

# Introduction to Computation Theory

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November 3, 2025

## **Abstract**

The lecture note of 2025 Fall Introduction to Computation Theory by professor 林智仁.

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# Chapter 0

## Basic Knowledge

### Lecture 1

#### 0.1 Mathematical Notions

2025-09-01

##### 0.1.1 Set & its operation

**Definition 0.1.1 (Set).** Omitted

**Definition (Sequence & Tuple).** Here are some definitions of basic containers

**Definition 0.1.2 (Sequence).** Sequence is the objects in order, which have two properties:

- Order:

$$(1, 2, 3) \neq (2, 1, 3)$$

- Repetition:

$$\text{Sequence : } (1, 2, 3) \neq (1, 1, 2, 3)$$

$$\text{Set : } \{1, 2, 3\} = \{1, 1, 2, 3\}$$

**Definition 0.1.3 (Tuple).** Finite sequence,  $(1, 2, 3)$  is a 3-tuple

**Definition 0.1.4 (Cartesian Product).** Here is the Cartesian Product between two sets. We define

$$A = \{1, 2\}, B = \{x, y\}$$

then,

$$A \times B = \{(1, x), (1, y), (2, x), (2, y)\}$$

### 0.1.2 Function & Relation

**Definition 0.1.5 (Function).** Function is a machine with single output.

**Definition (Equivalence Relations).** Here are the properties of Equivalence Relations.

**Definition 0.1.6 (reflexive).**

$$\forall x, xRx$$

**Definition 0.1.7 (symmetric).**

$$\forall x, y, xRy \iff yRx$$

**Definition 0.1.8 (transitive).**

$$xRy, yRz \implies xRz$$

**Example.**

$$i \equiv_7 j, \text{ if } 0 = i - j \pmod{7}$$

- Reflexive

$$i - i = 0 \pmod{7}$$

- Symmetric

$$i - j = 7a, j - i = -7a$$

- Transitive

$$i - j = 7a, j - k = 7b \implies i - k = 7(a + b)$$

### 0.1.3 String & Languages

**Definition (String & Languages).** Here is the definition of Language.

**Example (Alphabet).**

$$\{0, 1\}$$

**Example (String).**

$$01000$$

**Definition 0.1.9 (Language).** Set of Strings

$$L(A)$$

is the language of  $A$

## 0.2 Definitions, Theorems, and Proofs

- **Definition:** Introduce new concept.
- **Statement:** A sentence that is either true or false.
- **Theorem:** A statement that is true.
  - **Lemma:** A “helping” theorem.
  - **Corollary:** A theorem that follows easily from another theorem.

### 0.2.1 Proof by Construction

**Proposition 0.2.1.** Sum of degrees of every graph is even

**Proof.** Each edge contributes 2 nodes, so

$$\sum_{v \in V} \deg(v) = 2 \times |E|$$

Hence, the sum of degrees of every graph is even. ■

**Note.** The implication is the definition of graphs.

### 0.2.2 Proof by Contradiction

Assume the statement is false, then deduce a contradiction.

### 0.2.3 Proof by Induction

- Basis: Prove for  $n = 0$  or  $n = 1$  or some trivial case.
- Inductive Step: Assume true for  $n = k$  (Induction Hypothesis), prove for  $n = k + 1$ .

# Chapter 1

## Regular Languages

### 1.1 Deterministic Finite Automata (DFA)

- Automaton: single
- Automata: plural

**Definition 1.1.1** (Deterministic Finite Automata (DFA)). We define a DFA as a 5-tuple

$$(Q, \Sigma, \delta, q_0, F)$$

where

- $Q$ : Set of states (**Finite**)
- $\Sigma$ : Alphabet (i.e. set of input characters) (**Finite**)
- $\delta: Q \times \Sigma \rightarrow Q$ : Transition Function
- $q_0 \in Q$ : Start state
- $F \subset Q$ : Set of accept states

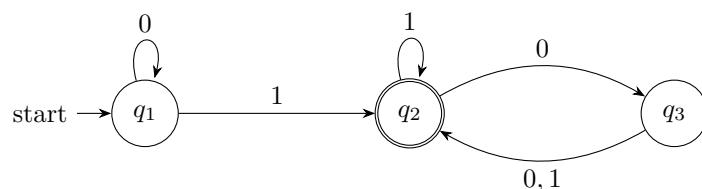


Figure 1.1: A state diagram

If we call this machine  $M$ , then we have.

$$M = (Q, \Sigma, \delta, q_0, F)$$

For the example given above,

$$Q = \{q_1, q_2, q_3\}$$

$$\Sigma = \{0, 1\}$$

$$q_0 = q_1$$

$$F = \{q_2\}$$

The  $\delta$  function:

	0	1
$q_1$	$q_1$	$q_2$
$q_2$	$q_3$	$q_2$
$q_3$	$q_2$	$q_2$

**Definition 1.1.2.** The language that recognize by a Machine  $M$  is denoted as

$$L(M) = A$$

We say  $A$  is recognized (accepted) by  $M$ .

### 1.1.1 Definition of Computation

Let,

- $M = (Q, \Sigma, \delta, q_0, F)$  be a finite automaton.
- $w = w_1, \dots, w_n$  be a string over  $\Sigma$ .

**Theorem 1.1.1.**  $M$  accepts  $w$  if  $\exists$  states  $r_0 \dots r_n$  such that

- (1)  $r_0 = q_0$
- (2)  $r_{i+1} = \delta(r_i, w_{i+1}), \quad i = [0, n - 1]$
- (3)  $r_n \in F$

**Definition 1.1.3 (Regular Language).** A language is regular if recognized by some automata.

### 1.1.2 Regular Operations

**Definition.** Assume  $A, B$  are given languages,

**Definition 1.1.4 (Union).**

$$A \cup B = \{w \mid w \in A \vee w \in B\}$$

**Definition 1.1.5 (Concatenation).**

$$A \circ B = \{w_1 w_2 \mid w_1 \in A, w_2 \in B\}$$

**Definition 1.1.6 (Kleene Star).**

$$A^* = \{w_1 \cdots w_k \mid k \geq 0, w_i \in A\}$$

which can also be defined as

$$\bigcup_{i=1}^{\infty} A_i = \{\varepsilon\} \cup A \cup A^2 \cup A^3 \cup \dots, \quad A^0 = \{\varepsilon\}, \quad A^n = \{wv \mid w \in A^{n-1}, v \in A\}$$

**Definition 1.1.7 (closed).** We say an operation  $R$  is closed if the following property holds if

$$x \in A, y \in A, \text{ then } xRy \in A$$

**Theorem 1.1.2.** Regular languages are closed under the union, concatenation, and Kleene star.

**Proof.** We define two machines as follows

$$M_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$$

$$M_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$$

if we union them, we can define a new machine

$$M_1 \cup M_2 = \begin{cases} M = (Q, \Sigma, \delta, q_0, F) \\ Q = \{(r_1, r_2) \mid r_1 \in Q_1, r_2 \in Q_2\} \\ \delta((r_1, r_2), a) = (\delta_1(r_1, a), \delta_2(r_2, a)) \\ q_0 = (q_1, q_2) \\ F = \{(r_1, r_2) \mid r_1 \in F_1 \text{ or } r_2 \in F_2\} \end{cases}$$

Hence, regular languages are closed under union. ■

## Lecture 2

### 1.2 Nondeterministic Finite Automata (NFA)

2025-09-08

First, we see a NFA that accept strings with 1 in 3rd position from the end,

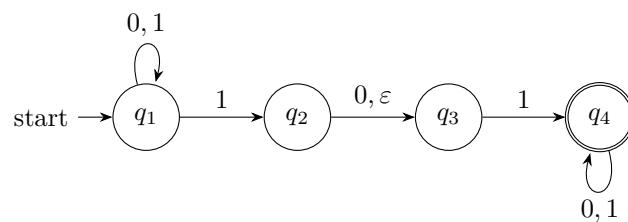


Figure 1.2: NFA machine

- $\delta$  is not a function, i.e.  $\delta(q_1, 1) = q_1$  or  $q_2$
- $\varepsilon$  between  $q_2, q_3$  means  $q_2$  can move to  $q_3$  without any input

We can transport NFA to DFA by some method, for example, for the above NFA we can have:

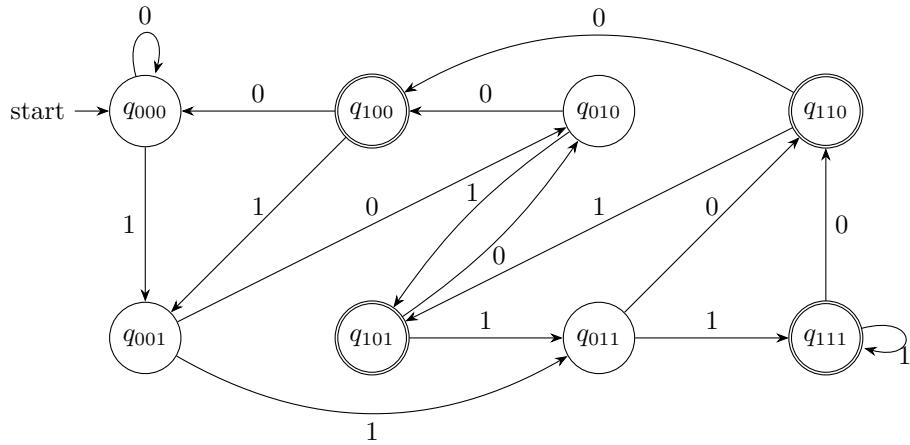


Figure 1.3: NFA machine transport to DFA

We can record it in three bits, it will be complicated.

**Definition 1.2.1 (power set).**

$$P(Q) = \{X | X \in Q\}$$

which contain all the  $2^{|Q|}$  combinations.

**Definition 1.2.2 (Nondeterministic Finite Automata (NFA)).** We define a NFA as a 5-tuple

$$M = (Q, \Sigma_\varepsilon, \delta, q_0, F)$$

where

- $Q$ : Set of states (**Finite**)
- $\Sigma_\varepsilon = \Sigma \cup \{\varepsilon\}$
- $\delta: Q \times \Sigma_\varepsilon \rightarrow P(Q)$
- $q_0 \in Q$
- $F \subset Q$

**Theorem 1.2.1.** We have  $w$

$$w = y_1 \cdots y_m \quad \text{where } y_i \in \Sigma_\varepsilon$$

A sequence  $r_0 \cdots r_m$  such that

- (1)  $r_0 = q_0$
- (2)  $r_{i+1} = \delta(r_i, y_{i+1}), \quad i = [0, n - 1]$
- (3)  $r_n \in F$

**Note.** So  $m$  may not be the original length (as  $y_i$  may be  $\varepsilon$ )

### 1.2.1 Equivalence of DFA and NFA

From DFA  $\Rightarrow$  NFA. Formally DFA is not an NFA due to  $\Sigma$  and  $\Sigma_\varepsilon$ . but we can easily handle this by adding

$$q_i, \varepsilon \rightarrow \emptyset$$

For  $\text{NFA} \Rightarrow \text{DFA}$ , we have the example on the slides on a graph.

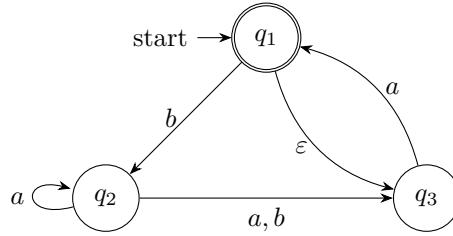


Figure 1.4: NFA example

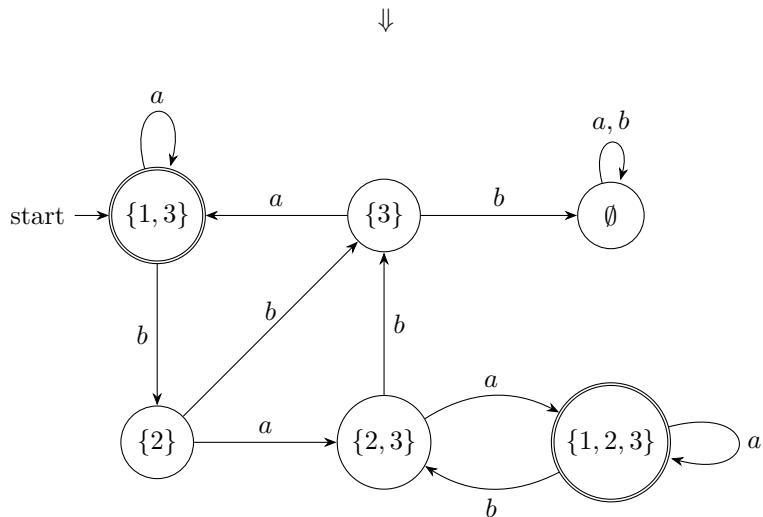


Figure 1.5: DFA conversion example

- Remove the states that are not reachable.
  - Remove the states that do not handle the  $\varepsilon$  transition. For example, the start state

$\{q_1\}$  wrong  $\rightarrow$   $\{q_1, q_3\}$  correct

### Definition 1.2.3.

$$E(\{q_0\}) = \{q_0\} \cup \{\text{states reached by } \varepsilon \text{ from } q_0\}$$

Then we can redefine the procedure formally.

**Theorem 1.2.2.** Given a NFA

$$M = (Q, \Sigma, \delta, q_0, F)$$

We can convert it to a DFA

$$M' = (Q', \Sigma, \delta', q'_0, F')$$

where

- $Q' = P(Q)$
- $q'_0 \in P(Q) = E(\{q_0\})$
- $F' = \{R \mid R \in Q', R \cap F \neq \emptyset\}$
- $\delta'$ :

$$\delta'(R, a) = \bigcup_{r \in R} E(\delta(r, a))$$

### 1.2.2 Closure under regular operations

We give two NFAs  $N_1, N_2$ ,

$$N_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$$

$$N_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$$

note that  $\varepsilon \notin \Sigma$ , and the graph of them are:

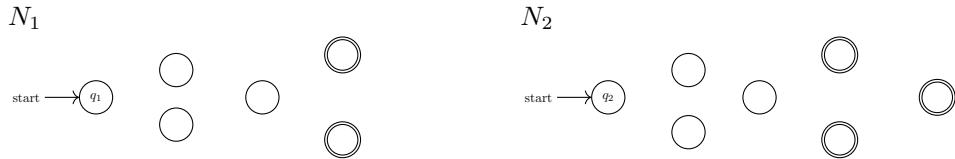


Figure 1.6:  $N_1, N_2$

- **Union:** We can construct the  $N_1 \cup N_2$  in

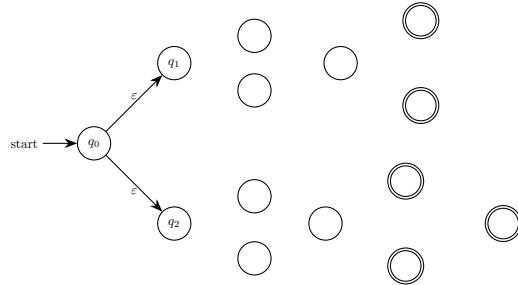


Figure 1.7:  $N_1 \cup N_2$

**Proposition 1.2.1** (Construction of Union). New NFA is

$$N_1 \cup N_2 = (Q, \Sigma, \delta, q_0, F)$$

where

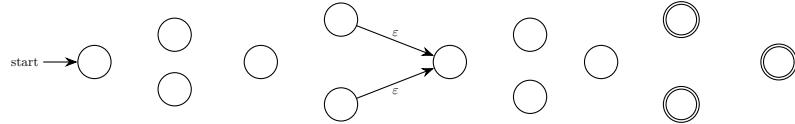
- $Q = Q_1 \cup Q_2 \cup \{q_0\}$

- $\delta$  :

$$\delta(q, a) = \begin{cases} \delta_1(q, a) & q \in Q_1 \\ \delta_2(q, a) & q \in Q_2 \\ \{q_1, q_2\} & q = q_0, a = \varepsilon \\ \emptyset & q = q_0, a \neq \varepsilon \end{cases}$$

- $F = F_1 \cup F_2$

- **Concatenation:** We can construct the  $N_1 \circ N_2$  in

Figure 1.8:  $N_1 \circ N_2$ 

**Proposition 1.2.2 (Construction of Concatenation).** New NFA is

$$N_1 \circ N_2 = (Q, \Sigma, \delta, q_0, F)$$

where

- $Q = Q_1 \cup Q_2$

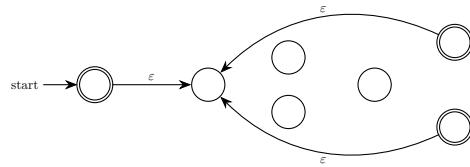
- $\delta$  :

$$\delta(q, a) = \begin{cases} \delta_1(q, a) & q \in Q_1, F_1 \\ \delta_2(q, a) & q \in Q_2 \\ \delta_1(q, \varepsilon) \cup \{q_2\} & q \in F_1, a = \varepsilon \\ \delta_1(q, \varepsilon) & q \in F_1, a \neq \varepsilon \end{cases}$$

- $q_0 = q_1$

- $F = F_2$

- **Kleene star:**  $N_1^*$  can also accept  $\{\emptyset\}$ , then we can construct the  $N_1^*$  in

Figure 1.9:  $N_1^*$ 

**Proposition 1.2.3 (Construction of Kleene Star).** New NFA is

$$N_1^* = (Q_1, \Sigma, \delta_1, q_0, F_1)$$

where

- $Q = Q_1 \cup \{q_0\}$

o  $\delta$ :

$$\delta(q, a) = \begin{cases} \delta_1(q, a) & q \in Q_1, F_1 \\ \delta_1(q, a) \cup \{q_1\} & q \in F_1, a = \varepsilon \\ \delta_1(q, \varepsilon) & q \in F_1, a \neq \varepsilon \\ \{q_1\} & q = q_0, a = \varepsilon \\ \emptyset & q = q_0, a \neq \varepsilon \end{cases}$$

o  $F = F_1 \cup \{q_0\}$

**Note.** Some operations are also closed under regular languages,

o **Intersection:**

$$A_1 \cap A_2$$

Use the product automaton (the same construction as for Union). A string is accepted if and only if the state is in the accept states of both  $N_1$  and  $N_2$  at the same time.

o **Set Difference:**

$$A_1 - A_2$$

Use the product automaton as well. A string is accepted if the state is in the accept states of  $N_1$  but *not* in the accept states of  $N_2$ .

o **Complement:**

$$A_1^c = \Sigma^* - A_1$$

Since  $\Sigma^*$  is regular and the class of regular languages is closed under set difference,  $A_1^c$  is also regular.

## Lecture 3

### 1.3 Regular expressions

2025-09-15

A regular expression is a tool to describe a language.

**Definition 1.3.1** (Regular expressions).  $R$  is a regular expression if it is one of the following expressions:

- (1)  $a$ , where  $a \in \Sigma$
- (2)  $\varepsilon$  ( $\varepsilon \notin \Sigma$ )
- (3)  $\emptyset$
- (4)  $R_1 \cup R_2$ , where  $R_1, R_2$  are regular expressions
- (5)  $R_1 \circ R_2$ , where  $R_1, R_2$  are regular expressions
- (6)  $R_1^*$ , where  $R_1$  is a regular expression

If there is no parentheses, we follow the order of:

Kleene star → Concatenation → Union

**Remark.**

$$R^+ = RR^*, \quad R^+ \cup \{\varepsilon\} = R^*$$

For  $\emptyset$  and  $\varepsilon$ , we have

- $\varepsilon$ : empty string
- $\emptyset$ : empty language (language without any string)

$$(0 \cup \varepsilon)1^* = 01^* \cup 1^*$$

$$(0 \cup \emptyset)1^* = 01^*$$

$$\emptyset 1^* = 1^* \emptyset = \emptyset$$

**Example.** Here are some examples,

- Strings that start and end with the same symbol:

$$0\Sigma^*0 \cup 1\Sigma^*1 \cup 0 \cup 1$$

- $(\Sigma\Sigma)^*$ : strings with even length
- $R \cup \emptyset = R$
- $R \circ \varepsilon = R$
- $\emptyset^* = \{\varepsilon\}$

Floating point numbers can also be represented by regular expressions. For example,

$$(+ \cup - \cup \varepsilon)(DD^* \cup DD^*.D^* \cup D^*.DD^*), \text{ where } D = \{0, \dots, 9\}$$

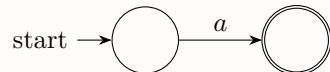
**Example.**

$$\begin{aligned} 72 &\in DD^* \\ 2.1 &\in DD^*.D^* \\ 7. &\in DD^*.D^* \\ .01 &\in D^*.DD^* \end{aligned}$$

**Lemma 1.3.1.** Language by a regular expression  $\implies$  Regular (described by an automaton)

**Proof.** The proof is by induction,

- $R = a \in \Sigma$  can be recognized by

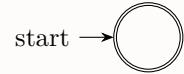


$$N = (\{q_1, q_2\}, \Sigma, \delta, q_1, \{q_2\})$$

$$\delta(q_1, a) = \{q_2\}$$

$$\delta(r, b) = \emptyset, r \neq q_1 \text{ or } b \neq a$$

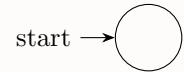
- $R = \epsilon$



$$N = (\{q_1\}, \Sigma, \delta, q_1, \{q_1\})$$

$$\delta(q_1, a) = \emptyset, \forall a$$

- $R = \emptyset$



$$N = (\{q\}, \Sigma, \delta, q, \emptyset)$$

$$\delta(r, a) = \emptyset, \forall r, a$$

- $R = R_1 \cup R_2, R = R_1 \circ R_2, R = R_1^*$  have proof by NFA.

■

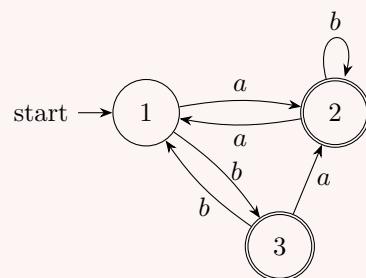
### 1.3.1 Convert a DFA to a regular expression

The idea is:

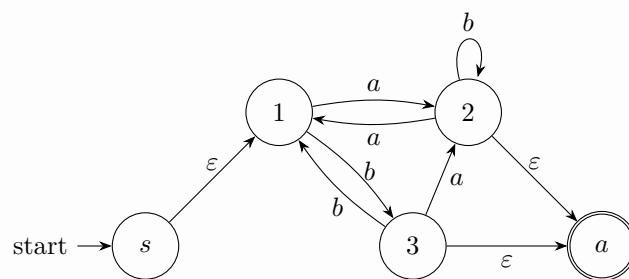
1° DFA  $\rightarrow$  GNFA

2° Remove states from GNFA until only the start and accept states.

**Question.** Convert the following DFA into regular expression.



**Answer.** First, convert to GNFA:



Next, is to remove the states one by one. We skip, so we can get the answer:

$$(a(aa \cup b)^*ab \cup b)((ba \cup a)(aa \cup b)^*ab \cup bb)^*((ba \cup a)(aa \cup b)^* \cup \varepsilon) \cup a(aa \cup b)^*$$

which is very complicated. ⊗

**Definition 1.3.2 (Generalized NFA(GNFA)).** We define a GNFA as a 5-tuple

$$G = (Q, \Sigma, \delta, q_{start}, q_{accept})$$

where

- $F$  is not a se, but a single accept state  $q_{accept}$
- $\delta$  function is:

$$(Q - \{q_{accept}\}) \times (Q - \{q_{start}\}) \rightarrow R$$

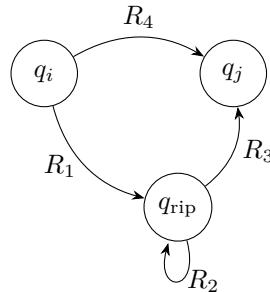
where  $R$  is all regular expressions over  $\Sigma$ .

- Two new states:

$$q_{start} \rightarrow q_0 \text{ with } \varepsilon$$

$$\text{any } q \in F \rightarrow q_{accept} \text{ with } \varepsilon$$

Consider  $q_{rip}$  is the state being removed



The new regular expression between  $q_i$  and  $q_j$  is

$$q_i \xrightarrow{(R_1)(R_2)^*(R_3) \cup (R_4)} q_j$$

We can write the whole process into an algorithm.

---

**Algorithm 1.1:** CONVERT( $G$ ) —State-Elimination from GNFA to RE

---

**Input:**  $G = (Q, \Sigma, \delta, q_s, q_a)$  a GNFA  
**Output:** A regular expression  $R$  for the language of  $G$

```

1  $k \leftarrow |Q|;$ 
2 ;                                     // number of states
3 if  $k = 2$  then
4   return  $\delta(q_s, q_a)$ ;           // the (single) edge label from  $q_s$  to  $q_a$ 
5 Choose any  $q_{rip} \in Q \setminus \{q_s, q_a\}$ ;
6  $Q' \leftarrow Q \setminus \{q_{rip}\}$ ;
7 Initialize  $\delta'$  as the restriction of  $\delta$  to  $Q' \times Q'$ ;
8 foreach  $q_i \in Q' \setminus \{q_a\}$  do
9   foreach  $q_j \in Q' \setminus \{q_s\}$  do
10     $R_1 \leftarrow \delta(q_i, q_{rip})$ ;
11     $R_2 \leftarrow \delta(q_{rip}, q_{rip})$ ;
12     $R_3 \leftarrow \delta(q_{rip}, q_j)$ ;
13     $R_4 \leftarrow \delta(q_i, q_j)$ ;
14     $\delta'(q_i, q_j) \leftarrow R_4 \cup (R_1 R_2^* R_3)$ ;
15  $G' \leftarrow (Q', \Sigma, \delta', q_s, q_a)$ ;
16 return CONVERT( $G'$ );

```

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## Lecture 4

### 1.4 Pumping lemma

2025-09-22

#### 1.4.1 Non regular language

Some languages cannot be recognized by DFA such as,

$$\{0^n 1^n \mid n \geq 0\}$$

We might remember #0 first, but # of possible  $n$ 's is  $\infty$ , so we have some method to prove that the language is non-regular.

**Theorem 1.4.1 (pumping lemma).** If  $A$  is regular,  $\exists p$  such that  $\forall s \in A, |s| \geq p$ ,

$\exists x, y, z$ , such that  $s = xyz$  and

$$1^\circ \quad \forall i \geq 0, xy^i z \in A$$

$$2^\circ \quad |y| > 0$$

$$3^\circ \quad |xy| \leq p$$

**Proof.** Skip, which is on the slides. ■

### 1.4.2 Example for Pumping Lemma

**Question.** Show that the language  $L = \{0^n 1^n \mid n \geq 0\}$  is not regular using the pumping lemma.

**Answer.** Now consider the string

$$s = 0^p 1^p$$

We know that  $|s| \geq p$ . By the lemma,  $s$  can be split into  $xyz$  such that

$$xy^i z \in B, \forall i \geq 0, |y| > 0, \text{ and } |xy| \leq p$$

1° If  $y = 0 \cdots 0$ , then

$$xy = 0 \cdots 0 \quad \text{and} \quad z = 0 \cdots 0 1 \cdots 1.$$

Thus,

$$xy^2 z : \#0 > \#1.$$

Hence  $xy^2 z \notin B$ , a contradiction.

2° If  $y = 1 \cdots 1$ , then similarly

$$xy^2 z \notin B \quad \text{as} \quad \#0 < \#1.$$

3° If  $y = 0 \cdots 0 1 \cdots 1$ , then

$$xy^2 z \notin B \quad \text{since it is not of the form } 0^* 1^*.$$

**Note.** Just pick one is sufficient to show the answer.

(\*)

**Question.** Show that the language  $C = \{w \mid \#0 = \#1\}$  is not regular using the pumping lemma.

**Answer.** We can use the situation in the previous example, consider

$$s = 0^p 1^p$$

We can't proof the third condition due to  $C = \{w \mid \#0 = \#1\}$  which just require the  $\#0 = \#1$ .

Then we can use the third condition

$$|xy| \leq p$$

which means  $y$  are strict into the first  $0^p$  we can only consider the first case.

$$|xy| \leq p \Rightarrow y = 0 \cdots 0 \text{ in } s = 0^p 1^p$$

Then,

$$xy^2 z \notin C$$

(\*)

**Lemma 1.4.1.** When using pumping lemma, we usually use contradiction, so we use

$$\forall p \exists s \in A, |s| \geq p, \left[ \forall x, y, z \left( (s = xyz \wedge |y| > 0 \wedge |xy| \leq p) \rightarrow \exists i \geq 0, xy^i z \notin A \right) \right].$$

Use the claim and the first, second condition to get the negation of the third condition.

**Question.**  $D = \{1^{n^2} \mid n \geq 0\}$  is not regular

**Answer.** We pick

$$s = 1^{p^2} \in D$$

Then, if  $s = xyz, |xy| \leq p, |y| > 0$ , we can get

$$p^2 < |xy^2z| \leq p^2 + p \leq (p+1)^2$$

hence,  $xy^2z \notin D$ .

⊗

# Chapter 2

## Context-Free Languages

### Lecture 5

#### 2.1 Context-Free Grammars (CFG)

2025-10-20

Which is more powerful, and can be used in compilers. A **Grammar** is a collection of substitution rules that describe the structure of a language.

**Example.** Consider a grammar  $G_1$ :

$$\begin{aligned}A &\rightarrow 0A1 \\A &\rightarrow B \\B &\rightarrow \#\end{aligned}$$

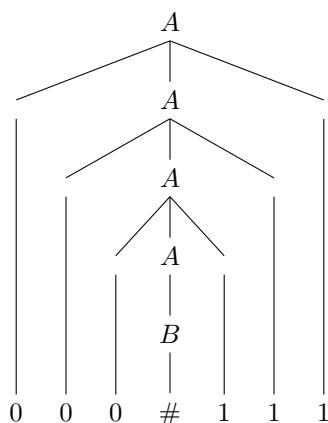
Here are the jargon terms:

- Each of one is called a **substitution rule**.
- **Variables** (non-terminals):  $A, B$  (Capital letters)
- **Terminals**:  $0, 1, \#$  (Lowercase letters, numbers, symbols)
- **Start variable**:  $A$  (the variable we start with)

The process of generating strings is called **derivation**.  $G_1$  generates  $000\#\underline{111}$  by

$$A \Rightarrow 0A1 \Rightarrow 00A11 \Rightarrow 000A111 \Rightarrow 000B111 \Rightarrow 000\#\underline{111}$$

We can show the derivation using a **parse tree**:



### 2.1.1 Definition of CFG

The language of grammar  $G$  is denoted by  $L(G)$ , for the language we discuss here,

$$L(G_1) = \{0^n \# 1^n \mid n \geq 0\}$$

Now we give the formal definition of CFG.

**Definition 2.1.1 (Context-Free Grammar).** We defined a CFG as a 4-tuple

$$G = (V, \Sigma, R, S)$$

where

- $V$ : Variables (Finite)
- $\Sigma$ : Terminals (Finite)
- $R$ : Rules:  

$$\text{Variables} \rightarrow \text{Strings of Variables and Terminals (including } \varepsilon\text{)}$$
- $S \in V$ : Start variable

For instance, for  $G_1$ ,

$$G_1 = (\{A, B\}, \{0, 1, \#\}, R, A)$$

where  $R$  is:

$$A \rightarrow 0A1 \mid B, \quad B \rightarrow \#$$

**Notation.** If  $u, v, w$  are strings and rule  $A \rightarrow w$  is applied, then we say

$$uAv \text{ yields } uwv$$

denoted as

$$uAv \Rightarrow uwv$$

**Notation.** If

$$u = v \text{ or } u \Rightarrow u_1 \Rightarrow \dots \Rightarrow u_k \Rightarrow v$$

then we write

$$v \xrightarrow{*} u$$

**Definition 2.1.2 (Language of a CFG).** The language generated by a CFG  $G$  with start variable  $S$  is

$$L(G) = \{w \in \Sigma^* \mid S \xrightarrow{*} w\}$$

### 2.1.2 Examples of CFGs

**Question.** Consider the grammar  $G_2 = (\{S\}, \{a, b\}, R, S)$ :

$$S \rightarrow aSb \mid SS \mid \varepsilon$$

What is  $L(G_2)$ ?

**Answer.** If we let  $a, b$  be the left and right parentheses respectively, then  $L(G_2)$  is the set of all balanced parentheses. (\*)

**Example.** Consider the grammar  $G_3 = (V, \Sigma, R, S)$  where

- $V = \{\langle \text{expr} \rangle, \langle \text{term} \rangle, \langle \text{factor} \rangle\}$
- $\Sigma = \{+, \times, (,), a\}$
- $R:$

$$\begin{aligned}\langle \text{expr} \rangle &\rightarrow \langle \text{term} \rangle + \langle \text{expr} \rangle \mid \langle \text{term} \rangle \\ \langle \text{term} \rangle &\rightarrow \langle \text{factor} \rangle \times \langle \text{term} \rangle \mid \langle \text{factor} \rangle \\ \langle \text{factor} \rangle &\rightarrow (\langle \text{expr} \rangle) \mid a\end{aligned}$$

Consider the string  $a + a \times a$ :

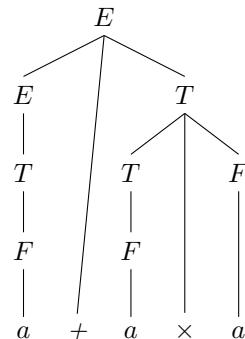


Figure 2.1: Parse tree of  $a + a \times a$

Consider the string  $(a + a) \times a$ :

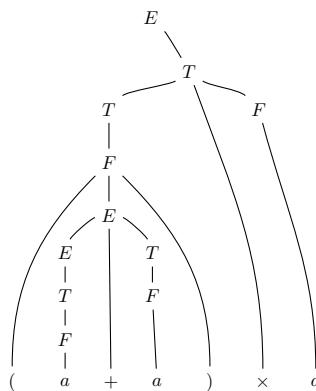


Figure 2.2: Parse tree of  $(a + a) \times a$

**Note.** The example above shows that CFGs can express operator precedence and associativity.

### 2.1.3 Design of CFGs

We can design CFGs in many methods. Here are some common patterns:

- Combining smaller parts:

**Example.**  $L(G) = \{a^n b^n \mid n \geq 0\} \cup \{b^n a^n \mid n \geq 0\}$

We can let the rule  $R$  be:

$$\begin{aligned} S_1 &\rightarrow aS_1b \mid \varepsilon \\ S_2 &\rightarrow bS_2a \mid \varepsilon \\ S &\rightarrow S_1 \mid S_2 \end{aligned}$$

- From DFA:

**Lemma 2.1.1.** For any regular language  $A$ , there exists a CFG  $G$  such that  $L(G) = A$ . The rules of CFG can be

$$\begin{aligned} R_i &\rightarrow aR_j \quad \text{for each transition } \delta(q_i, a) = q_j \\ R_i &\rightarrow \varepsilon \quad \text{if } q_i \in F \end{aligned}$$

The difference is that CFG allows the format

$$R_i \rightarrow aR_jb$$

But DFA only allows

$$R_i \rightarrow aR_j$$

where we treat  $R_i$  as the state and let  $\delta(R_i, a) = R_j$ .

### 2.1.4 Parse Trees and Ambiguity

If we let the rules of  $G_3$  be

$$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{expr} \rangle \mid \langle \text{expr} \rangle \times \langle \text{expr} \rangle \mid (\langle \text{expr} \rangle) \mid a$$

We can see the following two parse trees for  $a + a \times a$ :

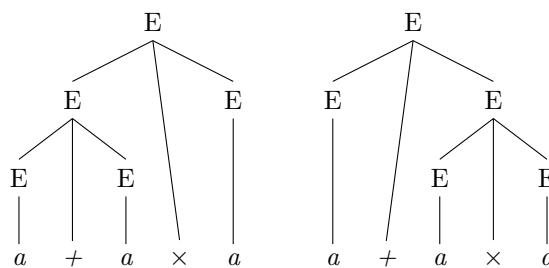


Figure 2.3: Two different parse trees for  $a + a \times a$  under ambiguous grammar

This is called **ambiguity**. A CFG is **ambiguous** if there exists some string with two or more different parse trees. The above  $G_3$  is **unambiguous**,  $G'_3$  with new rules is **ambiguous**.

However, an unambiguous grammar may also generate same parse tree but different derivations. Consider  $G_3$ :

- We can do derivation

$$\begin{aligned}\langle \text{expr} \rangle &\Rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle \\ &\Rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle \times \langle \text{factor} \rangle\end{aligned}$$

- We can also do derivation

$$\begin{aligned}\langle \text{expr} \rangle &\Rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle \\ &\Rightarrow \langle \text{term} \rangle + \langle \text{term} \rangle\end{aligned}$$

which is not considered ambiguous. So we have the following definition:

**Definition 2.1.3 (leftmost derivation).** A **leftmost derivation** is a derivation where at each step, the leftmost variable is replaced.

Then we can have the formal definition of ambiguity:

**Definition 2.1.4 (Ambiguous).**  $A$  is **ambiguous** if  $w \in A$  and there exists two or more different leftmost derivations for  $w$ .

**Definition 2.1.5 (Inherent Ambiguity).** A language is **inherently ambiguous** if it only has ambiguous grammars.

**Example.** Consider the language

$$L = \{a^i b^j c^k \mid i = j \text{ or } j = k\}$$

We can consider the string  $a^2b^2c^2$ . It can be generated by two different leftmost derivations. First we consider

$$S \Rightarrow S_1 \mid S_2$$

- Using  $i = j$ :

$$\begin{aligned}S_1 &\rightarrow AC \\ A &\rightarrow aAb \mid \epsilon \\ C &\rightarrow cC \mid \epsilon\end{aligned}$$

the derivation is

$$S_1 \Rightarrow AC \Rightarrow aAbC \Rightarrow aaAbbC \Rightarrow aabbC \Rightarrow aabbcC \Rightarrow aabbcc$$

- Using  $j = k$ :

$$\begin{aligned}S_2 &\rightarrow A'C' \\ A' &\rightarrow aA' \mid \epsilon \\ C' &\rightarrow bC'c \mid \epsilon\end{aligned}$$

the derivation is

$$S_2 \Rightarrow A'C' \Rightarrow aA'C' \Rightarrow aaA'bC'c \Rightarrow aabbC'cc \Rightarrow aabbcc$$

## Lecture 6

### 2.2 Chomsky Normal Form

2025-10-27

We want to simplify the structure of context-free grammars. One useful normal form is the Chomsky Normal Form (CNF).

**Definition 2.2.1 (Chomsky Normal Form).** A context-free grammar is in **Chomsky Normal Form** if all its production rules are of the form:

- $A \rightarrow BC$ , where  $A, B, C$  are non-terminal symbols and  $B, C$  are not the start symbol.
- $A \rightarrow a$ , where  $a \in \Sigma (\varepsilon \notin \Sigma)$
- $S \rightarrow \varepsilon$  is allowed, where  $S$  is the start symbol.

**Example.** Convert the following CFG to CNF:

$$\begin{aligned} S &\rightarrow ASA \mid aB \\ A &\rightarrow B \mid S \\ B &\rightarrow b \mid \varepsilon \end{aligned}$$

First, we add  $S_0$  as the new start symbol:

$$S_0 \rightarrow S \quad S \rightarrow ASA \mid aB \quad A \rightarrow B \mid S \quad B \rightarrow b \mid \varepsilon$$

Next, we remove the  $\varepsilon$ -productions  $B \rightarrow \varepsilon$ :

$$S_0 \rightarrow S \quad S \rightarrow ASA \mid aB \mid \textcolor{red}{a} \quad A \rightarrow B \mid \textcolor{red}{\varepsilon} \mid S \quad B \rightarrow b$$

Next, we remove the  $\varepsilon$ -productions  $A \rightarrow \varepsilon$ :

$$S_0 \rightarrow S \quad S \rightarrow ASA \mid aB \mid a \mid \textcolor{red}{AS} \mid \textcolor{red}{SA} \mid \textcolor{red}{S} \quad A \rightarrow B \mid S \quad B \rightarrow b$$

Next, we remove single production  $S \rightarrow S$ :

$$S_0 \rightarrow S \quad S \rightarrow ASA \mid aB \mid a \mid AS \mid SA \quad A \rightarrow B \mid S \quad B \rightarrow b$$

Next, we remove single production  $S_0 \rightarrow S$ :

$$S_0 \rightarrow \textcolor{red}{ASA} \mid \textcolor{red}{aB} \mid \textcolor{red}{a} \mid \textcolor{red}{AS} \mid \textcolor{red}{SA} \quad S \rightarrow ASA \mid aB \mid a \mid AS \mid SA \quad A \rightarrow B \mid S \quad B \rightarrow b$$

Next, we remove single production  $A \rightarrow B, A \rightarrow S$ :

$$S_0 \rightarrow ASA \mid aB \mid a \mid AS \mid SA \quad S \rightarrow ASA \mid aB \mid a \mid AS \mid SA \quad A \rightarrow \textcolor{red}{b} \mid \textcolor{red}{ASA} \mid \textcolor{red}{aB} \mid \textcolor{red}{a} \mid \textcolor{red}{AS} \mid \textcolor{red}{SA} \quad B \rightarrow b$$

Finally, we convert to CNF by introducing new variables for terminals and breaking down long productions:

$$\begin{aligned} S_0 &\rightarrow AA_1 \mid UB \mid a \mid AS \mid SA \\ S &\rightarrow AA_1 \mid UB \mid a \mid AS \mid SA \\ A &\rightarrow b \mid AA_1 \mid UB \mid a \mid AS \mid SA \\ A_1 &\rightarrow SA \\ B &\rightarrow b \\ U &\rightarrow a \end{aligned}$$

### 2.2.1 Procedure of Converting CFG to CNF

To convert any CFG to CNF, we can follow these steps:

1° **Add** a new start symbol  $S_0$  with the production

$$S_0 \rightarrow S$$

2° **Remove** all  $\varepsilon$ -productions, except for the start symbol, i.e.  $A \rightarrow \varepsilon$  ( $A \neq S_0$ ), for any

$$\dots \rightarrow uAv$$

add the production

$$\dots \rightarrow uv$$

3° **Remove** single productions of  $A \rightarrow B$  where  $A, B \in V/\{S\}$ .

$$A \rightarrow B, B \rightarrow \gamma \Rightarrow A \rightarrow \gamma$$

**Remark.**  $A \rightarrow \gamma$  can't be a unit rule previously removed.

4° **Convert** remaining productions to CNF:

$$A \rightarrow u_1 u_2 \dots u_k \quad u_i \in V \cup \Sigma$$

and

$$\text{if } k = 1, \text{ then } u_i \in \Sigma$$

Convert as follows:

$$A \rightarrow u_1 A_1$$

$$A_1 \rightarrow u_2 A_2$$

⋮

Replaced every terminal  $u_i \in \Sigma$  with a new variable  $U_i$ :

$$U_i \rightarrow u_i \quad u_i \in \Sigma$$

### 2.2.2 Infinite Loop in Converting

**Example.** Consider the grammar:

$$S \rightarrow B \mid \varepsilon$$

$$B \rightarrow S \mid \varepsilon$$

We first add a new start symbol:

$$S_0 \rightarrow S \quad S \rightarrow B \mid \varepsilon \quad B \rightarrow S \mid \varepsilon$$

Next, we remove the  $\varepsilon$ -productions:

$$S_0 \rightarrow S \mid \varepsilon \quad S \rightarrow B \quad B \rightarrow S \mid \varepsilon$$

Next, we remove the  $\varepsilon$ -productions again:

$$S_0 \rightarrow S \mid \varepsilon \quad S \rightarrow B \mid \varepsilon \quad B \rightarrow S$$

This process will continue indefinitely. The reason is  $S \rightarrow \varepsilon$  has been handled. So there is no need to add  $S \rightarrow \varepsilon$ .

## 2.3 Pushdown Automata

We now introduce the machine that recognizes context-free languages (CFL), called Pushdown Automata (PDA). PDA is a machine with a **stack**, which is a way to store previous states.

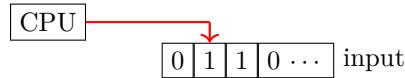


Figure 2.4: DFA or NFA

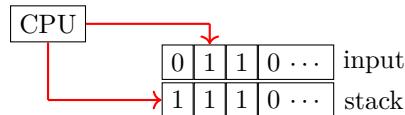


Figure 2.5: Pushdown Automata (PDA)

**Example.** Consider the language  $A = \{0^n 1^n \mid n \geq 0\}$ . We can design a PDA to recognize  $A$ :

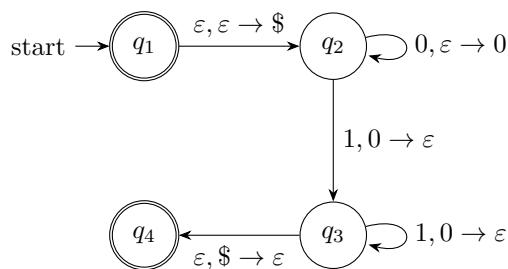


Figure 2.6: PDA for  $A = \{0^n 1^n \mid n \geq 0\}$

$\$$  is a special bottom stack symbol to indicate the initial state of the stack. The PDA works as follows:

- $q_2 \rightarrow q_2$ , put 0 into stack
- $q_2 \rightarrow q_3$  and  $q_3 \rightarrow q_3$ , read 1 and pop 0 up

If the input is 0011 which is same as  $\epsilon 0011 \epsilon$ , the process is as follows:

$q_1, \emptyset, \epsilon$   
 $q_2, \{\$\}, 0$   
 $q_2, \{0, \$\}, 0$   
 $q_2, \{0, 0, \$\}, 1$   
 $q_3, \{0, \$\}, 1$   
 $q_3, \{\$\}, \epsilon$   
 $q_4, \{\}$

**Notation.**  $\{\}$ : contents of the stack before processing the input character.

### 2.3.1 Formal definition of PDA

**Definition 2.3.1 (Pushdown Automata).** A **pushdown automaton** (PDA) is a 6-tuple

$$(Q, \Sigma, \Gamma, \delta, q_0, F)$$

, where

- $Q$ : States
- $\Sigma$ : Input alphabet
- $\Gamma$ : Stack alphabet
- $\delta$ : Transition function

$$Q \times \Sigma_\varepsilon \times \Gamma_\varepsilon \rightarrow \mathcal{P}(Q \times \Gamma_\varepsilon)$$

- $q_0 \in Q$ : Start state
- $F \subset Q$ : Set of accepting states

The definition of the above PDA for  $A = \{0^n 1^n \mid n \geq 0\}$  is as follows:

- $Q = \{q_1, q_2, q_3, q_4\}$
- $\Sigma = \{0, 1\}$
- $\Gamma = \{0, \$\}$
- $q_0 = q_1$
- $F = \{q_1, q_4\}$

For the transition function, we care about three things:

- Current state
- Current input
- **Top of the stack**

The transition function  $\delta$  works as follows:

	0	1	$\varepsilon$
0	$\$, \varepsilon$	$0, \$, \varepsilon$	$0, \$, \varepsilon$
$q_1$			$\{(q_2, \$)\}$
$q_2$	$\{(q_2, 0)\}$	$\{(q_3, \varepsilon)\}$	
$q_3$		$\{(q_3, \varepsilon)\}$	
$q_4$			$\{(q_4, \varepsilon)\}$

For example, we say the transition of  $q_2 \rightarrow q_3$  to be

$$\delta(q_2, 1, 0) = \{(q_3, \varepsilon)\}$$

### 2.3.2 Nondeterministic situation

**Example.** Design a PDA for the language  $B = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i = j \text{ or } j = k\}$ .

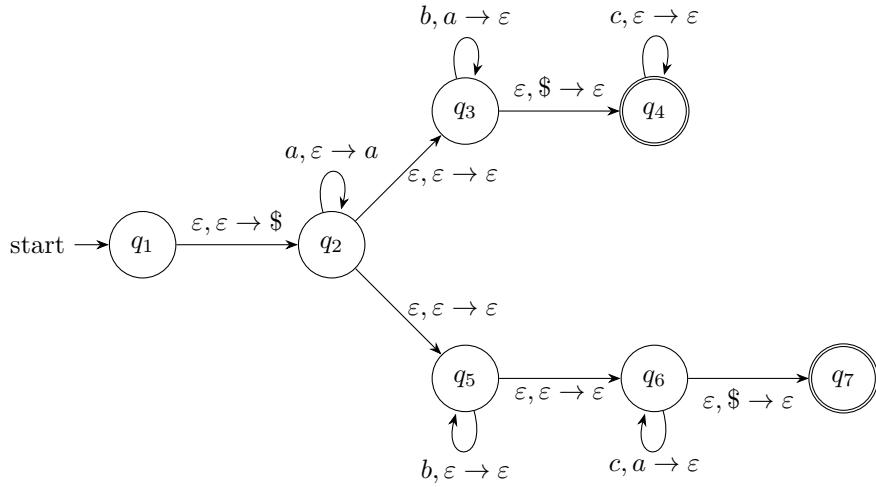
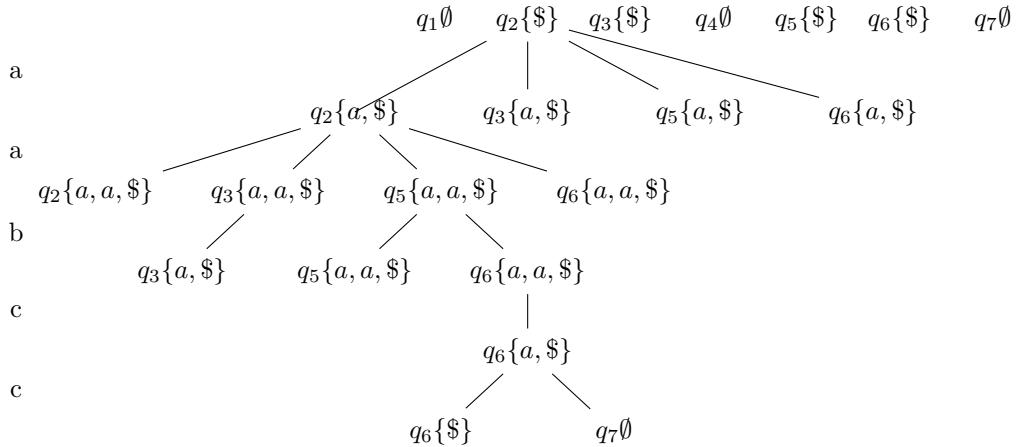


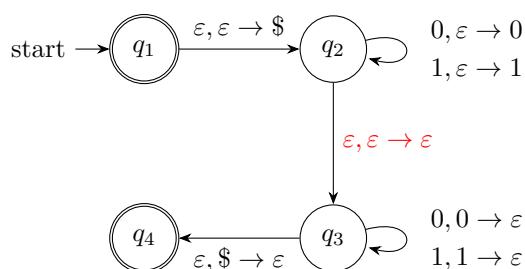
Figure 2.7: Nondeterministic PDA

We input  $a^2bc^2$ , to illustrate the process, we can build the following computation tree:



**Example.** Design a PDA for the language  $C = \{ww^R \mid w \in \{0, 1\}^*\}$ .

**Idea.** Symbols pushed to stack, nondeterministically guess middle is reached

Figure 2.8: PDA for  $C = \{ww^R \mid w \in \{0, 1\}^*\}$

### 2.3.3 Converting CFL to PDA

**Example.** Convert the CFG  $G$  to PDA that recognizes  $L(G)$ :

$$S \rightarrow aTb \mid b$$

$$T \rightarrow Ta \mid \varepsilon$$

**Idea.** For rule substitution, we replace the left-hand side variable with the right-hand side string i.e.

$$A \rightarrow \gamma \Rightarrow \text{pop } A \text{ from stack, push } \gamma \text{ to stack}$$

if there are multiple productions for  $A$ , we push them **in a reversed way**.

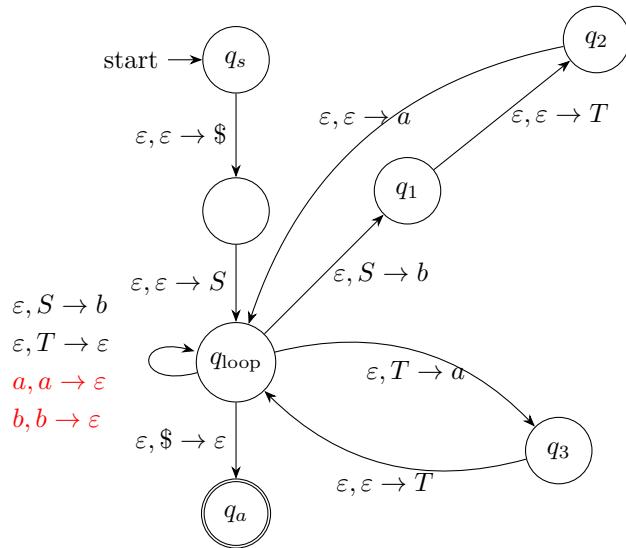


Figure 2.9: PDA for CFG  $G$

**Remark.** There are two transitions we must add to process the "input":

$$a, a \rightarrow \varepsilon$$

$$b, b \rightarrow \varepsilon$$

The procedure of converting CFG to PDA is as follows:

$$\begin{aligned} q_{\text{start}} &\xrightarrow{\varepsilon} q_{\text{loop}}, \{S, \$\} \xrightarrow{\varepsilon} q_1, \{b, \$\} \xrightarrow{\varepsilon} q_2, \{T, b, \$\} \\ &\xrightarrow{\varepsilon} q_{\text{loop}}, \{a, T, b, \$\} \xrightarrow{a} q_{\text{loop}}, \{T, b, \$\} \\ &\xrightarrow{\varepsilon} q_3, \{a, b, \$\} \xrightarrow{\varepsilon} q_{\text{loop}}, \{T, a, b, \$\} \\ &\xrightarrow{\varepsilon} q_3, \{a, a, b, \$\} \xrightarrow{\varepsilon} q_{\text{loop}}, \{T, a, a, b, \$\} \\ &\xrightarrow{\varepsilon} q_3, \{a, a, a, b, \$\} \xrightarrow{\varepsilon} q_{\text{loop}}, \{T, a, a, a, b, \$\} \\ &\xrightarrow{\varepsilon} q_{\text{loop}}, \{a, a, a, b, \$\} \xrightarrow{a} q_{\text{loop}}, \{a, a, b, \$\} \\ &\xrightarrow{a} q_{\text{loop}}, \{a, b, \$\} \xrightarrow{a} q_{\text{loop}}, \{b, \$\} \\ &\xrightarrow{b} q_{\text{loop}}, \{\$\} \xrightarrow{\varepsilon} q_{\text{accept}} \end{aligned}$$

**Proposition 2.3.1.** Even with a non-deterministic setting, we ensure that only strings generated by this CFG can be accepted by the PDA

- A string is accepted only if all characters are processed (this is part of the PDA definition!)
- We have  $\$$  to ensure that the stack is empty in the end

### 2.3.4 Converting PDA to CFL

**Lemma 2.3.1.** Language recognized by PDA  $\Rightarrow$  context free

**Note.** We need PDA to satisfy

- 1° Single start state
- 2° Stack empty before accepting
- 3° Each transition push or pop, but not both

**Idea.** For each pair of states  $p, q \in Q$  of a PDA  $P$ , we have  $A_{pq}$  and

$A_{pq}$  generates  $x \Rightarrow P$  from  $p$  with empty stack to  $q$  with empty stack, reading  $x$

First, we discuss how to handle transitions

$$\forall p, q, r \in Q, A_{pq} \rightarrow A_{pr}A_{rq}$$

We let the

- $x$ -axis: input string
- $y$ -axis: stack height

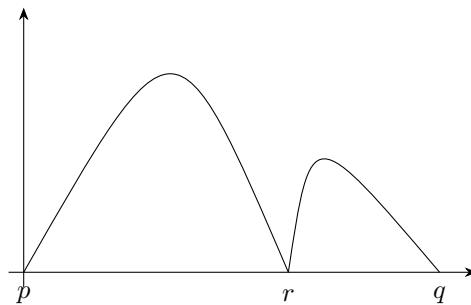


Figure 2.10: PDA transition  $A_{pq} \rightarrow A_{pr}A_{rq}$

If we can go

from  $p$  to  $r$  without changing stack

and

from  $r$  to  $q$  without changing stack

then we can do

from  $p$  to  $q$  without changing stack

Next, we have

$$\forall p, q, r, s \in Q, a, b \in \Sigma_\varepsilon, t \in \Gamma$$

If,

$$(r, t) \in \delta(p, a, \varepsilon) \text{ and } (q, \varepsilon) \in \delta(s, b, t)$$

we discuss how to handle transitions

$$A_{pq} \rightarrow aA_{rs}b$$

Then we have

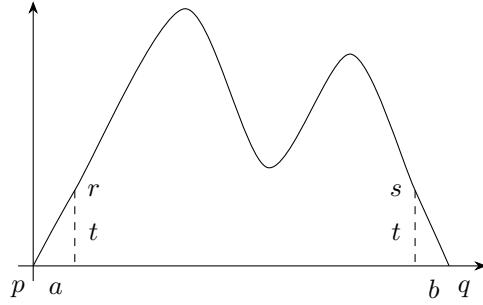


Figure 2.11: PDA transition  $A_{pq} \rightarrow aA_{rs}b$

Finally, we have the following base case:

$$\forall p \in Q, A_{pp} \rightarrow \varepsilon$$

To follow the condition (1°), we give a new example

**Example.** Consider the language  $L = \{0^n 1^n \mid n \geq 1\}$ .

Now  $q_1$  is not an accept state

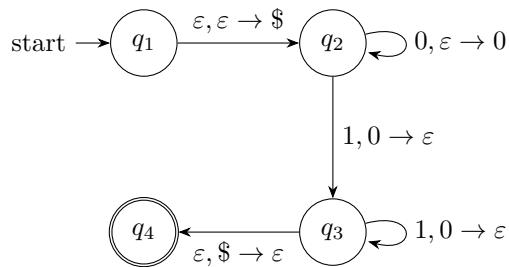


Figure 2.12: PDA for  $A = \{0^n 1^n \mid n \geq 1\}$

Consider two elements in  $\Gamma$

$$t_0 = \$, \quad t_1 = 0$$

- $t = \$$

p	r	s	q	t	a	b
1	2	3	4	\$	$\varepsilon$	$\varepsilon$

then we can get the rule

$$A_{14} \rightarrow A_{23}$$

- $t = 0$

p	r	s	q	t	a	b
2	2	2	3	0	0	1
2	2	3	3	0	0	1

then we can get the rules

$$\begin{aligned} A_{23} &\rightarrow 0A_{22}1 \\ A_{23} &\rightarrow 0A_{23}1 \end{aligned}$$

Other rules: 64 rules

$$\begin{aligned} A_{11} &\rightarrow A_{11}A_{11} \\ A_{11} &\rightarrow A_{12}A_{21} \\ A_{11} &\rightarrow A_{13}A_{31} \\ A_{11} &\rightarrow A_{14}A_{41} \\ &\vdots \end{aligned}$$

and

$$\begin{aligned} A_{11} &\rightarrow \varepsilon \\ A_{22} &\rightarrow \varepsilon \\ A_{33} &\rightarrow \varepsilon \\ A_{44} &\rightarrow \varepsilon \end{aligned}$$

### 2.3.5 Procedure of converting PDA to CFL

**Proposition 2.3.2.** Given a PDA

$$P = (Q, \Sigma, \Gamma, \delta, q_0, \{q_{accept}\})$$

We construct a CFG with variables

$$\text{var}(G) = \{A_{pq} \mid p, q \in Q\}$$

and start variable

$$S = A_{q_0 q_{accept}}$$

With rules

- 1° Single start state
- 2° Stack empty before accepting
- 3° Each transition push or pop, but not both

A new start  $q_s \rightarrow q_{s'}$  with  $\varepsilon, \varepsilon \rightarrow \$$ , and for any  $q \in F$ , we have  $\varepsilon, a \rightarrow \varepsilon$  back to  $q$ ,  $\forall a \in \Sigma$ . Then from any  $q \in F$ , we do  $\varepsilon, \$ \rightarrow \varepsilon$  to  $q_a$

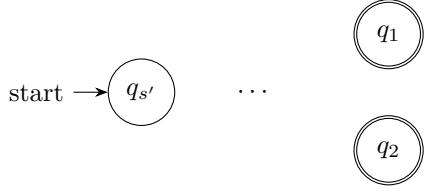
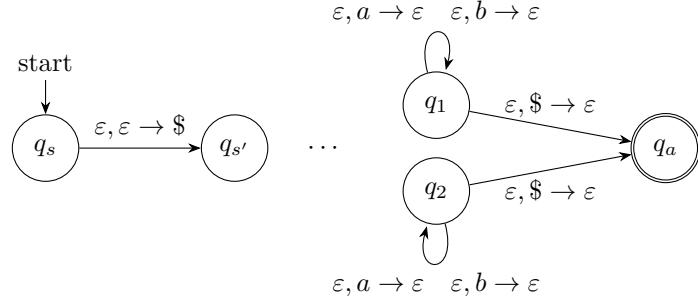


Figure 2.13: PDA with single accept state and empty stack before accepting

The new one will become



These is not enough to ensure condition (3°), we can do some modifications:

- To have each transition either push or pop (but not both), replace

$$q_1 \xrightarrow{a, a \rightarrow b} q_2$$

with the pair

$$q_1 \xrightarrow{a, a \rightarrow \varepsilon} q_3, \quad q_3 \xrightarrow{\varepsilon, \varepsilon \rightarrow b} q_2.$$

- Likewise, replace

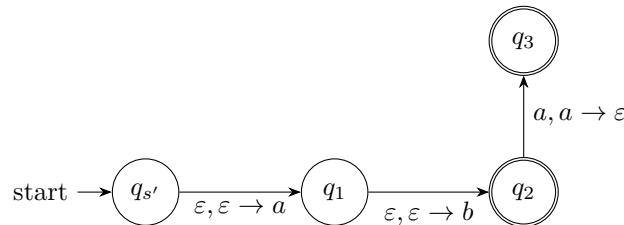
$$q_1 \xrightarrow{a, \varepsilon \rightarrow \varepsilon} q_2$$

with

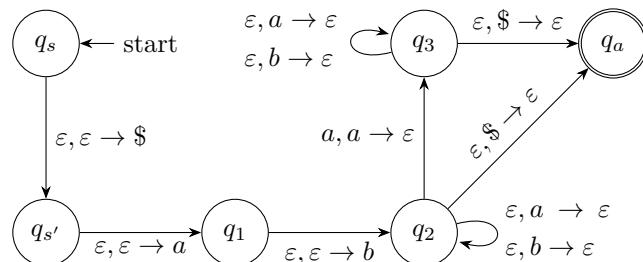
$$q_1 \xrightarrow{a, \varepsilon \rightarrow X} q_3, \quad q_3 \xrightarrow{\varepsilon, X \rightarrow \varepsilon} q_2,$$

where  $X$  is a fresh stack marker introduced for this simulation.

For another example, consider the PDA



After the modification, we have



The new PDA will accept the string  $a$  but the original PDA rejects it. Hence, we need to modify something else:

- A new start  $q_s \rightarrow q_{s'}$  with  $\varepsilon, \varepsilon \rightarrow \$$
- A **new state  $q_{\text{pop}}$**  that have  $\varepsilon, a \rightarrow \varepsilon$  back to  $q_{\text{pop}}$ ,  $\forall a$ .
- For  $q \in F$ , add a transition  $\varepsilon, \varepsilon \rightarrow \varepsilon$  from  $q$  to  $q_{\text{pop}}$
- Add a new accept state  $q_a$  and a transition  $\varepsilon, \$ \rightarrow \varepsilon$  from  $q_{\text{pop}}$  to  $q_a$

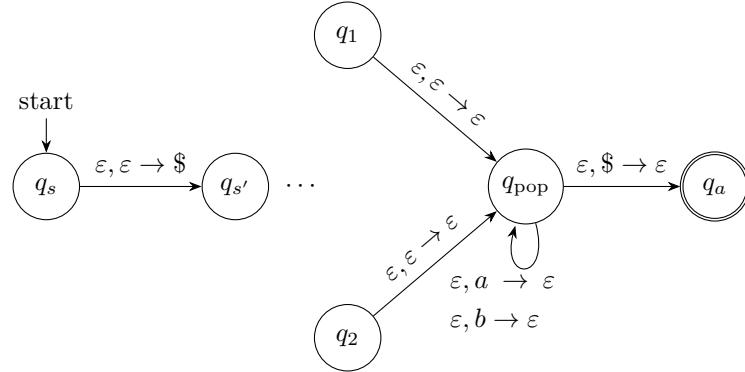


Figure 2.14: PDA with single accept state and empty stack before accepting

## 2.4 Deterministic Pushdown Automata

### Lecture 7

PDA is non-deterministic in general. However, there is a special class of PDA called **Deterministic** 2025-11-03 Pushdown Automata (DPDA). From Ch.1 we know

$$\text{DFA} \equiv \text{NFA}$$

but

$$\text{DPDA} \neq \text{PDA} \implies \text{CFL} \neq \text{DCFL}$$

**Definition 2.4.1** (Deterministic Pushdown Automaton (DPDA)). A deterministic pushdown automaton (DPDA) is a 6-tuple

$$M = (Q, \Sigma, \Gamma, \delta, q_0, F)$$

where

- $Q$ : States
- $\Sigma$ : Input alphabet
- $\Gamma$ : Stack alphabet
- $\delta$ : Transition function

$$Q \times \Sigma_\varepsilon \times \Gamma_\varepsilon \rightarrow (Q \times \Gamma_\varepsilon) \cup \{\emptyset\}$$

- $q_0 \in Q$ : Start state
- $F \subset Q$ : Set of accepting states

To build a DPDA, we first look at the difference between PDA and DPDA.

**As previously seen.** For PDA,

$$\delta : Q \times \Sigma_\varepsilon \times \Gamma_\varepsilon \rightarrow \mathcal{P}(Q \times \Gamma_\varepsilon)$$

**Note.** In DPDA, for  $\forall q \in Q, a \in \Sigma, x, \gamma \in \Gamma$ , at most and at least one of the following is true:

$$\delta(q, a, x) = (p, \gamma), \quad \delta(q, a, \varepsilon) = (p, \gamma), \quad \delta(q, \varepsilon, x) = (p, \gamma), \quad \delta(q, \varepsilon, \varepsilon) = (p, \gamma)$$

the rest must be  $\emptyset$ .

#### 2.4.1 Acceptance, Rejection of DPDA

The Rejection of DPDA is similar to PDA, which should only happen when

- Not end at an accept state after the last symbol.
- DPDA fails to read the input
  1. pop an empty stack
  2. Endless  $\varepsilon$ -transition

**Example.**  $L = \{0^n 1^n \mid n \geq 0\}$

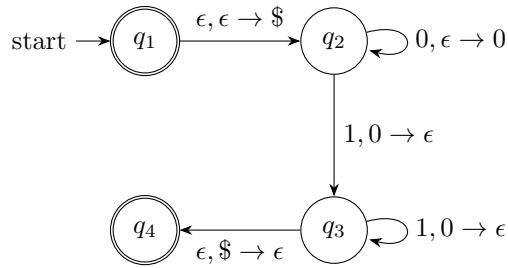


Figure 2.15: DPDA for  $L = \{0^n 1^n \mid n \geq 0\}$

The Transition function is defined as follows:

	0			1			$\epsilon$		
	0	\$	$\epsilon$	0	\$	$\epsilon$	0	\$	$\epsilon$
$q_1$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$	$(q_2, \$)$
$q_2$	$\emptyset$	$\emptyset$	$(q_2, 0)$	$(q_3, \epsilon)$	$q_r$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$
$q_3$	$q_r$	$\emptyset$	$\emptyset$	$(q_3, \epsilon)$	$\emptyset$	$\emptyset$	$\emptyset$	$(q_4, \epsilon)$	$\emptyset$
$q_4$	$q_r$	$q_r$	$\emptyset$	$q_r$	$q_r$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$
$q_r$	$q_r$	$q_r$	$\emptyset$	$q_r$	$q_r$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$

To find this transition table, for instance,

- consider the state  $q_1$ :

$$\delta(q_1, \epsilon, \epsilon) = (q_2, \$)$$

then we can implies that

$$\delta(q_1, a, \gamma) = \delta(q_1, a, \epsilon) = \delta(q_1, \epsilon, \gamma) = \emptyset, \quad \forall a \in \Sigma = \{0, 1\}, \gamma \in \Gamma = \{0, \$\}$$

- consider the state  $q_2$ :

$$\delta(q_2, 1, 0) = (q_3, \epsilon)$$

then we can implies that

	0			1			$\epsilon$		
	0	\$	$\epsilon$	0	\$	$\epsilon$	0	\$	$\epsilon$
$q_2$	$\emptyset$	$\emptyset$	$(q_2, 0)$	$(q_3, \epsilon)$	$\neq \emptyset$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$

due to

$$\delta(q_2, 1, \epsilon) = \delta(q_2, \epsilon, 0) = \delta(q_2, \epsilon, \epsilon) = \emptyset$$

Formally we have

$$\delta(q_2, 1, \$) = (q_r, \epsilon)$$

For the string 011, the computation of the DPDA is as follows:

$$q_1 \xrightarrow{\epsilon} q_2, \{\$\} \quad q_1 \xrightarrow{0} q_2, \{0, \$\} \quad q_1 \xrightarrow{1} q_3, \{\$\} \quad q_1 \xrightarrow{\epsilon} q_4, \emptyset$$

Follow the graph,

$$\delta(q_4, 1, \epsilon) \text{ and } \delta(q_4, \epsilon, \epsilon) = \emptyset$$

hence, the DPDA rejects 011.