

CSCI-UA 480.4: APS

Algorithmic Problem Solving

Linear Data Structures

Instructor: Joanna Klukowska

created based on materials for this class by
Bowen Yu and materials shared by the authors of
the textbook Steven and Felix Halim

Questions

- homework 1 questions ?

Basic Linear Data Structures

- list
 - array
 - linked
- stack
- queue

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Operations that determine performance of a data structure:

- insertion
- deletion
- query / find
- update / modify

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Most of the linear data structures are implemented in built-in libraries.

Lists

List as a Static Array

static array = fixed size, no need to resize it

- natively supported by both C/C++ and Java
- can be declared with appropriate size up-front if the problem specifies the maximum input size (HINT: use extra buffer for safety to avoid going out of bounds)
- can be multi-dimensional: 1D, 2D, 3D, ...

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- performance of operations
 - insert/delete in the back $O(1)$
 - insert/delete in the front or *middle* $O(N)$
 - update/modify $O(???)$
 - find $O(???)$

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 - update/modify: using index $O(1)$, if need to find, then see below
 - find $O(N)$ if not sorted, $O(\log N)$ if sorted

List as a Static Array

Java 1D and 2D static array

C/C++

List as a Dynamic/Resizable Array

- implementation provided by built-in classes:
 - `std::list` in C++ STL
 - `ArrayList` or `LinkedList` in Java (`ArrayList` is faster because it is unsynchronized)
- used when required size is not known at compile time

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performance of operations

use **amortized analysis**: average/amortized performance of a sequence of operations rather than each single operation
reason ???

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List as a Linked Structure (a Linked List)

- implementation provided by built-in classes:
 - in C++ STL
 - in Java
- rarely used due to poor performance of accessing elements
- (good exercise: implement your own in both Java and C++ to practice reference/pointer operations)
- performance of operations
 - insert/delete in the back/front $O(1)$ (assume a doubly linked)
 - insert/delete in the *middle* $O(N)$
 - update/modify: $O(N)$ (except for in the front/back)
 - find $O(N)$

List *Hybrids*

- dynamically allocated list that consists of short fixed sized arrays that are connected into a linked list

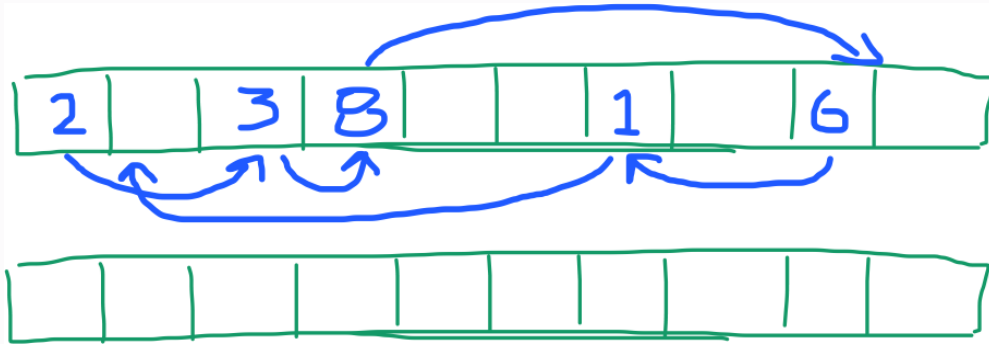


List *Hybrids*

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- linked list in which the "links" are provided by indexes in a separate array



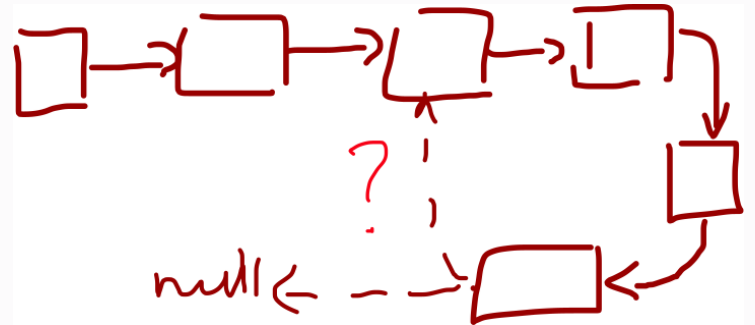
Challenge

Given a pointer/reference to a singly linked list, determine if it has a loop.

Determine the length of the loop if it exists.

Restrictions:

- Elements in the list are not unique.
- Do not use extra storage.



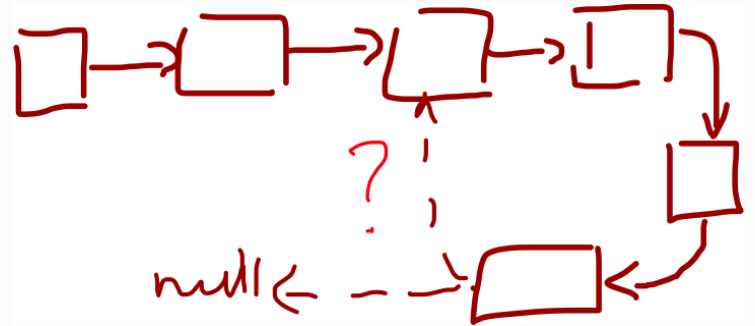
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Solution Rabbit and Turtle (two pointers)

- start two pointers at head
- advance rabbit two steps per iteration
- advance turtle one step per iteration
- if they meet (i.e., rabbit does not find the end of the list) then there is a loop and a note at which they meet is somewhere on the loop
- trace through the loop with one of them to determine the number of nodes in the loop

Challenge

- Can a singly linked list be circular?
- Can a doubly linked list be circular?
- Can a doubly linked list have a loop?

Stacks

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 - remove/pop $O(1)$ - removes from the top
 - top $O(1)$ - access the element on the top (optional)
 - empty $O(1)$ - determine if the stack is empty (optional)

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 - top $O(1)$ - access the element on the top (optional)
 - empty $O(1)$ - determine if the stack is empty (optional)
- used in many algorithms for solving problems
 - postfix, prefix calculations and conversions
 - graph algorithms

Challenge: Function Call Stack

What is the output of `foo()`?

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function foo() {
  console.log(1);
  foo();
  console.log(2);
}
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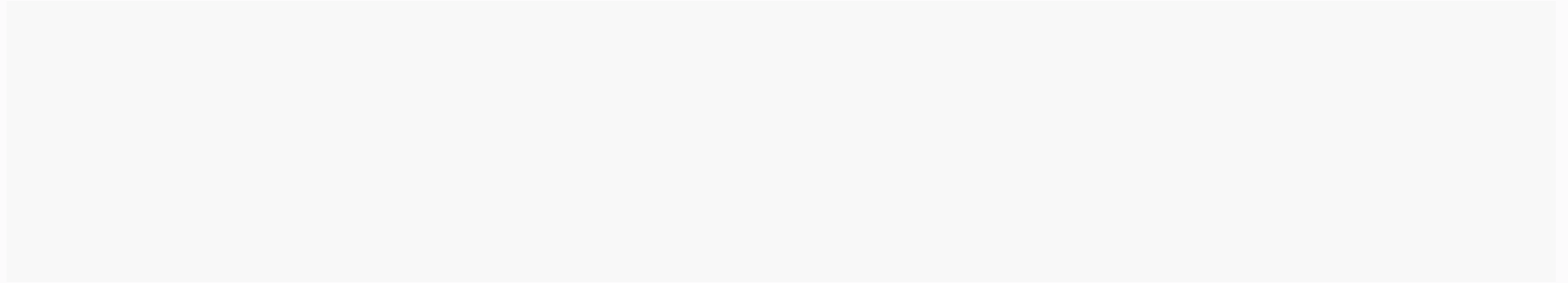
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def foo():  
    print(1)  
    bar()  
    print(2)  
  
def bar():  
    print(3)  
    foo()  
    print(4)  
  
foo()
```

Output

How about these functions? (Try it on your own after the class.)

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def foo():  
    print(1)  
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