

Distributed Algorithms 2020

Randomized coloring

 $(\Delta = maximum degree)$

Vertex coloring with Δ+1 colors

Week 4: deterministic, $O(\Delta + \log^* n)$

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Vertex coloring with Δ+1 colors

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Today: randomized, w.h.p. $O(\log n)$

Simplest possible idea:

everyone tries to pick
 a random free color

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- everyone tries to pick
 a random free color
- stop if successful

Pretty simple idea:

 nodes are active with probability 1/2

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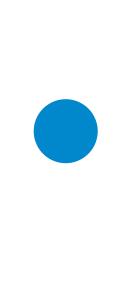
- nodes are active with probability 1/2
- only active nodes try to pick a random free color
- stop if successful

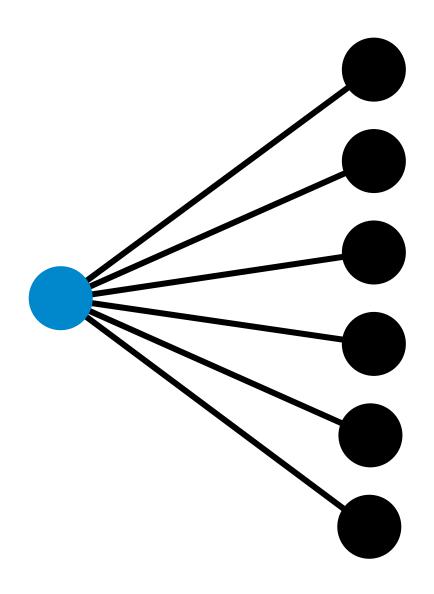
Corollary: The node is still running after T rounds with probability $\leq 0.75^{T}$

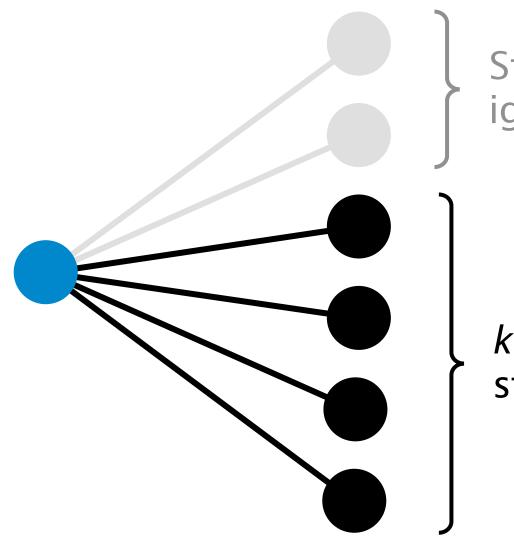
Corollary: The node will stop after $O(\log n)$ rounds w.h.p.

Corollary: All nodes will stop after O(log n) rounds w.h.p.

Proof...?

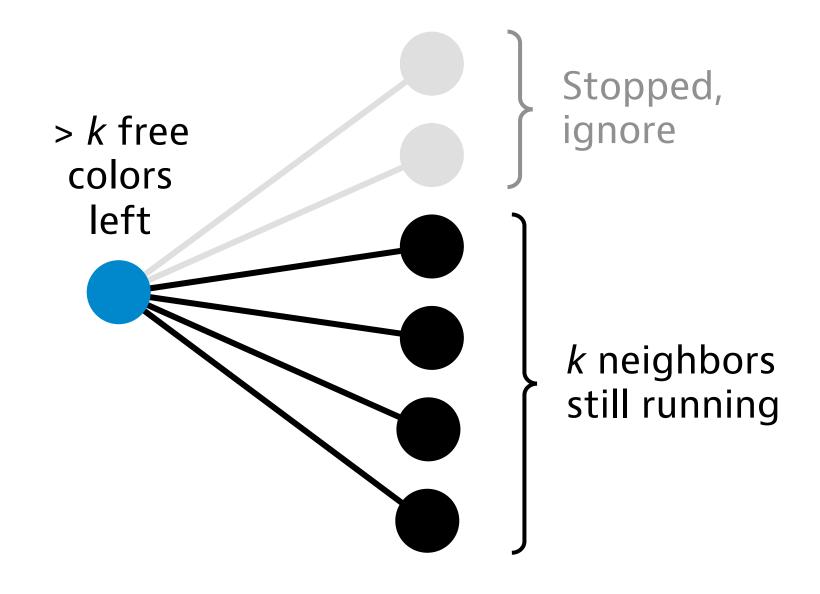


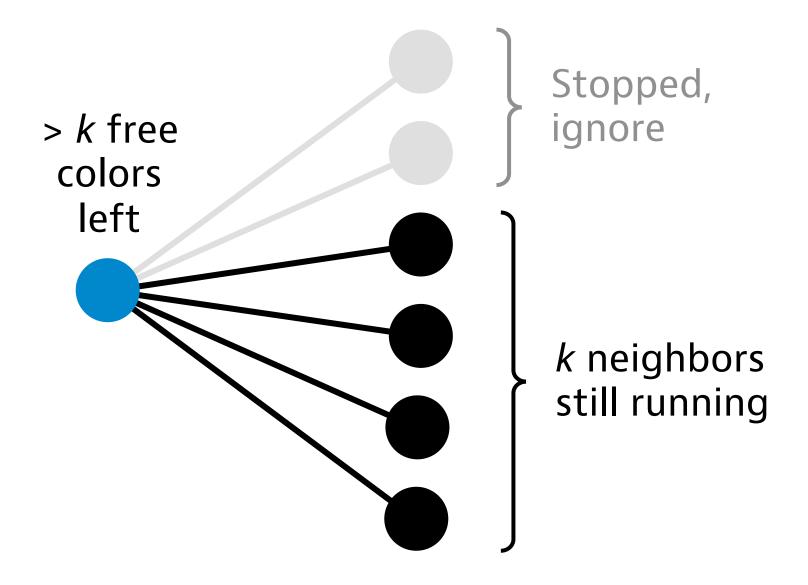


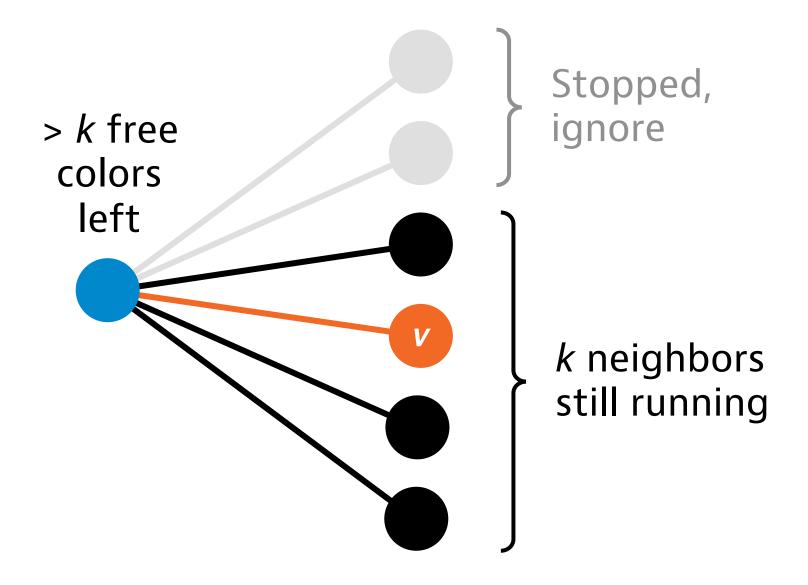


Stopped, ignore

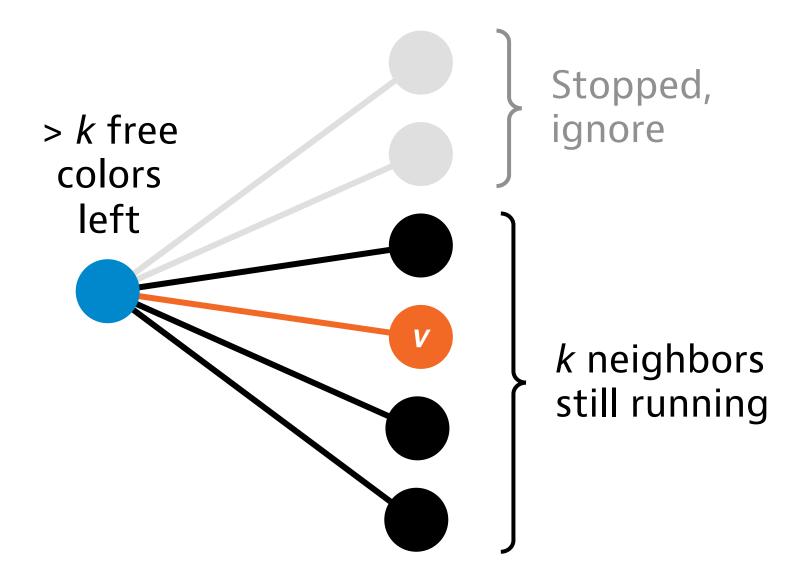
k neighbors still running





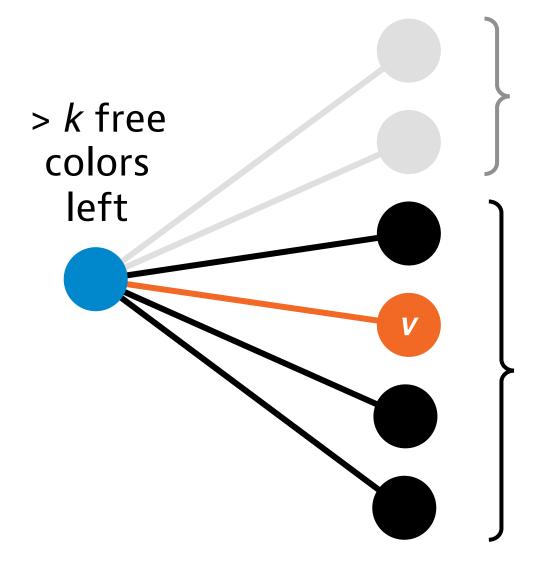


case 1: v is passive conflict probability 0



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case 2: v is active conflict probability < 1/k



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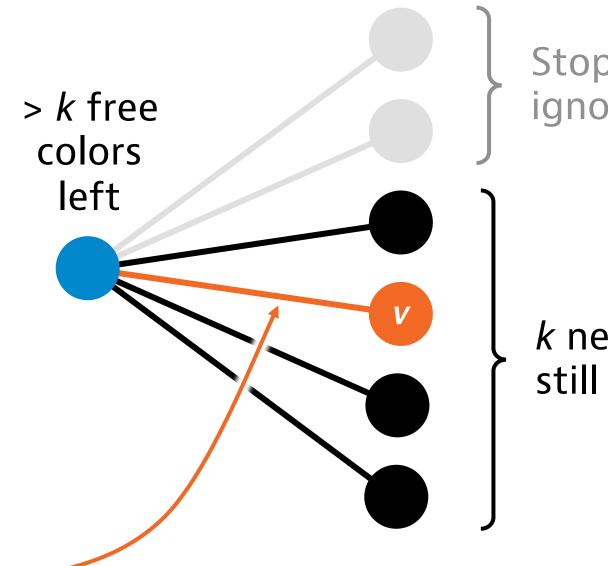
k neighbors still running

case 1: v is passive conflict probability 0

case 2: v is active conflict probability < 1/k

overall:

conflict probability < 1/(2k)



Stopped, ignore

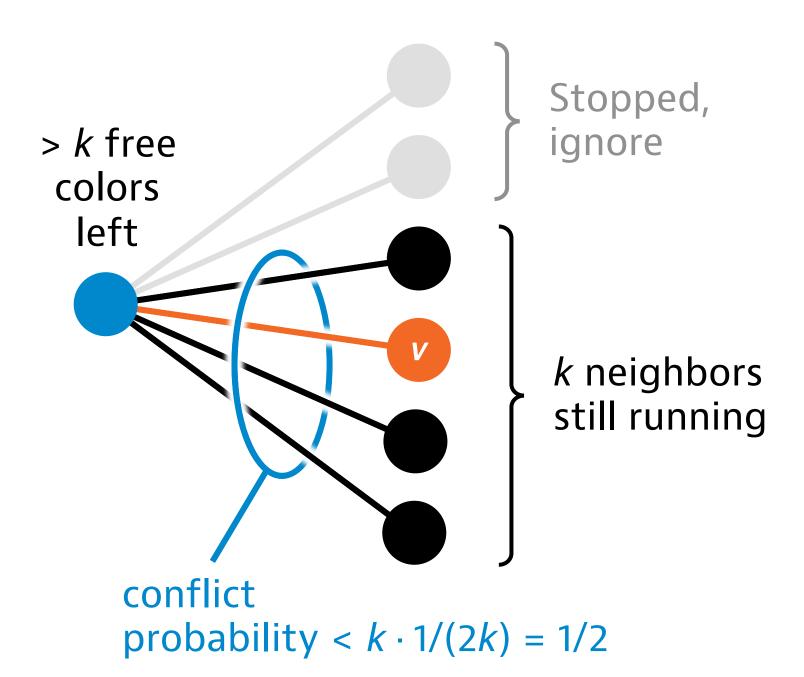
k neighborsstill running

case 1: v is passive conflict probability 0

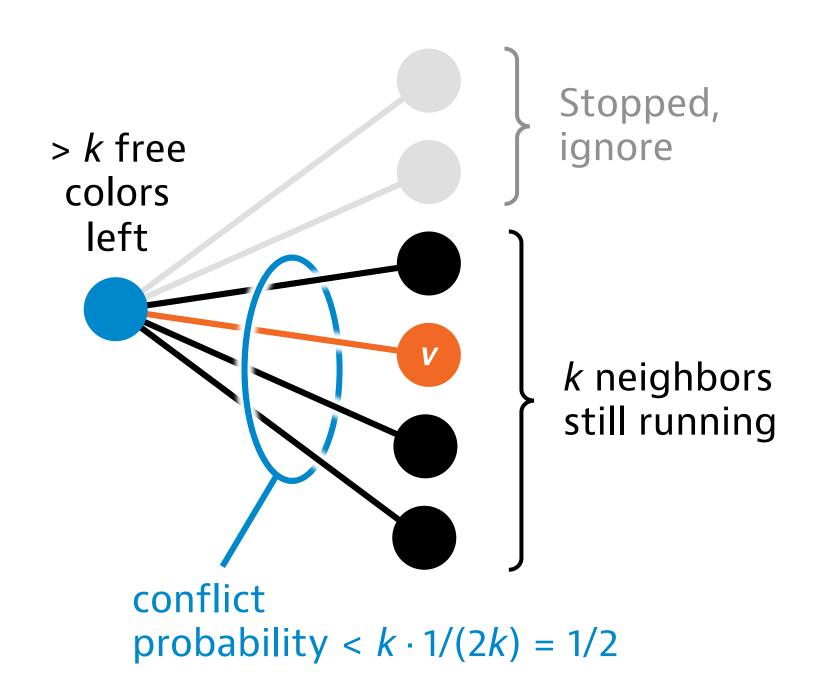
case 2: v is active conflict probability < 1/k

overall:

conflict probability < 1/(2k)

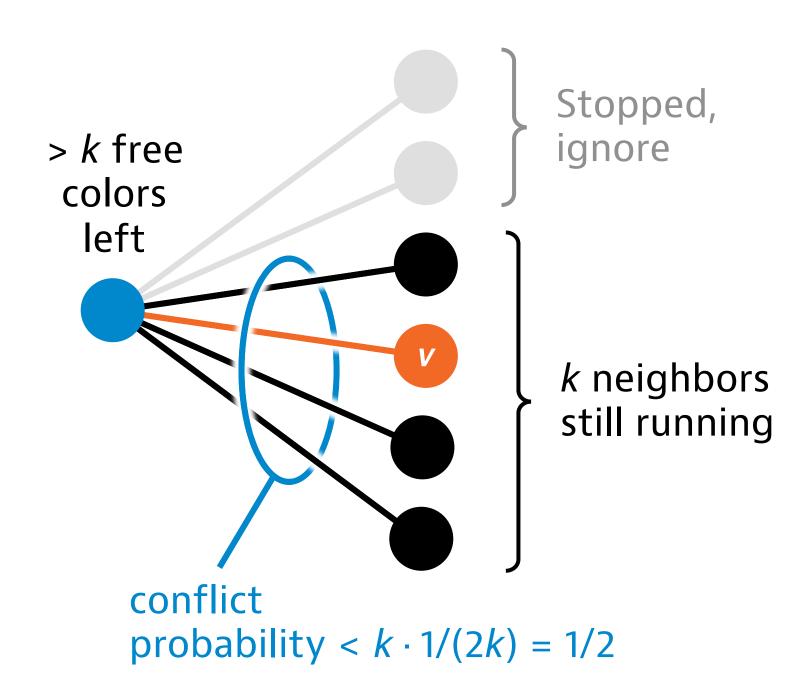


I can stop with probability > 1/2 if I'm active



I can stop with probability > 1/2 if I'm active

I'm active with probability 1/2



I can stop with probability > 1/2 if I'm active

I'm active with probability 1/2

Overall:

I can stop with probability > 1/4

