X: fruction invested in each asset

M: desired expected return on fund

o²: desired expected variance on fund

M= TX

 $M = r^T X$ asset covariance $\sigma^2 = x^T \sum_{x} x$

Problem: min o²
5.t. u fixel
frection bought 20

min
$$x^{T} \leq x$$

 $x^{T} \times x = x$
 $x \in \{x : x \neq 0 \ \forall i\}$

OR

max rTX X S.t. $X^T \leq X = \sigma^2$ $1^T X = 1$ $X_{i} \neq 0$ t

Trus action cost: Return $r_{X-C_{-}X_{-}-C_{+}X_{+}}$ $1_{X}^{T} + c_{-}^{T} x_{-} + c_{+}^{T} x_{+} = 1$ $x_{i} = \overline{x}_{i} + x_{+} = x_{-}$

Application: Facility location . Store 2

storez pol

min \(\frac{\x}{\infty} \) | | \(\times \) \(\cap \)

If p=2, + took cost 11x:- cllp, then C = 1 {x:

P=1 => roads are on grid